

Course Information

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| Programme | B.Tech. (Computer Science and Engineering) |
| Class, Semester | Final Year B. Tech., Sem VII |
| Course Code | 6CS451 |
| Course Name | Cryptography and Network Security Lab |
| PRN | 22510016 |

Experiment No. 08

Title – Implement the Diffie-Hellman Key Exchange algorithm for a given problem.

Objectives:

To implement the **Diffie-Hellman Key Exchange Algorithm** to enable two parties to securely share a secret key over an insecure communication channel. This shared key can then be used for symmetric encryption or decryption in secure communication.

Problem Statement:

In secure communications, two parties (commonly referred to as Alice and Bob) need to agree on a secret key that can be used for encrypting and decrypting messages. However, they must do this over a public channel where an attacker (Eve) might be listening.

Implement the **Diffie-Hellman Key Exchange Algorithm**, which allows Alice and Bob to securely compute a shared secret key without directly transmitting it over the insecure channel. The algorithm should:

1. Accept a large prime number p and a primitive root modulo p , g .
2. Allow each party to select a private key (a for Alice, b for Bob).
3. Compute the corresponding public keys:
 Alice computes $A = g^a \text{ mod } p$
 Bob computes $B = g^b \text{ mod } p$
4. Exchange public keys between Alice and Bob.
5. Compute the shared secret key:
 Alice computes $K = B^a \text{ mod } p$
 Bob computes $K = A^b \text{ mod } p$
6. Validate that both computed secret keys are equal, i.e., $K_{\text{Alice}} == K_{\text{Bob}}$.
7. Additionally, demonstrate the correctness of the algorithm with an example and optionally simulate an attacker attempting to derive the secret key without access to the private keys.

Equipment/Tools:

- Hardware: Computer System
 - Software: Python
 - Tools: Code Editor – VS Code
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Theory:

The Diffie-Hellman Key Exchange is a public-key cryptographic algorithm that enables two parties to establish a shared secret over an insecure channel. It is based on the difficulty of computing discrete logarithms in modular arithmetic.

Mathematical Steps:

1. Choose a large prime number p and a primitive root modulo p , g .
2. Alice chooses a private key a , and Bob chooses a private key b .
3. They compute their public keys:
 - o Alice: $A = g^a \text{ mod } p$
 - o Bob: $B = g^b \text{ mod } p$
4. Alice and Bob exchange public keys.
5. Each computes the shared secret:
 - o Alice: $K = B^a \text{ mod } p$
 - o Bob: $K = A^b \text{ mod } p$
6. Both results are equal:
$$K_{\text{Alice}} = K_{\text{Bob}} = g^{(ab)} \text{ mod } p$$

This shared key can now be used for symmetric encryption.

Procedure:

1. Input a large **prime number (p)** and a **primitive root (g)**.
 2. Input **private keys** for Alice and Bob (a and b).
 3. Compute **public keys**:
$$A = g^a \text{ mod } p$$

$$B = g^b \text{ mod } p$$
 4. Exchange **public keys**.
 5. Compute **shared secrets**:
$$\text{Alice: } K_{\text{Alice}} = B^a \text{ mod } p$$

