

Longest Common Subsequence

February 25, 2023

Bangladesh University of Engineering and Technology

What is LCS

Recursive Formulation

Using Dynamic Programming

What is LCS

Problem Definition

Given two strings `text1` and `text2`, return the length of their longest common subsequence

What is a sub-sequence?

Subsequence

A sub-sequence is a sequence that can be obtained by deleting elements from the original sequence without changing order of the elements.

Subsequence

A sub-sequence is a sequence that can be obtained by deleting elements from the original sequence without changing order of the elements.

Let `str = "ABCDE"` be a sequence

A B C D E F

Subsequence

A sub-sequence is a sequence that can be obtained by deleting elements from the original sequence without changing order of the elements.

Let `str = "ABCDE"` be a sequence

A B C D E F

Then some subsequence of `str` are : ABC

Subsequence

A sub-sequence is a sequence that can be obtained by deleting elements from the original sequence without changing order of the elements.

Let `str = "ABCDE"` be a sequence

A B C D E F

Then some subsequence of `str` are : ABC ACE

Subsequence

A sub-sequence is a sequence that can be obtained by deleting elements from the original sequence without changing order of the elements.

Let `str = "ABCDE"` be a sequence

A B C D E F

Then some subsequence of `str` are : ABC ACE BEF etc.

Longest Common Subsequence

From the two given sequences, we have to find a subsequence that is present in both of them and is the longest.

$$\begin{aligned} S_1 &= \text{ACCGGTCGAGTGC GCGGAAGCCGGCCGAA} \\ S_2 &= \text{GTCGTTCGGAATGCCGTTGCTCTGTAAA} \end{aligned}$$

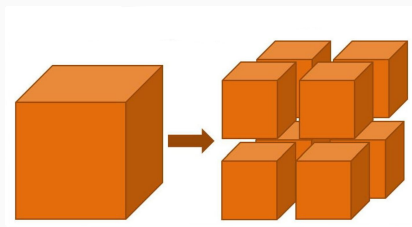
Longest Common Subsequence

$S_1 =$ ACCGGTCGAGTGCGCGGAAGCCGGCCGAA
 $S_2 =$ GTCGTTCGGAATGCCGTTGCTCTGTAAA

$S_3 =$ GTCGTTCGGAAGCCGGCCGAA

How to approach?

Break it down into smaller subproblems.



Breaking into smaller sub-problems

Suppose we have two sequences

X = A B C A D E

Y = A C B E

Breaking into smaller sub-problems

Suppose we have two sequences

X = A B C A D **E**

Y = A C B **E**

Breaking into smaller sub-problems

Suppose we have two sequences

X = A B C A D **E**

Y = A C B **E**

Longest Common Subsequence :E

Breaking into smaller sub-problems

Suppose we have two sequences

X = A B C A D E

Y = A C B E

Longest Common Subsequence :E

Breaking into smaller sub-problems

Suppose we have two sequences

X = A B C A D E

Y = A C B E

Longest Common Subsequence : _____E

Breaking into smaller sub-problems

Suppose we have two sequences

X = A B C A D **F**

Y = A C B E

Longest Common Subsequence : _____

Breaking into smaller sub-problems

Suppose we have two sequences

X = A B C A D **F**

Y = A C B **E**

Longest Common Subsequence : _____

Breaking into smaller sub-problems

Suppose we have two sequences

X = A B C A D **F**

Y = A C B **E**

Longest Common Subsequence : _____

Breaking into smaller sub-problems

Suppose we have two sequences

X = A B C A D F

Y = A C B E

Longest Common Subsequence : _____

Breaking into smaller sub-problems

Suppose we have two sequences

X = A B C A D F

Y = A C B E

Longest Common Subsequence : _____

Recursive Formulation

Recursion

If we have two sequences

$$X = [1 \cdot \cdot \cdot m] \qquad Y = [1 \cdot \cdot \cdot n]$$

Recursion

If we have two sequences

$$X = [\underbrace{1 \cdot \cdot \cdot a \cdot \cdot m}] \qquad Y = [\underbrace{1 \cdot \cdot b \cdot \cdot \cdot n}]$$

Let $LCS(a, b)$ denote the longest common subsequence of $x_1x_2\dots x_a$ and $y_1y_2\dots y_a$.

Recursion

If we have two sequences

$$X = [\underbrace{1 \cdot \cdot \cdot \cdot m}_{\text{red line}}] \qquad Y = [\underbrace{1 \cdot \cdot \cdot \cdot n}_{\text{red line}}]$$

Let $LCS(a, b)$ denote the longest common subsequence of $x_1x_2\dots x_a$ and $y_1y_2\dots y_a$.

