

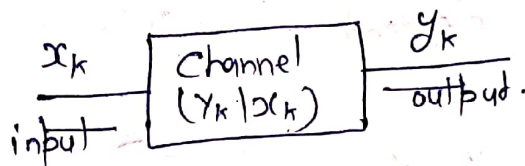


POORNIMA

COLLEGE OF ENGINEERING

DETAILED LECTURE NOTES

Discrete memoryless Channel:



The i/p and o/p are x_k and y_k discrete number. means those

are not vary with the continuous range

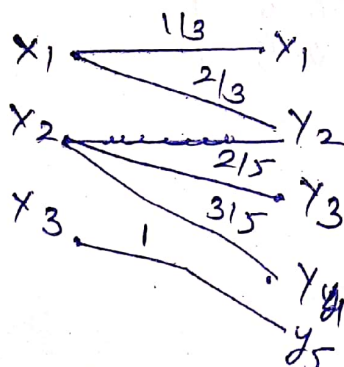
These values are not uniform.

Memoryless: it depends only present input, not depends on previous input. or pass element.

Channel matrix:

$$\left[\frac{Y_k}{X_k} \right] = \begin{matrix} & \begin{matrix} x_1 & x_2 & x_3 & \dots & x_m \end{matrix} \\ \begin{matrix} y_1 \\ y_2 \\ y_3 \\ \dots \\ y_n \end{matrix} & \begin{bmatrix} \frac{y_1}{x_1} & \frac{y_2}{x_1} & \frac{y_3}{x_1} & \dots & \frac{y_n}{x_1} \\ \frac{y_1}{x_2} & \frac{y_2}{x_2} & \frac{y_3}{x_2} & \dots & \frac{y_n}{x_2} \\ \frac{y_1}{x_3} & \frac{y_2}{x_3} & \frac{y_3}{x_3} & \dots & \frac{y_n}{x_3} \\ \dots & \dots & \dots & \dots & \dots \\ \frac{y_1}{x_m} & \frac{y_2}{x_m} & \frac{y_3}{x_m} & \dots & \frac{y_n}{x_m} \end{bmatrix} \end{matrix}$$

Lossless Channel

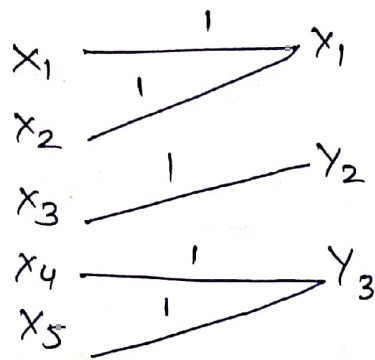


$$\begin{matrix} & \begin{matrix} y_1 & y_2 & y_3 & y_4 & y_5 \end{matrix} \\ \begin{matrix} x_1 \\ x_2 \\ x_3 \end{matrix} & \begin{bmatrix} \frac{1}{3} & \frac{2}{3} & 0 & 0 & 0 \\ 0 & 0 & \frac{2}{5} & \frac{3}{5} & 0 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix} \end{matrix}$$

$$\sum R_i = 1$$

→ every column has one & non zero element

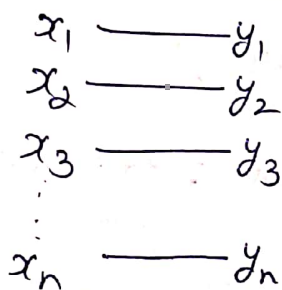
2. Deterministic Channel:



$$\begin{matrix} & y_1 & y_2 & y_3 \\ \begin{matrix} x_1 \\ x_2 \\ x_3 \\ x_4 \\ x_5 \end{matrix} & \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \\ 0 & 0 & 1 \end{bmatrix} \end{matrix}$$

every element is one. → Every Row has one non zero value.

