# ITS202: Algorithms and Data Structure Graphs

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#### Graphs

#### Graph

A graph is a set of vertices and a collection of edges that each connect a pair of vertices.

### Graphs

#### Example of Graph: Facebook



## **Graph Applications**

graph	vertex	edge	
communication	telephone, computer	fiber optic cable	
circuit	gate, register, processor	wire	
mechanical	joint rod, beam, spri		
financial	stock, currency	transactions	
transportation	intersection	street	
internet	class C network	connection	
game	board position	legal move	
social relationship	person	friendship	
neural network	neuron	synapse	
protein network	protein	protein-protein interaction	
molecule	atom	bond	

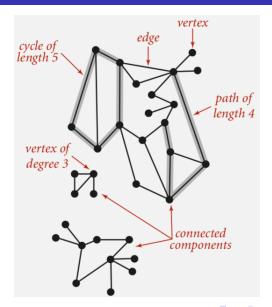
#### Graph Terminology

Path: Sequence of vertices connected by edges.

Cycle: Path whose first and last vertices are the same.

Two vertices are connected if there is a path between them.

## Graph Terminology



# Some graph-processing problems

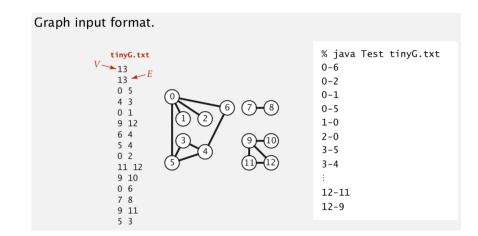
problem	description	
s-t path	Is there a path between s and t?	
shortest s-t path	What is the shortest path between s and t?	
cycle	Is there a cycle in the graph?	
Euler cycle	Is there a cycle that uses each edge exactly once?	
Hamilton cycle	Is there a cycle that uses each vertex exactly once?  Is there a way to connect all of the vertices?	
connectivity		
biconnectivity	Is there a vertex whose removal disconnects the graph?	
planarity	Can the graph be drawn in the plane with no crossing edges?	
graph isomorphism	Do two adjacency lists represent the same graph?	

## Undirected Graph: Graph API

public class	Graph	
	Graph(int V)	create an empty graph with V vertices
	Graph(In in)	create a graph from input stream
void	addEdge(int v, int w)	add an edge v-w
Iterable <integer></integer>	adj(int v)	vertices adjacent to v
int	V()	number of vertices
int	E()	number of edges

# Undirected Graph: Graph API

## Graph API: Sample Client



# Graph API: Sample Client

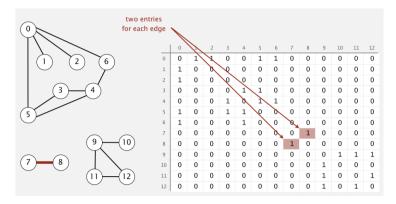
```
// degree of vertex v in graph G
public static int degree(Graph G, int v)
    int degree = 0;
    for (int w : G.adj(v))
       degree++;
    return degree;
```

#### **Graph Representation**

- Adjacency-matrix graph representation
- Adjacency-list graph representation

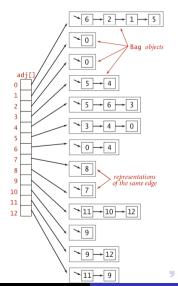
# Graph Representation: Adjacency-matrix graph representation

Maintain a two-dimensional V-by-V boolean array; for each edge v-w in graph: adj[v][w] = adj[w][v] = true.



# Graph Representation: Adjacency-list graph representation

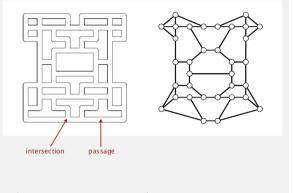
Maintain vertex-indexed array of lists.



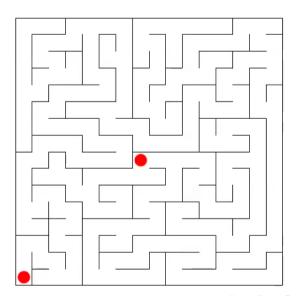
#### Maze Exploration

#### Maze graph.

- Vertex = intersection.
- Edge = passage.



