

Resume : Description Logic

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C is the set of concept names and R the set of role names. Every concept names is a concept description (CD).

Formal notation :

$C \sqcap D \rightarrow$ conjunction

$C \sqcup D \rightarrow$ disjunction

$\neg C \rightarrow$ negation

$\exists r.C \rightarrow$ existential restriction

$\forall r.C \rightarrow$ value restriction

An interpretation $I = (\Delta, \cdot)$ such that $\Delta \neq \emptyset$, (Δ is called the domain of the interpretation) and with the following :

$$A \in C \rightarrow A^I \subseteq \Delta$$

$$r \in R \rightarrow r^I \subseteq \Delta \times \Delta$$

$$\top^I = \Delta$$

$$\perp^I = \Delta$$

$$(C \sqcap D)^I = C^I \cap D^I$$

$$(C \sqcup D)^I = C^I \cup D^I$$

$$(\neg C)^I = \Delta \setminus C^I$$

$$(\exists r.C)^I = \{d \in \Delta \mid \exists e \in \Delta \text{ with } (d, e) \in r^I \text{ and } e \in C^I\}$$

$$(\forall r.C)^I = \{d \in \Delta \mid \forall e \in \Delta, \text{ if } (d, e) \in r^I, \text{ then } e \in C^I\}$$

We call C^I the extension of C in I and $b \in \Delta^I$ an r -filler of a in I if $(a, b) \in r^I$.

Lemma 1. *Let I be an interpretation, C, D concepts and r a role. Then*

$$\top^I = (C \sqcup \neg C)^I$$

$$\perp^I = (C \sqcap \neg C)^I$$

$$(\neg \neg C)^I = C^I$$

$$(\neg(C \sqcap D))^I = (\neg C \sqcup \neg D)^I$$

$$(\neg(C \sqcup D))^I = (\neg C \sqcap \neg D)^I$$

$$(\neg(\exists r.C))^I = (\forall r.\neg C)^I$$

$$(\neg(\forall r.C))^I = (\exists r.\neg C)^I$$

Because the \sqcup operator can be tricky sometimes, we use the following relation to replace the \sqcup by \sqcap : $C \sqcup D \rightarrow \neg(\neg C \sqcap \neg D)$.

ALC TBoxes

For C and D possibly compound *ALC* concepts, an expression of form $C \sqsubseteq D$ is called an *ALC* general concept inclusion (GCI). We use $C \equiv D$ as an abbreviation for $C \sqsubseteq D, D \sqsubseteq C$. A finite set of GCI is called an *ALC TBox* and noted τ .

Lemma 2. *If $\tau \in \tau'$ for two TBox τ and τ' , then each model of τ' is also a model of τ .*

Example of a TBox T_{ex} :

$$\begin{aligned} T_{ex} = \{ & Course \sqsubseteq \neg Person, \\ & UGC \sqsubseteq Course, \\ & PGC \sqsubseteq Course, \\ & Teacher \equiv Person \sqcap \exists teaches.Course, \\ & \exists teaches.T \sqsubseteq Person, \\ & Student \equiv Person \sqcap \exists attends.Course, \\ & \exists attends.T \sqsubseteq Person \} \end{aligned}$$

ALC ABoxes

Let I be a set of individual names disjoint from R and C . For $a, b \in I$ individual names, C a possibly compound *ALC* concept, and $r \in R$ a role name, an expression of the form $a : C$ is called an *ALC* concept assertion and $(a, b) : r$ is called an *ALC* role assertion.

A finite set of concept and role assertion is called an *ALC ABox*. An interpretation function \cdot^I is additionally required to map every individual name $a \in I$ to an element a^I in Δ^I . An interpretation I satisfies a concept assertion $a : C$ if $a^I \in C^I$ and a role assertion $(a, b) : r$ if $(a^I, b^I) \in r^I$.

An interpretation that satisfies each concept assertion and each role assertion in a *ABox* A is called a model of A .

Example of a *ABox* A_{ex} :

$$\begin{aligned} A_{ex} = \{ & Mary : Person, \\ & CS600 : Course, \\ & Ph456 : Course \sqcap PGC, \\ & Hugo : Person, \\ & Betty : Person \sqcap Teacher, \\ & (Mary, CS600) : teaches, \\ & (Hugo, Ph456) : teaches, \\ & (Betty, Ph456) : attends, \\ & (Mary, Ph456) : attends \} \end{aligned}$$

Then we can create an interpretation I of this $ABox$ which is a model:

$$\begin{aligned}
\Delta^I &= \{h, m, c6, p4\}, \\
Mary^I &= m \\
Betty^I &= Hugo^I = h, \\
CS600^I &= c6, \\
Ph456^I &= p4, \\
Person^I &= \{h, m, c6, p4\}, \\
Teacher^I &= \{h, m\}, \\
Course^I &= \{c6, p4\}, \\
PGC^I &= \{p4\}, \\
UGC^I &= \{cs\}, \\
Student^I &= \emptyset, \\
teaches^I &= \{(m, c6), (h, p4)\}, \\
attends^I &= \{(h, p4), (m, p4)\}
\end{aligned}$$

Note that isn't a model of the previous example of a $TBox$.