

Course Information

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| Programme       | B.Tech. (Computer Science and Engineering) |
| Class, Semester | Final Year B. Tech., Sem VII               |
| Course Code     | 6CS451 C                                   |
| Course Name     | Cryptography and Network Security Lab      |
| PRN             | 22510016                                   |

**Experiment No. 04**

**Title -** Implementation of Chinese Remainder Theorem (CRT)

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**Objectives:**

To understand and implement the Chinese Remainder Theorem (CRT) for solving systems of simultaneous congruences with pairwise coprime moduli.

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**Problem Statement:**

Given a system of simultaneous congruences:

$$x \equiv a_1 \pmod{n_1}$$

$$x \equiv a_2 \pmod{n_2}$$

$$\dots$$

$$x \equiv a_k \pmod{n_k}$$

where the moduli  $n_1, n_2, \dots, n_k$  are pairwise coprime positive integers, find the smallest non-negative integer  $x$  that satisfies all these congruences simultaneously.

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**Equipment/Tools:**

Computer or laptop

Python 3.x

Code editor/IDE

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**Theory:**

The Chinese Remainder Theorem states that if  $n_1, n_2, \dots, n_k$  are pairwise coprime, then the system of congruence's has a unique solution modulo  $N$ , where  $N = n_1 * n_2 * \dots * n_k$ .

Let  $M_i = N/n_i$ , and  $y_i$  be the modular inverse of  $M_i$  modulo  $n_i$ .

Then, the solution is:

$$x = (\sum a_i * M_i * y_i) \pmod{N}$$

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**Procedure:**

Input the arrays of remainders (a) and moduli (n).

Compute N = product of all moduli.

For each congruence, compute  $M_i = N/n_i$ .

Find  $y_i$  = modular inverse of  $M_i \text{ mod } n_i$  using Extended Euclidean Algorithm.

Compute  $x = \sum (a_i * M_i * y_i)$ .

Take  $x \text{ mod } N$  to get the final answer.

Verify the result with all congruences.

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**Steps:** Solve the system:

$$x \equiv 2 \pmod{3}$$

$$x \equiv 3 \pmod{4}$$

$$x \equiv 1 \pmod{5}$$

$$N = 60$$

$$M_1 = 20 \rightarrow y_1 = 2$$

$$M_2 = 15 \rightarrow y_2 = 3$$

$$M_3 = 12 \rightarrow y_3 = 3$$

$$x = (2 \cdot 2 + 3 \cdot 3 + 1 \cdot 3) \pmod{60} = 11$$

Final Answer:  $x = 11$

```
from math import gcd
from functools import reduce

def egcd(a, b):
    if b == 0:
        return (a, 1, 0)
    g, x1, y1 = egcd(b, a % b)
    return(g,y1,x1 - (a//b)*y1)

def modinv(a, m):
    g, x, _ = egcd(a, m)
    if g != 1:
        raise ValueError("Inverse does not exist")
    return x % m

def crt_pairwise_coprime(a, n):
    N=reduce(lambda x,y:x*y, n, 1)
    x = 0
    for ai, ni in zip(a, n):
        Mi= N//ni
        yi = modinv(Mi, ni)
        x+= ai*Mi *yi
    return x % N, N
```

```

print("Example 1:")
print("x ≡ 2 (mod 3), x ≡ 3 (mod 4), x ≡ 1 (mod 5)")
print("Solution:", crt_pairwise_coprime([2, 3, 1], [3, 4, 5]))
print()

print("Example 2:")
print("x ≡ 1 (mod 2), x ≡ 2 (mod 3), x ≡ 3 (mod 5)")
print("Solution:", crt_pairwise_coprime([1, 2, 3], [2, 3, 5]))
print()

print("Example 3:")
print("x ≡ 2 (mod 5), x ≡ 3 (mod 7), x ≡ 2 (mod 9)")
print("Solution:", crt_pairwise_coprime([2, 3, 2], [5, 7, 9]))
print()

print("Example 4:")
print("x ≡ 3 (mod 4), x ≡ 4 (mod 7), x ≡ 2 (mod 9)")
print("Solution:", crt_pairwise_coprime([3, 4, 2], [4, 7, 9]))
print()

print("Example 5:")
print("x ≡ 1 (mod 5), x ≡ 4 (mod 11), x ≡ 6 (mod 17)")
print("Solution:", crt_pairwise_coprime([1, 4, 6], [5, 11, 17]))
print()

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PS C:\Users\Aishw\OneDrive\Documents\Final Year\CNS Practicals\Assignment_4_23520001> python crt.py
Example 1:
x ≡ 2 (mod 3), x ≡ 3 (mod 4), x ≡ 1 (mod 5)
Solution: (11, 60)

Example 2:
x ≡ 1 (mod 2), x ≡ 2 (mod 3), x ≡ 3 (mod 5)
Solution: (23, 30)

Example 3:
x ≡ 2 (mod 5), x ≡ 3 (mod 7), x ≡ 2 (mod 9)
Solution: (227, 315)

Example 4:
x ≡ 3 (mod 4), x ≡ 4 (mod 7), x ≡ 2 (mod 9)
Solution: (11, 252)

Example 5:
x ≡ 1 (mod 5), x ≡ 4 (mod 11), x ≡ 6 (mod 17)
Solution: (686, 935)

```

***Observations:***

The solution is unique modulo N.

The Extended Euclidean Algorithm is efficient for finding modular inverses.

The method is much faster than brute force checking.

***Conclusion:***

The Chinese Remainder Theorem provides a systematic and efficient way to solve simultaneous congruences when moduli are pairwise coprime. The implementation works correctly and gives the smallest non-negative solution.

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