

Analysis of Graphs Operations

Let n = # vertices, e = # edges

But typically e is much less than n !!!

2 graph implementations: **adjacency matrix** or **adjacency list**

	Adj. List	Adj. Matrix
Adding or removing an edge	$O(e)$	$O(1)$
Checking whether particular edge exists	$O(e)$	$O(1)$
Processing each edge of particular vertex	$O(e)$	$O(n)$

	<u>Adj. List</u>	<u>Adj. Matrix</u>
<u>Depth-First Search (DFS)</u>		
initialize visited[n] to false;		
for each vertex v do	$O(n)$	$O(n)$
if (visited[v] == false)		
DFS(v);		
void DFS(vertex v) {		
visited[v] = true;		
for each neighbor w of v do	$O(e)$	$O(n)$
if (visited[w] == false)		
DFS(w);		
}		
	So $O(n * e)$	So $O(n^2)$

<u>Breadth-First Search (BFS)</u>		
initialize visited[n] to false;		
for each vertex v do		
if (visited[v] == false)		
BFS(v);		
void BFS(vertex v) {		
visited[v] = true;		
enqueue vertex v into rear of queue;		
while (queue is not empty) {	$O(n)$	$O(n)$
dequeue vertex u from front of queue;		
for each neighbor w of u do	$O(e)$	$O(n)$
if (visited[w] == false) {		
visited[w] = true;		
enqueue w into rear of queue;		
}		
}		
}		
}	So $O(n * e)$	So $O(n^2)$

Shortest-Paths Problem

- Assume each edge has a (positive value) weight associated with it
- Represent with adjacency matrix as $A[v][w]$ = weight for edge (v, w) , or ∞ if no such edge
- Want to find the “cost” of the **shortest path from a source vertex to every other vertex**

// This returns 1D array D where $D[i]$ = cost of shortest path from source to vertex i

Dijkstra’s Algorithm:

```
for each vertex v do {
    D[ v ] = A[ start ][ v ];
    visited[ v ] = false;
}

D[ start ] = 0; visited[ start ] = true;

for i = 1..(n - 1) do // where n = # vertices
{
    v = unvisited vertex with smallest value in D[ ];
    visited[ v ] = true;

    for each neighbor w of v do
        D[ w ] = min(D[ w ], D[ v ] + A[ v ][ w ]);
}
```

Example:

	0	1	2	3	4
0	∞	10	3	20	∞
1	∞	∞	∞	5	∞
2	∞	2	∞	∞	15
3	∞	∞	∞	∞	11
4	∞	∞	∞	∞	∞