Abstract State Machines

Verification and Final Thoughts

Verification of ASM

Three methods:

- 1. Reviews
- 2. Testing
- 3. Formal Proofs

Review of ASMs

- Similar to code reviews
- Natural to the system
- Knowledge base

Testing ASMs

- ASMs can be executed
 - AsmL
 - XASM
- Testing is performed earlier

Formal Proofs of ASMs

- ASMs are a Formal Description
- Verification of ASMs

Example

Prove the definition of Kruskal's algorithm for finding a minimum spanning tree terminates.

```
if CurrentMode = initial then
    do-forall p: Node
         \mathsf{Label}(p) := p
    enddo
    CurrentMode := run
elseif CurrentMode = run then
    choose e: Edge: Eligible(e) and ((\forall f: Edge) Eligible(f) \Rightarrow Weight(f) \geq Weight(e))
         Tree(e) := true
         choose p, q: Node: \{p, q\} = \text{Endpoints}(e)
             do-forall r: Node: Label(r) = Label(p)
                  \mathsf{Label}(r) := \mathsf{Label}(q)
             enddo
         endchoose
    ifnone CurrentMode := done
    endchoose
endif
```

- Prove the algorithm finds a spanning tree
 - Graph of *n* nodes
 - Algorithm creates edges between nodes
- Termination:
 - All nodes with the same label
 - n 1 edges

- At each state:
 - \circ Label $\rightarrow p$

 - Set of nodes labelled $p → L_p$ Number of nodes in $L_p → n_p$ Edges connected to nodes in $L_p → t_p$
- Each set of nodes in a non-empty L_n
 - $\circ t_p = n_p 1$
- Loop until only one non-empty L_n

- Now prove the claim by induction
- At the first step each node is separately labelled
 - \circ $n_n = 1$
 - \circ $t_{n} = 0$
 - Člaim holds true so far.

- Edge e chosen between two labels
 - \circ (p and q)
- All nodes labelled p are re-labelled q
- L'_q is the new set of nodes labelled q
- The proof holds true
 - \circ L'_n is now empty
 - All other labels remain the same
- Number of non-empty labels decreases by 1

Final Thoughts

- ASMs not limited to software
- Multi-agent ASMs
 - Distributed Systems
 - Real-time Systems