

Introduction to Apache ZooKeeper™



http://zookeeper.apache.org/

Who am I?



- Saurav Haloi
- Engineer at Symantec
- Work in Hadoop & Distributed System
- FOSS enthusiast

What is a Distributed System?



A distributed system consists of multiple computers that communicate through a computer network and interact with each other to achieve a common goal.

- Wikipedia

Fallacies of Distributed Computing



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- The network is reliable.
- Latency is zero.
- Bandwidth is infinite.
- The network is secure.
- Topology doesn't change.
- There is one administrator.
- Transport cost is zero.
- The network is homogeneous.

Reference: http://en.wikipedia.org/wiki/Fallacies_of_Distributed_Computing

Coordination in a distributed system



- Coordination: An act that multiple nodes must perform together.
- Examples:
 - Group membership
 - Locking
 - Publisher/Subscriber
 - Leader Election
 - Synchronization
- Getting node coordination correct is very hard!





Introducing ZooKeeper



ZooKeeper allows distributed processes to coordinate with each other through a shared hierarchical name space of data registers.

- ZooKeeper Wiki

ZooKeeper is much more than a distributed lock server!

What is ZooKeeper?



- An open source, high-performance coordination service for distributed applications.
- Exposes common services in simple interface:
 - naming
 - configuration management
 - locks & synchronization
 - group services

... developers don't have to write them from scratch

Build your own on it for specific needs.

ZooKeeper Use Cases

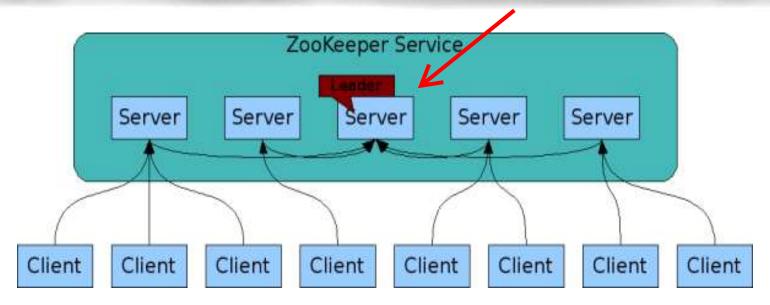


- Configuration Management
 - Cluster member nodes bootstrapping configuration from a centralized source in unattended way
 - Easier, simpler deployment/provisioning
- Distributed Cluster Management
 - Node join / leave
 - Node statuses in real time
- Naming service e.g. DNS
- Distributed synchronization locks, barriers, queues
- Leader election in a distributed system.
- Centralized and highly reliable (simple) data registry

The ZooKeeper Service

Leader





- ZooKeeper Service is replicated over a set of machines
- All machines store a copy of the data (in memory)
- A leader is elected on service startup
- Clients only connect to a single ZooKeeper server & maintains a TCP connection.
- Client can read from any Zookeeper server, writes go through the leader & needs majority consensus.

Image: https://cwiki.apache.org/confluence/display/ZOOKEEPER/ProjectDescription

The ZooKeeper Data Model



- ZooKeeper has a hierarchal name space.
- Each node in the namespace is called as a ZNode.
- Every ZNode has data (given as byte[]) and can optionally have children.

```
parent: "foo"
|-- child1: "bar"
|-- child2: "spam"

`-- child3: "eggs"

`-- grandchild1: "42"
```

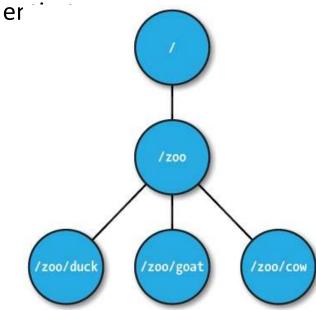
ZNode paths:

- canonical, absolute, slash-separated
- no relative references.
- names can have Unicode characters

ZNodes



- Maintain a stat structure with version numbers for data changes, ACL changes and timestamps.
- Version numbers increases with changes
- Data is read and written in its



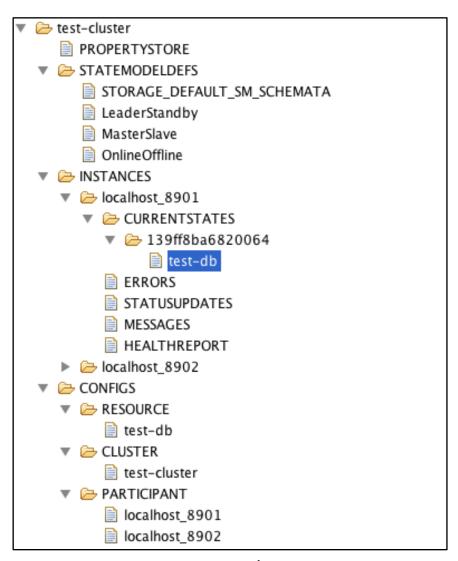


Image: http://helix.incubator.apache.org/Architecture.html

ZNode Types



- Persistent Nodes
 - exists till explicitly deleted
- Ephemeral Nodes
 - exists as long as the session is active
 - can't have children
- Sequence Nodes (Unique Naming)
 - append a monotonically increasing counter to the end of path
 - applies to both persistent & ephemeral nodes

ZNode Operations



Operation	Туре
create	Write
delete	Write
exists	Read
getChildren	Read
getData	Read
setData	Write
getACL	Read
setACL	Write
sync	Read

ZNodes are the main entity that a programmer access.

