

## Sorting

- Arrange array elements to be in ascending (or descending) order by key
- Sorting algorithms differ by:
  - **ease of implementation**
  - **efficiency (in best, worst, and average cases)**
  - **adaptability for memory limitations** (i.e., internal vs. external sorting)

### Bubblesort

- Imagine array turned vertically (A[0] at “top”, A[n-1] at “bottom”)
- Make repeated passes over the array, from bottom to top
- If 2 adjacent elements out of order, then reverse them (i.e., “lighter” one moves up)
- On 0<sup>th</sup> pass, “lightest” element rises all the way to the top
- On i<sup>th</sup> pass, no element will have to rise above position i

```

for (i = 0; i < n - 1; i++)           // n-2 times
  for (j = n - 1; j >= i + 1; j--)    // n-1 times, n-2 times, ..., 1 time
    if (A[j].key < A[j - 1].key)      // constant time
      swap(A[j], A[j - 1]);          // constant time

```

<u>When</u>	<u>What j loop will do</u>	<u># times if stmt will be executed</u>
i = 0	j = (n-1)..1	n-1
i = 1	j = (n-1)..2	n-2
i = 2	j = (n-1)..3	n-3
...		
i = n-2	j = (n-1)..(n-1)	n - (n - 1) = 1

So  $T(n) = (n-1) + (n-2) + (n-3) + \dots + (n - (n - 1)) = ((n - 1) * n) - \sum_{i=1}^{n-1} i = O(n^2)$

*...What if array is already in sorted order???*

### Selection Sort

- On the i<sup>th</sup> pass, select the lowest key among A[i], ..., A[n-1], and swap it with A[i]

```

for (i = 0; i < n-1; i++) {           // n-2 times
  min = A[i].key; minPos = i;
  for (j = i + 1; j < n; j++)         // n-1 times, n-2 times, ..., 1 time
    if (A[j].key < min) {
      min = A[j].key;
      minPos = j;
    }
  if (i != minPos) swap(A[i], A[minPos]);
}

```

This algorithm is also  $O(n^2)$  *...Again, what if array is already in sorted order???*

### Insertion Sort

- On the  $i^{\text{th}}$  pass, insert  $A[i]$  into its rightful place among  $A[0], A[1], \dots, A[i-1]$
- After that  $A[0], A[1], \dots, A[i]$  are in sorted order

for ( $i = 1; i < n; i++$ )

    insert  $A[i]$  into the elements before  $A[i]$  (which are already in sorted order)

This algorithm is also  $O(n^2)$  ...*What if array is already in sorted order???*

### Characteristics of Quadratic Time Sorting Algorithms

- Simple to implement
- Not very efficient, but O.K. if  $n$  not too big
- Require all elements of array to be in memory at one time