2. Operating System Structures

ECE30021/ITP30002 Operating Systems

Agenda

- Operating-system services
- Interfaces for users and programmers
- Components and their interconnections
- Virtual Machines
- Design, implementation, generation
- System boot

Objectives

- To describe the services an operating system provides to users, processes, and other systems
- To discuss the various ways of structuring an operating system
- To explain how operating systems are installed and customized and how they boot

- Operating systems provide an environment for execution of programs and services to programs and users
- One set of operating-system services provides functions that are helpful to the user:
 - User interface Almost all operating systems have a user interface (UI).
 - □ Varies between Command-Line (CLI), Graphics User Interface (GUI), Batch
 - Program execution The system must be able to load a program into memory and to run that program, end execution, either normally or abnormally (indicating error)
 - I/O operations A running program may require I/O, which may involve a file or an I/O device
 - **File-system manipulation** The file system is of particular interest. Programs need to read and write files and directories, create and delete them, search them, list file Information, permission management.

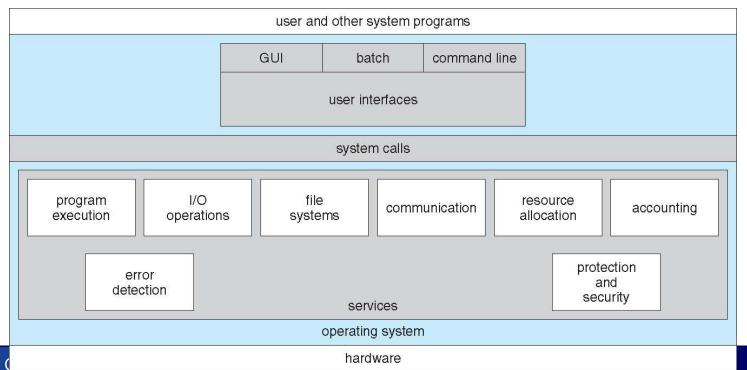
- **Communications** Processes may exchange information, on the same computer or between computers over a network
 - Communications may be via shared memory or through message passing (packets moved by the OS)
- Error detection OS needs to be constantly aware of possible errors
 - May occur in the CPU and memory hardware, in I/O devices, in user program
 - □ For each type of error, OS should take the appropriate action to ensure correct and consistent computing
 - Debugging facilities can greatly enhance the user's and programmer's abilities to efficiently use the system

- Another set of OS functions exists for ensuring the efficient operation of the system itself via resource sharing
 - Resource allocation When multiple users or multiple jobs running concurrently, resources must be allocated to each of them
 - Many types of resources
 Some (such as CPU cycles, main memory, and file storage)
 may have special allocation code,
 others (such as I/O devices) may have general request and release code
 - Accounting To keep track of which users use how much and what kinds of computer resources

- Protection and security The owners of information stored in a multiuser or networked computer system may want to control use of that information, concurrent processes should not interfere with each other
 - Protection involves ensuring that all access to system resources is controlled
 - □ Security of the system from outsiders requires user authentication, extends to defending external I/O devices from invalid access attempts
 - If a system is to be protected and secure, precautions must be instituted throughout it.

- Services for user
 - User interface
 - Program execution
 - I/O operation
 - File-system manipulation
 - Communications
 - Error detection

- Functions for efficient operation of system itself
 - Resource allocation
 - Accounting
 - Protection and security



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OS User Interface - CLI

- CLI or command interpreter allows direct command entry
 - Sometimes implemented in kernel, sometimes by systems program
 - Sometimes multiple flavors implemented shells
 - Primarily fetches a command from user and executes it
 - □ Sometimes commands built-in, sometimes just names of programs
 - □ If the latter, adding new features doesn't require shell modification

Command Line Interpreter (CLI)

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sd1	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0	0	
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fd0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0	0	
sd0	0.6	0.0	38.4	0.0	0.0	0.0	8.2	0	0	
sd1	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0	0	
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OS User Interface - GUI

- User-friendly desktop metaphor interface
 - Usually mouse, keyboard, and monitor
 - Icons represent files, programs, actions, etc
 - Various mouse buttons over objects in the interface cause various actions (provide information, options, execute function, open directory (known as a folder)
 - Invented at Xerox PARC
- Many systems now include both CLI and GUI interfaces
 - Microsoft Windows is GUI with CLI "command" shell
 - Apple Mac OS X is "Aqua" GUI interface with UNIX kernel underneath and shells available
 - Unix and Linux have CLI with optional GUI interfaces (CDE, KDE, GNOME)

The Mac OS X GUI



Touchscreen Interfaces

- Touchscreen devices require new interfaces
 - Mouse not possible or not desired
 - Actions and selection based on gestures
 - Virtual keyboard for text entry



GUI (3D Desktop)