Our goal is to find $LCS(m, n)$

Recursive Formulation

$$X = [1 \ 2 \ 3 \ \dots \ m] \quad Y = [1 \ 2 \ 3 \ \dots \ n]$$

Recursive Formulation

$$X = [1 \ 2 \ 3 \ \dots \ m] \quad Y = [1 \ 2 \ 3 \ \dots \ n]$$

When $X[m] == Y[n]$

Recursive Formulation

$$X = [\boxed{1 \ 2 \ 3 \ \dots} \ m] \quad Y = [\boxed{1 \ 2 \ 3 \ \dots} \ n]$$

When $X[m] == Y[n]$

$$LCS(m, n) = LCS(m - 1, n - 1) + 1$$

Recursive Formulation

$$X = [1 \ 2 \ 3 \ \dots \ m] \quad Y = [1 \ 2 \ 3 \ \dots \ n]$$

When $X[m] \neq Y[n]$

Recursive Formulation

$$X = [\boxed{1 \ 2 \ 3 \ \dots \ m}] \quad Y = [\boxed{1 \ 2 \ 3 \ \dots} \ n]$$

When $X[m] \neq Y[n]$

$$LCS(m, n) = \max\{LCS(m, n - 1),$$

Recursive Formulation

$$X = [\boxed{1 \ 2 \ 3 \ \dots} \ m] \quad Y = [\boxed{1 \ 2 \ 3 \ \dots \ n}]$$

When $X[m] \neq Y[n]$

$$LCS(m, n) = \max\{LCS(m, n - 1), LCS(m - 1, n)\}$$

Recursive Formulation

$$X = [1 \ 2 \ 3 \ \cdots \ m] \quad Y = [1 \ 2 \ 3 \ \cdots \ n]$$

The recursion solution:

$$LCS(m, n) = \begin{cases} LCS(m-1, n-1) + 1 & X[m] == Y[n] \\ \max\{LCS(m-1, n), LCS(m, n-1)\} & X[m] \neq Y[n] \end{cases}$$

What about the base case?

Recursive Formulation

If $m == 0$ then $LCS(m,n) = 0$

If $n == 0$ then $LCS(m,n) = 0$

The final recursion solution:

$$LCS(m, n) = \begin{cases} 0 & m == 0 || n == 0 \\ LCS(m - 1, n - 1) + 1 & X[m] == Y[n] \\ \max\{LCS(m - 1, n), LCS(m, n - 1)\} & X[m] \neq Y[n] \end{cases}$$

Recursion

$$LCS(m, n) = \begin{cases} 0 & \text{if } m = 0 \text{ or } n = 0 \\ LCS(m-1, n-1) + 1 & \text{if } X[m] = Y[n] \\ \text{MAX}(LCS(m-1, n), LCS(m, n-1)) & \text{if } X[m] \neq Y[n] \end{cases}$$

```
def LCS(X[1..m], Y[1..n]) :  
    if m = 0 or n = 0  
        return 0
```

Recursion

$$LCS(m, n) = \begin{cases} 0 & \text{if } m = 0 \text{ or } n = 0 \\ LCS(m-1, n-1) + 1 & \text{if } X[m] = Y[n] \\ \text{MAX}(LCS(m-1, n), LCS(m, n-1)) & \text{if } X[m] \neq Y[n] \end{cases}$$

```
def LCS(X[1..m], Y[1..n]) :  
    if m = 0 or n = 0  
        return 0  
  
    if X[m] = Y[n] :  
        result = LCS(m-1, n-1) + 1
```


Recursion

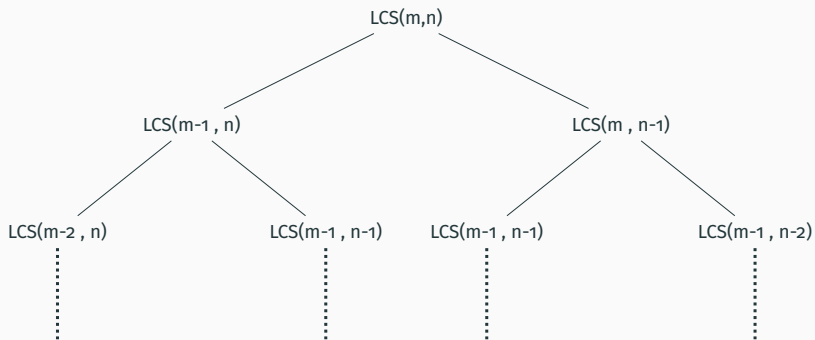
$$LCS(m, n) = \begin{cases} 0 & \text{if } m = 0 \text{ or } n = 0 \\ LCS(m-1, n-1) + 1 & \text{if } X[m] = Y[n] \\ \text{MAX}(LCS(m-1, n), LCS(m, n-1)) & \text{if } X[m] \neq Y[n] \end{cases}$$

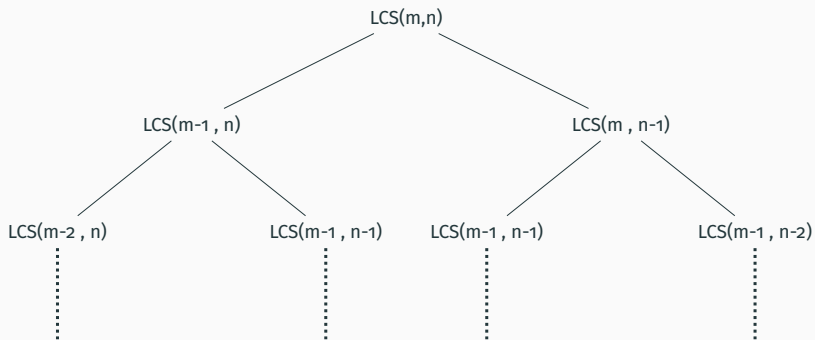
```
def LCS(X[1..m], Y[1..n]) :  
    if m = 0 or n = 0  
        return 0  
  
    if X[m] = Y[n] :  
        result = LCS(m-1, n-1) + 1  
  
    else :  
        result = max(LCS(m-1, n), LCS(m, n-1))
```