3. Noiseless Channel

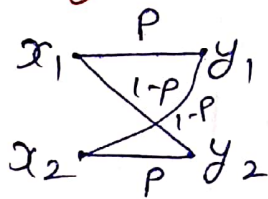


$$\begin{matrix} & y_1 & y_2 & \dots & y_n \\ \begin{matrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{matrix} & \begin{bmatrix} 1 & 0 & \dots & 0 \\ 0 & 1 & \dots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & \dots & \dots & 1 \end{bmatrix} \end{matrix}$$

→ Every Column has one non zero value.

only diagonal terms are coming. So any intermediate signal is there.

4. Binary Symmetric Channel



$$\begin{matrix} & y_1 & y_2 \\ \begin{matrix} x_1 \\ x_2 \end{matrix} & \begin{bmatrix} p & 1-p \\ 1-p & p \end{bmatrix} \end{matrix}$$

$$x_1 = x_2 = 0.5$$

$$\begin{bmatrix} 0.5 & 0.5 \end{bmatrix} \begin{bmatrix} p & 1-p \\ 1-p & p \end{bmatrix}$$

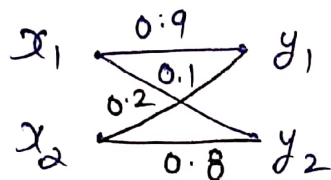
$$= \begin{bmatrix} 0.5p + 0.5(1-p) \\ 0.5(1-p) + 0.5p \end{bmatrix} = \begin{bmatrix} 0.5 \\ 0.5 \end{bmatrix}$$



POORNIMA

COLLEGE OF ENGINEERING

DETAILED LECTURE NOTES



$$\begin{bmatrix} 0.9 & 0.1 \\ 0.2 & 0.8 \end{bmatrix}$$

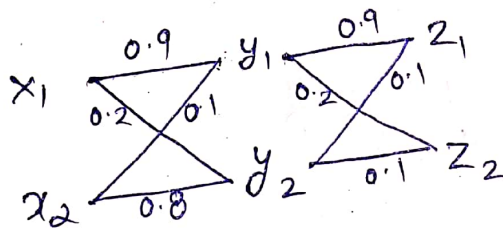
$$= \begin{bmatrix} 0.5 & 0.5 \end{bmatrix} \begin{bmatrix} 0.9 & 0.1 \\ 0.2 & 0.8 \end{bmatrix} = \begin{bmatrix} 0.55 & 0.45 \end{bmatrix}$$

$$= \begin{bmatrix} 0.55 \\ 0.45 \end{bmatrix}$$

Joint probability of X_1 & Y_2

$$\begin{bmatrix} 0.5 & 0 \\ 0 & 0.5 \end{bmatrix} \begin{bmatrix} 0.9 & 0.1 \\ 0.2 & 0.8 \end{bmatrix}$$

Casading of channels



$$\begin{bmatrix} 0.9 & 0.1 \\ 0.2 & 0.8 \end{bmatrix}_{2 \times 2} \begin{bmatrix} 0.9 & 0.1 \\ 0.2 & 0.8 \end{bmatrix}_{2 \times 2}$$

$$\text{Channel Matrix} = \text{Channel Matrix} \left(\frac{Y}{X} \right) \times \text{Channel Matrix} \left(\frac{Z}{Y} \right)$$

(2x2)

Symmetric Channel \rightarrow Represented by channel matrix in which the second and subsequent Row contain the same element as that in the first Row. but in different order.

$$P(B|A) = \begin{matrix} & \begin{matrix} b_1 & b_2 & \dots & b_s \end{matrix} \\ \begin{matrix} a_1 \\ a_2 \\ \vdots \\ a_n \end{matrix} & \begin{bmatrix} p_1 & p_2 & \dots & p_s \\ p_2 & p_3 & \dots & p_1 \\ \vdots & \vdots & \ddots & \vdots \\ p_3 & \dots & \dots & p_1 \end{bmatrix} \end{matrix}$$

Channel Capacity for all the channels:

1. Binary Symmetric Channel $C = 1 - h$

where $h = P \log_2 \frac{1}{P} + P \log_2 \frac{1}{P}$

2. Noiseless Channel =

$\log_2 n$ bits/sec.

