CS390 Computational Game Theory and Mechanism Design July 10, 2013

Problem Set 4

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Problem 1 We'll first prove that $\forall \varepsilon, \sigma =$

Problem 2 (Collaborated with YYY and ZZZ.)

The root is player 6, and all the other players are his children and the terminated nodes. For $\forall i, i \neq 6$, suppose player 6 choose the strategy $\sigma = aH + bT$, $0 \leq a, b \leq 1$. Then player i choose the strategy $\sigma' = a'H + b'T$, $0 \leq a', b' \leq 1$. We have

$$u_i = (a-b)(a'-b')$$

Suppose that a > b, to maximize u_i , we'll make a' = 1, b' = 0. Similarly when b > a, a' = 0, b' = 1.

And for $a=b=\frac{1}{2},\,a',b'$ will take arbitrary value. Therefore,

$$T(aH+bT,H)=1\text{a ; b}$$

$$T(aH+bT,T)=1\text{a ; b}$$

$$T(aH+bT,a'H+b'T)=1\text{a = b, }0\leq a',b'\leq 1$$

$$T(X,Y)=0\text{otherwise}$$

For player 6, when $a \neq b$, it's obvious that σ is not the best response to his children's choices.

So we'll only consider $a = b = \frac{1}{2}$ to find a $z_1, z_2 \dots z_5$ that σ being the best response to them.

Clearly, if the opponents of player 6 has higher probability to choose Head or Tail. He'll choose the other one, so he can get more utility, then a pure strategy is his best response. So only when all his opponents have the same probability to choose either sides of the coin, σ is the best choice for him.

Therefore, $T(\frac{1}{2}H + \frac{1}{2}T) = 1$. The witness list is all $(z_1, z_2, z_2, z_3, z_4, z_5)$, sum of probability that all the children choose Head is equal to sum of probability that all the children choose Tail.

Then we'll find the Nash Equilibrium top-down.

$$\sigma_i = a_i H + b_i T, 0 \le a_i, b_i \le 1, i \ne 6$$

$$\sigma_6 = \frac{1}{2} H + \frac{1}{2} T$$

$$\Sigma_{i=1}^5 a_i = \Sigma_{i=1}^5 b_i$$

Problem 3 (No collaborator.)

Your solution goes here.

References

- [1] M. J. Osborne and A. Rubinstein. A course in game theory. MIT Press, 1994.
- [2] N. Nisan, T. Roughgarden, E. Tardos, and V. Vazirani (eds). *Algorithmic game theory*. Cambridge University Press, 2007. (Available at http://www.cambridge.org/journals/nisan/downloads/Nisan_Non-printable.pdf.)