



Theory of Machines and Languages

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Turing Machines as Language Accepters

□ Example

For $\Sigma = \{a, b\}$, design a Turing machine that accepts

$$L = \{a^n b^n : n \geq 1\}.$$

$Q = \{q_0, q_1, q_2, q_3, q_4\}$, $F = \{q_4\}$, $\Sigma = \{a, b\}$, $\Gamma = \{a, b, x, y, \square\}$. The transitions can be broken into several parts.

$$\begin{array}{lll} \delta(q_0, a) = (q_1, x, R), & \delta(q_2, y) = (q_2, y, L), & \delta(q_0, y) = (q_3, y, R), \\ \delta(q_1, a) = (q_1, a, R), & \delta(q_2, a) = (q_2, a, L), & \delta(q_3, y) = (q_3, y, R), \\ \delta(q_1, y) = (q_1, y, R), & \delta(q_2, x) = (q_0, x, R). & \delta(q_3, \square) = (q_4, \square, R). \\ \delta(q_1, b) = (q_2, y, L) \end{array}$$

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□ Example (Cont.)

- The particular input $aabb$ gives the following successive instantaneous descriptions:

$$\begin{aligned}
 q_0 a a b b &\vdash x q_1 a b b \vdash x a q_1 b b \vdash x q_2 a y b \\
 &\vdash q_2 x a y b \vdash x q_0 a y b \vdash x x q_1 y b \\
 &\vdash x x y q_1 b \vdash x x q_2 y y \vdash x q_2 x y y \\
 &\vdash x x q_0 y y \vdash x x y q_3 y \vdash x x y y q_3 \square \\
 &\vdash x x y y \square q_4 \square.
 \end{aligned}$$

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□ Example

- Design a Turing machine that accepts $L = \{a^n b^n c^n : n \geq 1\}$

