Overview

- Nondeterminism
- 2 Markov Decision Processes
- Probabilities in MDPs
- Policies
 - Finite-memory policies
- **5** Summary

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Nondeterminism

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http://moves.rwth-aachen.de/teaching/ws-1819/movep18/

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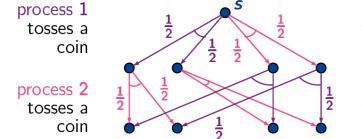
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Nondetermini

Randomness and concurrency

Markov chains are not appropriate for modeling randomized distributed systems, since they cannot adequately model the interleaving behavior of the concurrent processes.



process 2

tosses a coin

process 1 tosses a coin

Nondeterminism

The use of nondeterminism

- Concurrency scheduling of parallel components
 - ▶ in randomised distributed algorithms, several components run partly autonomously and interact asynchronously
- Abstraction
 - ▶ partition state space of a DTMC in similar (but not bisimilar) states
 - replace probabilistic branching by a nondeterministic choice
- Unknown environments
 - interaction with unknown environment
 - example: security in which the environment is an unknown adversary

Beware

Nondeterminism is not the same as a uniform distribution!

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Markov Decision Processe

Markov decision process (MDP)

Markov decision processes

- ▶ In MDPs, both nondeterministic and probabilistic choices coexist.
- ▶ MDPs are transition systems in which in any state a nondeterministic choice between probability distributions exists.
- ▶ Once a probability distribution has been chosen nondeterministically, the next state is selected probabilistically—as in DTMCs.
- ► Any MC is thus an MDP in which in any state the probability distribution is uniquely determined.

Randomized distributed algorithms are typically appropriately modeled by MDPs, as probabilities affect just a small part of the algorithm and nondeterminism is used to model concurrency between processes by means of interleaving.

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Markov Decision Proce

Markov decision process (MDP)

Markov decision process

An MDP \mathcal{M} is a tuple (S, Act, P, ι_{init} , AP, L) where

- ▶ S is a countable set of states with initial distribution $\iota_{\text{init}}: S \to [0,1]$
- ► Act is a finite set of actions
- ▶ $P: S \times Act \times S \rightarrow [0, 1]$, transition probability function such that:

for all
$$s \in S$$
 and $\alpha \in Act$: $\sum_{s' \in S} \mathbf{P}(s, \alpha, s') \in \{0, 1\}$

▶ *AP* is a set of atomic propositions and labeling $L: S \to 2^{AP}$.

Markov decision process (MDP)

Markov decision process

An MDP \mathcal{M} is a tuple $(S, Act, \mathbf{P}, \iota_{\text{init}}, AP, L)$ where

- ▶ S, $\iota_{\text{init}}: S \rightarrow [0,1]$, AP and L are as before, i.e., as for DTMCs, and
- ► Act is a finite set of actions
- ▶ **P** : $S \times Act \times S \rightarrow [0, 1]$, transition probability function such that:

for all $s \in S$ and $\alpha \in Act$: $\sum_{s' \in S} \mathbf{P}(s, \alpha, s') \in \{0, 1\}$

Enabled actions

Let $Act(s) = \{ \alpha \in Act \mid \exists s' \in S. \mathbf{P}(s, \alpha, s') > 0 \}$ be the set of enabled actions in state s. We require $Act(s) \neq \emptyset$ for any state s.

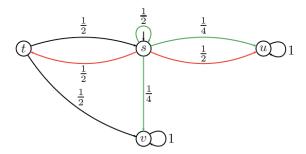
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Markov Decision Processes

An example MDP



- Initial distribution: $\iota_{\text{init}}(s) = 1$ and $\iota_{\text{init}}(t) = \iota_{\text{init}}(u) = \iota_{\text{init}}(u) = 0$
- ▶ Set of enabled actions in state s is $Act(s) = \{ \alpha, \beta \}$ where
 - ▶ $P(s, \alpha, s) = \frac{1}{2}$, $P(s, \alpha, t) = 0$ and $P(s, \alpha, u) = P(s, \alpha, v) = \frac{1}{4}$ ▶ $P(s, \beta, s) = P(s, \beta, v) = 0$, and $P(s, \beta, t) = P(s, \beta, u) = \frac{1}{2}$
- $Act(t) = \{ \alpha \}$ with $P(t, \alpha, s) = P(t, \alpha, u) = \frac{1}{2}$ and 0 otherwise

Markov decision process (MDP)

Markov decision process

An MDP \mathcal{M} is a tuple $(S, Act, \mathbf{P}, \iota_{\text{init}}, AP, L)$ where