Recursion

$$LCS(m, n) = \begin{cases} 0 & \text{if } m = 0 \text{ or } n = 0 \\ LCS(m-1, n-1) + 1 & \text{if } X[m] = Y[n] \\ \max(LCS(m-1, n), LCS(m, n-1)) & \text{if } X[m] \neq Y[n] \end{cases}$$

```
def LCS(X[1..m], Y[1..n]) :  
    if m = 0 or n = 0  
        return 0  
  
    if X[m] = Y[n] :  
        result = LCS(m-1, n-1) + 1  
  
    else :  
        result = max(LCS(m-1, n), LCS(m, n-1))  
  
    return result
```





So the time complexity is $O(2^n)$

Memoized Algorithm

$$LCS(m, n) = \begin{cases} 0 & \text{if } m = 0 \text{ or } n = 0 \\ LCS(m-1, n-1) + 1 & \text{if } X[m] = Y[n] \\ \text{MAX}(LCS(m-1, n), LCS(m, n-1)) & \text{if } X[m] \neq Y[n] \end{cases}$$

```
def LCS(X[1..m], Y[1..n]) :  
    if m = 0 or n = 0  
        return 0  
    if LCS_memo[m][n] != -INF:  
        return LCS_memo[m][n]  
    else if X[m] = Y[n] :  
        result = LCS(m-1, n-1) + 1  
    else :  
        result = max(LCS(m-1, n), LCS(m, n-1))  
    LCS_memo[m][n] = result  
    return result
```

Memoized Algorithm

$$LCS(m, n) = \begin{cases} 0 & \text{if } m = 0 \text{ or } n = 0 \\ LCS(m-1, n-1) + 1 & \text{if } X[m] = Y[n] \\ \text{MAX}(LCS(m-1, n), LCS(m, n-1)) & \text{if } X[m] \neq Y[n] \end{cases}$$

```
def LCS(X[1..m], Y[1..n]) :  
    if m = 0 or n = 0  
        return 0  
    if LCS_memo[m][n] != -INF:  
        return LCS_memo[m][n]  
    else if X[m] = Y[n] :  
        result = LCS(m-1, n-1) + 1  
    else :  
        result = max(LCS(m-1, n), LCS(m, n-1))  
    LCS_memo[m][n] = result  
    return result
```

Time complexity is $O(mn)$

Using Dynamic Programming

Dynamic Programming

		0	1	2	3	4	5	6
	y_j	B	D	C	A	B	A	
0	x_i							
1	A							
2	B							
3	C							
4	B							
5	D							
6	A							
7	B							

If $m=0$ or $n=0$

$$LCS(m,n) = 0$$

Else if $X[m] = Y[n]$

$$LCS(m,n) = LCS(m-1, n-1) + 1$$

Else

$$LCS(m,n) = \max(LCS(m-1,n), LCS(m, n-1))$$

Dynamic Programming

		0	1	2	3	4	5	6
		Y_j	B	D	C	A	B	A
0	X_i	0	0	0	0	0	0	0
1	A	0						
2	B	0						
3	C	0						
4	B	0						
5	D	0						
6	A	0						
7	B	0						

If $m=0$ or $n=0$

$$\text{LCS}(m,n) = 0$$

Else if $X[m] = Y[n]$

$$\text{LCS}(m,n) = \text{LCS}(m-1, n-1) + 1$$

Else

$$\text{LCS}(m,n) = \max(\text{LCS}(m-1,n), \text{LCS}(m, n-1))$$

Dynamic Programming

		0	1	2	3	4	5	6
		y_j	B	D	C	A	B	A
0	x_i	0	0	0	0	0	0	0
1	A	0						
2	B	0						
3	C	0						
4	B	0						
5	D	0						
6	A	0						
7	B	0						

If $m=0$ or $n=0$

$$\text{LCS}(m,n) = 0$$

Else if $X[m] = Y[n]$

$$\text{LCS}(m,n) = \text{LCS}(m-1, n-1) + 1$$

Else

$$\text{LCS}(m,n) = \max(\text{LCS}(m-1,n), \text{LCS}(m, n-1))$$

Dynamic Programming

		0	1	2	3	4	5	6
		y_j	B	D	C	A	B	A
0	x_i	0	0	0	0	0	0	0
1	A	0	0					
2	B	0						
3	C	0						
4	B	0						
5	D	0						
6	A	0						
7	B	0						

If $m=0$ or $n=0$

$$\text{LCS}(m,n) = 0$$

Else if $X[m] = Y[n]$

$$\text{LCS}(m,n) = \text{LCS}(m-1, n-1) + 1$$

Else

$$\text{LCS}(m,n) = \max(\text{LCS}(m-1,n), \text{LCS}(m, n-1))$$

Dynamic Programming

		0	1	2	3	4	5	6
	y_j	B	D	C	A	B	A	
0	x_i	0	0	0	0	0	0	0
1	A	0	0					
2	B	0						
3	C	0						
4	B	0						
5	D	0						
6	A	0						
7	B	0						

If $m=0$ or $n=0$

$$\text{LCS}(m,n) = 0$$

Else if $X[m] = Y[n]$

$$\text{LCS}(m,n) = \text{LCS}(m-1, n-1) + 1$$

Else

$$\text{LCS}(m,n) = \max(\text{LCS}(m-1,n), \text{LCS}(m, n-1))$$

Dynamic Programming

		0	1	2	3	4	5	6
	y_j	B	D	C	A	B	A	
0	x_i	0	0	0	0	0	0	0
1	A	0	0	0				
2	B	0						
3	C	0						
4	B	0						
5	D	0						
6	A	0						
7	B	0						

