SS 2015

Solution hints for exercise sheet 1

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Algorithmic Game Theory (IN2239)

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Exercise 2 Transitivity (H)

Let \succeq be a complete preference relation.

Show that \succsim is transitive if and only if \succ and \sim are both transitive.

Solution hints:

- \Rightarrow : Assume \succeq is transitive. Consider arbitrary x,y,z such that $x \sim y \sim z$. Then, $x \succeq y$ as well as $y \succeq z$. Hence by transitivity of \succeq , $x \succeq z$. Moreover, also $z \succeq y$ and $y \succeq x$. Again by transitivity of \succeq , $z \succeq x$. By definition of \sim , it now follows that $x \sim z$, as desired. Now, consider arbitrary x,y,z such that $x \succ y \succ z$. Then again, $x \succeq y \succeq z$ and by transitivity of \succeq , $x \succeq z$. Assume for contradiction that $z \succeq x$. By transitivity of \succeq , we get $z \succeq y$ which contradicts $y \succ z$. Therefore $x \succ z$.
- \Leftarrow : Assume that both \succ and \sim are transitive. Consider arbitrary x, y, z and assume that $x \succsim y \succsim z$. Assume for a contradiction that it is not the case that $x \succsim z$. By completeness, then, $z \succsim x$. Hence, $z \succ x$. We distinguish four cases:
 - i) x > y and y > z,
 - ii) $x \succ y$ and $y \sim z$,
 - iii) $x \sim y$ and $y \succ z$,
 - iv) $x \sim y$ and $y \sim z$.
 - If i), by transitivity of \succ , $x \succ z$ and hence $x \succsim z$, a contradiction. If iv), by transitivity of \sim , $x \sim z$ and hence $x \succsim z$, a contradiction. If ii), both $z \succ x$ and $x \succ y$. Hence, by transitivity of \succ , $z \succ y$, which contradicts $y \sim z$. Finally, if (iii), $y \succ z$ and $z \succ x$, by transitivity of \succ , $y \succ x$, contradicting $x \sim y$.

Exercise 3 Extending rational preferences (H)

Let X, Y be disjoint sets and $R \subseteq X \times X$, $R' \subseteq Y \times Y$ complete and transitive relations on X and Y, respectively. Show that there is a complete and transitive extension of $R \cup R'$ on $X \cup Y$. Solution hints: Consider $R'' = R \cup R' \cup (X \times Y)$. R'' is obviously complete.

For transitivity, consider arbitrary $x, y, z \in X \cup Y$ and, for contradiction, assume that x R'' y R'' z but not x R'' z. Either $z \in X$ or $z \in Y$. If the former, also $x \in X$ and $y \in X$. But then x R y R z and by transitivity of R on X, also x R z. It follows that x R'' z as well, a contradiction.

In the latter case, also $x \in Y$ (because of zR''x). Hence, $y \in Y$ as well. Then, x R' y R' z and by transitivity of R' on Y, moreover, x R' z. By definition of R'', also x R'' z, a contradiction.

Exercise 4 Representing preferences by utility functions (T)

- (a) Show that rational preferences over a countable set A of alternatives can be represented by a utility function $u: A \to [-1, 1]$.
- (b) Let $A = [0, 1] \times [0, 1]$. Define lexicographic preferences \succeq over A such that for all $x, y \in A$,

$$x \succeq y \text{ iff } x_1 > y_1 \text{ or both } x_1 = y_1 \text{ and } x_2 \geqslant y_2.$$

Show that this preference relation cannot be represented by a utility function.

Solution hints: For the solution of this exercise, we will use that

- ullet Q is dense in \mathbb{R} , i.e., for any two reals, there is a rational number between them, and
- the cardinality of \mathbb{R} is strictly greater than the cardinality of \mathbb{Q} .
- (a) In the lecture, the statement was made for a *finite* number of alternatives, we proof the stronger statement for a *countable* number of alternatives. First, enumerate the alternatives a_1, \ldots, a_n . This can be done as A is countable. We will subsequently assign to each a_i a utility $u(a_i)$ in \mathbb{R} , starting with a_1 .