- \triangleright S, $\iota_{\text{init}}: S \rightarrow [0, 1]$, AP and L are as before, i.e., as for DTMCs, and
- ► Act is a finite set of actions
- ▶ **P** : $S \times Act \times S \rightarrow [0, 1]$, transition probability function such that:

for all
$$s \in S$$
 and $\alpha \in Act$: $\sum_{s' \in S} \mathbf{P}(s, \alpha, s') \in \{0, 1\}$

If |Act(s)|=1 for any state s, then the nondeterministic choice in any state is over a singleton set. In this case, $\mathcal M$ is a DTMC. Vice versa, a DTMC is an MDP such that |Act(s)|=1 for all s.

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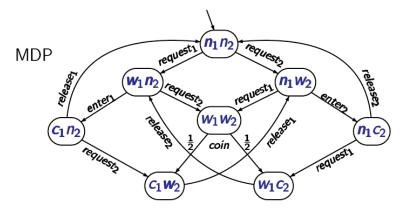
Markov Decision Proce

Example: randomized mutual exclusion

- 2 concurrent processes \mathcal{P}_1 , \mathcal{P}_2 with 3 phases:
 - n_i noncritical actions of process \mathcal{P}_i
 - w_i waiting phase of process \mathcal{P}_i
 - c_i critical section of process \mathcal{P}_i
- competition of both processes are waiting
- resolved by a randomized arbiter who tosses a coin

Randomized mutual exclusion

- interleaving of the request operations
- competition if both processes are waiting
- randomized arbiter tosses a coin if both are waiting



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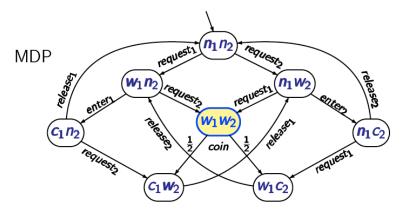
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Markov Decision Processes

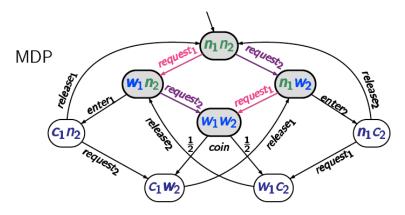
Randomized mutual exclusion

- interleaving of the request operations
- competition if both processes are waiting
- randomized arbiter tosses a coin if both are waiting



Randomized mutual exclusion

- interleaving of the request operations
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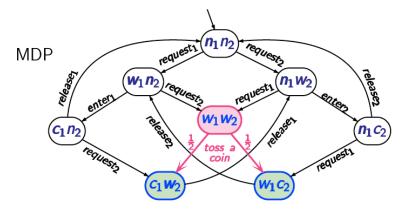
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Markov Decision Proce

Randomized mutual exclusion

- interleaving of the request operations
- competition if both processes are waiting
- randomized arbiter tosses a coin if both are waiting



Intuitive operational behavior

Intuitive operational MDP behavior

- 1. A stochastic experiment according to ι_{init} yields starting state s_0 with probability $\iota_{\text{init}}(s_0) > 0$.
- 2. On entering state s, a non-deterministic choice among Act(s) determines the next action $\alpha \in Act(s)$, say.
- 3. The next state t is randomly chosen with probability $P(s, \alpha, t)$.
- 4. If t is the unique α -successor of s, then almost surely t is the successor after selecting α , i.e., $\mathbf{P}(s, \alpha, t) = 1$.
- 5. Continue with step 2.

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Probabilities in MDPs

Paths in an MDP

State graph

The *state graph* of MDP \mathcal{M} is a digraph G = (V, E) with V are the states of \mathcal{M} , and $(s, s') \in E$ iff $\mathbf{P}(s, \alpha, s') > 0$ for some $\alpha \in Act$.

Paths

An infinite *path* in an MDP $\mathcal{M} = (S, Act, P, \iota_{\text{init}}, AP, L)$ is an infinite sequence $s_0 \alpha_1 s_1 \alpha_2 s_2 \alpha_3 \ldots \in (S \times Act)^{\omega}$, written as

$$\pi = s_0 \xrightarrow{\alpha_1} s_1 \xrightarrow{\alpha_2} s_2 \xrightarrow{\alpha_3} \ldots$$

such that $P(s_i, \alpha_{i+1}, s_{i+1}) > 0$ for all $i \ge 0$. Any finite prefix of π that ends in a state is a *finite path*.

Let $Paths(\mathcal{M})$ denote the set of paths in \mathcal{M} , and $Paths^*(\mathcal{M})$ the set of finite prefixes thereof.

