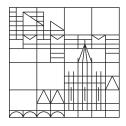
## Minimization of Stochastic Dynamical Systems

A Comparison of Models & Methods

## M.Sc. Project Report

by Fabian Klopfer

Universität Konstanz



Faculty of Sciences

Department of Computer and Information Science

Modelling of Complex Self-Organizing Systems Group

**Supervisor:** Dr. Stefano Tognazzi

**Professor:** Prof. Dr. Tatjana Petrov

Konstanz, 2020

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## 1. Introduction

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# A. Software Requirement Specification & Design Documentation

#### A.1. Introduction

#### A.1.1. Purpose

The purpose of this Software Requirements Specification is to describe the features, constraints and demands of a tool for model reduction and minimization of stochastic dynamical systems in detail. This document is intended for both the supervisors and the developers of the tool.

#### A.1.2. Scope

This software system shall implement a tool for model reduction and minimization of stochastic dynamical systems, that provides support for differential equations and automata, equivalent benchmark examples for each, at least one algorithm for the reduction of each — namely [29] and [10]. An optional feature is the implementation of a conversion algorithm between automata and differential equations, when theoretically possible.

#### A.1.3. Definitions, Acronyms, Abbreviations

**Differential equation** (DE) An equation relating a function to its derivative with respect to one or more independent variable.

**Weighted automata** (WA) An automaton consisting of states and a transition function which associates a weight to each transition

**Possible DE extensions** Extensions of the differential equation framework. For details on the specifics, see section 2.2.3

**Isomorphism** A structure-preserving, revertable, i.e. bijective morphism or mapping from a domain to a codomain

#### A.1.4. Overview

In the second chapter several conditions, assumptions and circumstances will be mentioned, that help charachterizing the software's special use case. In the thrid chapter the concrete requirements are listed.

#### A.2. Overall description

#### A.2.1. Product Perspective

The software system shall be the practical and evaluative basis for the underlying theory, described in the previous report. Especially it shall provide an empirical basis for the analysis of equivalence of reduction techniques for DEs and the minimization of WAs, i.e. if reducing a DE using an algorithm, does one arrive at the same complexity as when minimising a weighted automaton representing the same system as the DE. For that it shall offer scalable benchmark examples, that is an example DE with an equivalent WA. Further it shall provide a base for implementing further algorithms or new methods to be developed (e.g. if an isomorphism is found between DEs and WAs a conversion tool).

#### A.2.2. Product Functions

The system shall take a DE or a WA as input and apply one selected algorithm onto those and then output the minimized model. Further it shall be able to compare the results of a minimisation in terms of model complexity and equivalence, i.e. answer the question weather DE reduction and WA minimisation yield different representations of the same model or if one model is more complex.

#### A.2.3. User Characteristics

For now the only user is going to be the developer in the course of the master project and thesis. When successful enough this may or may not change. This implies that maintainability and extensibility is important, if it's ever handed-over to another developer or even extended to be usable by non-developers.

#### A.2.4. Constraints

- Run time is crucial for these algorithms to be feasible. The system shall not impose significant overhead and the implementations shall be as fast as possible
- Memory safety is important as inputs may induce infinite memory requirements that must not execute if so (else the system will crash)
- Correctness is the most important aspect of the implementation as the results will be unusable
  if minor errors are introduced that corrupt the reduction or minimization processes as well
  as comparisons and benchmarks

#### A.2.5. Assumptions and Dependencies

to be set when dependencies are fixed

#### A.3. Requirement Specification

#### A.3.1. Software System requirements

The system shall be implemented with focus on the following traits:

- S 1. Correctness
- S 2. Extensibility
- S 3. Maintainability
- S 4. Lightweightness
- S 5. As fast as possible
- S 6. Memory safe for all inputs
- S 7. Cross-compileable.

#### A.3.2. External Requirements

The system shall:

E 1. to be set when dependencies are fixed

#### A.3.3. Functional requirements

#### F 1. Control flow system

The system shall

- 1. provide a mechanism for applying the methods depending on the specified model, i.e. apply WA validation checks, algorithm and evaluation for WA input and the respective equivalent for DEs.
- 2. provide a file based input and output mechanism.
- 3. provide a user interface based input and output mechanism.
- 4. validate the inputs before algorithm execution

#### F 2. User interface

The system shall

- 1. offer 3 different methods to be applied: WA minimization, DE reduction and Benchmarking.
- 2. provide a way to specify which input method is to be used, either file-based or user interface-based.
- 3. provide a way to specify which output method is to be used.
- 4. provide the possibility to select an algorithm from a set of available ones for both models.
- 5. display the input, the output and summaries and statistics provided by the model specific modules

#### F 3. WA minimization

The system shall

- 1. specify an input format for WA.
- 2. provide a validation machanism to ensure the well-formedness of the input.
- 3. implement the algorithm specified in [29].
- 4. summarize the results and statistics on the difference between input and output.

#### F 4. DE reduction

The system shall

1. specify an input format for arbitrary DEs and possible extensions, specifically stochastic

differential, functional differential, integro-differential and differential-algebraic equations.