If $m=0$ or $n=0$

$$\text{LCS}(m,n) = 0$$

Else if $X[m] = Y[n]$

$$\text{LCS}(m,n) = \text{LCS}(m-1, n-1) + 1$$

Else

$$\text{LCS}(m,n) = \max(\text{LCS}(m-1,n), \text{LCS}(m, n-1))$$

Dynamic Programming

		0	1	2	3	4	5	6
		y_j	B	D	C	A	B	A
0	x_i	0	0	0	0	0	0	0
1	A	0	0	0				
2	B	0						
3	C	0						
4	B	0						
5	D	0						
6	A	0						
7	B	0						

If $m=0$ or $n=0$

$$\text{LCS}(m,n) = 0$$

Else if $X[m] = Y[n]$

$$\text{LCS}(m,n) = \text{LCS}(m-1, n-1) + 1$$

Else

$$\text{LCS}(m,n) = \max(\text{LCS}(m-1,n), \text{LCS}(m, n-1))$$

Dynamic Programming

		0	1	2	3	4	5	6
		y_j	B	D	C	A	B	A
0	x_i	0	↑ 0	↑ 0	↑ 0	0	0	0
1	A	0	0	0	0			
2	B	0						
3	C	0						
4	B	0						
5	D	0						
6	A	0						
7	B	0						

If $m=0$ or $n=0$

$$\text{LCS}(m,n) = 0$$

Else if $X[m] = Y[n]$

$$\text{LCS}(m,n) = \text{LCS}(m-1, n-1) + 1$$

Else

$$\text{LCS}(m,n) = \max(\text{LCS}(m-1,n), \text{LCS}(m, n-1))$$

Dynamic Programming

		0	1	2	3	4	5	6
		y_j	B	D	C	A	B	A
0	x_i	0	↑ 0	↑ 0	↑ 0	0	0	0
1	A	0	0	0	0			
2	B	0						
3	C	0						
4	B	0						
5	D	0						
6	A	0						
7	B	0						

If $m=0$ or $n=0$

$$LCS(m,n) = 0$$

Else if $X[m] = Y[n]$

$$LCS(m,n) = LCS(m-1, n-1) + 1$$

Else

$$LCS(m,n) = \max(LCS(m-1,n), LCS(m, n-1))$$

Dynamic Programming

		0	1	2	3	4	5	6
	y_j	B	D	C	A	B	A	
0	x_i	0	0	0	0	0	0	
1	A	0	0	0	0	1		
2	B	0						
3	C	0						
4	B	0						
5	D	0						
6	A	0						
7	B	0						

If $m=0$ or $n=0$

$$LCS(m,n) = 0$$

Else if $X[m] = Y[n]$

$$LCS(m,n) = LCS(m-1, n-1) + 1$$

Else

$$LCS(m,n) = \max(LCS(m-1,n), LCS(m, n-1))$$

Dynamic Programming

		0	1	2	3	4	5	6
		y_j	B	D	C	A	B	A
0	x_i	0	0	0	0	0	0	0
1	A	0	0	0	0	1	1	
2	B	0						
3	C	0						
4	B	0						
5	D	0						
6	A	0						
7	B	0						

If $m=0$ or $n=0$
 $LCS(m,n) = 0$

Else if $X[m] = Y[n]$
 $LCS(m,n) = LCS(m-1, n-1) + 1$

Else

$LCS(m,n) = \max(LCS(m-1,n), LCS(m, n-1))$

Dynamic Programming

		0	1	2	3	4	5	6
		y_j	B	D	C	A	B	A
0	x_i	0	0	0	0	0	0	0
1	A	0	0	0	0	1	1	
2	B	0						
3	C	0						
4	B	0						
5	D	0						
6	A	0						
7	B	0						

If $m=0$ or $n=0$
 $LCS(m,n) = 0$

Else if $X[m] = Y[n]$
 $LCS(m,n) = LCS(m-1, n-1) + 1$

Else

$LCS(m,n) = \max(LCS(m-1,n), LCS(m, n-1))$

Dynamic Programming

		0	1	2	3	4	5	6
	y_j	B	D	C	A	B	A	
0	x_i	0	0	0	0	0	0	
1	A	0	0	0	0	1	1	
2	B	0						
3	C	0						
4	B	0						
5	D	0						
6	A	0						
7	B	0						

If $m=0$ or $n=0$

$$LCS(m,n) = 0$$

Else if $X[m] = Y[n]$

$$LCS(m,n) = LCS(m-1, n-1) + 1$$

Else

$$LCS(m,n) = \max(LCS(m-1,n), LCS(m, n-1))$$

Dynamic Programming

		0	1	2	3	4	5	6
	y_j	B	D	C	A	B	A	
0	x_i	0	0	0	0	0	0	
1	A	0	0	0	0	1	1	
2	B	0						
3	C	0						
4	B	0						
5	D	0						
6	A	0						
7	B	0						

