

# Automata Lecture 1

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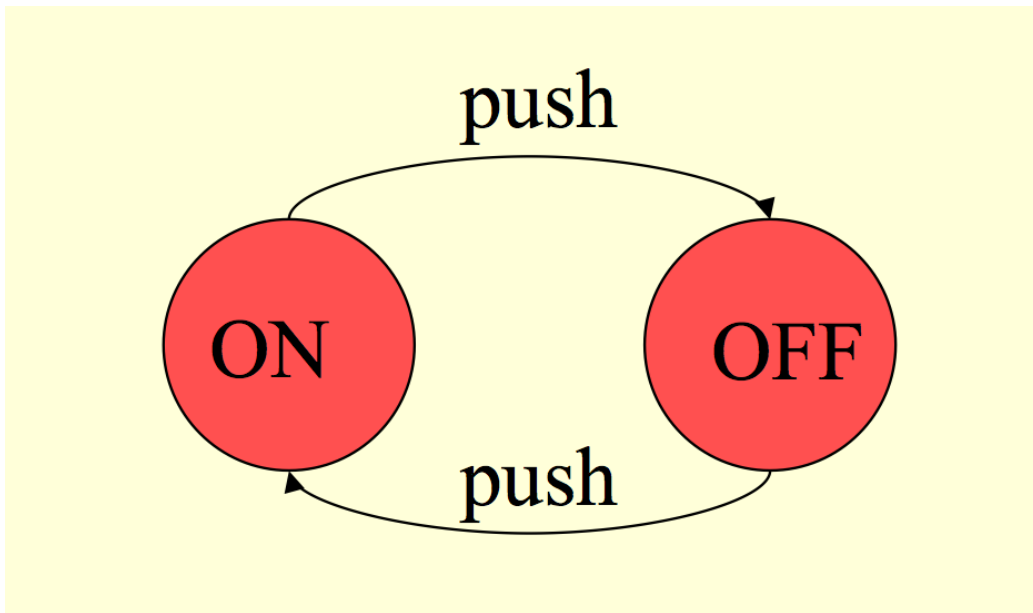
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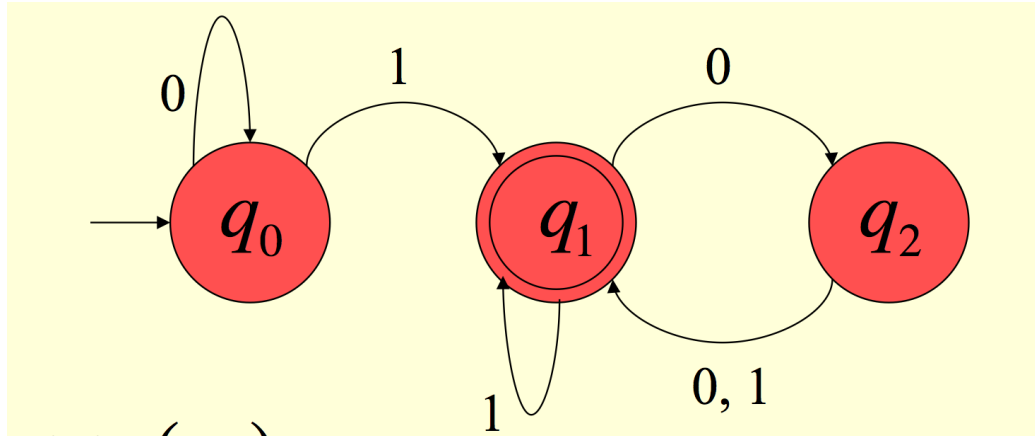
# 1 Finite state Machines FA (finite Automaton)

## 1.1 Example light switch



- It captures, using a finite number of states, the behaviour of a system
- When designing an FA, you must remember only the important parts, and discard the unimportant

## 1.2 Example 2



- The above picture is called a state transition diagram
- Start state  $q(0)$
- finish or accept state  $q(1)$
- Input characters trigger transitions between states
- The automaton receives a string , processes it, and produces an output: accept or reject

if we put 1101:

$q_0 - q_1 - q_1 - q_2 - q_1$

Accepted

## 2 Sets

A set is a **group** of objects represented as a unit.

Objects are called its elements or members.

**Example:**  $\{1,5,4\}$   $\{1,3,2,\dots\}$   $\{a,b,c,d\}$

### 3 Tuples and sequences

- A sequence is a list of **objects** in some **order**.
- example (1,2,6)
- a tuple is a finite sequence
- **k-tuple** sequence of k objects
  - (1,2,6) is a 3 tuple
  - (1,3) is a pair is a pair (2 tuple)
- The cartesian product of A and B.  $A \times B$ .  
is the set of all pairs where the first element is a member of A and the second is a member of B.

### 4 Function

A function is an object that sets an input output relationship.

$f(a) = b$ .

**Domain** The set of possible inputs to the function.

**Range** The outputs of a function.

### 5 Strings and Languages

An alphabet is a **finite set** and its elements are called symbols.

A string over an alphabet is a finite sequence of symbols from the alphabet

**length** of a string is the number of symbols that it contains.

$|w| = n$  if w contains n symbols

and a **language** is a set of strings.

