# L21 - Topological Sort/Job Sequences

Tuesday, June 9, 2020 8:20 AM

An application of Directed Acyclic Graph (DAG) i.e. a directed graph without cycle.

Topological Sorting of a DAG is Linear Ordering of its vertices such that for an edge (u, v), the vertex u appears before vertex v.

### TOPOLOGICAL-SORT(G)

- 1 call DFS(G) to compute finishing times v.f for each vertex v
- 2 as each vertex is finished, insert it onto the front of a linked list
- 3 return the linked list of vertices

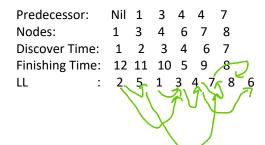
st Predecessor: Nil 1 3 4 4 7 Nil 2
Node: 1 3 4 6 7 8 2 5
Discovery Time: 1 2 3 4 6 7 13 14
Finishing Time: 12 11 10 5 9 8 16 15

#### DFS(G)

- 1. for each vertex  $k \in G.V$
- 2. u. *color* = WHITE
- 3.  $k.\pi = NIL$
- 4. Linked\_list LL
- 5. *time* = 0
- 6. for each vertex  $k \in G.V$
- 7. if k. *color* == WHITE
- 8. DFS-VISIT(G; k)

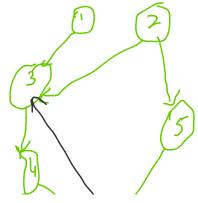
#### DFS-VISIT(G; k)

- 1. time = time + 1 // white vertex u has just been discovered
- 2. k. *d* = *time* // k. d means the time when vertex k was first **discovered**.
- 3. k. color = GRAY
- 4. for each  $l \in G.Adj[k]$  // explore edge (k; l)
- 5. if l. *color* == WHITE
- 6.  $I.\pi = k$
- 7. DFS-VISIT(G; I)
- 8. k. color = BLACK // blacken k; it is finished
- 9. *time* = *time* + 1



### Adjacency List Representation

1	3	
2	3	5
3	4	
4	6	7
5	7	
6		
7	8	
8		



- 7. DFS-VISIT(G; I)
- 8. k. color = BLACK // blacken k; it is finished
- 9. *time* = *time* + 1
- 10. k. f = time //finish time
- 11. LL.Insert\_at\_Head(k)

# Time Complexity of Topological Sorting O(|V|+|E|)

Predecessor: Node: Discovery Time: Finish Time:

 Predecessor:
 Nil
 3
 4
 4
 7
 Nil
 1
 2

 Node:
 3
 4
 6
 7
 8
 1
 2
 5

 Discovery Time:
 1
 2
 3
 5
 6
 11
 13
 14

 Finish Time:
 10
 9
 4
 8
 7
 12
 16
 15

2 5 1 3 4 7 8 6