$$\text{Bob: } K_{\text{Bob}} = A^b \text{ mod } p$$
 6. Verify that both secret keys are **equal**.
 7. Simulate an attacker (Eve) trying to compute the secret without private keys.
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Steps:

```
def is_prime(n):
    if n < 2:
        return False
    for i in range(2, int(n**0.5) + 1):
        if n % i == 0:
            return False
    return True

def diffie_hellman(p, g, a, b):
    if not is_prime(p):
        print("Error: p must be a prime number.")
        return
    if g <= 1 or g >= p:
        print("Error: g must be between 2 and p-1.")
        return
    if not (1 <= a < p - 1) or not (1 <= b < p - 1):
        print("Error: Private keys (a, b) must be between 1 and p-1.")
        return
    print("\n--- Input is valid ---")
    print(f"Prime number (p): {p}")
    print(f"Primitive root (g): {g}")
    print(f"Alice's private key (a): {a}")
    print(f"Bob's private key (b): {b}")
    A = pow(g, a, p)
    B = pow(g, b, p)
    print("\n--- Public Keys ---")
    print(f"Alice's Public Key (A) = {A}")
    print(f"Bob's Public Key (B) = {B}")

    K_Alice = pow(B, a, p)
    K_Bob = pow(A, b, p)

    print("\n--- Shared Secret Computation ---")
    print(f"Alice's Shared Secret (K_Alice) = {K_Alice}")
    print(f"Bob's Shared Secret (K_Bob) = {K_Bob}")

if K_Alice == K_Bob:
    print("\nShared secret key successfully established!")
    print(f"Final Shared Secret Key = {K_Alice}")
else:
    print("\nKeys do not match. Error in computation!")
print()

print("==== Diffie-Hellman Key Exchange ===")

p = int(input("Enter a large prime number (p): "))
g = int(input("Enter a primitive root modulo p (g): "))
a = int(input("Enter Alice's private key (a): "))
b = int(input("Enter Bob's private key (b): "))

diffie_hellman(p, g, a, b)
```

```

    === Diffie-Hellman Key Exchange ===
Enter a large prime number (p): 23
Enter a primitive root modulo p (g): 5
Enter Alice's private key (a): 6
Enter Bob's private key (b): 15

--- Input is Valid ---
Prime number (p): 23
Primitive root (g): 5
Alice's private key (a): 6
Bob's private key (b): 15

--- Public Keys ---
Alice's Public Key (A) = 8
Bob's Public Key (B) = 19

--- Shared Secret Computation ---
Alice's Shared Secret (K_Alice) = 2
Bob's Shared Secret (K_Bob) = 2

Shared secret key successfully established!
Final Shared Secret Key = 2

```

```

    === Diffie-Hellman Key Exchange ===
Enter a large prime number (p): 29
Enter a primitive root modulo p (g): 2
Enter Alice's private key (a): 19
Enter Bob's private key (b): 13

--- Input is Valid ---
Prime number (p): 29
Primitive root (g): 2
Alice's private key (a): 19
Bob's private key (b): 13

--- Public Keys ---
Alice's Public Key (A) = 26
Bob's Public Key (B) = 14

--- Shared Secret Computation ---
Alice's Shared Secret (K_Alice) = 10
Bob's Shared Secret (K_Bob) = 10

Shared secret key successfully established!
Final Shared Secret Key = 10

```

```

    === Diffie-Hellman Key Exchange ===
Enter a large prime number (p): 21
Enter a primitive root modulo p (g): 5
Enter Alice's private key (a): 6
Enter Bob's private key (b): 15
Error: p must be a prime number.

```

```

    === Diffie-Hellman Key Exchange ===
Enter a large prime number (p): 23
Enter a primitive root modulo p (g): 5
Enter Alice's private key (a): 0
Enter Bob's private key (b): 15
Error: Private keys (a, b) must be between 1 and p-1.

```

Observations and Conclusion:

- **Valid Inputs Produce Matching Shared Keys:**
 - When a prime number p , a valid primitive root g , and private keys a and b within the allowed range are used, both Alice and Bob compute the **same shared secret key**.
 - Example:
 - $p=23, g=5, a=6, b=15 \rightarrow$ Shared key = 2
 - $p=29, g=2, a=19, b=13 \rightarrow$ Shared key = 10
 - This demonstrates the correctness of the Diffie-Hellman key exchange algorithm.
- **Invalid Prime Number:**
 - Inputting a non-prime p (e.g., $p=21$) triggers an **error**: "Error: p must be a prime number."
 - This ensures that the algorithm only works in a proper modular arithmetic group, preserving security.
- **Invalid Private Keys:**
 - Inputting a private key outside the range $1 \leq \text{key} \leq p-1$ (e.g., $a=0$) triggers an **error**: "Error: Private keys (a, b) must be between 1 and $p-1$."
 - This prevents invalid computation and maintains proper key generation.
- **Algorithm Behavior:**
 - The program correctly identifies and handles **wrong inputs**, preventing generation of insecure or meaningless shared keys.
 - Public keys are correctly computed using modular exponentiation.
 - Shared secret computation is consistent for valid inputs.

The Diffie-Hellman Key Exchange algorithm successfully allows two parties to securely generate a shared secret over an insecure channel when inputs are valid.

Input validation is crucial: the prime number p , the primitive root g , and private keys a and b must meet specific criteria to ensure correctness and security.

The program demonstrates that invalid inputs are detected, preventing computation of insecure or incorrect keys.

This implementation provides a secure foundation for further use of the shared key in symmetric encryption and secure communication.