Goal. Explore every intersection in the maze.



Goal. Systematically traverse a graph.

DFS (to visit a vertex v)

Mark v as visited.

Recursively visit all unmarked vertices w adjacent to v.

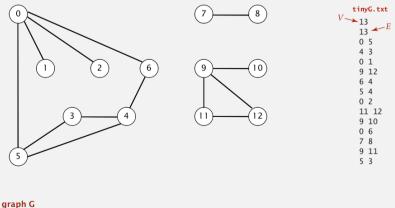
#### Typical applications

- Find all vertices connected to a given source vertex.
- Find a path between two vertices.

## Depth-first Search Demo

#### To visit a vertex v:

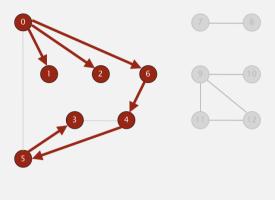
- Mark vertex v as visited.
- Recursively visit all unmarked vertices adjacent to v.



## Depth-first Search Demo

#### To visit a vertex v:

- Mark vertex v as visited.
- Recursively visit all unmarked vertices adjacent to  $\nu$ .



v	marked[]	edgeTo[]
0	Т	_
1	Т	0
2	Т	0
3	Т	5
4	Т	6
5	Т	4
6	Т	0
7	F	-
8	F	-
9	F	-
10	F	-
11	F	-
12	F	-

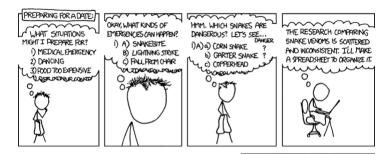
vertices reachable from 0

## Depth-first Search: Data Structure

#### Data structures.

- Boolean array marked[] to mark visited vertices.
- Integer array edgeTo[] to keep track of paths. (edgeTo[w] == v) means that edge v-w taken to visit w for first time .

## Depth-first search application: preparing for a date







I REALLY NEED TO STOP USING DEPTH-FIRST SEARCHES.

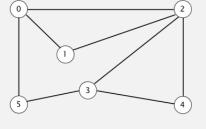
#### Depth-first search application: preparing for a date

Maze Application that we discussed

#### Repeat until queue is empty:



- Remove vertex v from queue.
- Add to queue all unmarked vertices adjacent to v and mark them.

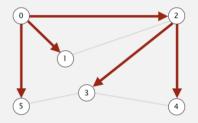




graph G

#### Repeat until queue is empty:

- Remove vertex v from queue.
- Add to queue all unmarked vertices adjacent to v and mark them.



v	edgeTo[]	distTo[]
0	-	0
1	0	1
2	0	1
3	2	2
4	2	2
5	0	1

done

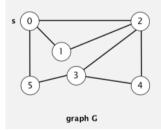
#### Algorithm

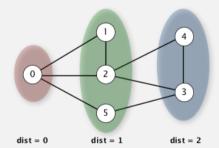
#### BFS (from source vertex s)

Put s onto a FIFO queue, and mark s as visited. Repeat until the queue is empty:

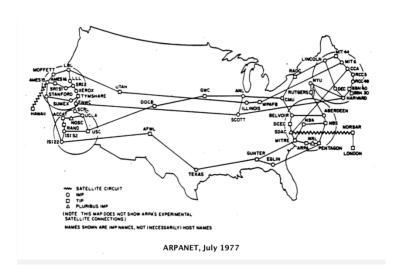
- remove the least recently added vertex v
- add each of v's unvisited neighbors to the queue,
   and mark them as visited.

Proposition. In any connected graph G, BFS computes shortest paths from s to all other vertices in time proportional to E+V.





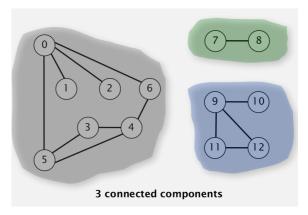
#### Breadth-first search application: routing



#### **Connected Componets**

Def. Vertices v and w are connected if there is a path between them.

Def. A connected component is a maximal set of connected vertices.



## **Graph Traversal Summary**

problem	BFS	DFS	time
path between s and t	~	~	E + V
shortest path between s and t	~		E + V
connected components	~	V	E + V