



# **Array as a Data Structures**



# Array as a Data Structure

## ADT array

- Objects Elements of the same type arranged in a sequence. An associated index has finite ordinal type. There is an one-to-one correspondence between the values of the index and the array elements.
- Operations
  - (1) `store_array (a,i,e)` -- store `e`'s value in the `i`th element of array `a`
  - (2) `retrieve_array (a,i)` -> `e` -- return the value of the `i`th element of array `a`



# Array as a Data Structure...

- **Design**

The required no. of memory locations are statically allocated consecutively.

- **Implementation**

Built into the language.

What are the constraints ?



# Higher Dimensional Arrays

- Two- and higher-dimensional arrays are extension of the same ADT.
- There are two design choices :
  1. **Row-major** – Elements are stored such that the last index increases most rapidly.
  2. **Column-major** – Elements are stored such that the first index increases most rapidly.
- Whenever there is a reference to an array element, the compiler generates codes for calculating the physical address of the element. This address is then used to access the element exactly like a simple variable.



# Polynomials – Application of Array

- Operations
  - Is-zero – returns true if polynomial is zero.
  - Coef – returns the coeff. of a specified exponent.
  - add - add two polynomials
  - mult - multiply two polynomials
  - Cmult - multiply a polynomial by a const.
  - degree – returns the degree of the polynomial
- Representation decisions
- 1. Exponents should be unique and be arranged in decreasing order.
- 2. Storage alternatives ?



# Sparse Matrix

It is natural to represent matrices as 2-d arrays. But for sparse matrices this involves wastage of a lot of memory space. The operations like transpose, add, multiply also takes a lot of time.

- An alternative approach :

Store the nonzero elements of a sparse matrix in the form of 3-tuples (i, j, val) in an array.

i = row-position

j = column position

val = value at position (i, j) [nonzero].

[The first triple contains (m, n, t), for a m X n matrix having t non-zero values]

- This representation reduces the space requirement. Algorithms developed on this representation may be faster than the first approach.



# Lists

The most obvious application of arrays is in representing lists of elements of the same type.

Some of the operations performed on lists :

- 1. find out the length of a list
- 2. read the list from either direction
- 3. retrieve the  $i^{\text{th}}$  element
- 4. store a new value into the  $i^{\text{th}}$  position
- 5. insert a new element at position  $i$
- 6. delete the element at position  $i$
- 7. search the list for a specified value
- 8. sort the list in some order on the value of the elements.



# Summary

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- Abstract Data Type for Array defined
- Design options of Array Data Structure and its Implementation discussed
- Application of array to implement other data structures discussed