**Example:** Language E (english) belongs to the set alphabet  $\{a,b,c,d,e,f,\dots,x,y,z\}$

$\Sigma^k$  is the set of strings of length k, each of whose symbols are in  $\Sigma$

$\Sigma^1 = \{1, 2\}$

$\Sigma^2 = \{00, 01, 10, 11\}$

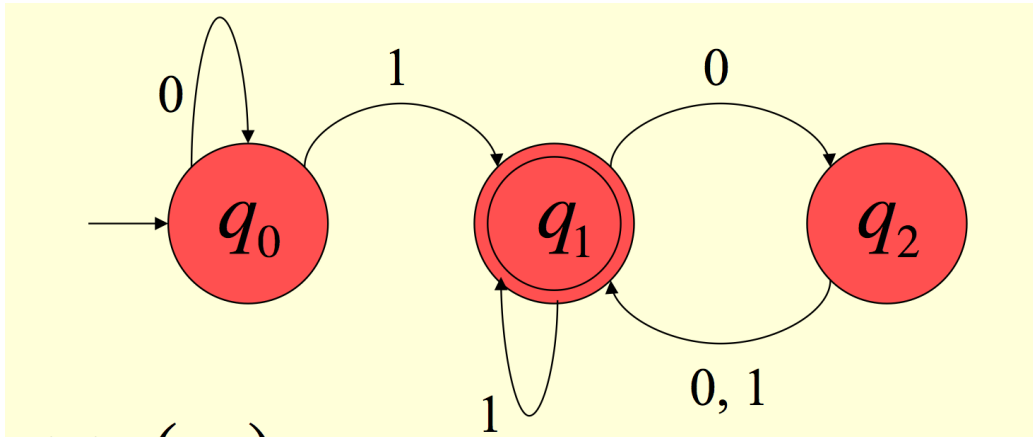
$\Sigma^1 = \{0, 1\}$

$\Sigma^0 = \{\varepsilon\}$

$\varepsilon$  is the string of length 0. the set of all strings over alphabet  $\Sigma$  is

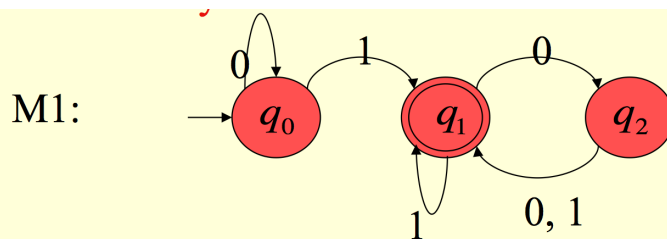
$\Sigma^* = \Sigma^1 + \Sigma^2 + \dots + \Sigma^n$

## Back to FA



### 5.1 Formal Definition of Finite Automaton

- 5 tuple  $(Q, \Sigma, \delta, q_0, F)$
- $\Sigma$  is an **alphabet** (the set of input characters).
- $\delta : Q \times \Sigma \rightarrow Q$  is the **transition function**
- $q_0 \in Q$  is the start state and
- $F \subseteq Q$  is the set of **accept states**



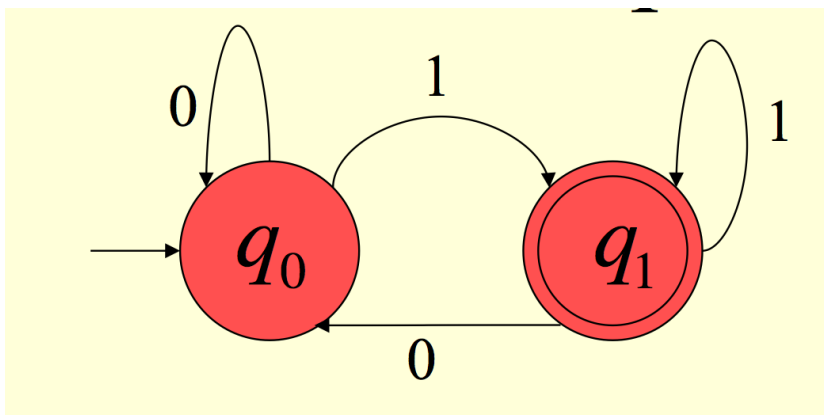
$$M_1 = (\{q_0, q_1, q_2\}, \{0, 1\}, \delta, q_0, \{q_1\})$$

$\delta:$		0	1
$q_0$		$q_0$	$q_1$
$q_1$		$q_2$	$q_1$
$q_2$		$q_1$	$q_1$

## 5.2 Language of machine M

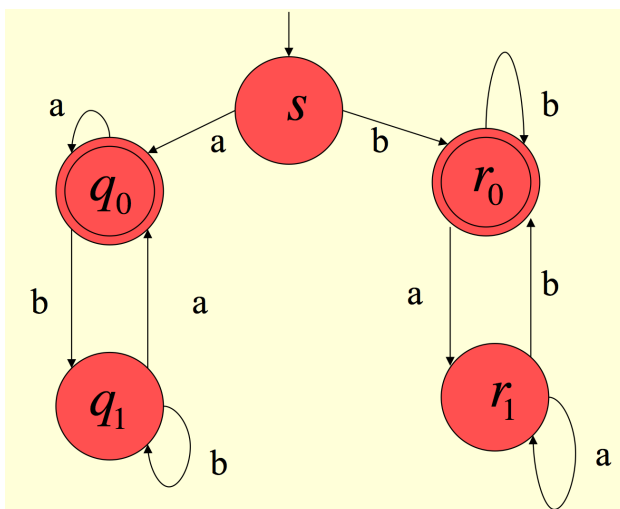
- If  $A$  is the set of all strings that machine  $M$  accepts, we say that  $A$  is the language of machine  $M$
- We say that  $M$  recognises  $A$
- A machine may accept several strings, and it recognises one language

### 5.2.1 Other Example M2



It accepts only strings which doesn't end with 0.

### 5.2.2 M4



accepts only string which begin and ends with same element.