

Binary search

- Sorted array)

```
while (left <= right)
    // if (a[mid] < search-element)
        mid = left + (right - left) / 2;
        if (a[mid] == search-ele)
            return mid;
        if (search-ele < a[mid])
            left = mid + 1;
        else if (search-ele > a[mid])
            right = mid - 1;
```

Recursive approach

```
public static int binarySearch(int array[], int left,
    int right, int item) {
    if (left <= right)
    {
        int mid = left + (right - left) / 2;
        if (array[mid] == item)
            return mid;
        if (array[mid] > item)
            return binarySearch(array, left, mid - 1, item);
        else
            return binarySearch(array, mid + 1, right, item);
    }
    else
        return -1;
}
```

Time complexity

Space complexity

Best : $O(1)$

$O(1)$

Average : $O(\log n)$

Worst : $O(\log n)$

Sorting

in-place sorting

not-in-place sorting

Bubble sort, selection
sort, insertion
Heap

Merge

Quick

Bubble sort

- 1 We will take starting two elements from the list.
- 2 we will compare those elements
- 3 If These elements are found in unsorted order, we will sort them.
- 4 else we will compare next to elements
- 5 Repeat until we get sorted array

pass 1 28 6 4 2 24 → 6 28 4 2 24
6 28 4 2 24 → 6 4 28 2 24
6 4 28 2 24 → 6 4 2 28 24
6 4 2 28 24 → 6 4 2 24 28

pass 2 6 4 2 24 28 → 4 6 2 24 28
4 6 2 24 28 → 4 2 6 24 28
4 2 6 24 28 → 4 2 6 24 28

Algo

```
for (int i=0; i<len-1; i++) {  
    for (int j=0; j<len-1-i; j++) {  
        if (a[j] > a[j+1])
```

```
        {  
            int temp = a[j];  
            a[j] = a[j+1];  
            a[j+1] = temp;
```

```
        }  
    }
```


Best: $O(n)$

Average = $O(n^2)$

Worst = $O(n^2)$

Space complexity

$O(1)$

In optimized bubble sort, use a boolean variable swapped, initialize with 'false' for every iteration. If it remains false at end of iteration i.e., no swapping occurs in a pass (it indicates that array is already sorted) then you can break the loop.

Selection Sort - At any pt of implementation, we'll have array divided into a sorted array part on the left and unsorted array part on the right.

1. We find smallest ele on unsorted part of array
2. Replace it with first item on unsorted part of array
3. Sorted arr is by one element.

Example:
78 50 10 44 8 20 → 8 | 50 10 44 72 20
→ 8 10 | 50 44 72 20
→ 8 10 20 | 44 72 50
→ 8 10 20 44 | 72 50
→ 8 10 20 44 50 72

Algo

```
for (i = 0; i < len - 1; i++) { min = i
```

```
    for (j = i + 1; j < len; j++)
```

```
        if (a[j] < a[min])
```

```
            min = j
```

```
    swap(a[min], a[i])
```

Time complexity

$O(n^2)$

Insertion Sort

1 3 5 8 9 2 4 7

5 3 1 9 8 2 4 7

2 5

3 5 1 9 8 2 4 7

1 3 5 9 8 2 4 7

1 3 5 9 8 2 4 7

1 3 5 8 9 2 4 7

1 2 3 5 8 9 4 7

1 2 3 4 5 8 9 7

1 2 3 4 5 7 8 9

```
for(i=1; i<len; i++)
```

```
    int key = arr[i]
```

```
    j = i - 1
```

```
    while(j >= 0 && arr[j] > key)
```

```
    { arr[j+1] = arr[j]
```

```
      j = j - 1
```

```
    }
```

```
    arr[j+1] = key
```

Space complexity

$O(1)$

Time complexity

$O(n^2)$

$i=6$ $j=5$
 $9 \ 2 \ 2$
 $1 \ 3 \ 5 \ 8 \ 9 \ 2$
 $1 \ 3 \ 5 \ 8 \ 9$
 $1 \ 3 \ 5 \ 5 \ 8 \ 9$
 $1 \ 3 \ 3 \ 5 \ 8 \ 9$
 $1 \ 3 \ 3 \ 5 \ 8 \ 9$
 $1 \ 2 \ 3 \ 5 \ 8 \ 9$

