

CSE 214

Algorithm

"Assignment-II"

SUBMITTED BY:

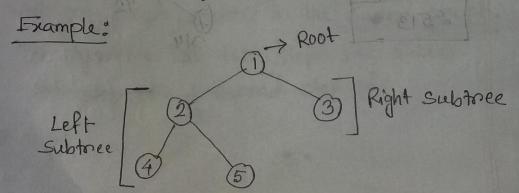
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1. Full Tree Traversal 1 10 101

DFS: DFS teachnique used for traversing tree or Graph. Here backtracking is used for traversal. In the first traversal the deepest node is visited and then backtracks to its parent node if no st sibling of that not exist.

In a tree, we can simply begin from a node then traverse its adjacent without carring about cycles.

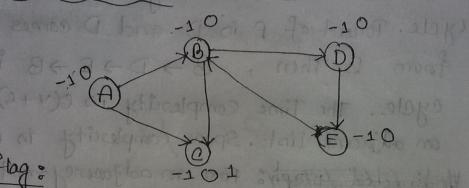


Therefore, DFS Traversal of this tree will be:

- a) Inondero (Left, Root, Right): 425 13
- b) Preorder (Root, Left, Right): 12453
- c) Postonders (Left, Right, Poot): 45231

planted at the property of the bank of

D Dinected Graph:



-1 = Unvisited

0 = Vinited and in stack

1 = Visited and Poped out from Stack

01	Stack	CI
1		450
1		V.
-	F	
t	D	
1	<u>B</u>	-
	-	

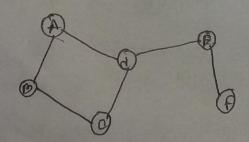
Vinited set	vertext	Parent
ABODE	2 A Hers	-
algrinous no a	0	A
	C	13
(4)	D	18
	E	

Last vinited note E, Now, E > B, that edge Should be find out the flag of O and B is already in this stack. So, that edge E to B is a cycle. Parent of E is D. and D comes out from B. Then, B > D > E > B is a cycle. The Time Complexity is O(V+e) for an adjacent List. Space complexity is O(V).

3. Component Finding

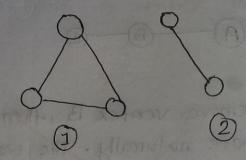
Component is a is a subgreaph any two ventices are connected by paths and no ventices are connected with any others ventices in the super graph.

Here is an areample:



This is a graph. Every node is connected with each others. This graph has one connected components.

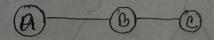
Again here is at an example.



This is also a greaph. But Every node is not connected with each others. This graph has two connected components.

Anticulation Points Frinding

- Anticulation point is such a mock on ventex in the graph such that if we semove that ventex then the graph devided into different components.



If we remove vertex B, then the edges also remove naturally. The we get two Components in the graph A and C. Then there are no paths to from A to C. A, C components are disconnected. So, Removing point B the graph devided into diffount components. B vertex 15 called as Articulation points.

5. Topological Sovot (DFS)

Key Points:

- 1. It is a linear ordering of its ventices such that for every edge uv for vertex u to V, U comes before vertex v in the ordering.
- 2. Graph Should be Dissected Acyclic Graph.
- 3. Every Directed Acycle Graph will have atteast one topological ordering.

After Pop! 1) levois los pologol . ? topological sort order: 20416537

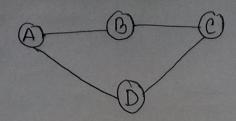
Total Running time: O(n+e)

6. Strongly connected components

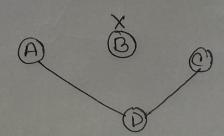
In a greaph, Connected component is a sub graph in which any two vertices we connected by Paths and no vertices are is connected with any others verotices in the supero grouph. This algorithm has two steps DFS GOBFOS. 18th step DFS & Second step BFS.

19t Steps	1110
[6]	19 417 516 11/12
4	213
3	6 13/16
1	3/8 F 14/15

1st OFS Provens finished. 2nd step: Revense travense. After pop: Strongly connected boyol15 mora effort Camponento 13/14 2513



If, we remove B verstex. Then the graph will be.



B in not Apticulation point because the graph Still connected remains. There is a Path from vertex A to vertex C through vertex D. So. B can't be the Apticulation poins. In this graph has no Apticulation point.