

Coverage of Removing Single, Pair, and Triple Operations in Code Refactoring

Let $I(A)$ denote the quantity of inconsistencies associated with dependencies connected to node A . It is defined as:

$$I(A) = \sum_{i=1}^{l_A} r_{A_i} S_{A_i} = \sum_{i=1}^{l_A} r_{A_i} H(A, A(i), d)$$

Where:

- A denotes the node.
- A_i represents the dependency connected to node A , expressed as either $A_i = (A, A(i))$ or $(A(i), A)$.
- $A(i)$ denotes the other endpoint of dependency A_i with $1 \leq i \leq l_A$.
- l_A is the count of dependencies connected to node A .
- r_{A_i} is the weight of dependency A_i , which corresponds to the number of dependencies between A and $A(i)$, and is considered a constant.
- $S_{A_i} = H(A, A(i), d)$, H is the function that assesses whether dependency A_i is consistent with the architecture design, where 0 represents consistency and 1 represents inconsistency, which is equivalent to S_{A_i} . d indicates the direction of the dependency, with values of 1 or -1, representing a dependency or inverse dependency relationship.

Step 1 : Removing one refactoring operation (single)

Upon the removal of operation A , the impact is only on the inconsistency quantity of dependencies connected to node A . If operation A leads to a change for node A from A to A' , the change in inconsistency quantity, $\Delta I(A)$, is given by:

$$\begin{aligned} \Delta I(A) &= \sum_{i=1}^{l_A} r_{A_i} \Delta S_{A_i} = \sum_{i=1}^{l_A} r_{A_i} \Delta H(A, A(i), d) \\ &= \sum_{i=1}^{l_A} r_{A_i} [H(A', A(i), d) - H(A, A(i), d)] \end{aligned}$$

Analogous expressions can be derived for $\Delta I(B)$ and $\Delta I(C)$.

$$\begin{aligned}
\Delta I(B) &= \sum_{i=1}^{l_B} r_{B_i} \Delta S_{B_i} = \sum_{i=1}^{l_B} r_{B_i} \Delta H(B, B(i), d) \\
&= \sum_{i=1}^{l_B} r_{B_i} [H(B', B(i), d) - H(B, B(i), d)] \\
\Delta I(C) &= \sum_{i=1}^{l_C} r_{C_i} \Delta S_{C_i} = \sum_{i=1}^{l_C} r_{C_i} \Delta H(C, C(i), d) \\
&= \sum_{i=1}^{l_C} r_{C_i} [H(C', C(i), d) - H(C, C(i), d)]
\end{aligned}$$

Step 2: Removing two refactoring operation concurrently (pair)

Assuming dependencies exist between nodes A and B , which are characterized as follows:

- Dependencies connecting A to B are the l_A -th and $l_A - 1$ -th dependencies in the set of dependencies connected to A ,
- Dependencies connecting B to A are the $l_B - 2$ -th and $l_B - 3$ -th dependencies in the set of dependencies connected to B

The inconsistency quantity connected to nodes A and B can be formulated as:

$$\begin{aligned}
I(A, B) &= \sum_{i=1}^{l_A-2} r_{A_i} S_{A_i} + \sum_{i=1}^{l_B-4} r_{B_i} S_{B_i} + \sum_{i=l_B-1}^{l_B} r_{B_i} S_{B_i} + r_{A_{l_A}} H(A, B, 1) \\
&\quad + r_{A_{l_A-1}} H(A, B, -1)
\end{aligned}$$

And the change in inconsistency quantity due to the concurrent removal of operations A, B, C is expressed as:

$$\begin{aligned}
\Delta I(A, B) &= \sum_{i=1}^{l_A-2} r_{A_i} \Delta S_{A_i} + \sum_{i=1}^{l_B-4} r_{B_i} \Delta S_{B_i} + \sum_{i=l_B-1}^{l_B} r_{B_i} \Delta S_{B_i} + r_{A_{l_A}} \Delta H(A, B, 1) \\
&\quad + r_{A_{l_A-1}} \Delta H(A, B, -1)
\end{aligned}$$

Similarly,

$$\begin{aligned}
\Delta I(B, C) &= \sum_{i=1}^{l_B-2} r_{B_i} \Delta S_{B_i} + \sum_{i=1}^{l_C-4} r_{C_i} \Delta S_{C_i} + \sum_{i=l_C-1}^{l_C} r_{C_i} \Delta S_{C_i} + r_{B_{l_B}} \Delta H(B, C, 1) \\
&\quad + r_{B_{l_B-1}} \Delta H(B, C, -1)
\end{aligned}$$

$$\begin{aligned}\Delta I(C, A) = & \sum_{i=1}^{l_C-2} r_{C_i} \Delta S_{C_i} + \sum_{i=1}^{l_A-4} r_{A_i} \Delta S_{A_i} + \sum_{i=l_A-1}^{l_A} r_{A_i} \Delta S_{A_i} + r_{C_{l_C}} \Delta H(C, A, 1) \\ & + r_{C_{l_C-1}} \Delta H(C, A, -1)\end{aligned}$$

Step 3: Removing three refactoring operation concurrently (triple)

Assuming dependencies exist between nodes A , B and C , which are characterized as follows:

- Dependencies connecting A to B are the l_A -th and $l_A - 1$ -th dependencies in the set of dependencies connected to A ,
- Dependencies connecting A to C are the $l_A - 2$ -th and $l_A - 3$ -th dependencies in the set of dependencies connected to A ,
- Dependencies connecting B to C are the l_B -th and $l_B - 1$ -th dependencies in the set of dependencies connected to B ,
- Dependencies connecting B to A are the $l_B - 2$ -th and $l_B - 3$ -th dependencies in the set of dependencies connected to B ,
- Dependencies connecting C to A are the l_C -th and $l_C - 1$ -th dependencies in the set of dependencies connected to C ,
- Dependencies connecting C to B are the $l_C - 2$ -th and $l_C - 3$ -th dependencies in the set of dependencies connected to C .

