

Data Modeling and Databases: Homework 3

Sokolov Maxim
email: somal1996@rambler.ru

Innopolis University

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1 Task 1

I translated the ER model to following relation data model:

- Treatment (PID, code dosage, indications)
- Medicine (code, price)
- Patient (PID, name, address, code)
- Doctor (EID, speciality, E_address, salary, E_name)
- Trainee (EID, speciality, E_address, salary, E_name)
- Permanent (EID, speciality, E_address, salary, E_name)
- Visiting (EID, speciality, E_address, salary, E_name)
- Rooms (Room_id, period)
- Nurse (EID, shift, E_address, salary, E_name, Room_id)
- Employee (EID, E_address, salary, E_name)
- Receptionist (EID, E_address, salary, E_name)
- Record (RecordN, appointment, Patient_id, EID)
- attends (PID, EID)
- assigned (PID, Room_id)

2 Task 2

- a) $(\sigma_{Room_id=107}(Nurse) \bowtie Employee) \cup ((\sigma_{Room_id=107}(assigned) \bowtie attends) \bowtie Employee)$
- b) $Nurse - (((\sigma_{E_name="Dr.Alex"}(Doctor)) \bowtie attends) \bowtie assigned) \bowtie Nurse)$
- c) $(Employee \bowtie_{salary > docs.salary} \rho_{docs}(Employee \bowtie Doctors)) \bowtie Employee$
Second variant:
 $\sigma_{salary > \min(Doctor.salary)}(Employee)$
- d) $assigned \bowtie Rooms$

3 Task 3

- a) $\sigma_{E2.salary = \max(E2.salary)} \rho_{E2}(\sigma_{E1.salary < \max(E1.salary)} \rho_{E1}(\sigma_{Employees.salary < \max(Employees.salary)}(Employees)))$
Result: $\{(3, Alex, \$60,000)\}$
- b) $((\sigma_{range \leq distance}(\pi_{distance}(\sigma_{flight\# = 100}(Flights)) \bowtie Aircraft) \bowtie Certified) \bowtie Employees)$
Result: $\{(3, Alex, \$60,000), (4, Sam, \$40,000)\}$
- c) $((((\sigma_{salary > 70000}(Employees)) \bowtie Certified) \bowtie Aircraft) \bowtie_{distance \leq range} Flights) \bowtie Flights)$
Result: $\{(112, Moscow, Kazan, 600), (111, Kazan, Moscow, 600), (300, Kazan, St. Petersburg, 1000)\}$

4 Task 4

$$0 \leq |R \bowtie S| \leq m * n$$

Minimal number of tuples in $|R \bowtie S|$ is 0. It will be if R and S have several common domains, but there are not have intersected tuples. Maximal number of tuples is $m*n$. It will be if two relations have common domains which contains single number.

5 Task 5

$$\pi_N(R - S) \neq \pi_N(R) - \pi_N(S).$$

Example:

$R = \{(1,1), (2,2), (3,3)\}$, where $R(N,A)$.

$S = \{(1,1), (2,3), (4,3)\}$, where $S(N,A)$.

Then $\pi_N(R - S) = \pi_N(\{(2,2), (3,3)\}) = \{(2), (3)\}$,

but $\pi_N(R) - \pi_N(S) = \{(1), (2), (3)\} - \{(1), (2), (4)\} = \{(3)\}$.

6 Task 6

- a) In this exercise we can see amount of domain: $R \bowtie S$ (A, B, C, D) have 4 domains, but $(\sigma_{R.C=S.C}(R \times S)(R.A, R.B, R.C, S.C, S.D))$ have 5 domains. So these relations are not equivalent.
- b) For solving this problem I modelled all actions on relations (include tuples): join and projection. Let $R = \{(a_1, b_1, c_1), \dots, (a_n, b_n, c_n)\}$, where $|R| = n$; $S = \{(c'_1, d_1), \dots, (c'_n, d_n)\}$, where $|S| = m$.
We know that R and S have one intersection by domain "C". And "C" will define $R \bowtie S$. So we can divide all domain "C" in R and in S on two part: c_i with pair in S.C domain and without. And so we can change order of tuples in relations using following condition: $c_i = c'_i$, where $1 \leq i \leq k \leq \min(m, n)$. So if $i \leq k$ than c_i have pair, if $i \geq k$ than have not pair. It's main think. Now we will model all actions using our manipulation:

- $R \bowtie S = \{(a_i, b_i, c_i, d_i)\}$, where $1 \leq i \leq k$
 $\pi_c(R \bowtie S) = \{(c_i)\}$, where $1 \leq i \leq k$
- $\pi_c(R) = \{(c_i)\}$, where $1 \leq i \leq n$
 $\pi_c(S) = \{(c'_j)\}$, where $1 \leq j \leq m$
 $\pi_c(R) \bowtie \pi_c(S) = \{(c_i)\}$, where $1 \leq i \leq k$
- In the end, we have following:
 $\pi_c(R \bowtie S) = \{(c_i)\} = \pi_c(R) \bowtie \pi_c(S)$