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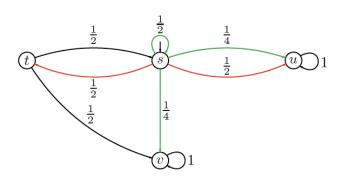
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Probabilities in MDI

Paths in MDPs



$$s \xrightarrow{\alpha} s \xrightarrow{\alpha} s \xrightarrow{\beta} t \xrightarrow{\alpha} s \xrightarrow{\beta} u \dots$$

$$s \xrightarrow{\beta} t \xrightarrow{\alpha} s \xrightarrow{\beta} t \xrightarrow{\alpha} s \dots$$

Probabilities in MDPs

- ▶ For DTMCs, a set of infinite paths is equipped with a σ -algebra and a probability measure that reflects the intuitive notion of probabilities for paths.
- ▶ Due to the presence of nondeterminism, MDPs are not augmented with a unique probability measure.
- Example: suppose we have two coins: a fair one, and a biased one, say $\frac{1}{6}$ for heads and $\frac{5}{6}$ for tails. We select nondeterministically one of the coins, and are interested in the probability of obtaining tails. This, however, is not specified! This also applies if we select one of the two coins repeatedly.
- ▶ Reasoning about probabilities of sets of paths of an MDP relies on the resolution of nondeterminism. This resolution is performed by a policy. A policy chooses in any state s one of the actions $\alpha \in Act(s)$.

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Policies

Policies

Policy

Let $\mathcal{M} = (S, Act, \mathbf{P}, \iota_{\text{init}}, AP, L)$ be an MDP. A *policy* for \mathcal{M} is a function $\mathfrak{S}: S^+ \to Act$ such that $\mathfrak{S}(s_0 s_1 \dots s_n) \in Act(s_n)$ for all $s_0 s_1 \dots s_n \in S^+$.

The path

$$\pi = s_0 \xrightarrow{\alpha_1} s_1 \xrightarrow{\alpha_2} s_2 \xrightarrow{\alpha_3} \dots$$

is called a \mathfrak{S} -path if $\alpha_i = \mathfrak{S}(s_0 \dots s_{i-1})$ for all i > 0.

For any scheduler, the actions are omitted from the *history* $s_0 s_1 \dots s_n$. This is not a restriction as for any sequence $s_0 s_1 \dots s_n$ the relevant actions α_i are given by $\alpha_{i+1} = \mathfrak{S}(s_0 s_1 \dots s_i)$. Hence, the scheduled action sequence can be constructed from prefixes of the path at hand.

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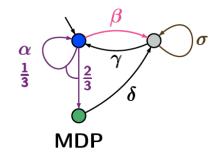
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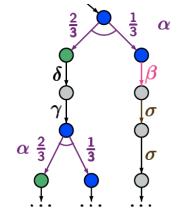
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Policie

Induced Markov chain





Each policy induces an infinite DTMC. States are finite prefixes of paths in the MDP.

¹Also called scheduler, strategy or adversary.

Induced DTMC of an MDP by a policy

DTMC of an MDP induced by a policy

Let $\mathcal{M} = (S, Act, \mathbf{P}, \iota_{\text{init}}, AP, L)$ be an MDP and \mathfrak{S} a policy on \mathcal{M} . The DTMC induced by \mathfrak{S} , denoted $\mathcal{M}_{\mathfrak{S}}$, is given by

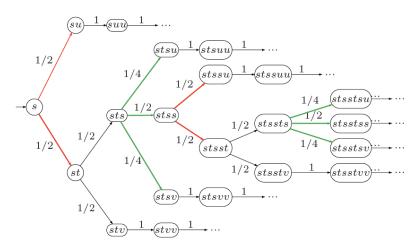
$$\mathcal{M}_{\mathfrak{S}} = (S^+, \mathbf{P}_{\mathfrak{S}}, \iota_{\text{init}}, AP, L')$$

where for $\sigma = s_0 s_1 \dots s_n$: $\mathbf{P}_{\mathfrak{S}}(\sigma, \sigma s_{n+1}) = \mathbf{P}(s_n, \mathfrak{S}(\sigma), s_{n+1})$ and $L'(\sigma) = L(s_n).$

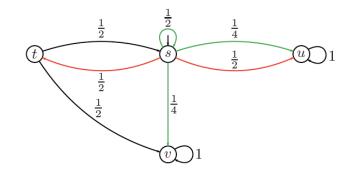
 $\mathcal{M}_{\mathfrak{S}}$ is infinite, even if the MDP \mathcal{M} is finite. Intuitively, state $s_0 s_1 \dots s_n$ of DTMC $\mathcal{M}_{\mathfrak{S}}$ represents the configuration where the MDP \mathcal{M} is in state s_n and $s_0 s_1 \dots s_{n-1}$ stands for the history. Since policy \mathfrak{S} might select different actions for finite paths that end in the same state s, a policy as defined above is also referred to as history-dependent.

