

Theory of Machines and Languages

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□ A function f with domain D is said to be **Turing-computable** or just **computable** if there exists some Turing machine $M = (Q, \Sigma, \Gamma, \delta, q_0, \square, F)$ such that

$$q_0w \stackrel{*}{\vdash}_M q_f f(w), \qquad q_f \in F,$$

for all $w \in D$.

□ All the common mathematical functions, no matter how complicated, are Turing-computable

Example

 $\delta(q_3, \square) = (q_4, \square, R).$

- For two positive integers x and y, design a Turing machine that computes x + y
 - O Positive integer x is represented by $w(x) \in \{1\}^+$, such that |w(x)| = x
 - We want to design a Turing machine for performing the computation:

$$q_0w(x) 0w(y) \stackrel{*}{\vdash} q_fw(x+y) 0$$

$$\begin{split} \delta\left(q_{0},1\right) &= \left(q_{0},1,R\right),\\ \delta\left(q_{0},0\right) &= \left(q_{1},1,R\right),\\ \delta\left(q_{1},1\right) &= \left(q_{1},1,R\right),\\ \delta\left(q_{1},1\right) &= \left(q_{2},\square,L\right),\\ \delta\left(q_{2},1\right) &= \left(q_{3},0,L\right),\\ \delta\left(q_{3},1\right) &= \left(q_{3},1,L\right), \end{split} \qquad \begin{array}{l} q_{0}111011 \vdash 1q_{0}1011 \vdash 111q_{0}011\\ \vdash 1111q_{1}11 \vdash 111111q_{1}1 \vdash 111111q_{1}1\\ \vdash 11111q_{2}1 \vdash 11111q_{3}10\\ & \vdash q_{3}\square 1111110 \vdash q_{4}111110. \end{split}$$

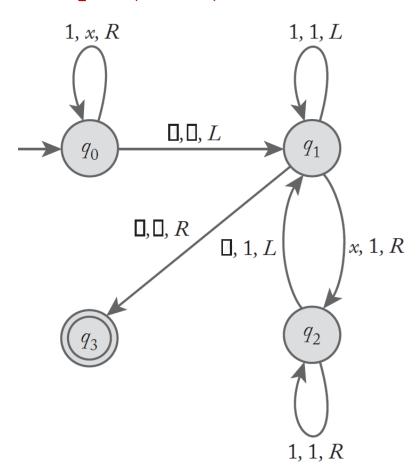
Example

> Design a Turing machine that copies strings of 1's

$$q_0w \stackrel{*}{\vdash} q_fww$$

- 1. Replace every 1 by an x
- 2. Find the rightmost x and replace it with 1
- 3. Travel to the right end of the current nonblank region and create a 1 there
- 4. Repeat Steps 2 and 3 until there are no more x's

□ Example (Cont.)



$$q_0 11 \vdash xq_0 1 \vdash xxq_0 \Box \vdash xq_1 x$$
 $\vdash x1q_2 \Box \vdash xq_1 11 \vdash q_1 x11$
 $\vdash 1q_2 11 \vdash 11q_2 1 \vdash 111q_2 \Box$
 $\vdash 11q_1 11 \vdash 1q_1 111$
 $\vdash q_1 1111 \vdash q_1 \Box 1111 \vdash q_3 1111.$

Turing Machines

■ Exercise: Construct Turing machines that will accept the following language

$$L = \left\{ a^n b^{2n} : n \ge 1 \right\}.$$

