



Theory of Machines and Languages

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Three Basic Concepts

2. Grammars

$$\begin{array}{l} S \rightarrow Ab, \\ A \rightarrow aAb, \\ A \rightarrow \lambda. \end{array} \quad L = \{a^n b^{n+1} : n \geq 0\}$$

$$\begin{array}{l} S \rightarrow SS, \\ S \rightarrow \lambda, \\ S \rightarrow aSb, \\ S \rightarrow bSa \end{array} \quad L = \{w : n_a(w) = n_b(w)\}$$

Three Basic Concepts

2. Grammars

- Two grammars G_1 and G_2 are equivalent if they generate the same language, that is, if $L(G_1) = L(G_2)$

Consider the grammar $G_1 = (\{A, S\}, \{a, b\}, S, P_1)$, with P_1 consisting of the productions

$$\begin{aligned} S &\rightarrow aAb|\lambda, \\ A &\rightarrow aAb|\lambda. \end{aligned}$$



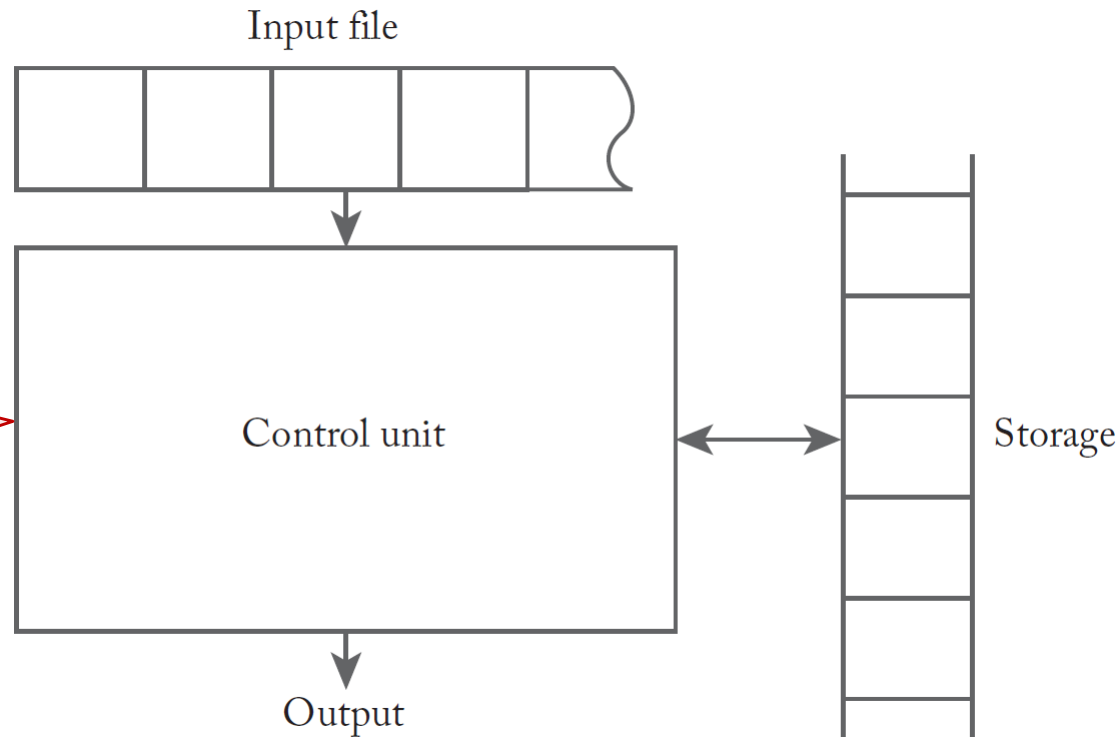
$$L(G_1) = \{a^n b^n : n \geq 0\}.$$

Three Basic Concepts

3. Automata

- An automaton is an abstract model of a digital computer

Control unit can be in any one of a finite number of internal states



Three Basic Concepts

□ Example

Find a grammar that generates the language

$$L = \{ww^R : w \in \{a, b\}^+\}.$$

Find grammars for the following languages on $\Sigma = \{a\}$.