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Example induced DTMC



Induced DTMC for a policy that alternates between selecting red and green.



Consider a policy that alternates between selecting red and green, starting with red.

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MDP paths versus paths in the induced DTMC

There is a one-to-one correspondence between the \mathfrak{S} -paths of the MDP \mathcal{M} and the paths in the Markov chain $\mathcal{M}_{\mathfrak{S}}$.

For \mathfrak{S} -path $\pi = s_0 \xrightarrow{\alpha_1} s_1 \xrightarrow{\alpha_2} \dots$, the corresponding path in DTMC $\mathcal{M}_{\mathfrak{S}}$ is:

$$\pi^{\mathfrak{S}} = \widehat{\pi}_0 \, \widehat{\pi}_1 \, \widehat{\pi}_2 \dots$$
 where $\widehat{\pi}_n = s_0 \, s_1 \dots s_n$.

Vice versa, for a path $\widehat{\pi}_0\,\widehat{\pi}_1\,\widehat{\pi}_2\,\dots$ in the DTMC $\mathcal{M}_{\mathfrak{S}},\,\widehat{\pi}_0=s_0$ for some state s_0 such that $\iota_{\text{init}}(s_0) > 0$ and, for each n > 0, $\widehat{\pi}_n = \widehat{\pi}_{n-1} s_n$ for some state s_n in the MDP \mathcal{M} such that $\mathbf{P}(s_{n-1},\mathfrak{S}(\widehat{\pi}_{n-1}),s_n)>0$. Hence:

$$s_0 \xrightarrow{\mathfrak{S}(\widehat{\pi}_0)} s_1 \xrightarrow{\mathfrak{S}(\widehat{\pi}_1)} s_2 \xrightarrow{\mathfrak{S}(\widehat{\pi}_2)} \dots$$

is a \mathfrak{S} -path in \mathcal{M} .

Probability measure on MDP

Probability measure on MDP

Let $Pr_{\mathfrak{S}}^{\mathcal{M}}$, or simply $Pr^{\mathfrak{S}}$, denote the probability measure $Pr^{\mathcal{M}_{\mathfrak{S}}}$ associated with the DTMC $\mathcal{M}_{\mathfrak{S}}$.

This measure is the basis for associating probabilities with events in the MDP \mathcal{M} . Let, e.g., $P \subseteq (2^{AP})^{\omega}$ be an ω -regular property. Then $Pr^{\mathfrak{S}}(P)$ is defined as:

$$Pr^{\mathfrak{S}}(P) = Pr^{\mathcal{M}_{\mathfrak{S}}}(P) = Pr_{\mathcal{M}_{\mathfrak{S}}} \{ \pi \in Paths(\mathcal{M}_{\mathfrak{S}}) \mid trace(\pi) \in P \}.$$

Similarly, for fixed state s of \mathcal{M} , which is considered as the unique starting state,

$$Pr^{\mathfrak{S}}(s \models P) = Pr_{s}^{\mathcal{M}_{\mathfrak{S}}} \{ \pi \in Paths(s) \mid trace(\pi) \in P \}$$

where we identify the paths in $\mathcal{M}_{\mathfrak{S}}$ with the corresponding \mathfrak{S} -paths in \mathcal{M} .

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Finite-memory policies

- Finite-memory policies (shortly: fm-policies) are a generalisation of positional policies.
- ▶ The behavior of an fm-policy is described by a deterministic finite automaton (DFA).
- ightharpoonup The selection of the action to be performed in the MDP $\mathcal M$ depends on the current state of \mathcal{M} (as before) and the current state (called mode) of the policy, i.e., the DFA.

Positional policy

Positional policy

Let \mathcal{M} be an MDP with state space S. Policy \mathfrak{S} on \mathcal{M} is *positional* (or: *memoryless*) iff for each sequence $s_0 s_1 \dots s_n$ and $t_0 t_1 \dots t_m \in S^+$ with $s_n = t_m$:

$$\mathfrak{S}(s_0 s_1 \ldots s_n) = \mathfrak{S}(t_0 t_1 \ldots t_m).$$

In this case, \mathfrak{S} can be viewed as a function $\mathfrak{S}: S \to Act$.

