



Theory of Machines and Languages

Fatemeh Deldar

1403-1404

Context-Free Languages

Context-Free Grammars

- A grammar $G = (V, T, S, P)$ is said to be **context-free** if all productions in P have the form

$$A \rightarrow x,$$

where $A \in V$ and $x \in (V \cup T)^*$.

A language L is said to be context-free if and only if there is a context-free grammar G such that $L = L(G)$.

- **Every regular grammar is context-free, so a regular language is also a context-free one**
- **The family of regular languages is a proper subset of the family of context-free languages**

Context-Free Grammars

□ Example

$$\begin{aligned} S &\rightarrow aSa, \\ S &\rightarrow bSb, \\ S &\rightarrow \lambda, \end{aligned}$$

A context-free
grammar



$$L(G) = \{ww^R : w \in \{a,b\}^*\}$$

A context-free
language

□ Example

$$\begin{aligned} S &\rightarrow abB, \\ A &\rightarrow aaBb, \\ B &\rightarrow bbAa, \\ A &\rightarrow \lambda, \end{aligned}$$



$$L(G) = \{ab(bbaa)^n bba(ba)^n : n \geq 0\}$$

Context-Free Grammars

□ Example

$$L = \{a^n b^m : n \neq m\}$$

Grammar

$$S \rightarrow AS_1 | S_1 B,$$

$$S_1 \rightarrow aS_1 b | \lambda,$$

$$A \rightarrow aA | a,$$

$$B \rightarrow bB | b.$$

□ Example

$$S \rightarrow aSb | SS | \lambda.$$

Language

$$L = \{w \in \{a, b\}^* : n_a(w) = n_b(w) \text{ and } n_a(v) \geq n_b(v), \\ \text{where } v \text{ is any prefix of } w\}.$$

Context-Free Grammars

- Find context-free grammars for the following languages

$$L = \{a^n b^m : 2n \leq m \leq 3n\} \Rightarrow S \rightarrow aSbb \mid aSbbb \mid \lambda$$

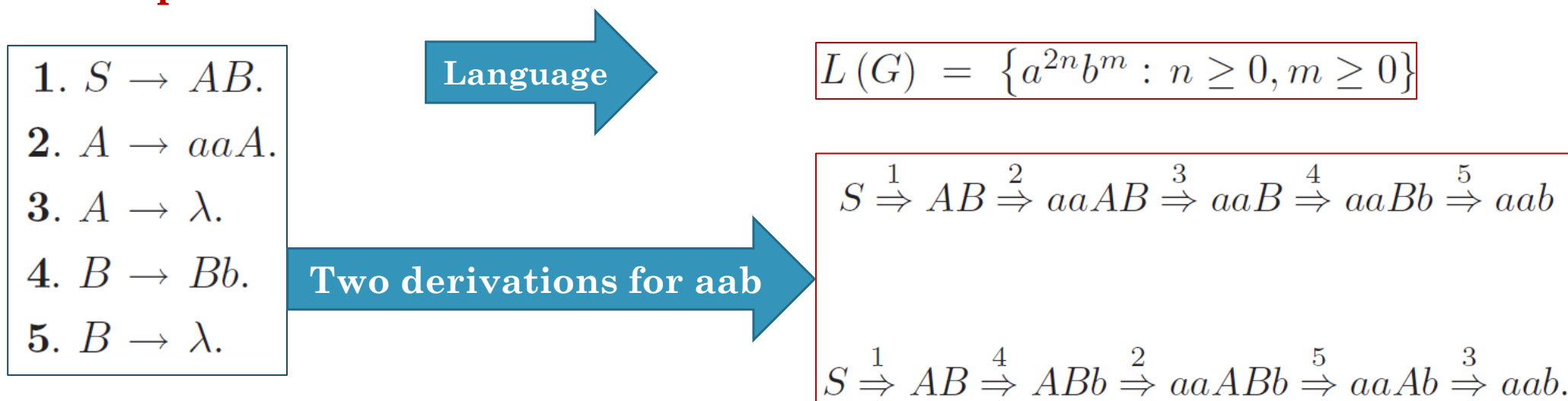
$$L = \{a^n b^m c^k : n = m \text{ or } m \leq k\} \Rightarrow \begin{array}{l} S \rightarrow AC \mid DBC \\ A \rightarrow aAb \mid \lambda \\ B \rightarrow bBc \mid \lambda \\ C \rightarrow cC \mid \lambda \\ D \rightarrow aA \mid \lambda \end{array}$$

$$L = \{uvwv^R : u, v, w \in \{a, b\}^+, |u| = |w| = 2\} \Rightarrow \begin{array}{l} S \rightarrow AB \\ A \rightarrow aa \mid ab \mid ba \mid bb \\ B \rightarrow aBa \mid bBb \mid aAa \mid bAb \end{array}$$

Leftmost and Rightmost Derivations

- In a grammar that is not linear, a derivation may involve sentential forms with more than one variable
 - In such cases, we have a choice in the order in which variables are replaced

□ Example



Leftmost and Rightmost Derivations

□ Leftmost derivation

- In each step the leftmost variable in the sentential form is replaced

□ Rightmost derivation

- In each step the rightmost variable in the sentential form is replaced

□ Example

$$\begin{aligned} S &\rightarrow aAB, \\ A &\rightarrow bBb, \\ B &\rightarrow A|\lambda. \end{aligned}$$

Leftmost derivation of the string *abbbb*

$$S \Rightarrow aAB \Rightarrow abBbB \Rightarrow abAbB \Rightarrow abbBbbB \Rightarrow abbbbB \Rightarrow abbbb$$

$$S \Rightarrow aAB \Rightarrow aA \Rightarrow abBb \Rightarrow abAb \Rightarrow abbBbb \Rightarrow abbbb$$

Rightmost derivation of the string *abbbb*