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## Using QuickCheck to test C APIs

2015-06-15 - Magnus Therning

Last year at ICFP I attended the <u>tutorial on QuickCheck</u> with John Hughes. We got to use the Erlang implementation of QuickCheck to test a C API. Ever since I've been planning to do the same thing using Haskell. I've put it off for the better part of a year now, but then Francesco Mazzoli wrote about <u>inline-c</u> (<u>Call C functions from Haskell</u> without bindings and I found the motivation to actually start writing some code.

#### The general idea

Many C APIs are rather stateful beasts so to test it I

- 1. generate a sequence of API calls (a program of sorts),
- 2. run the sequence against a model,
- 3. run the sequence against the real implementation, and
- 4. compare the model against the real state each step of the way.

#### The C API

To begin with I hacked up a simple implementation of a stack in C. The "specification" is

```
/**
  * Create a stack.
  */
void *create();

/**
  * Push a value onto an existing stack.
  */
void push (void *, int);

/**
  * Pop a value off an existing stack.
  */
int pop(void *);
```

Using inline-c to create bindings for it is amazingly simple:

```
{-# LANGUAGE QuasiQuotes #-}
{-# LANGUAGE TemplateHaskell #-}
```

```
module CApi
    where

import qualified Language.C.Inline as C
import Foreign.Ptr

C.include "stack.h"

create :: IO (Ptr ())
create = [C.exp| void * { create() } |]

push :: Ptr () -> C.CInt -> IO ()
push s i = [C.exp| void { push($(void *s), $(int i)) } |]

pop :: Ptr () -> IO C.CInt
pop s = [C.exp| int { pop($(void *s)) } |]
```

In the code below I import this module qualified.

### Representing a program

To represent a sequence of calls I first used a custom type, but later realised that there really was no reason at all to not use a wrapped list:

```
newtype Program a = P [a]
  deriving (Eq, Foldable, Functor, Show, Traversable)
```

Then each of the C API functions can be represented with

```
data Statement = Create | Push Int | Pop
    deriving (Eq, Show)
```

## Arbitrary for Statement

My implementation of Arbitrary for Statement is very simple:

```
instance Arbitrary Statement where
    arbitrary = oneof [return Create, return Pop, liftM Push ar
    shrink (Push i) = Push <$> shrink i
    shrink _ = []
```

That is, arbitrary just returns one of the constructors of Statement, and shrinking only returns anything for the one constructor that takes an argument, Push.

## Prerequisites of Arbitrary for Program Statement

I want to ensure that all Program Statement are valid, which means I need to define the model for running the program and functions for checking the precondition of a statement as well as for updating the model (i.e. for running the Statement).

Based on the <u>C API</u> above it seems necessary to track creation, the contents of the stack, and even if it isn't explicitly mentioned it's probably a good idea to track the popped value. Using <u>record</u> (Record is imported as R, and Record.Lens as RL) I defined it like this:

```
type ModelContext = [R.r| { created :: Bool, pop :: Maybe Int,
```

Based on the rather informal specification I coded the pre-conditions for the three statements as

```
preCond :: ModelContext -> Statement -> Bool
preCond ctx Create = not $ RL.view [R.1| created |] ctx
preCond ctx (Push _) = RL.view [R.1| created |] ctx
preCond ctx Pop = RL.view [R.1| created |] ctx
```

#### That is

- Create requires that the stack hasn't been created already.
- Push i requires that the stack has been created.
- Pop also requires that the stack has been created.

Furthermore the "specification" suggests the following definition of a function for running a statement:

(This definition assumes that the model satisfies the pre-conditions, as can be seen in the use of tail.)

#### Arbitrary for Program Statement

With this in place I can define Arbitrary for Program Statement as follows.

```
instance Arbitrary (Program Statement) where
    arbitrary = liftM P $ ar baseModelCtx
    where
        ar m = do
            push <- liftM Push arbitrary
        let possible = filter (preCond m) [Create, Pop,
        if null possible
            then return []
        else do
            s <- oneof (map return possible)
        let m' = modelRunStatement m s
            frequency [(499, liftM2 (:) (return s)</pre>
```

The idea is to, in each step, choose a valid statement given the provided model and

cons it with the result of a recursive call with an updated model. The constant 499 is just an arbitrary one I chose after running arbitrary a few times to see how long the generated programs were.

For shrinking I take advantage of the already existing implementation for lists:

### Some thoughts so far

I would love making an implementation of Arbitrary s, where s is something that implements a type class that contains preCond, modelRunStatement and anything else needed. I made an attempt using something like

```
class S a where
   type Ctx a :: *

baseCtx :: Ctx a
  preCond :: Ctx a -> a -> Bool
   ...
```

However, when trying to use <code>baseCtx</code> in an implementation of <code>arbitrary I</code> ran into the issue of injectivity. I'm still not entirely sure what that means, or if there is something I can do to work around it. Hopefully someone reading this can offer a solution.

#### Running the C code

When running the sequence of Statement against the C code I catch the results in

```
type RealContext = [r| { o :: Ptr (), pop :: Maybe Int } |]
```

Actually running a statement and capturing the output in a RealContext is easily done using inline-c and record:

## **Comparing states**

Comparing a ModelContext and a RealContext is easily done:

#### Verifying a Program Statement

With all that in place I can finally write a function for checking the validity of a program:

```
validProgram :: Program Statement -> IO Bool
validProgram p = and <$> snd <$> mapAccumM go (baseModelCtx, ba
where
    runSingleStatement mc rc s = realRunStatement rc s >>=
    go (mc, rc) s = do
        ctxs@(mc', rc') <- runSingleStatement mc rc s
        return (ctxs, compCtx mc' rc')</pre>
```

(This uses mapAccumM from an earlier post of mine.)

### The property, finally!

To wrap this all up I then define the property

```
prop_program :: Program Statement -> Property
prop_program p = monadicIO $ run (validProgram p) >>= assert
```

and a main function

```
main :: IO ()
main = quickCheck prop_program
```

Edit 2015-07-17: Adjusted the description of the pre-conditions to match the code.

Tags: haskell, quickcheck, testing

← mapAccum in monad Systemd watchdog ⇒

#### Jesse McDonald

2015-06-15

I believe the problem with baseCtx is that there could be two or more instances of S which have the same definition for Ctx a, which means that the compiler can't infer a

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(and thus the instance to use) from Ctx a. One option would be to add a Proxy a parameter to baseCtx to select the instance. Another would be to make S a multiparameter type class with a one-to-one functional dependency between the parameters, like so:

```
class S a ctx | a -> ctx, ctx -> a where
  baseCtx :: ctx
...
```

#### **Magnus Therning**

2015-06-16

I should probably point out that I'm using version 0.3.1.2 of record.

The library seems to have seen dramatic changes as of version 0.4.0.0.

#### Leave a comment

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