The inconsistency quantity connected to nodes A , B and C can be formulated as

$$\begin{aligned}I(A, B, C) = & \sum_{i=1}^{l_A-4} r_{A_i} S_{A_i} + \sum_{i=1}^{l_B-4} r_{B_i} S_{B_i} + \sum_{i=1}^{l_C-4} r_{C_i} S_{C_i} \\ & + r_{A_{l_A}} H(A, B, 1) + r_{A_{l_A-1}} H(A, B, -1) \\ & + r_{B_{l_B}} H(B, C, 1) + r_{B_{l_B-1}} H(B, C, -1) \\ & + r_{C_{l_C}} H(C, A, 1) + r_{C_{l_C-1}} H(C, A, -1)\end{aligned}$$

And the change in inconsistency quantity due to the concurrent removal of operations A , B and C is expressed as

$$\begin{aligned}\Delta I(A, B, C) = & \sum_{i=1}^{l_A-4} r_{A_i} \Delta S_{A_i} + \sum_{i=1}^{l_B-4} r_{B_i} \Delta S_{B_i} + \sum_{i=1}^{l_C-4} r_{C_i} \Delta S_{C_i} \\ & + r_{A_{l_A}} \Delta H(A, B, 1) + r_{A_{l_A-1}} \Delta H(A, B, -1) \\ & + r_{B_{l_B}} \Delta H(B, C, 1) + r_{B_{l_B-1}} \Delta H(B, C, -1)\end{aligned}$$

$$+r_{c_{l_C}}\Delta H(C,A,1)+r_{c_{l_C-1}}\Delta H(C,A,-1)$$

Step 4: Scenario Coverage Confirmation

To prove whether the scenario of removing single and pair refactoring operations can cover the scenario of removing triple refactoring operations. It can be calculated as follows:

$$\begin{aligned}
& \Delta I(A,B)+\Delta I(B,C)+\Delta I(C,A)-\Delta I(A,B,C) \\
&= \left(\sum_{i=1}^{l_A-2} r_{A_i} \Delta S_{A_i} + \sum_{i=1}^{l_B-4} r_{B_i} \Delta S_{B_i} + \sum_{i=l_B-1}^{l_B} r_{B_i} \Delta S_{B_i} \right) + \left(\sum_{i=1}^{l_B-2} r_{B_i} \Delta S_{B_i} + \sum_{i=1}^{l_C-4} r_{C_i} \Delta S_{C_i} \right. \\
&\quad \left. + \sum_{i=l_C-1}^{l_C} r_{C_i} \Delta S_{C_i} \right) + \left(\sum_{i=1}^{l_C-2} r_{C_i} \Delta S_{C_i} + \sum_{i=1}^{l_A-4} r_{A_i} \Delta S_{A_i} + \sum_{i=l_A-1}^{l_A} r_{A_i} \Delta S_{A_i} \right) \\
&\quad - \left(\sum_{i=1}^{l_A-4} r_{A_i} \Delta S_{A_i} + \sum_{i=1}^{l_B-4} r_{B_i} \Delta S_{B_i} + \sum_{i=1}^{l_C-4} r_{C_i} \Delta S_{C_i} \right) \\
&= \sum_{i=1}^{l_A-2} r_{A_i} \Delta S_{A_i} + \sum_{i=l_B-1}^{l_B} r_{B_i} \Delta S_{B_i} + \sum_{i=1}^{l_B-2} r_{B_i} \Delta S_{B_i} + \sum_{i=l_C-1}^{l_C} r_{C_i} \Delta S_{C_i} + \sum_{i=1}^{l_C-2} r_{C_i} \Delta S_{C_i} \\
&\quad + \sum_{i=l_A-1}^{l_A} r_{A_i} \Delta S_{A_i} \\
&= \sum_{i=1}^{l_A} r_{A_i} \Delta S_{A_i} + \sum_{i=1}^{l_B} r_{B_i} \Delta S_{B_i} + \sum_{i=1}^{l_C} r_{C_i} \Delta S_{C_i} \\
&= \Delta I(A) + \Delta I(B) + \Delta I(C)
\end{aligned}$$

Therefore,

$$\Delta I(A,B,C) = \Delta I(A,B) + \Delta I(B,C) + \Delta I(C,A) + \Delta I(A) + \Delta I(B) + \Delta I(C)$$

If removing operations A, B, and C individually does not cause any inconsistency, and removing in pairs does not cause any inconsistency, that is $\Delta I(A) = \Delta I(B) = \Delta I(C) = \Delta I(A,B) = \Delta I(B,C) = \Delta I(C,A) = 0$. Then $\Delta I(A,B,C) = 0$

This result demonstrates that if the individual removal of operations A, B and C does not introduce any inconsistencies, and removing in pairs does not cause any inconsistency, then their removal in triple will also not affect the inconsistency quantity.