Quick sort in Java

1. choose pivot element
2. move ele less than pivot in left partition
3. move ele greater than pivot to right "
4. partition index is discovered at the end

pivot = high

70 90 10 30 50 20 60

10 20 30 50

70 90

10 20 30

70

70 90 10 30 50 20 60

70 90 10 30 50 20 60

70 90 10 30 50 20 60

70 90 10 30 50 20 60

70 90 10 30 50 20 60

70 90 10 30 50 20 60

70 90 10 30 50 20 60

70 90 10 30 50 20 60

70 90 10 30 50 20 60

70 90 10 30 50 20 60

70 90 10 30 50 20 60

70 90 10 30 50 20 60

70 90 10 30 50 20 60

70 90 10 30 50 20 60

70 90 10 30 50 20 60

70 90 10 30 50 20 60

70 90 10 30 50 20 60

70 90 10 30 50 20 60

70 90 10 30 50 20 60

Time complexity

$O(n \log n)$ ← Best, Average

$O(n^2)$ ← If pivot is chosen as first element.

now swap ($s_0 + 1$ & pivot)

10 30 50 20 60 90 70

10 30 50 20 60 90 70

10 30 50 20 60 90 70

10 30 50 20 60 90 70

10 30 50 20 60 90 70

10 30 50 20 60 90 70

10 30 50 20 60 90 70

10 30 50 20 60 90 70

10 30 50 20 60 90 70

10 30 50 20 60 90 70

Space complexity

$O(1)$

Auxiliary space complexity $O(\log n)$

↑
function call stack

Quick sort

```
main(    )
```

```
{
```

```
int a[] = {1, 2, 3, 4, 5, 6, 7, 8, 9};
```

```
int size = a.length;
```

```
quickSort(a, 0, size-1)
```

```
}
```

```
quickSort(int a[], int low, int high)
```

```
{
```

```
if (low < high) {
```

```
int indexPI = partition(arr, low, high)
```

```
quickSort(arr, low, indexPI-1);
```

```
quickSort(arr, indexPI+1, high);
```

```
}
```

```
partition(int a[], int low, int high)
```

```
{
```

```
pivot = arr[high]
```

```
swapIndex = low-1
```

```
for (j=low, j<=high-1; j++)
```

```
{ if (arr[j] < pivot)
```

```
{
```

```
swapIndex++
```

```
swap(arr[swapIndex], arr[j])
```

```
}
```

```
swap(arr[swapIndex+1], arr[high]);
```

```
return swapIndex+1;
```

```
}
```


Merge Sort

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main()

{
 mergeSort(a, 0, size-1)

}

mergeSort(a , left , right)

{

 int mid;

 if (left < right) {

 mid = (left + (right - left) / 2);

 mergeSort(a, left, mid);

 mergeSort(a, mid+1, right);

 merge(a, left, mid, right);

 }

}

merge (arr , left , mid , right)

{

 int i, j, k, n1, n2

 n1 = mid - left + 1 // length of left subarray

 n2 = right - left mid

 L[n1] , R[n2]

 for(i=0; i<n1; i++)

 L[i] = arr[left+i]

 for(j=0; j<n2; j++)

 R[j] = arr[mid+1+j]

 i=0, j=0 // starting index of L, R merged subarray)
 k = left //

 while(i < n1 && j < n2) {

 if (L[i] <= R[j]) {

 arr[k] = L[i]

 i++

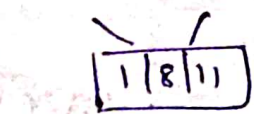
 } else {

 arr[k] = R[j]; j++

 }

 k++

 }



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```

while (i < n) {
    arr[i] = L[i];
    i++;
    k++;
}

```

```

while (j < n) {
    arr[j] = R[j];
    j++;
    k++;
}

```

} // merge method close

Time complexity

Space complexity

Best } $O(n \log n)$
 Worst }
 Average }

$O(n)$

Recursions