ZooKeeper Shell



```
[zk: localhost:2181(CONNECTED) 0] help
                                             [zk: localhost:2181(CONNECTED) 1] Is /
ZooKeeper -server host:port cmd args
                                             [hbase, zookeeper]
   connect host:port
                                              [zk: localhost:2181(CONNECTED) 2] ls2 /zookeeper
    get path [watch]
    Is path [watch]
                                              [quota]
    set path data [version]
                                              cZxid = 0x0
    rmr path
                                              ctime = Tue Jan 01 05:30:00 IST 2013
   delquota [-n|-b] path
                                              mZxid = 0x0
    quit
                                              mtime = Tue Jan 01 05:30:00 IST 2013
    printwatches on off
                                              pZxid = 0x0
    create [-s] [-e] path data acl
                                              cversion = -1
    stat path [watch]
                                              dataVersion = 0
                                              aclVersion = 0
   close
                                              ephemeralOwner = 0x0
   Is2 path [watch]
                                              dataLength = 0
    history
                                              numChildren = 1
    listquota path
    setAcl path acl
   getAcl path
                                              [zk: localhost:2181(CONNECTED) 3] create /test-znode HelloWorld
                                             Created /test-znode
    sync path
    redo cmdno
                                              [zk: localhost:2181(CONNECTED) 4] Is /
    addauth scheme auth
                                              [test-znode, hbase, zookeeper]
    delete path [version]
                                              [zk: localhost:2181(CONNECTED) 5] get /test-znode
   setquota -n | -b val path
                                             HelloWorld
```

ZNode Watches



- Clients can set watches on znodes:
 - NodeChildrenChanged
 - NodeCreated
 - NodeDataChanged
 - NodeDeleted
- Changes to a znode trigger the watch and ZooKeeper sends the client a notification.
- Watches are one time triggers.
- Watches are always ordered.
- Client sees watched event before new znode data.
- Client should handle cases of latency between getting the event and sending a new request to get a watch.

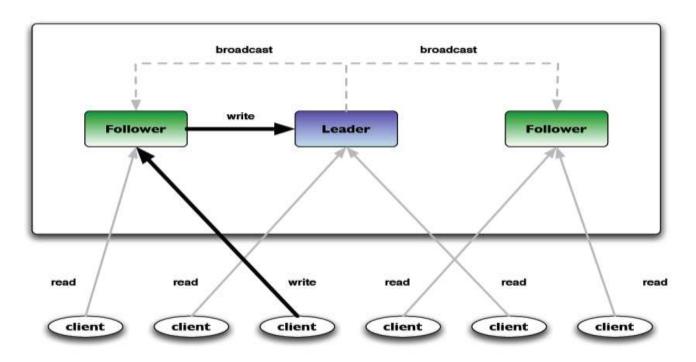
API Synchronicity



- API methods are sync as well as async
- Sync: exists("/test-cluster/CONFIGS", null);
- Async:

ZNode Reads & Writes





- Read requests are processed locally at the ZooKeeper server to which the client is currently connected
- Write requests are forwarded to the leader and go through majority consensus before a response is generated.

Image: http://www.slideshare.net/scottleber/apache-zookeeper

Consistency Guarantees



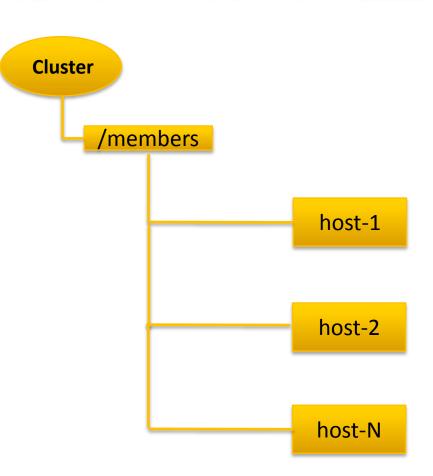
- Sequential Consistency: Updates are applied in order
- Atomicity: Updates either succeed or fail
- Single System Image: A client sees the same view of the service regardless of the ZK server it connects to.
- Reliability: Updates persists once applied, till overwritten by some clients.
- Timeliness: The clients' view of the system is guaranteed to be up-to-date within a certain time bound. (Eventual Consistency)

Recipe #1: Cluster Management



Each Client Host i, i:=1 .. N

- 1. Watch on /members
- Create /members/host-\${i} as ephemeral nodes
- 3. Node Join/Leave generates alert
- Keep updating /members/host-\${i}
 periodically for node status
 changes
 (load, memory, CPU etc.)