If $m=0$ or $n=0$

$$LCS(m,n) = 0$$

Else if $X[m] = Y[n]$

$$LCS(m,n) = LCS(m-1, n-1) + 1$$

Else

$$LCS(m,n) = \max(LCS(m-1,n), LCS(m, n-1))$$

Dynamic Programming

		0	1	2	3	4	5	6
	y_j		B	D	C	A	B	A
0	x_i	0	0	0	0	0	0	0
1	A	0	0	0	0	1	1	1
2	B	0	1					
3	C	0						
4	B	0						
5	D	0						
6	A	0						
7	B	0						

If $m=0$ or $n=0$
 $LCS(m,n) = 0$

Else if $X[m] = Y[n]$
 $LCS(m,n) = LCS(m-1, n-1) + 1$

Else

$LCS(m,n) = \max(LCS(m-1,n), LCS(m, n-1))$

Dynamic Programming

		0	1	2	3	4	5	6
	y_j		B	D	C	A	B	A
0	x_i	0	0	0	0	0	0	0
1	A	0	0	0	0	1	1	1
2	B	0	1					
3	C	0						
4	B	0						
5	D	0						
6	A	0						
7	B	0						

If $m=0$ or $n=0$
 $LCS(m,n) = 0$

Else if $X[m] = Y[n]$
 $LCS(m,n) = LCS(m-1, n-1) + 1$

Else

$LCS(m,n) = \max(LCS(m-1,n), LCS(m, n-1))$

Dynamic Programming

		0	1	2	3	4	5	6
	y_j	B	D	C	A	B	A	
0	x_i	0	↑ 0	↑ 0	↑ 0	0	0	0
1	A	0	↓ 0	0	0	↑ 1	← 1	1
2	B	0	↖ 1	← 1	← 1	↑ 1	↖ 2	
3	C	0						
4	B	0						
5	D	0						
6	A	0						
7	B	0						

If $m=0$ or $n=0$
 $LCS(m,n) = 0$

Else if $X[m] = Y[n]$
 $LCS(m,n) = LCS(m-1, n-1) + 1$

Else

$LCS(m,n) = \max(LCS(m-1,n), LCS(m, n-1))$

Dynamic Programming

		0	1	2	3	4	5	6
	y_j	B	D	C	A	B	A	
0	x_i	0	0	0	0	0	0	0
1	A	0	0	0	0	1	1	1
2	B	0	1	1	1	1	2	
3	C	0						
4	B	0						
5	D	0						
6	A	0						
7	B	0						

If $m=0$ or $n=0$
 $LCS(m,n) = 0$

Else if $X[m] = Y[n]$
 $LCS(m,n) = LCS(m-1, n-1) + 1$

Else

$LCS(m,n) = \max(LCS(m-1,n), LCS(m, n-1))$

Dynamic Programming

		0	1	2	3	4	5	6
	y_j	B	D	C	A	B	A	
0	x_i	0	0	0	0	0	0	
1	A	0	0	0	0	1	1	
2	B	0	1	1	1	1	2	
3	C	0						
4	B	0						
5	D	0						
6	A	0						
7	B	0						

If $m=0$ or $n=0$

$$LCS(m,n) = 0$$

Else if $X[m] = Y[n]$

$$LCS(m,n) = LCS(m-1, n-1) + 1$$

Else

$$LCS(m,n) = \max(LCS(m-1,n), LCS(m, n-1))$$

Dynamic Programming

		0	1	2	3	4	5	6
	y_j	B	D	C	A	B	A	
0	x_i	0	0	0	0	0	0	
1	A	0	0	0	0	1	1	
2	B	0	1	1	1	1	2	
3	C	0						
4	B	0						
5	D	0						
6	A	0						
7	B	0						

If $m=0$ or $n=0$

$$LCS(m,n) = 0$$

Else if $X[m] = Y[n]$

$$LCS(m,n) = LCS(m-1, n-1) + 1$$

Else

$$LCS(m,n) = \max(LCS(m-1,n), LCS(m, n-1))$$

Dynamic Programming

		0	1	2	3	4	5	6
	y_j	B	D	C	A	B	A	
0	x_i	0	↑ 0	↑ 0	↑ 0	0	0	0
1	A	0	↓ 0	0	0	↑ 1	← 1	1
2	B	0	↑ 1	← 1	← 1	1	↑ 1	← 2
3	C	0	↑ 1	↑ 1				
4	B	0						
5	D	0						
6	A	0						
7	B	0						

If $m=0$ or $n=0$

$$LCS(m,n) = 0$$

Else if $X[m] = Y[n]$

$$LCS(m,n) = LCS(m-1, n-1) + 1$$

Else

$$LCS(m,n) = \max(LCS(m-1,n), LCS(m, n-1))$$

Dynamic Programming

		0	1	2	3	4	5	6
	Y_j	B	D	C	A	B	A	
0	X_i	0	0	0	0	0	0	0
1	A	0	0	0	0	1	1	1
2	B	0	1	1	1	1	2	2
3	C	0	1	1	2	2	2	2
4	B	0						
5	D	0						
6	A	0						
7	B	0						

If $m=0$ or $n=0$

$$\text{LCS}(m,n) = 0$$

Else if $X[m] = Y[n]$

$$\text{LCS}(m,n) = \text{LCS}(m-1, n-1) + 1$$

Else

$$\text{LCS}(m,n) = \max(\text{LCS}(m-1,n), \text{LCS}(m, n-1))$$

Dynamic Programming

		0	1	2	3	4	5	6
	y_j	B	D	C	A	B	A	
0	x_i	0	0	0	0	0	0	
1	A	0	0	0	0	1	1	
2	B	0	1	1	1	1	2	
3	C	0	1	1	2	2	2	
4	B	0	1	1	2	2	3	
5	D	0						
6	A	0						
7	B	0						

