1. Division 
$$\frac{x}{y}$$
 computer  $g: ves$  us  $f(f(x)/f(y))$ 

=  $(x(1-\delta_1)/y(1-\delta_2))(1-\delta_3)$ 

=  $\frac{x}{y} \frac{(1-\delta_1)(1-\delta_3)}{(1-\delta_2)}$ 

=  $\frac{x}{y} \frac{(1-\delta_1-\delta_2)}{(1-\delta_2)} \frac{(1+\delta_2)}{(1+\delta_2)}$ 

=  $\frac{x}{y} \frac{(1-\delta_1-\delta_3)}{(1-\delta_2)} \frac{(1+\delta_2)}{(1+\delta_2)}$ 

=  $\frac{x}{y} \frac{(1-\delta_1-\delta_3)+\delta_2-\delta_3\delta_2-\delta_3\delta_2}{(1-\delta_2\delta_2)}$ 

=  $\frac{x}{y} \frac{(1-\delta_1-\delta_3)+\delta_2-\delta_3\delta_2-\delta_3\delta_2}{(1-\delta_1+\delta_2)}$ 

=  $\frac{x}{y} \frac{(1-\delta_1-\delta_3)+\delta_2-\delta_3\delta_2-\delta_3\delta_2}{(1-\delta_1+\delta_2)}$ 

where  $|S_1| \leq 3 \in$ 

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2. a) We have LyPylzPzlzPzl,P,A= U We need PAILU where P= Py Pz Pz Pi L= [-1 [-1 L-1 L-1] Ly Pulz Pz Lz Pz L, P, A = LyPyL3P3 L2P2L,P2/P2P1)A \since P-1=P; = Ly Py L3(3 L2 P3)(F3 P2 L, P2 P3)(P3 P2 P1) A(P; P; = I = Ly/2 L3 Py/2 P3 L2 P3 Py/2 P3 P2 C1 P2 P3 Py/P4 P3 P2 P1) A Let 23 = (Pyl3Py) (2 - PyP3 L2P3Py C, = Pr. P3 P2 L, P2P3 P4 ad P2 P4P3 P2P, We have Ly Lz Lz L, PA-U > PA- [, 2] 27 Ly U Let L= 2, 2, 2, Ly => PA = 1. U  $\sum_{i} = (P_{i}P_{3}P_{2}L_{1}P_{2}P_{3}P_{4})^{-1}$ SIALL PIZP = P4 P2P2 L1 P2P2P4

b) 
$$L_{1} = P_{1}P_{2}P_{2}L_{1}P_{2}P_{3}P_{4}$$

$$L_{1} = \begin{bmatrix} 1 & 0 & 0 & 0 & 0 & 0 \\ M_{21} & 0 & 0 & 0 & 0 & 0 \\ M_{31} & 0 & 0 & 0$$

$$P_{4}P_{3}P_{2}L_{1}^{-1}P_{2}P_{3}P_{4}$$
  $-\frac{1}{2}P_{3}P_{4}$   $-\frac{1}{2}P_{4}P_{4}$   $-\frac{1}{2}P_{4}$   $-\frac{1}$ 

We have 
$$L_2P_2L_1P_1H=U$$
 $\Leftrightarrow L_2P_2L_1P_2P_2P_1H=U$ 
 $\Leftrightarrow L_2\hat{L}_1P_2P_2P_1H=U$ 
 $\Leftrightarrow P_2P_1H=U$ 
 $\Leftrightarrow P_2P_1H=U$ 

We have 
$$Ax=b \Leftrightarrow PAx=Pb \Leftrightarrow LUX=b$$

Let  $Ux=d$ 

S. We  $Ld=Pb$ 

$$\begin{bmatrix} 1 & 0 & 0 & 0 \\ -1 & 1 & 0 & 0 \\ -1 & 2 & 1 & 0 \end{bmatrix} = \begin{bmatrix} 24 & 0 & 0 \\ 24 & 1 & 0 & 0 \\ -1 & 2 & 1 & 0 \end{bmatrix}$$

$$dz=24$$

$$dz=20$$

$$dz=12$$

$$X_2=3$$

$$X_3=3$$

$$X_2=3$$

$$X_1=1$$

$$X_1=1$$

$$X_1=1$$

$$X_2=3$$

$$X_1=1$$

$$X_1=1$$

$$X_1=1$$

$$X_2=3$$

$$X_1=1$$

$$X_1=1$$

$$X_1=1$$

$$X_1=1$$

$$X_1=1$$

$$X_2=3$$

$$X_1=1$$

c) When you have screal linear systems to solve which differ only in the right hand side, you can save much computation by factoring A once and then using the factorization to solve each of the systems. For example, this situation arises in the iterative improvement algorithm.

4. Aenember: Premoltiplying swaps rows, postawltiplying

Swaps volumns

Ax=b 

PAX=Pb

PAQQ=Pb

where Qx=y

Lly=Pb

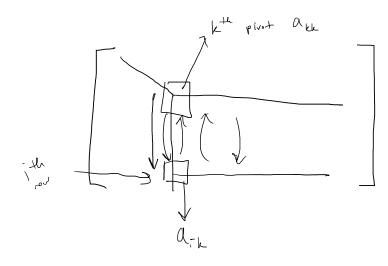
Ld=Pb

Solve Ld=Pb

for d

Solve Ld=Pb

Compute X=Qy



- if Lont pivot and pivot element is very small, multipliers become very large in this stage since pivot element is in the denominator of each multiplier if the multipliers are large, each element in lewer right submatrix (out get significantly amplified in this stage. Not good since round of this stage. Not good since round of the ecror is directly proportional to the largest elements that occur during factorization

- if pivot, multiplies are gravanteed to be

{ I in magnitude in this stage and

{ lenents in lover right submatrix will be

{ lenents in lover right of round off error is reduced.

} amped. Thus, growth of round off error is reduced.

$$\sqrt{1+x} - \sqrt{1-x} \times \sqrt{1+x} + \sqrt{1-x}$$

$$=\frac{1+x-(1-x)}{\sqrt{1+x}+\sqrt{1-x}}$$

and 
$$||Ax||_1 \leq ||A||_1 ||x||_2 = \sum_{i=1}^{M} |A_i|_i$$

Note that if we choose X=ex= [] ck+1 on the unit vector

Note this is romalized

then  $\|Ax\|_1$  is maximized  $\|Ax\|_1 = \sum_{i=1}^{\infty} |\hat{Z}| \alpha_{i} \times |\hat{Z}| = \sum_{i=1}^{\infty} |\alpha_{i} \times |$ 

exice  $\frac{e^{x}-1}{e^{x}}$  cancellation when  $x \to 0$   $\frac{e^{x}-1}{e^{x}} = \frac{x^{2}}{n!} + \frac{x^{3}}{3!} + \dots$   $\frac{e^{x}-1}{e^{x}} = \frac{x^{2}+x^{3}}{3!} = \frac{x^{2}-1}{2!}$