

Prob. Modelling for Computer Scientists
Exercise/Lecture Notes

① Repetition of slides of previous lecture

Notation $\Diamond \varphi$ (Finally / Eventually)

$$\square \sim G(G \circ L, H_y)$$

G is noted by S' to not coherency with Global operator

* is not by * to not compile with future operator

$$S_2 T$$

PROOF

(2) $\Box \Diamond S' = \{ s_0 s_1 s_2 \dots \mid \forall i \geq 0 \exists j \geq i \ s_j \in S' \}$ is uncountable

$\Diamond \Box S' = \{ s_0 s_1 s_2 \dots \mid \exists i \geq 0 \forall j \geq i \ s_j \in S' \}$

$= \bigcup_{i=0}^{\infty} \left(\bigcap_{j=i}^{\infty} \{ s_0 s_1 \dots s_i \dots \mid s_j \in S' \} \right)$ \square

$$C_g((*)^{*}g_j)^{ES}$$

(3) Reachability computation \rightarrow

→ I. Multiply Transition Matrix

II. Recursive computation $\left(\sum P(s|t) \cdot x_t + \sum P(s, y) \right)$

III. Linear eq. solution + pruning

$$\text{Pruning } S = S_1 \cup S_2$$

(4) Constrained reachability $\rightarrow \frac{I_0}{I}$ with pruning the states from S
 $\Pr(\bar{S} \cup S')$

~~I~~ with Transient computations

$$P^n(y,t) \Leftrightarrow \Pr(X_{m+n} = + \mid X_m = s) = P$$

$\hookrightarrow P_r^D(\diamond^{\leq n} S') = P_r^{D[G]}(\diamond^n S') = I$ *something $<$ something*
 \hookrightarrow initial distr.

Ⓢ quiz (ask Matej if you did not know answer)

⑥ try assignment 3 Exercise 1

① Solution for original HWAS2 exercise 1

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Show that two sets of paths are measurable \Leftrightarrow you are able to compute γ .

a) the set of paths in state in state s_{init} and remaining forever in states from A

5, remaining forever in states A and passing through a state B after exactly 3

$$a_1 = \square A = \{ s_0 s_1 s_2 \dots \mid \forall i: s_i \in A \} = \overline{0A} \Rightarrow \cancel{0A} \quad P_r(\overline{0A}) = 1 - P_r(0A)$$



$$\text{Free } \diamond G = \{s_0 s_1 \dots \mid \exists i s_i \in G\} = \bigcup_{k=0}^{\infty} \{ \text{cylinder}(s_0 \dots s_k) \mid s_0 \dots s_k \in G, s_k \in G \}$$

$$\Pr(\Diamond \bar{A}) = \sum_{k=0}^{\infty} \Pr(\text{Cylinder}(s_0 \dots s_k))$$

$\diamond \text{BU} \bigcup_{k=0}^{\infty} \{CH(s_0, \dots, s_k) \mid s_5 \in B, s_0, \dots, s_{k-1} \notin A, s_k \in A\}$

Assignment 2 Exer 2.

~~Example for b)~~

$$a) P_{03}^{(3)} = \Pr(X_3 = s_3 | X_0 = s_0) = P_{03}^3 \quad \text{TRANSIENT}$$

$$= \Pr(\Diamond^3 \{s_3\}) = \Pr(\{Cyl(q_0, q_1, q_2, q_3) | q_0 = s_0, q_3 = s_3\})$$

$$b) \pi_0 = [1/4, 1/4, 1/4, 1/4]$$

$$\Pr(X_3 = s_3 | X_0 \sim \pi_0) = \pi_0 P^{(3)}$$

! In PRISM add dummy states s.t. $\forall s \in S \quad P(\text{dummy}, s) = 1/4 \wedge P_s = \Pr(X_4 = s_4 | X_0 = \text{dummy})$