If $m=0$ or $n=0$

$$LCS(m,n) = 0$$

Else if $X[m] = Y[n]$

$$LCS(m,n) = LCS(m-1, n-1) + 1$$

Else

$$LCS(m,n) = \max(LCS(m-1,n), LCS(m, n-1))$$

Dynamic Programming

		0	1	2	3	4	5	6
	Y_j	B	D	C	A	B	A	
0	X_i	0	0	0	0	0	0	0
1	A	0	0	0	0	1	1	1
2	B	0	1	1	1	1	2	2
3	C	0	1	1	2	2	2	2
4	B	0	1	1	2	2	3	3
5	D	0	1	2	2	2	3	3
6	A	0						
7	B	0						

If $m=0$ or $n=0$

$$LCS(m,n) = 0$$

Else if $X[m] = Y[n]$

$$LCS(m,n) = LCS(m-1, n-1) + 1$$

Else

$$LCS(m,n) = \max(LCS(m-1,n), LCS(m, n-1))$$

Dynamic Programming

		0	1	2	3	4	5	6
	Y_j	B	D	C	A	B	A	
0	X_i	0	0	0	0	0	0	0
1	A	0	0	0	0	1	1	1
2	B	0	1	1	1	1	2	2
3	C	0	1	1	2	2	2	2
4	B	0	1	1	2	2	3	3
5	D	0	1	2	2	2	3	3
6	A	0	1	2	2	3	3	4
7	B	0						

If $m=0$ or $n=0$

$$\text{LCS}(m,n) = 0$$

Else if $X[m] = Y[n]$

$$\text{LCS}(m,n) = \text{LCS}(m-1, n-1) + 1$$

Else

$$\text{LCS}(m,n) = \max(\text{LCS}(m-1,n), \text{LCS}(m, n-1))$$

Dynamic Programming

		0	1	2	3	4	5	6
	Y_j	B	D	C	A	B	A	
0	X_i	0	0	0	0	0	0	0
1	A	0	0	0	0	1	1	1
2	B	0	1	1	1	1	2	2
3	C	0	1	1	2	2	2	2
4	B	0	1	1	2	2	3	3
5	D	0	1	2	2	2	3	3
6	A	0	1	2	2	3	3	4
7	B	0	1	2	2	3	4	4

If $m=0$ or $n=0$

$$LCS(m,n) = 0$$

Else if $X[m] = Y[n]$

$$LCS(m,n) = LCS(m-1, n-1) + 1$$

Else

$$LCS(m,n) = \max(LCS(m-1,n), LCS(m, n-1))$$

Dynamic Programming

		0	1	2	3	4	5	6
	Y_j	B	D	C	A	B	A	
0	X_i	0	0	0	0	0	0	0
1	A	0	0	0	0	1	1	1
2	B	0	1	1	1	1	2	2
3	C	0	1	1	2	2	2	2
4	B	0	1	1	2	2	3	3
5	D	0	1	2	2	2	3	3
6	A	0	1	2	2	3	3	4
7	B	0	1	2	2	3	4	4

If $m=0$ or $n=0$

$LCS(m,n) = 0$

Else if $X[m] = Y[n]$

$LCS(m,n) = LCS(m-1, n-1) + 1$

Else

$LCS(m,n) = \max(LCS(m-1,n), LCS(m, n-1))$

Simulated Annealing (SA)

		0	1	2	3	4	5	6
	y_j	B	D	C	A	B	A	
0	x_i	0	0	0	0	0	0	0
1	A	0	0	0	0	1	1	1
2	B	0	1	1	1	1	2	2
3	C	0	1	1	2	2	2	2
4	B	0	1	1	2	2	3	3
5	D	0	1	2	2	2	3	3
6	A	0	1	2	2	3	3	4
7	B	0	1	2	2	3	4	4

```
def dp_lcs(x[1..m], y[1..n]):
```

```
    for i = 0 to m :
```

```
        L[i, 0] = 0
```

```
    for j = 0 to n :
```

```
        L[0, j] = 0
```

```
    for i = 1 to m :
```

```
        for j = 1 to n :
```

```
            if x[i] == y[j] :
```

```
                L[i, j] = L[i-1, j-1] + 1
```

```
            else if L[i-1, j] >= L[i, j-1] :
```

```
                L[i, j] = L[i-1, j]
```

```
            else :
```

```
                L[i, j] = L[i, j-1]
```

```
    return L[m, n]
```

**What if we want to find the
subsequence too?**

Going back

		0	1	2	3	4	5	6
	y_j	B	D	C	A	B	A	
0	x_i	0	↑ 0	↑ 0	↑ 0	0	0	0
1	A	0	↖ 0	0	0	↑ 1	← 1	↖ 1
2	B	0	↑ 1	← 1	← 1	1	↑ 2	← 2
3	C	0	↑ 1	↑ 1	↑ 2	↑ 2	← 2	↑ 2
4	B	0	↑ 1	↖ 1	↑ 2	↑ 2	↑ 3	← 3
5	D	0	↑ 1	↑ 2	↑ 2	↑ 2	↑ 3	↖ 3
6	A	0	↑ 1	↑ 2	↑ 2	↑ 3	↑ 3	↑ 4
7	B	0	↑ 1	↑ 2	↑ 2	↑ 3	↖ 4	↑ 4

