

# L23/24 - Dijkstra's Algorithm (Single Source Shortest Path) (Greedy Algorithm)

Tuesday, June 16, 2020 7:30 AM

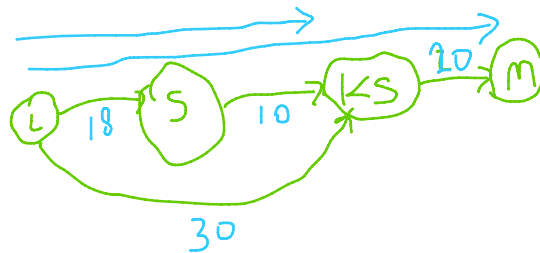
Shortest Distance: Sum of weights on edges from source to destination vertices.

Pre-Conditions:

1. All edge weights must be positive.
2. Directed weighted Graphs

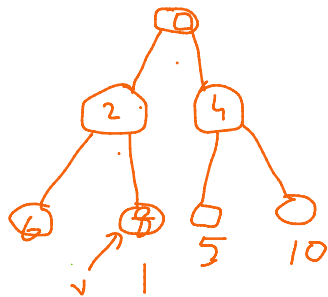
Constraints

Principle of Optimality:



INITIALIZE-SINGLE-SOURCE( $G, s$ )

1. for each vertex  $v \in G.V$
2.  $v.d = \infty$
3.  $v.\pi = \text{NIL}$
- 4.
5.  $s.d = 0$



DIJKSTRA ( $G, w, s$ )

1. INITIALIZE-SINGLE-SOURCE( $G, s$ )
2.  $X = \emptyset$
3.  $H = G.V$
4. While  $H \neq \emptyset$
5.  $u = \text{EXTRACT-MIN}(H)$
6.  $X = X \cup \{u\}$
7. for each vertex  $v \in G.Adj[u]$
8. if  $v.d > u.d + w(u, v)$
9.  $v.d = u.d + w(u, v)$
10.  $v.\pi = u$

$O(|V|)$

$O(|V| * \log |V|)$

$O(|V|)$

$O(|V| * \log |V|)$

$O(|E|)$

$O(|E|)$

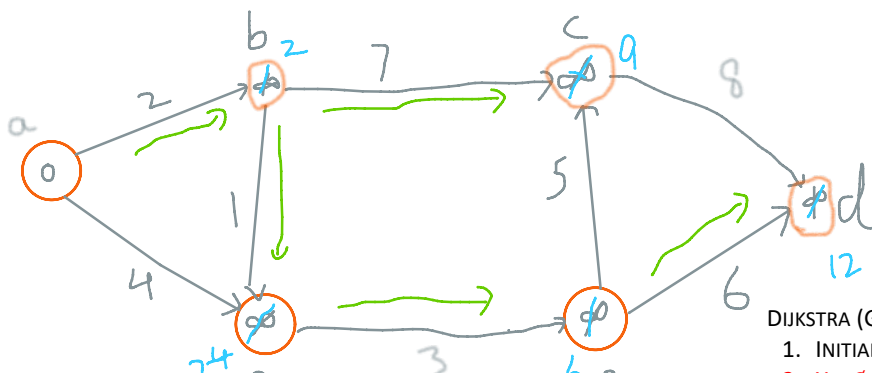
$O(|E| * \log(|V|))$

$O(|E|)$

$$O((|V| + |E|) * \log(|V|))$$

Adjacency List

a b|2 e|4  
b c|7 e|1  
c d|1  
d  
e f|3  
f c|5 d|6



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1. INITIALIZE-SINGLE-SOURCE( $G, s$ )

$O(|V|)$



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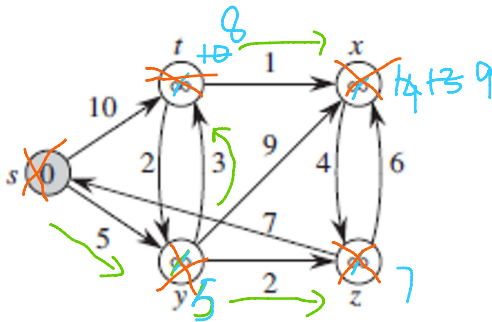
$O(|E|)$

