

Key Indexed Counting

Sonam Wangmo

Gyalpozhing College of Information Technology

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Key Indexed Counting

- A simple method for sorting that is effective whenever the keys are small integers.
- There are mainly four steps in processing key-indexed counting sort.
 - 1 Compute frequency counts.
 - 2 Transform counts to indices.
 - 3 Distribute the data.
 - 4 Copy back.

Example: Inputs are given below

Table 1: array a[]

Anderson	2
Brown	3
Davis	3
Garcia	4
Harris	1
Jackson	3
Johnson	4
Jones	3
Martin	1
Martinez	2
Miller	2
Moore	1
Robinson	2
Smith	4
Taylor	3

1. Compute frequency counts

The first step is to count the frequency of occurrence of each key value, using an int array `count[]`. For each item, we use the key to access an entry in `count[]` and increment that entry. If the key value is r , we increment `count[r+1]`.

```
1
2     for (i = 0; i < N; i++)
3         count[a[i].key() + 1]++;
```

2.Transform counts to indices

Next, we use `count[]` to compute, for each key value, the starting index positions in the sorted order of items with that key.

```
1
2     for (int r = 0; r < R; r++)
3         count[r+1] += count[r];
```

3. Distribute the data

With the `count[]` array transformed into an index table, we accomplish the actual sort by moving the items to an auxiliary array `aux[]`.

```
1
2     for (int i = 0; i < N; i++)
3         aux[count[a[i].key()]++] = a[i]
```

4. Copy back

Since we accomplished the sort by moving the items to an auxiliary array, the last step is to copy the sorted result back to the original array.

```
1      for (int i = 0; i < N; i++)  
2      a[i] = aux[i]
```

Key-indexed counting ($a[i].key$ is an int in $[0, R)$).

```
1      int N = a.length;
2      String[] aux = new String[N];
3      int[] count = new int[R+1];
4      // Compute frequency counts.
5      for (int i = 0; i < N; i++)
6          count[a[i].key() + 1]++;
7      // Transform counts to indices.
8      for (int r = 0; r < R; r++)
9          count[r+1] += count[r];
10     // Distribute the records.
11     for (int i = 0; i < N; i++)
12         aux[count[a[i].key()]++] = a[i];
13     // Copy back.
14     for (int i = 0; i < N; i++)
15         a[i] = aux[i];
```


Q. Do key indexing counting sort for the given array.

i	a[i]
0	d
1	a
2	c
3	f
4	f
5	b
6	d
7	b
8	f
9	b
10	e
11	a

use a for 0
b for 1
c for 2
d for 3
e for 4
f for 5

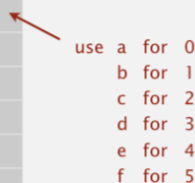


Figure 1: a[i]

Running Time of Key Indexed Counting

- key-indexed counting does no compares (it accesses data only through `key()`). When R is within a constant factor of N , we have a linear-time sort.
- $O(N + R)$ R can be characters of alphabets, Hence we ignore R .
- $O(N)$