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Probabilistic Modelling for Computer Scientists

Exercise lecture 1 notes by  
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① see photo sorry, I did not know that I should do this

(sigma)

def 2.1.  $\sigma$ -algebra  $(\Omega, \mathcal{F})$  complement

2.  $A \in \mathcal{F} \rightarrow A^c \in \mathcal{F}$

1.  $\Omega \in \mathcal{F}$

3.  $A_1, A_2, A_3, \dots \in \mathcal{F} \rightarrow \bigcup_{i=1}^{\infty} A_i \in \mathcal{F}$

power set function

$2^{\Omega} = \mathcal{P}(\Omega) = \{A \mid A \subseteq \Omega\}$

Ex 2.2  $\Omega = \{1, 2, 3\}$

$2^{\Omega} = \{\emptyset, \{1\}, \{2\}, \{3\}, \{1, 2\}, \{1, 3\}, \{2, 3\}, \{1, 2, 3\}\}$

② Homework Exercise 1 ( $\Omega$  countably infinite,  $\mathcal{F}_{\Omega}$  - smallest class of subsets of  $\Omega$ ,  $\mathcal{F}_{\Omega} \subseteq 2^{\Omega}$ )

(i)  $A$  is finite  $\Rightarrow A \in \mathcal{F}_{\Omega}$

(ii)  $A \in \mathcal{F}_{\Omega} \Rightarrow A^c \in \mathcal{F}_{\Omega}$

a, Solution  $\Omega = \{1, 2, 3, \dots\} = \mathbb{N}$

by given for  $A = \Omega$

! does not relate with (i)

no  $\Omega = \mathbb{N}$

! SOLUTION  $A = \{1, 3, 5, 7, \dots\} \rightarrow$  counterexample that does not need to be in  $\mathcal{F}_{\Omega}$  1

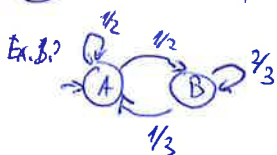
since  $\mathcal{F}_{\Omega}$  should be smallest set it is not  $\Rightarrow$  is not trivial.

b, yes it is trivial  $\wedge \mathcal{F}_{\Omega} = 2^{\Omega}$

c, no because it does not contain 1a counterexample, but  $\sigma$ -algebra does

because  $\mathcal{F}_{\Omega}$  contains each odd number  $\Rightarrow A = \bigcup \{\text{odd number}\} \Rightarrow$  (iii) does not hold for  $\mathcal{F}_{\Omega}$

③ DTMC examples transient



def 3.1.  $\mathcal{D} = (S, P, \mu_{init}, AP, L)$

$(\{A, B\}, \begin{pmatrix} 1/2 & 1/2 \\ 1/3 & 2/3 \end{pmatrix}, (1, 0), AP, L)$

def 3.2.  $Pr(X_1 = s' \mid X_0 = s) = P_{(s, s')}^{(1)}$

$Pr(X_n = s' \mid X_0 = s) = P_{(s, s')}^{(n)}$

def 3.3 Chapman-Kolmogorov relations

$Pr(X_n = s' \mid X_0 = s) = P_{(s, s')}^{(n)} = \sum_{s'' \in S} P(X_n = s', X_i = s'', X_0 = s)$

$= \sum P(X_i = s'', X_0 = s) \cdot P(X_n = s' \mid X_i = s'', X_0 = s) = P(X) \cdot P(Y \mid X)$

$= \sum_{s'' \in S} P_{(s, s'')}^{(i)} \cdot P_{(s'', s')}^{(n-i)}$

def 3.4

$\pi(n)$ : a row vector denoting transient prob distr. at time  $n$

def.  $\pi(0) = \mu_{init}$

$\pi(1) = \mu_{init} \cdot P = \text{EXAMPLE } (1, 0) \begin{pmatrix} 1/2 & 1/2 \\ 1/3 & 2/3 \end{pmatrix} = (1/2 \ 1/2)$

$\pi(2) = \pi(1) \cdot P = \pi(0) \cdot P \cdot P =$

$\pi(0) \cdot P^2$

conditional probability  $P(X \mid Y) = \frac{P(X, Y)}{P(Y)}$   
 $\Rightarrow P(X \cap Y) = P(Y) \cdot P(X \mid Y)$   
↳ memory less prop

def. 3.5 stationary distribution  $\tilde{\pi}$

$$\pi(n+1) = \pi(n) \cdot P = \pi(n) = \tilde{\pi} \quad \text{or} \quad \tilde{\pi} P = \tilde{\pi}$$

EXAMPLE  $[\tilde{\pi}_A \quad \tilde{\pi}_B] \begin{bmatrix} 1/2 & 1/2 \\ 1/3 & 2/3 \end{bmatrix} = [\pi_A \quad \pi_B]$

$$\frac{1}{2}\pi_A + \frac{1}{3}\pi_B = \pi_A \quad \text{and} \quad \pi_A = \frac{2}{3}\pi_B$$

$$\frac{1}{2}\pi_A + \frac{2}{3}\pi_B = \pi_B$$

def 3.6  $\begin{matrix} \in \mathbb{R}^{n \times n} \\ \in \mathbb{R}^{n \times 1} \\ \in \mathbb{R} \end{matrix}$  Eigen value and Eigen vector

$$M N = \lambda N$$

3.7 Each stochastic matrix has a eigenvalue equal 1 and eigen vector of ones =  $\begin{pmatrix} 1 \\ 1 \\ \vdots \\ 1 \end{pmatrix}$

Lemma  $\begin{bmatrix} P \end{bmatrix} \begin{pmatrix} 1 \\ 1 \\ 1 \\ 1 \end{pmatrix} = \begin{pmatrix} 1 \\ 1 \\ 1 \\ 1 \end{pmatrix} \Rightarrow N = \begin{pmatrix} 1 \\ 1 \\ 1 \\ 1 \end{pmatrix} \wedge \lambda = 1$

additional iHf  
Perron-T

theorem

def 3.8  $P$  is stochastic matrix iff each row is a distribution

$$\begin{pmatrix} p_{11} & p_{1n} \\ & p_{nn} \\ & p_{nn} \\ & p_{nn} \end{pmatrix} \rightarrow p_{ij}$$

$$\Rightarrow \sum_{i \in \{1, \dots, n\}} p_{ij} = 1$$

Lemma 3.9  $P^{(n)}$  is stochastic matrix if  $P$  is stoch. matrix