3. Deterministic Channel = $\log_2 S$ bit/sec.

4. Symmetric Channel = $\log_2 S - h$

Noiseless Channel

$$I(A, B) = H(A) - H(A|B)$$

$$H(A|B) = 0$$

So

$$I(A, B) = H(A)$$

Channel Capacity

$$C = \max(I(A, B)) \text{ n/s}$$

$$C = \max[H(A)] \text{ n/s}$$

Deterministic Channel

$$I(A, B) = H(B) - H(B|A)$$

$$H(B|A) = 0$$

$$I(A, B) = H(B)$$

Channel Capacity $C = \max[I(A, B)] \times \text{n/s}$

$$C = \max[H(B)] \text{ n/s}$$

Q. For the channel matrix given below find the Channel Capacity

$P(b_j/a_i) =$

	b_1	b_2	b_3
a_1	$1/2$	$1/3$	$1/6$
a_2	$1/3$	$1/6$	$1/2$
a_3	$1/6$	$1/2$	$1/3$

Channel Capacity = $\log_2 S - h$

$$h = H(B|A) = \sum_{j=1}^3 \frac{1}{2} \log_2 2 + \frac{1}{3} \log_2 3 + \frac{1}{6} \log_2 6$$

$$h = 1.459 \text{ b/sym}$$

$$C = \log_2 3 - 1.459$$

$$C = 0.1258 \text{ b/sec}$$

$S = 3$ with 3 sequence.



POORNIMA

COLLEGE OF ENGINEERING

DETAILED LECTURE NOTES

PAGE NO.

Channel Example:

Q4 A bsc has the following noise matrix with source probabilities of $p(x_1) = 2/3$ & $p(x_2) = 1/3$

$$P(y|x) = \begin{matrix} & y_1 & y_2 \\ x_1 & \begin{bmatrix} 3/4 & 1/4 \end{bmatrix} \\ x_2 & \begin{bmatrix} 1/4 & 3/4 \end{bmatrix} \end{matrix} \left\{ \begin{array}{l} \text{also represent} \\ \text{by } \overline{P} \text{ and } P \\ \begin{bmatrix} \overline{P} & P \\ P & \overline{P} \end{bmatrix} \end{array} \right.$$

- (i) Determine $H(x)$, $H(y)$
- (ii) find channel capacity C
- (iii) find channel efficiency and Redundancy

Ans

$$(i) H(x) = \sum_{i=1}^2 p(x_i) \log_2 \left(\frac{1}{p(x_i)} \right)$$

$$= \frac{2}{3} \log_2 \left(\frac{3}{2} \right) + \frac{1}{3} \log_2 (3)$$

$$\boxed{H(x) = 0.9183} \text{ bits}$$

$$(ii) H(y) = P(x,y) = \begin{matrix} & y_1 & y_2 \\ x_1 & \begin{bmatrix} 1/2 & 1/6 \end{bmatrix} \\ x_2 & \begin{bmatrix} 1/12 & 1/4 \end{bmatrix} \end{matrix} \Rightarrow P(y) = \begin{matrix} y_1 = \frac{1}{2} + \frac{1}{12} = \frac{7}{12} \\ y_2 = \frac{1}{6} + \frac{1}{4} = \frac{5}{12} \end{matrix}$$

