

# Minimization of Weighted Automata

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# Introduction

In the last presentation . . .

- two models for stochastic dynamical systems were considered:  
Weighted Automata (WA) and Differential Equations (DE)
- an example for modelling a CRN's dynamics in both models was given:
  - DE Solving Chemical Master Equation
  - WA Monte Carlo CTMC

## Goals specified

1. Implement minimization algorithm for weighted automata [1]. ✓
2. Implement model reduction algorithm for ODEs [2].
3. Develop reproducible benchmarks
4. Write report

## What has been done so far

- Software Requirement Specification & Software Design Document
- Random Basis Minimal WA Construction Algorithm by Kiefer/Schützenberger [1]
- Execution of example by Matlab script and hand
- Implementation of minimization & equivalence algorithm, interfaces, TUI, CLI, tests

# The Weighted Automaton Minimization Algorithm [1] I

Weighted Automaton  $A = (n, \Sigma, \alpha, \mu, \eta)$ , where

- $n$  the number of states
- $\Sigma$  the input alphabet
- $\alpha$  the initial vector with a non-zero value for all starting states
- $\mu$  the set of transition matrices, one per input character
- $\eta$  the final vector with non-zero values for all ending states

## The Weighted Automaton Minimization Algorithm [1] II

Reduction to Schwarz-Zippel Lemma (Polynomial Identity Testing) provide bounds on correctness  $\frac{n}{K}$

Author claims complexity is in  $\mathcal{NC}$ , this is not correct (c.f. later)

In the following slides use forward reduction, the backwards reduction is similar

## The Weighted Automaton Minimization Algorithm [1] III

- Find a basis  $F$  of the prefix space using random vectors  $r_i$ 
  - Add the vectors of all prefix words up to length  $n$  together and multiply this vector by  $n$  different factors yielding  $\{v_1, \dots, v_n\}$
  - Factors are derived by random vectors and structure of prefixes
  - Base is then the maximally linear independent subset of  $\{\alpha, v_1, \dots, v_n\}$
- Use basis to do Schützenberger Construction [3]:  $\vec{A} = (\vec{n}, \Sigma, \vec{\alpha}, \vec{\mu}, \vec{\eta})$  With
  - $\vec{\mu} = \vec{F} \mu \vec{F}^{-1}$  or  $\vec{F} \mu = \vec{\mu} \vec{F}$
  - $\vec{\alpha} = e_1$
  - $\vec{\eta} = \vec{F} \eta$
  - $\vec{n} = \text{rank}(\vec{\mu})$

# The Weighted Automaton Minimization Algorithm: Pseudo Code I

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**Algorithm 1:** minimize(WeightedAutomaton WA)

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**Input:** A weighted automata WA

**Parameters:** K setting the maximal random number

**Output:** A minimal version of WA

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**begin**

    List<Matrix> randVs;

    WeightedAutomaton minWA;

    randVs  $\leftarrow$  generate\_random\_vectors(WA, K);

    minWA  $\leftarrow$  forward\_reduction(WA, randVs);

    randVs  $\leftarrow$  generate\_random\_vectors(minWA, K);

    minWA  $\leftarrow$  backward\_reduction(minWA, randVs);

**return** minWA;

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Array `randVs` has size  $A \cdot n$  and matrices  $r_i \in \{1, \dots, K \cdot n\}^{\Sigma \times n}$



# The Weighted Automaton Minimization Algorithm: Pseudo Code II

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**Algorithm 2:** forward\_reduction(WA, randVs)

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**Input:** A weighted automata WA, random vectors randVs

**Output:** WA transformed by a random minimal forward space base

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**begin**

List<Vector> rhoVectors  $\leftarrow$  calculate\_rho\_vectors(WA, randVs);

Matrix  $\vec{F} \leftarrow \text{vstack}(\mathbf{A}.\alpha, \text{rhoVectors}[0], \dots, \text{rhoVectors}[n - 1]);$

int  $\vec{n} \leftarrow \text{rank}(\vec{F});$

$\vec{F} \leftarrow \vec{F}[:, \vec{n} - 1, :];$

RowVector  $\vec{\alpha} \leftarrow \text{standard\_basis}(\vec{n})[0, :];$

Vector  $\vec{\eta} \leftarrow \vec{F} \cdot \eta;$

**foreach**  $\mu_i \in \mathbf{WA}.\mu$  **do**

└  $\vec{\mu}_i \leftarrow \text{Solver}(\vec{F}^T).\text{solve}((\vec{F} \cdot \mu_i)^T)^T;$