Recipe #2: Leader Election



- 1. A znode, say "/svc/election-path"
- 2. All participants of the election process create an ephemeral-sequential node on the same election path.
- 3. The node with the smallest sequence number is the leader.
- 4. Each "follower" node listens to the node with the next lower seq. number
- 5. Upon leader removal go to election-path and find a new leader, or become the leader if it has the lowest sequence number.
- Upon session expiration check the election state and go to election if needed

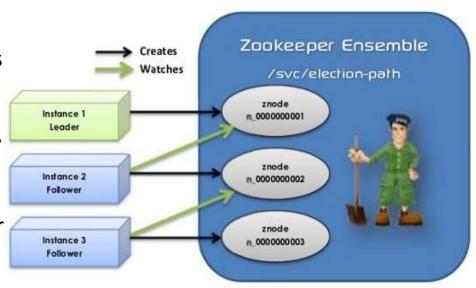


Image: http://techblog.outbrain.com/2011/07/leader-election-with-zookeeper/

Recipe #3: Distributed Exclusive Lock



Assuming there are N clients trying to acquire a lock

- Clients creates an ephemeral, sequential znode under the path /Cluster/_locknode_
- Clients requests a list of children for the lock znode (i.e. _locknode_)
- The client with the least ID according to natural ordering will hold the lock.
- Other clients sets watches on the znode with id immediately preceding its own id
- Periodically checks for the lock in case of notification.
- The client wishing to release a lock deletes the node, which triggering the next client in line to acquire the lock.

```
ZK
---Cluster
  +---config
  +---memberships
  +--- locknode
    +---host1-3278451
    +---host2-3278452
    +---host3-3278453
    \---hostN-3278XXX
```

Language Bindings



- ZooKeeper ships client libraries in:
 - Java
 - Θ C
 - Perl
 - Python
- Community contributed client bindings available for Scala, C#, Node.js, Ruby, Erlang, Go, Haskell

https://cwiki.apache.org/ZOOKEEPER/zkclientbindings.html

A few points to remember



- Watches are one time triggers
 - Continuous watching on znodes requires reset of watches after every events / triggers
- Too many watches on a single znode creates the "herd effect" causing bursts of traffic and limiting scalability
- If a znode changes multiple times between getting the event and setting the watch again, carefully handle it!
- Keep session time-outs long enough to handle long garbage-collection pauses in applications.
- Set Java max heap size correctly to avoid swapping.
- Dedicated disk for ZooKeeper transaction log

Who uses ZooKeeper?



Companies:

- Yahoo!
- Zynga
- Rackspace
- LinkedIn
- Netflix
- and many more...

Projects in FOSS:

- Apache Map/Reduce (Yarn)
- Apache HBase
- Apache Solr
- Neo4j
- Katta
- and many more...

Reference: https://cwiki.apache.org/confluence/display/ZOOKEEPER/PoweredBy

ZooKeeper In Action @Twitter



- Used within Twitter for service discovery
- How?
 - Services register themselves in ZooKeeper
 - Clients query the production cluster for service "A" in data center "XYZ"
 - An up-to-date host list for each service is maintained
 - Whenever new capacity is added the client will automatically be aware
 - Also, enables load balancing across all servers.



Reference: http://engineering.twitter.com/

References



- The Chubby lock service for loosely-coupled distributed systems Google Research (7th USENIX Symposium on Operating Systems Design and Implementation (OSDI), {USENIX} (2006))
- ZooKeeper: Wait-free coordination for Internet-scale systems Yahoo Research (USENIX Annual Technology Conference 2010)
- Apache ZooKeeper Home: http://zookeeper.apache.org/
- Presentations:
 - http://www.slideshare.net/mumrah/introduction-to-zookeeper-trihug-may-22-2012
 - http://www.slideshare.net/scottleber/apache-zookeeper
 - https://cwiki.apache.org/confluence/display/ZOOKEEPER/ZooKeeperPresentations

Interesting Reads



- The Google File System
- The Hadoop Distributed File System
- MapReduce: Simplified Data Processing on Large Clusters
- Bigtable: A Distributed Storage System for Structured Data
- PNUTS: Yahoo!'s Hosted Data Serving Platform
- Dynamo: Amazon's Highly Available Key-value Store
- Spanner: Google's Globally Distributed Database
- Centrifuge: Integrated Lease Management and Partitioning Cloud Services (Microsoft)
- ZAB: A simple totally ordered broadcast protocol (Yahoo!)
- Paxos Made Simple by Leslie Lamport.
- Eventually Consistent by Werner Vogel (CTO, Amazon)
- http://www.highscalability.com/



Questions?



Thank You!

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