[for getting the probability of y_1 when occur of x_1]

$$x_1 = \frac{2}{3} \times \frac{3}{4} = \frac{1}{2}$$

$$P(y) = \left[\frac{7}{12} \quad \frac{5}{12} \right]$$

$$H(y) = \sum_{j=1}^2 P(y_j) \log_2 \left(\frac{1}{P(y_j)} \right)$$

$$= \frac{7}{12} + \log_2 \frac{12}{7} + \frac{5}{12} \log_2 \frac{12}{5}$$

$$H(y) = 0.9799 \text{ bits}$$

2 Channel Capacity

$$C = 1 - h$$

$$\text{where } h = P \log_2 \frac{1}{P} + P \log_2 \frac{1}{P}$$

$$C = 1 - \left[\frac{3}{4} \log_2 \frac{4}{3} + \frac{1}{4} \log_2 4 \right]$$

$$C = 1 - 0.813$$

$$C = 0.1887 \text{ bits}$$

$$(iii) \quad n_{\text{channel}} = \frac{I(x, y)}{C}$$

$$I(x, y) = H(x) - H(x|y)$$

$$H(x|y) = H(x, y) - H(y)$$

$$\rightarrow H(x) + H(y|x)$$

$H(x, y) \rightarrow$ check all the probabilities when x is applied



POORNIMA

COLLEGE OF ENGINEERING

DETAILED LECTURE NOTES

PAGE NO.

Qm $H(x, y) = H(x) + H(y|x)$

$$H(x, y) = 2 \times \frac{3}{4} \log_2 \frac{4}{3} + 2 \times \frac{1}{4} \log_2 4$$
$$= 1.623$$

$$H(x|y) = H(x, y) - H(y)$$

$$H(x|y) = 0.749$$

$$I(x, y) = 0.9183 - 0.749$$

$$I(x, y) = 0.1686$$

(iii) $\eta_{en} = I(x, y) / C$

$$= 0.1686 / 0.1887$$

$$\eta = 89.35\%$$

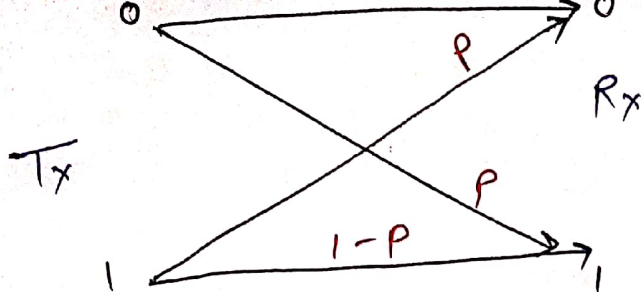
iv) Redundancy $= \gamma = 1 - \eta_{en}$

$$= 1 - 0.89$$

$$\gamma = 10.65\%$$

Example of binary Symmetric Channel (BSC)

A channel can be visualized for transmitting 1's and 0's from the transmitter and Receiver. The defined input symbol shown in the figure. A BSC channel flips 1 to 0 and equal probability.



A BSC channel

$$P(Y=0 | X=0) = 1-P$$

$$P(Y=0 | X=1) = P$$

$$P(Y=1 | X=1) = 1-P$$

$$P(Y=1 | X=0) = P$$

$$\begin{aligned} P(Y=0) &= P(X=0) \cdot P(Y=0 | X=0) + P(X=1) \cdot P(Y=0 | X=1) \\ &= 0.5(1-P) + 0.5(P) \end{aligned}$$

$$= 0.5$$

$$\begin{aligned} P(Y=1) &= P(X=0) \cdot P(Y=1 | X=0) + P(X=1) \cdot P(Y=1 | X=1) \\ &= 0.5(P) + 0.5(1-P) \end{aligned}$$

$$= 0.5$$

$$I(x_0; y_0) = I(0; 0)$$

$$= \log_2 \left(\frac{P(Y=0 | X=0)}{P(Y=0)} \right) = \log_2 \left(\frac{1-P}{0.5} \right)$$

$$\boxed{I(x_0; y_0) = \log_2 2(1-P)}$$

$$I(x_1; y_0) = I(1; 0)$$

$$\log_2 \left(\frac{P(Y=0 | X=1)}{P(Y=0)} \right) = \log_2 \left(\frac{P}{0.5} \right)$$

$$\boxed{I(x_1; y_0) = \log_2 2P}$$