**return**  $(\vec{n}, \mathbf{WA}.\Sigma, \vec{\alpha}, \vec{\mu}, \vec{\eta});$

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# The Weighted Automaton Minimization Algorithm: Pseudo Code III

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**Algorithm 3:** calculate\_rho\_forward\_vectors(WA, randVs)

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**Input:** A weighted automata WA, random vectors randVs

**Output:** Candidate vectors as base of the Forward space

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**begin**

```
List<(Word, Vector)> words ← generate_words_forward(WA, WA.n);
for i = 0 to randVs.length - 1 do
    for j = 0 to words.length - 1 do
        vi += words[j].vector * get_word_factor(words[j].word, randVs[i]);
return { v0, . . . , vrandVs.length-1 };
```

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**Algorithm 4:** get\_word\_factor(WA, randVs)

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**Input:** Word w, randVs random vector

**Output:** WA transformed by a random minimal forward space base

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**begin**

```
result ← 1;
for i = 0 to do
    result *= randVs(word[i], i);
```

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# The Weighted Automaton Minimization Algorithm: Pseudo Code IV

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**Algorithm 5:** generate\_words\_forward(WA, randVs)

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**Input:** A weighted automata WA, length of words k

**Output:** List of Tuples of word and corresponding vector

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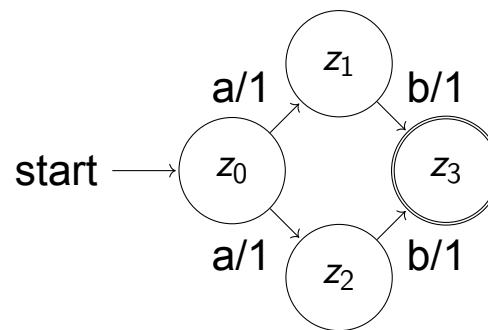
```
begin
  List<Word, Vector> result;
  if k == 1 then
    result = ;
    foreach  $\mu_i \in WA.\mu$  do
      vect  $\leftarrow$  WA. $\alpha \cdot \mu_i$ ;
      if !vect.isZero() then
        result.add(i, vect);
  else
    result = generate_words_forward(WA, k - 1);
    for (word, wVector)  $\in$  result do
      if word.length == k - 1 then
        foreach  $\mu_i \in WA.\mu$  do
          vect  $\leftarrow$  wVector. $\cdot \mu_i$ ;
          if !vect.isZero() then
            newWord  $\leftarrow$  word;
            newWord.append(i);
            result.add(newWord, vect);
  return result;
```

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# The Weighted Automaton Minimization Algorithm: Example I

Weighted automaton  $\mathcal{A} = (n, \Sigma, \mu, \alpha, \eta)$  with

- $n = 3,$
  - $\Sigma = \{a, b\},$
  - $\alpha = (1, 0, 0, 0),$
  - $\eta = (0, 0, 0, 1),$
- $\mu(a) = \begin{pmatrix} 0 & 1 & 1 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix},$
  - $\mu(b) = \begin{pmatrix} 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{pmatrix}.$



## The Weighted Automaton Minimization Algorithm: Example II

$$r^{(1)} = \begin{pmatrix} 9 & 5 & 5 & 7 \\ 6 & 11 & 2 & 1 \end{pmatrix}; \quad r^{(2)} = \begin{pmatrix} 2 & 3 & 1 & 2 \\ 12 & 3 & 9 & 4 \end{pmatrix} \quad r^{(3)} = \begin{pmatrix} 2 & 7 & 9 & 10 \\ 1 & 11 & 2 & 6 \end{pmatrix} \quad r^{(4)} = \begin{pmatrix} 4 & 5 & 2 & 10 \\ 5 & 9 & 5 & 5 \end{pmatrix}$$