Construction hypothesis At step n, we have that $u(a_1), \ldots, u(a_n)$ are already fixed in a way s.t. for all $i, j \in \{1, \ldots, n\}$

$$u(\alpha_i)\not\in\{-1,1\} \text{ and }$$

$$u(\alpha_i)\geqslant u(\alpha_j) \text{ iff } \alpha_i\succsim\alpha_j.$$

Construction Basis:

$$u(a_1) = 0$$

Construction Step: Assume construction hypothesis holds for n. Need to fix a value for $u(a_{n+1})$.

- If $a_{n+1} \sim a_k$ for some $k \in \{1, \dots, n\}$, set $u(a_{n+1} = u(a_k))$.
- If not, consider

$$\begin{split} A_n^+ = & \{u(\alpha_i) \mid i \in \{1,\dots,n\} \text{ and } \alpha_{n+1} \succsim \alpha_i\} \cup \{-1\} \text{ and } \\ A_n^- = & \{u(\alpha_i) \mid i \in \{1,\dots,n\} \text{ and } \alpha_i \succsim \alpha_{n+1}\} \cup \{1\} \end{split}$$

We claim that A_n^+ and A_n^- form discrete intervals of [-1,1] that do not overlap, i.e., for all $\mathfrak{u}(\mathfrak{a}^+) \in A_n^+$, $\mathfrak{u}(\mathfrak{a}^-) \in A_n^-$ we have $\mathfrak{u}(\mathfrak{a}^+) < \mathfrak{u}(\mathfrak{a}^-)$. If this was not the case, there would be $\mathfrak{u}(\mathfrak{a}^+) \in A_n^+$ and $\mathfrak{u}(\mathfrak{a}^-) \in A_n^-$ with $\mathfrak{u}(\mathfrak{a}^-) < \mathfrak{u}(\mathfrak{a}^+)$. But this implies (by construction hypothesis) that $\mathfrak{a}^+ \succ \mathfrak{a}^-$. Also, by definition of A_n^- , $\mathfrak{a}^- \succ \mathfrak{a}_{n+1}$. By transitivity, we get $\mathfrak{a}^+ \succ \mathfrak{a}_{n+1}$, violating $\mathfrak{a}^+ \in A_n^+$, a contradiction. Given that either $|A_n^+| = 1$ or $|A_n^-| = 1$, the discrete intervals do not overlap by the hypothesis. Thus, the claim holds. Therefore, there is "space" between A^+ and A^- from where we can choose the utility for \mathfrak{a}_{n+1} , e.g.,

$$\mathfrak{u}(\alpha_{n+1}) = \frac{\max A_n^+ + \min A_n^-}{2}.$$

Now, we need to check whether this utility function does actually represent the preferences. To this end, let $b,c\in A$. By the enumeration, there are k and ℓ such that $b=a_k$ and $c=a_\ell$. W.l.o.g k>l. It is straightforward to check that the construction determined $u(a_\ell)$ in a way that $u(a_\ell)\geqslant u(a_k)$ if and only if $a_\ell\succsim a_k$. This concludes the proof.

(b) The high-level idea is as follows: For each value in the first dimension, we need to reserve some space in the reals for a utility representation. Reserving space uncountably many times is problematic because in each such space lies a rational number, entailing a contradiction.

Assume that there is a utility representation of \succeq . Observe that for all $a \in [0,1]$, we have $(a,1) \succ (a,0)$, implying u(a,1) > u(a,0).

As \mathbb{Q} is dense in \mathbb{R} , we can find a $q(a) \in \mathbb{Q} \cap [\mathfrak{u}(a,0),\mathfrak{u}(a,1)]$ for every $a \in [0,1]$. This defines a function $q:[0,1] \to \mathbb{Q}$. If we can show that q is injective, we have a contradiction because the cardinality of [0,1] is strictly larger than the cardinality of \mathbb{Q} . (This function q has no additional meaning, we just use it for the contradiction.)

Assume that q was not injective, i.e., there exist distinct $b, c \in [0, 1]$ such that q(b) = q(c). W.l.o.g., c > b. But

$$\bigg(\mathfrak{u}(\mathfrak{b},0)\leqslant\bigg)q(\mathfrak{b})\leqslant\mathfrak{u}(\mathfrak{b},1)<\mathfrak{u}(\mathfrak{c},0\leqslant q(\mathfrak{c})\bigg(\leqslant\mathfrak{u}(\mathfrak{c},1)\bigg).$$

A contradiction and q is indeed injective. As said before, such a function cannot exist and therefore the initial assumption on the existence of such a utility function was wrong and we proved that no such utility function exists.