$O(|E| * \log |V|)$

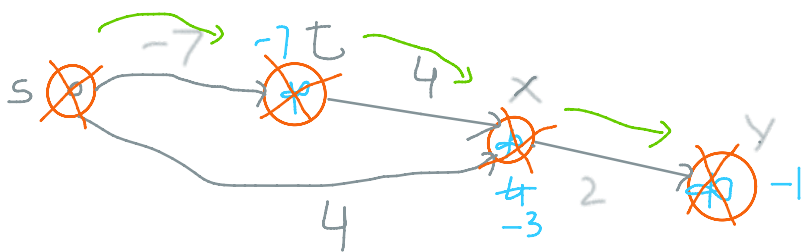
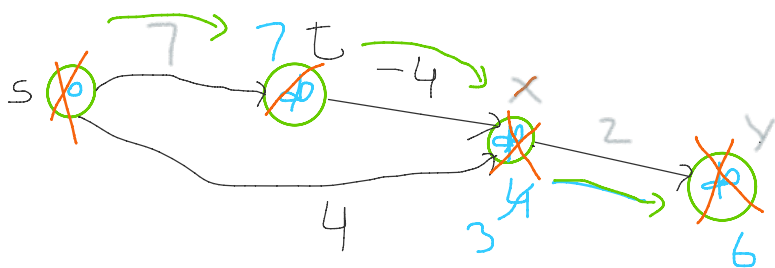
$O(|E|)$

| Predecessor Node | Nil  | a        | b        | f        | a <sup>b</sup> | e        |
|------------------|------|----------|----------|----------|----------------|----------|
| Selected         | d(a) | d(b)     | d(c)     | d(d)     | d(e)           | d(f)     |
|                  | 0    | $\infty$ | $\infty$ | $\infty$ | $\infty$       | $\infty$ |
| a                | 0    | 2        | $\infty$ | $\infty$ | 4              | $\infty$ |
| b                | 0    | 2        | 9        | $\infty$ | 3              | $\infty$ |
| e                | 0    | 2        | 9        | $\infty$ | 3              | 6        |
| f                | 0    | 2        | 9        | 12       | 3              | 6        |
| c                | 0    | 2        | 9        | 12       | 3              | 6        |
| d                | 0    | 2        | 9        | 12       | 3              | 6        |

Build the Distance Table for the Directed Weighted Graph shown below:



| Predecessor Node | Nil | s        | t        | y        | z        |
|------------------|-----|----------|----------|----------|----------|
| Selected         | s.d | t.d      | y.d      | x.d      | z.d      |
| s                | 0   | $\infty$ | $\infty$ | $\infty$ | $\infty$ |
| y                | 0   | 8        | 5        | 14       | 7        |
| z                | 0   | 8        | 5        | 13       | 7        |
| t                | 0   | 8        | 5        | 9        | 7        |
| x                | 0   | 8        | 5        | 9        | 7        |



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3.  $H = G.V$

$|V|$

$O(|V| \cdot \log |V|)$

4. While  $H \neq \emptyset$

$O(|V|)$

5.  $u = \text{EXTRACT-MIN}(H)$

$O(|V| \cdot \log |V|)$

6.  $X = X \cup \{u\}$

- 
- ```

graph TD
    5((5)) --- 8((8))
    5 --- 12((12))
    8 --- 10((10))
    8 --- 7((7))
    12 --- 15((15))

```

$$\begin{aligned} & O(2(|V| \times \log |V|)) + O(|E| \times \log(|V|)) \\ &= O((|V| + |E|) \times \log |V|) \\ &= O(|E| \times \log |V|) \end{aligned}$$