- 2. provide a validation machanism to ensure the well-formedness of the expression.
- 3. implement the algorithm specified in [10]
- 4. summarize the results and statistics on the difference between inputs and outputs

#### F 5. Benchmark The system shall

- 1. provide example WA and DEs, which are
  - 1. theoretically proven to be equivalent, see section 4.1
  - 2. scalable by a parameter, defining the number of states or the number of variables
- 2. implement a method to compare the results in terms of
  - 1. run time,
  - 2. memory
  - 3. resulting model complexity

#### F 6. Conversion or Equivalence The system shall

- 1. be able to convert the benchmark examples from one model to the other or to prove that instances of distinct models are equivalent
- 2. (Optional) be able to convert all instances of one model to the other model

#### A.4. Proposed Software Architecture

#### A.4.1. Overview

**Main Module** Entry point of the process, handles for the control flow.

**User Interface** Lets the user control the main module.

**Model Interface** Defines a model including format.

**Reduction Method Interface** Defines a reduction method for a specific model.

**Conversion/Equivalence Interface** Defines the conversion or equivalence notion from a specific model to a specific model.

**Benchmark Module** Defines benchmark examples and generation processes for those.

#### A.4.2. Subsystem decomposition

**Main Module** The main module is the main entry point of the process, responsible for executing the user interface, the algorithms and utilities from the model specific modules, the conversions, as well as dealing with benchmark flows and loading the necessary examples. There are two mandatory basic control flows and one optional one that depends on the theoretical feasibility that is to be proven:

- 1. Standard model reduction for a single model with user specified input
- 2. Benchmark flow using the provided examples, resulting in two sequential model reduction executions. They can not be parallelized, as they would infere with the caches and slow each other down. After the execution it converts one instance to the other model preserving model complexity or checks if they are equivalent.

**User Interface** The User Interface let's the user control the main module, provides means of input to the main module and displays the output from the main module. Also contains usage information.

A text-based user interface is to be implemented (also called command-line interface). The user interface must contain a menu letting the user specify what control flow is to be executed, thus what task the user wants to perform. It must be able to either read the data from file or from the standard input of the terminal and display the input, output and summary on the standard output of the terminal, aswell as possible errors.

**Model Interface** The model interface specifies the details of a model including the format, input validation functionality, concrete instances will hold an array of function pointers of the reduction method interface type for the concrete model. Depending on the task that the user wants to perform either one or another instance is used to execute the control flow. The data flow from the user via the main module is passed to the model interface instance which does the actual processing using the defined functions and one of the functions from the array holding the model reduction method pointers. Instances of this interface must also specify how to summarize and compute statistics over the result, i.e. the differences in the model and in model complexity. To be implemented are instances for DEs and WAs.

**Reduction Method Interface** This interface is used to implement concrete reduction methods for a specific model. For each models there may be multiple reduction methods. To be implemented are the reduction methods specified in [10, 29].

**Conversion/Equivalence Interface** Using the model interfaces, implementations of this interface shall provide a method to convert an instance of a model to another model or to prove the

equivalence of two instances of different models. As of now this is a question to be further investigated in terms of theory. One open question is if there exists an isomorphism between the to be implemented methods and if so how to do the conversion or equivalence check without also transforming the model complexity (e.g. from DEs to WAs, how not to produce more or less states than there are implicitly in the DE representation). Two options are viable for implementation structurally:

- 1. Attach the conversions to the model interface as function pointer holding one function for each convertable model
- 2. Use a struct holding function pointers for equivalence/conversion in both directions

#### A.4.3. Interfaces

The AccessWebService interface provides three functions: getAccessToken, accessValid and access. getAccessToken is used to obtain access tokens. accessValid checks, if the provided access token is valid. This function is used by the MSS to authenticate users. The function access is used to execute operations on the database. In before a authentification check is made and it is tested if the user is allowed to execute the given operation.

```
Interface AccessWebService {
public string accesstoken getAccessToken(string username, string password) {
    require username.length > 0 && password.length > 0;
    ensure database.getUser(username).accesstoken == accesstoken &&
    database.getUser(username).expirationDate;
    ensure containsLogEntry(user, system.getCurrentDate());
}
public boolean valid accessValid(string username, string accesstoken) {
    require database.existsEntry(username);
    ensure valid == database.getUser(username).isIsValidToken(accesstoken);
}
public Response r access(string username,
                            string accesstoken, Operation op) {
    require database.entryExists(username);
    ensure iff accessValid(username, accesstoken) then
    iff database.getUser(username).isLocked() then
        operation == new PasswordChangeRequiredResponse();
    else
        iff database.getUser().getPrivilegeLvl() > getPrivilegeLvl(operation) then
        operation is executed and r is the response of the operation;
        r == new AccessDeniedResponse();
    ensure if accesstoken is not valid operation is not executed and r is null;
}
}
```

#### A.4.4. Data and Control Flow

#### A.4.5. Boundary conditions