$$\alpha\mu(a)r_i + \alpha\mu(a)\mu(b)r_i = (0, 1, 1, 0)r_i + (0, 0, 0, 2)r_i$$

$$v_1 = 9 \cdot (0, 1, 1, 0) + 9 \cdot 11 \cdot (0, 0, 0, 2) = (0, 9, 9, 198)$$

$$v_2 = 2 \cdot (0, 1, 1, 0) + 2 \cdot 3 \cdot (0, 0, 0, 2) = (0, 2, 2, 12)$$

$$v_3 = 2 \cdot (0, 1, 1, 0) + 2 \cdot 11 \cdot (0, 0, 0, 2) = (0, 2, 2, 44)$$

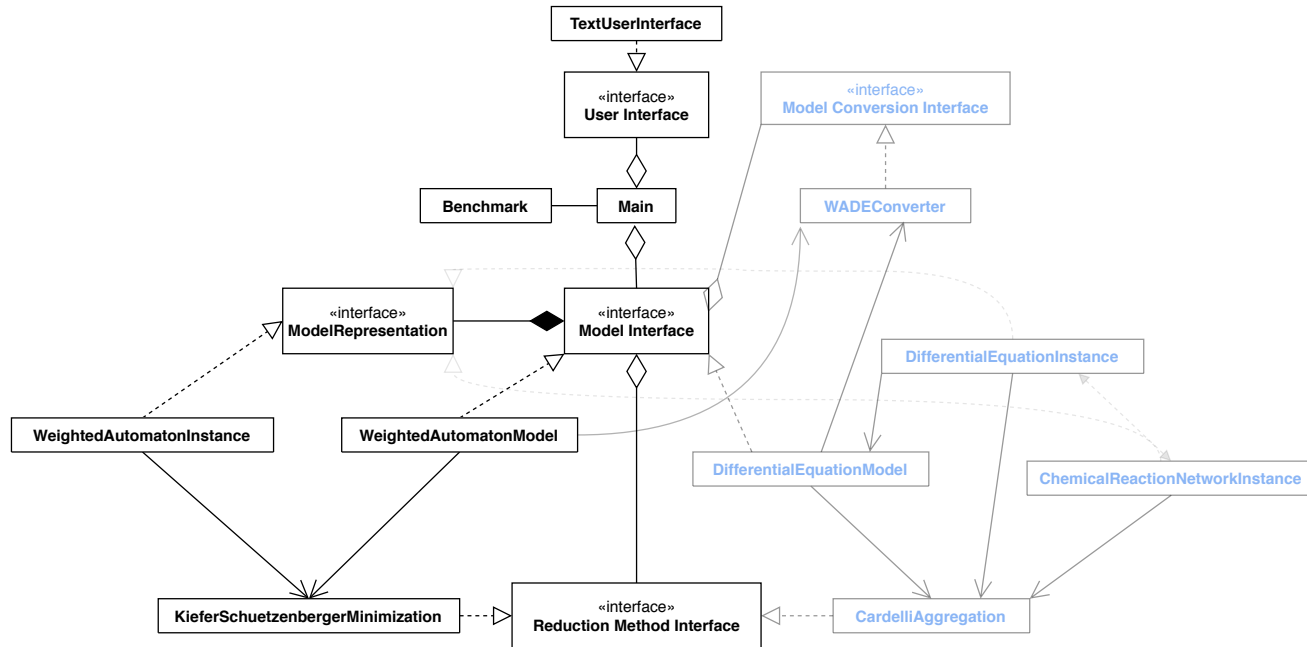
$$v_4 = 4 \cdot (0, 1, 1, 0) + 4 \cdot 9 \cdot (0, 0, 0, 2) = (0, 4, 4, 72)$$

$$\vec{F} = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 9 & 9 & 198 \\ 0 & 2 & 2 & 12 \end{pmatrix}$$

## The Weighted Automaton Minimization Algorithm: Example III

$$\begin{aligned}\vec{F}\mu(\sigma) &= \vec{\mu}(\sigma)\vec{F} \equiv \vec{\mu}(\sigma) = \vec{F}\mu(\sigma)\vec{F}_R^{-1} \\ \vec{\mu}(a) &= \begin{pmatrix} 0 & \frac{-1}{24} & \frac{11}{16} \\ 0 & 0 & 0 \\ 0 & 0 & 0 \end{pmatrix} & \vec{\mu}(b) &= \begin{pmatrix} 0 & 0 & 0 \\ 0 & \frac{1}{8} & \frac{-9}{16} \\ 0 & \frac{1}{36} & \frac{-1}{8} \end{pmatrix} \\ \vec{\eta} = \vec{F}\eta &= \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 9 & 9 & 198 \\ 0 & 2 & 2 & 12 \end{pmatrix} \cdot (0, 0, 0, 1)^T = (0, 198, 12)^T \\ \Rightarrow \vec{\mathcal{A}} &= (3, \{a, b\}, \vec{\mu}, (1, 0, 0), (0, 198, 12)^T)\end{aligned}$$

# Implementation Details I






## Up Next

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# Bibliography

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-  M. P. Schützenberger, “On the definition of a family of automata”, *Information and control*, vol. 4, no. 2-3, pp. 245–270, 1961.