If $m=0$ or $n=0$
 $LCS(m,n) = 0$

Else if $X[m] = Y[n]$
 $LCS(m,n) = LCS(m-1, n-1) + 1$

Else
 $LCS(m,n) = \max(LCS(m-1,n), LCS(m, n-1))$

Going back

		0	1	2	3	4	5	6
	y_j	B	D	C	A	B	A	
0	x_i	0	↑ 0	↑ 0	↑ 0	0	0	0
1	A	0	↖ 0	0	0	↑ 1	← 1	↖ 1
2	B	0	↑ 1	← 1	← 1	1	↑ 2	← 2
3	C	0	↑ 1	↑ 1	↑ 2	↑ 2	← 2	↑ 2
4	B	0	↑ 1	↖ 1	↑ 2	↑ 2	↑ 3	← 3
5	D	0	↑ 1	↑ 2	↑ 2	↑ 2	↑ 3	↖ 3
6	A	0	↑ 1	↑ 2	↑ 2	↑ 3	↑ 3	↖ 4
7	B	0	↑ 1	↑ 2	↑ 2	↑ 3	↑ 4	4

If $m=0$ or $n=0$
 $LCS(m,n) = 0$

Else if $X[m] = Y[n]$
 $LCS(m,n) = LCS(m-1, n-1) + 1$

Else
 $LCS(m,n) = \max(LCS(m-1,n), LCS(m,n-1))$

Going back

		0	1	2	3	4	5	6
	y_j	B	D	C	A	B	A	
0	x_i	0	↑ 0	↑ 0	↑ 0	0	0	0
1	A	0	↖ 0	0	0	↑ 1	← 1	↖ 1
2	B	0	↑ 1	← 1	← 1	1	↑ 2	← 2
3	C	0	↑ 1	↑ 1	↑ 2	↑ 2	← 2	↑ 2
4	B	0	↑ 1	↖ 1	↑ 2	↑ 2	↑ 3	← 3
5	D	0	↑ 1	↑ 2	↑ 2	↑ 2	↑ 3	↑ 3
6	A	0	↑ 1	↑ 2	↑ 2	↑ 3	↑ 3	↖ 4
7	B	0	↑ 1	↑ 2	↑ 2	↑ 3	↑ 4	↑ 4

If $m=0$ or $n=0$
 $LCS(m,n) = 0$

Else if $X[m] = Y[n]$
 $LCS(m,n) = LCS(m-1, n-1) + 1$

Else
 $LCS(m,n) = \max(LCS(m-1,n), LCS(m,n-1))$

Going back

		0	1	2	3	4	5	6
	y_j	B	D	C	A	B	A	
0	x_i	0	↑ 0	↑ 0	↑ 0	0	0	0
1	A	0	0	0	0	↑ 1	← 1	↖ 1
2	B	0	↑ 1	↑ 1	1	1	↑ 2	← 2
3	C	0	↑ 1	↑ 1	↑ 2	↑ 2	↑ 2	↑ 2
4	B	0	↑ 1	↑ 1	↑ 2	↑ 2	↑ 3	← 3
5	D	0	↑ 1	↑ 2	↑ 2	2	↑ 3	↖ 3
6	A	0	↑ 1	↑ 2	↑ 2	↑ 3	3	↑ 4
7	B	0	1	2	2	3	4	4

If $m=0$ or $n=0$
 $LCS(m,n) = 0$

Else if $X[m] = Y[n]$
 $LCS(m,n) = LCS(m-1, n-1) + 1$

Else
 $LCS(m,n) = \max(LCS(m-1,n), LCS(m, n-1))$

Going back

		0	1	2	3	4	5	6
	y_j	B	D	C	A	B	A	
0	x_i	0	↑ 0	↑ 0	↑ 0	0	0	0
1	A	0	↖ 0	0	0	↑ 1	← 1	↖ 1
2	B	0	↑ 1	← 1	← 1	1	↑ 2	← 2
3	C	0	↑ 1	↑ 1	↑ 2	↑ 2	↑ 2	↑ 2
4	B	0	↑ 1	↑ 1	↑ 2	↑ 2	↖ 3	↑ 3
5	D	0	↑ 1	↑ 2	↑ 2	2	↑ 3	↑ 3
6	A	0	↑ 1	↑ 2	↑ 2	↑ 3	↑ 3	↖ 4
7	B	0	↑ 1	↑ 2	↑ 2	3	↑ 4	↑ 4

If $m=0$ or $n=0$
 $LCS(m,n) = 0$

Else if $X[m] = Y[n]$
 $LCS(m,n) = LCS(m-1, n-1) + 1$

Else
 $LCS(m,n) = \max(LCS(m-1,n), LCS(m,n-1))$

Going back

		0	1	2	3	4	5	6
	Y_j	B	D	C	A	B	A	
0	X_i	0	↑ 0	↑ 0	↑ 0	0	0	0
1	A	0	0	0	0	↑ 1	← 1	1
2	B	0	↑ 1	↑ 1	1	1	↑ 2	← 2
3	C	0	↑ 1	↑ 1	↑ 2	2	2	2
4	B	0	↑ 1	↑ 1	↑ 2	↑ 2	↑ 3	← 3
5	D	0	↑ 1	↑ 2	↑ 2	2	↑ 3	↑ 3
6	A	0	↑ 1	↑ 2	↑ 2	↑ 3	3	↑ 4
7	B	0	1	2	2	3	4	4