- (a) $L = \{w : |w| \bmod 3 > 0\}.$
- (b) $L = \{w : |w| \bmod 3 = 2\}.$
- (c) $w = \{|w| \bmod 5 = 0\}.$

Three Basic Concepts

□ Example

Let $\Sigma = \{a, b\}$. For each of the following languages, find a grammar that generates it.

(a) $L_1 = \{a^n b^m : n \geq 0, m < n\}$.

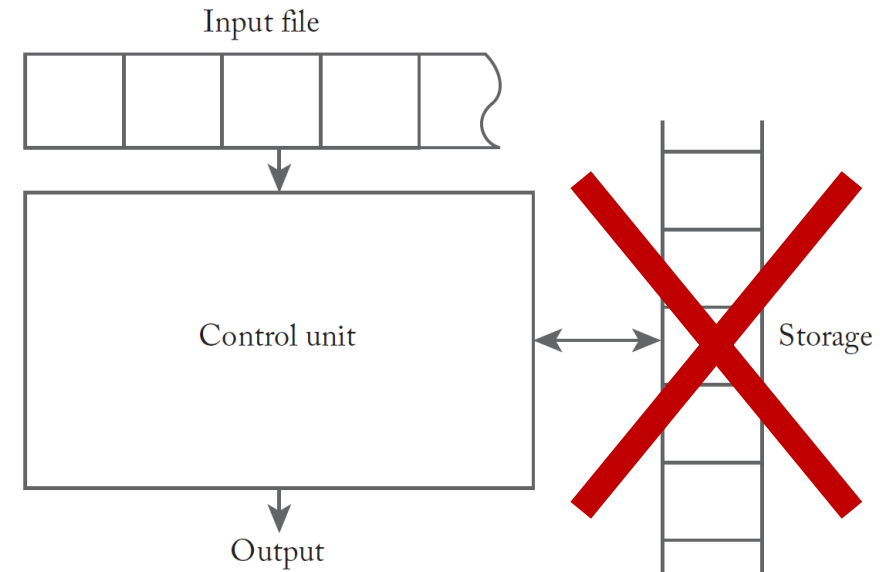
(b) $L_2 = \{a^{3n} b^{2n} : n \geq 2\}$.

Finite Automata

Finite Automata

□ This type of automaton is characterized by:

1. Having no temporary storage
2. An input file cannot be rewritten
3. A finite amount of information can be retained in the control unit by placing the unit into a specific state
 - Since the number of such states is finite, a finite automaton can only deal with situations in which the information to be stored at any time is strictly bounded



Deterministic Finite Acceptor

A **deterministic finite acceptor** or **dfa** is defined by the quintuple

$$M = (Q, \Sigma, \delta, q_0, F) ,$$

where

Q is a finite set of **internal states**,

Σ is a finite set of symbols called the **input alphabet**,

$\delta : Q \times \Sigma \rightarrow Q$ is a total function called the **transition function**,

$q_0 \in Q$ is the **initial state**,

$F \subseteq Q$ is a set of **final states**.

- ❑ The transitions from one internal state to another are governed by the transition function δ
 - **For example**, if $\delta(q_0, a) = q_1$ then if the dfa is in state q_0 and the current input symbol is a , the dfa will go into state q_1

Deterministic Finite Acceptor

- Transition graphs are used to visualize and represent finite automata
 - Vertices represent states
 - Edges represent transitions
 - The labels on the vertices are the names of the states
 - The labels on the edges are the current values of the input symbol

□ Example

$$M = (\{q_0, q_1, q_2\}, \{0, 1\}, \delta, q_0, \{q_1\})$$

$$\begin{array}{ll} \delta(q_0, 0) = q_0, & \delta(q_0, 1) = q_1, \\ \delta(q_1, 0) = q_0, & \delta(q_1, 1) = q_2, \\ \delta(q_2, 0) = q_2, & \delta(q_2, 1) = q_1. \end{array}$$