Policy \mathfrak{S} is positional if it always selects the same action in a given state. This choice is independent of what has happened in the history, i.e., which path led to the current state.

Finite-memory policy

Finite-memory policy

Let \mathcal{M} be an MDP with state space S and action set Act.

A finite-memory policy \mathfrak{S} for \mathcal{M} is a tuple $\mathfrak{S} = (Q, act, \Delta, start)$ with:

- Q is a finite set of modes.
- ▶ $\Delta : Q \times S \rightarrow Q$ is the transition function,
- $act: Q \times S \rightarrow Act$ is a function that selects an action $act(q, s) \in Act(s)$ for any mode $q \in Q$ and state $s \in S$ of \mathcal{M} ,
- ightharpoonup start : S o Q is a function that selects a starting mode for state $s \in S$.

An MDP under a finite-memory policy

The behavior of an MDP \mathcal{M} under fm-policy $\mathfrak{S} = (Q, act, \Delta, start)$ is:

- ▶ Initially, a starting state s_0 is randomly determined according to the initial distribution ι_{init} , i.e., $\iota_{\text{init}}(s_0) > 0$.
- ▶ The fm-policy \mathfrak{S} initializes its DFA to the mode $q_0 = start(s_0) \in Q$.
- ▶ If \mathcal{M} is in state s and the current mode of \mathfrak{S} is q, then the decision of \mathfrak{S} , i.e., the selected action, is $\alpha = act(q, s) \in Act(s)$.
- ► The policy changes to mode $\Delta(q, s)$, while \mathcal{M} performs the selected action α and randomly moves to the next state according to the distribution $\mathbf{P}(s, \alpha, \cdot)$.

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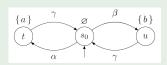
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Policie

Positional versus fm-policies

Positional policies are insufficient for ω -regular properties

Consider the MDP:



Positional policy \mathfrak{S}_{α} always chooses α in state s_0 Positional policy \mathfrak{S}_{β} always chooses β in state s_0 . Then:

$$Pr_{\mathfrak{S}_{\alpha}}(s_0 \models \Diamond a \wedge \Diamond b) = Pr_{\mathfrak{S}_{\alpha}}(s_0 \models \Diamond a \wedge \Diamond b) = 0.$$

Now consider fm-policy $\mathfrak{S}_{\alpha\beta}$ which alternates between selecting α and β . Then: $Pr_{\mathfrak{S}_{\alpha\beta}}(s_0 \models \Diamond a \land \Diamond b) = 1$.

Thus, the class of positional policies is insufficiently powerful to characterise minimal (or maximal) probabilities for ω -regular properties.

Finite-memory policies

Relation fm-policy to definition policy

An fm-policy $\mathfrak{S} = (Q, act, \Delta, start)$ is identified with policy, $\mathfrak{S}' : Paths^* \to Act$ which is defined as follows.

- 1. For the starting state s_0 , let $\mathfrak{S}'(s_0) = act(start(s_0), s_0)$.
- 2. For path fragment $\hat{\pi} = s_0 s_1 \dots s_n$ let

$$\mathfrak{S}'(\widehat{\pi}) = act(q_n, s_n)$$

where $q_0 = start(s_0)$ and $q_{i+1} = \Delta(q_i, s_i)$ for $0 \le i \le n$.

Positional policies can be considered as fm-policies with just a single mode.

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Polici

Other kinds of policies

- ► Counting policies that base their decision on the number of visits to a state, or the length of the history (i.e., number of visits to all states)
- Partial-observation policies that base their decision on the trace $L(s_0) \ldots L(s_n)$ of the history $s_0 \ldots s_n$.
- ▶ Randomised policies. This is applicable to all (deterministic) policies. For instance, a randomised positional policy $\mathfrak{S}: S \to Dist(Act)$, where Dist(X) is the set of probability distributions on X, such that $\mathfrak{S}(s)(\alpha) > 0$ iff $\alpha \in Act(s)$. Similar can be done for fm-policies and history-dependent policies etc..
- ► There is a strict hierarchy of policies, showing their expressiveness (black board).

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Summ

Summary

Important points

- 1. An MDP is a model exhibiting non-determinism and probabilities.
- 2. Non-determinism is important for e.g., randomized distributed algorithms.
- 3. Policies are functions that select enabled actions in states.
- 4. A policy on an MDP induces an infinite DTMC, even if the MDP is finite.
- 5. Probability measures on MDP paths are defined using induced DTMC paths.
- 6. A positional policy selects in a state always the same action.

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