If $m=0$ or $n=0$
 $LCS(m,n) = 0$

Else if $X[m] = Y[n]$
 $LCS(m,n) = LCS(m-1, n-1) + 1$

Else
 $LCS(m,n) = \max(LCS(m-1,n), LCS(m,n-1))$

Going back

		0	1	2	3	4	5	6
	Y_j	B	D	C	A	B	A	
0	X_i	0	↑ 0	↑ 0	↑ 0	0	0	0
1	A	0	0	0	0	↑ 1	← 1	1
2	B	0	↑ 1	↑ 1	1	1	↑ 2	← 2
3	C	0	↑ 1	↑ 1	2	↑ 2	2	2
4	B	0	↑ 1	↑ 1	2	↑ 2	3	↑ 3
5	D	0	↑ 1	↑ 2	2	2	3	↑ 3
6	A	0	↑ 1	↑ 2	2	↑ 3	3	4
7	B	0	1	2	2	3	4	4

If $m=0$ or $n=0$
 $LCS(m,n) = 0$

Else if $X[m] = Y[n]$
 $LCS(m,n) = LCS(m-1, n-1) + 1$

Else
 $LCS(m,n) = \max(LCS(m-1,n), LCS(m, n-1))$

Going back

		0	1	2	3	4	5	6
	y_j	B	D	C	A	B	A	
0	x_i	0	↑ 0	↑ 0	↑ 0	0	0	0
1	A	0	0	0	0	↑ 1	← 1	1
2	B	0	↑ 1	← 1	1	1	↑ 2	← 2
3	C	0	↑ 1	↑ 1	← 2	2	2	2
4	B	0	↑ 1	↑ 1	↑ 2	↑ 2	← 3	3
5	D	0	↑ 1	↑ 2	↑ 2	2	↑ 3	3
6	A	0	↑ 1	↑ 2	↑ 2	↑ 3	3	4
7	B	0	1	2	2	3	4	4

If $m=0$ or $n=0$
 $LCS(m,n) = 0$

Else if $X[m] = Y[n]$
 $LCS(m,n) = LCS(m-1, n-1) + 1$

Else
 $LCS(m,n) = \max(LCS(m-1,n), LCS(m, n-1))$

Going back

		0	1	2	3	4	5	6
	y_j		B	D	C	A	B	A
0	x_i	0	↑ 0	↑ 0	↑ 0	0	0	0
1	A	0	0	0	0	↑ 1	← 1	↖ 1
2	B	0	↑ 1	← 1	1	1	↑ 2	↑ 2
3	C	0	↑ 1	↑ 1	↖ 2	↑ 2	↑ 2	↑ 2
4	B	0	↑ 1	↑ 1	↑ 2	↑ 2	↖ 3	↑ 3
5	D	0	↑ 1	↑ 2	↑ 2	2	↑ 3	↑ 3
6	A	0	↑ 1	↑ 2	↑ 2	↑ 3	↑ 3	↖ 4
7	B	0	1	2	2	3	4	4

If $m=0$ or $n=0$
 $LCS(m,n) = 0$

Else if $X[m] = Y[n]$
 $LCS(m,n) = LCS(m-1, n-1) + 1$

Else
 $LCS(m,n) = \max(LCS(m-1,n), LCS(m, n-1))$

Going back

		0	1	2	3	4	5	6
	y_j		B	D	C	A	B	A
0	x_i	0	0	0	0	0	0	0
1	A	0	0	0	0	1	1	1
2	B	0	1	1	1	1	2	2
3	C	0	1	1	2	2	2	2
4	B	0	1	1	2	2	3	3
5	D	0	1	2	2	2	3	3
6	A	0	1	2	2	3	3	4
7	B	0	1	2	2	3	4	4

If $m=0$ or $n=0$
 $LCS(m,n) = 0$

Else if $X[m] = Y[n]$
 $LCS(m,n) = LCS(m-1, n-1) + 1$

Else
 $LCS(m,n) = \max(LCS(m-1,n), LCS(m, n-1))$

Find LCS String

Function LCS-Generate-String($X[1..m]$, $Y[1..n]$):

Initialization

for $i \leftarrow 0$ to m **do**

$L[i][0] \leftarrow 0$ /* Initialization */

end

for $j \leftarrow 0$ to n **do**

$L[0][j] \leftarrow 0$ /* Initialization */

end

Find LCS String

Function LCS-Generate-String($X[1\dots m]$, $Y[1\dots n]$):

Matching Characters

```
for i ← 1 to m do
  for j ← 1 to n do
    /* If current characters match, add 1 to LCS */
    if  $X[i] = Y[j]$  then
       $L[i][j] \leftarrow L[i-1][j-1] + 1$ 
       $D[i][j] = 1$ ;
    end
    else
      /* If they don't match, take the maximum of the LCS */
      if  $L[i-1][j] \geq L[i][j-1]$  then
         $L[i][j] \leftarrow L[i-1][j]$ 
         $D[i][j] = 2$ ;
      end
      else
         $L[i][j] \leftarrow L[i][j-1]$ 
         $D[i][j] = 3$ ;
      end
    end
  end
end
```


THANK YOU