$$c) \lim_{n \rightarrow \infty} \Pr(X_n = s_3) = 0$$

$$\hat{\pi} P = \hat{\pi}$$

$$P^T \hat{\pi}^T = (\hat{\pi}^T)^T \quad \text{transposed}$$

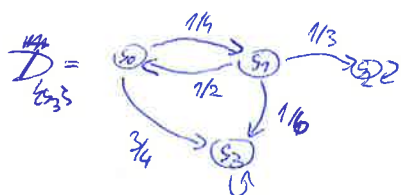
$$d) \Pr(\Diamond^{\leq 3} s_3) = \Pr(\{s_0 s_3, s_0 s_1 s_3, s_0 s_1 s_0 s_3\}) = \frac{3}{4} + \frac{1}{4} \cdot \frac{1}{6} + \frac{1}{4} \cdot \frac{1}{2} \cdot \frac{3}{4} = \frac{85}{96} \approx 0.8854$$

in 1 step in 2 steps in 3 steps

Theorem by Example TRANSFORMATION

Recall $D = \text{ORIGINAL}$

$$\Pr(\Diamond^{\leq 3} \{s_3\}) = \Pr(\Diamond_{D \setminus s_3}^{\leq 3} \{s_3\})$$



$$e) \Pr(\Diamond \{s_3\})$$

$S_0 = \{s_2\}$ states that ^{does not} reach s_3 $X_{s_0} = 0$
 $S_1 = \{s_0, s_1\}$ states that reach s_3 $X_{s_1} = 1$

by complement

$$\Pr(\Diamond \{s_3\}) = \Pr(\Diamond \{s_0, s_1, s_2\}) =$$

$$\sum_{k=0}^{\infty} Cyl(s_0(s_1 s_0)^k s_1 s_2) =$$

$$\sum_{k=0}^{\infty} \frac{1}{4} \cdot \frac{1}{3} \cdot \left(\frac{1}{2} \cdot \frac{1}{4}\right)^k = \frac{1}{12} \cdot \frac{1}{1 - 1/8} = \frac{2}{21}$$

$$\frac{1}{2 \cdot 2^k} \cdot \frac{1}{2} = \frac{1}{2^{k+2}}$$

$$X_{s_0} = \frac{3}{4} + \frac{1}{4} X_{s_1}$$

$$X_{s_1} = \frac{1}{6} + \frac{1}{2} X_{s_0}$$

$$X_{s_1} = \frac{1}{6} + \frac{1}{2} \cdot \left(\frac{3}{4} + \frac{1}{4} X_{s_1}\right)$$

$$X_{s_1} = \frac{1}{6} + \frac{3}{8} + \frac{1}{8} X_{s_1}$$

$$\frac{7}{8} X_{s_1} = \frac{13}{24} \quad X_{s_1} = \frac{13}{21} \quad \text{WHAT?}$$

$$X_{s_0} = \frac{3}{4} + \frac{1}{4} \cdot \frac{13}{21} = \frac{19}{21}$$

$$X_{s_0} = \frac{3}{4} + \frac{1}{4} \cdot \left(\frac{1}{6} + \frac{1}{2} X_{s_0}\right) = \frac{3}{4} + \frac{1}{24} + \frac{1}{8} X_{s_0}$$

$$S_0 = \frac{28}{7} \cdot \left(\frac{3 \cdot 24 + 4}{24 \cdot 4}\right) = \frac{28}{7 \cdot 12} = \frac{19}{21}$$

also for iterative m.

$$\begin{bmatrix} X_{s_0} \\ X_{s_1} \end{bmatrix} = \begin{bmatrix} 0 & 1/4 \\ 1/2 & 0 \end{bmatrix} \begin{bmatrix} X_{s_0} \\ X_{s_1} \end{bmatrix} + \begin{bmatrix} 3/4 \\ 1/6 \end{bmatrix}$$

$$(A - I)x = -b$$

$$\begin{bmatrix} -1 & 1/4 \\ 1/2 & -1 \end{bmatrix} \begin{bmatrix} X_{s_0} \\ X_{s_1} \end{bmatrix} = -\begin{bmatrix} 3/4 \\ 1/6 \end{bmatrix}$$