

You are given a string `word`, and an integer `numFriends`.

Alice is organizing a game for her `numFriends` friends. There are multiple rounds in the game, where in each round:

- `word` is split into `numFriends` **non-empty** strings, such that no previous round has had the **exact** same split.
- All the split words are put into a box.

Find the **lexicographically largest** string from the box after all the rounds are finished.

Example 1:

Input: `word = "dbca"`, `numFriends = 2`

Output: `"dbc"`

Explanation:

All possible splits are:

- `"d"` and `"bca"`.
- `"db"` and `"ca"`.
- `"dbc"` and `"a"`.

Example 2:

Input: `word = "gggg"`, `numFriends = 4`

Output: `"g"`

Explanation:

The only possible split is: `"g"`, `"g"`, `"g"`, and `"g"`.

"lexicographically Largest" → which appear later in the dict

Ex : "a b c" < "a b d"

"a" = "a" → move forward
"b" = "b" → move forward
"c" < "d" → d appears later in dict than c
∴ "abc" < "abd"

Ex : "abcd" < "c"
"a" < "c" → c comes after a
∴ "abcd" < "c" in dict.
"c" comes later in dict than "abcd".

⇒ To find lexicographically largest start comparing strings character wise.

Ex : "ebcde" > "e" x

"e" = "e"

"b" > " " → Empty string
no more letters/
chars left.

∴ "ebcde" appears after "e"
in dict.

Thought Process

word = "d b c a", numFriends = 2 .
↓

⇒ look for largest char (that comes later in alphabet)

Ex: "d" in dbca .

how to distribute among friends ?

Ex :

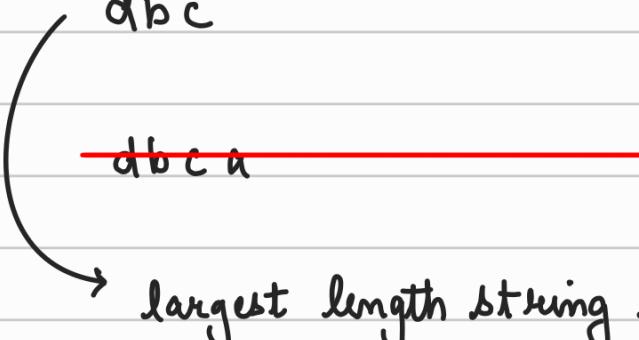
⇒ we know we have to start with "largest char"
i.e. "d"

Friend 1

d

db

dbc



Friend 2 .

bca

ca

a

" " (NOT ALLOWED)

"we gave the last friend "one char" & gave all

"the other chars" to the first friend "

Ex : 3 friends .

⇒ we start with "d"

"dbc a"

friend 1

friend 2

friend 3

db .

c

a

NOTE : Give one - one char to the remaining friends
whatever chars are left add it with "d"
& give it to first friend .

This way we will get our lexicographically largest string .

⇒ if we have numFriends .

⇒ Fr 1 .

(numFriends - 1)

give 1 char each .

whatever remaining
give to fr 1 to
add with largest
char .

length of lexicographically largest string =

word.length() - (numFr - 1) chars ;

Ex: "dbca"

↓
4.

giving 1 char each
to remaining friends .

Max length substring we
would get .

Ex: "ea**bcde**", numFriends = 3 .

↓

Total length = 6

F₁

(6 - 2)

4
 $((\text{word.len} - (\text{NumFr} - 1))$

F₂

1 char

2

F₃

1 char

2 (NumFr - 1)

↓

largest length substring we can get .

"eabc"

"d"

"e"

another ex :

" a b c e d "
↑
 $F_K = 3$.

F₁

3 ch

F₂

1 ch

F₃.

1 ch.

ed → lexicographically largest

c

after "e" in og string
there was only "d"

left

although we can make 3 ch
largest substring with "e"
but take whatever is
available it will still be
lexicographically largest.

e → idx = 3.

longest possible = $n - (\text{num}F_K - 1)$;
(starting with largest
char)

substring (i , 3)
possible ↓

index of largest
char

↗
its not possible
in this string
as "ed" is only available.

∴ add a check ,

[jitna mil huye utna hi lelo]

longest possible = $n - (\text{numFr} - 1)$;

take possible length = $\min(\text{longest possible}, \underline{n-i})$;
(CanTake length) char avail.. to
substr(i , take possible length); the right side
of it

Ex: "abc^ed"

takepossible length = 2
substring [i , 2] → length
starting index .
"ed"

Ex : $5 - 3 = 2$

"abced"

$i \rightarrow \text{idx of}$
 i^{th} char .

(in this ex, i is idx
of largest char)

Ex : $\begin{matrix} 0 & 1 & 2 & 3 & 4 & 5 \\ \text{e} & \text{a} & \text{b} & \text{c} & \text{e} & \text{d} \end{matrix}$ Friends = 3 .
 i i

largest char at two indexes → 0 & 4 .

$$\begin{aligned}\text{longest Possible} &= n - (\text{Fr} - 1) \\ &= 6 - (3 - 1) \\ &= 4 .\end{aligned}$$

if we start from $i = 0$

we can make longest possible substring.

"eabc"

if we start from $i = 4$

we get, "ed"

NOTE : "e~~a~~bc" < "e~~d~~"

∴ lexicographically largest \rightarrow "ed"

what to do for this kind of edge case ??

1st way

\Rightarrow find largest char "e"

\Rightarrow check in which indices it is present in $\{0, 4\}$

\Rightarrow go to each i & check if we are able "eabc" "ed"
to make "largest possible" length
substring or not.

\Rightarrow then compare b/w which is
bigger. (lexicographically)
"eabc" < "ed"
"ed"

2nd way

→ go to every index & try to make "longest possible" length substring.

66 ^{0 1 2 3 4 5}
 e a b c e d
 |
 i

"eabc" \Rightarrow best answer.

↓
66 ^{0 1 2 3 4 5}
 e a b c e d
 |
 i

"abce" \rightarrow still the above is best answer.

↓

66 ^{0 1 2 3 4 5}
 e a b c e d
 |
 i

"bcde" \rightarrow still the

- - -

↓

66 ^{0 1 2 3 4 5}
 e a b c e d
 |
 i

\nearrow longest possible.

can't take len = 4, doesn't matter
take whatever length is possible.

"ced" \rightarrow still

- - -

↓
"e a b c e d"
i

"ed" → better than prev
last answer
∴ update it.

↓
"e a b c e d"
i

"d" → still
.....

↓
"e a b c e d"
i

i oob

Answer found.

longestPossible = $n - (\text{numFriends} - 1)$;

for (i=0; i<n; i++) { } $\rightarrow O(n)$

canTakeLength = min (longestPossible, n-i);

result = max (result, word.substr(i, canTakeLength))
 \downarrow
 $O(n)$

}
return result;

T.C = $O(n^2)$

S.C = $O(1)$

optimisation
↓

$O(n)$.

\hookrightarrow edge case if no of friends = 1 give the whole word to it.

```
1 class Solution {
2 public:
3     string answerString(string word, int numFriends) {
4         int n = word.length();
5         if( numFriends == 1) { }  $\hookrightarrow$  edge case if no of friends = 1 give the whole word to it.
6             return word;
7
8         string result;
9
10        int longestPossible = n - (numFriends - 1);
11
12        for(int i = 0; i < n; i++) {
13            int canTakeLength = min(longestPossible, n-i);
14
15            result = max(result, word.substr(i, canTakeLength));
16        }
17
18        return result;
19    }
20};
```