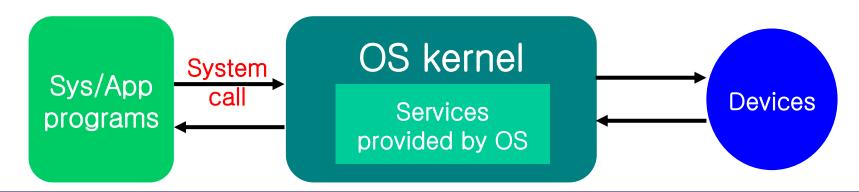
Programming Interfaces

- System calls
 - Primitive programming interface provided through interrupt
 - System-call interface
 - Connection between program language and OS
 Ex) implementations of open(), close(), ...
- API (Application Programming Interface)
 - High-level programming interface
 Ex) MessageBox(...);
 - Ex) Win32 API, POSIX API, Java API

System Calls

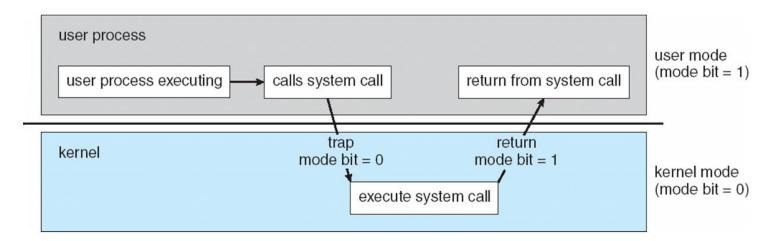
- System call: the mechanism used by an application program to request service from OS kernel
 - "Function calls to OS kernel available through interrupt"
 - Generally, provided as interrupt handlers written in C/C++ or assembly.
 - A mechanism to transfer control safely from lesser privileged modes to higher privileged modes.

Ex) POSIX system calls: open, close, read, write, fork, kill, wait, ...



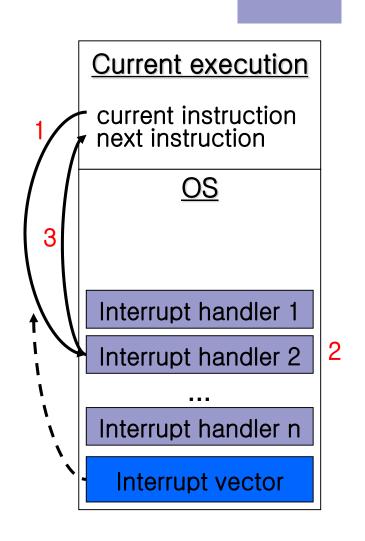
Dual Mode Operation

- User mode
 - User defined code (application)
 - Privileged instructions, which can cause harm to other system, are prohibited
 - Privileged instruction can be invoked only through OS system call
- Kernel mode (supervisor mode, system mode, privileged mode)
 - OS code
 - Privileged instructions are permitted



Interrupt Mechanism

- Interrupt handling
 - 1. CPU stops current work and transfers execution to interrupt handler
 - Interrupt vector: table of interrupt handlers for each types interrupt
 - 2. Interrupt is handled by corresponding handler
 - 3. Return to the interrupted program
 - Before interrupt handler is invoked, necessary information should be saved (return address, state)



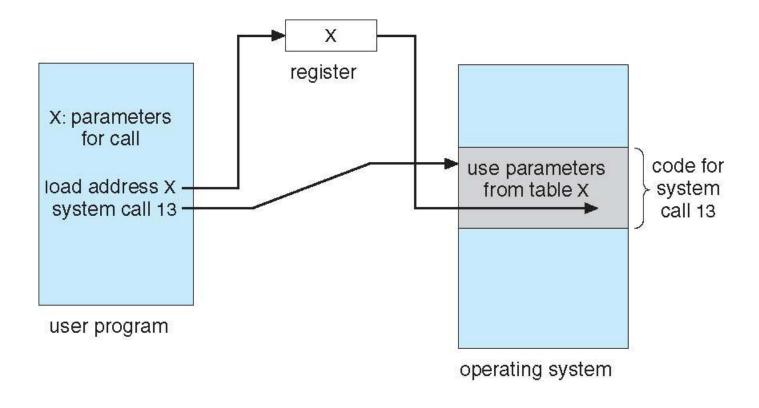
Parameter Passing in System Call

- Often, more information is required than simply identity of desired system call
 - Exact type and amount of information vary according to OS and call
- Three general methods used to pass parameters to the OS
 - Simplest: pass the parameters in registers
 - □ In some cases, may be more parameters than registers
 - Parameters stored in a block, or table, in memory, and address of block passed as a parameter in a register
 - □ This approach taken by Linux and Solaris
 - Parameters placed, or pushed, onto the stack by the program and popped off the stack by the operating system
 - Some OS prefer the block or stack method because they do not limit the number or length of parameters being passed

Parameter Passing in System Call

- Internally, system call is serviced through interrupt
 - Additional information can be necessary
- Parameter passing methods
 - Register (simple information)
 - Address of block (large information)
 - System stack

Passing of parameters as a table



Types of System Calls

- Process control
- File management
- Device management
- Information maintenance
- Communication



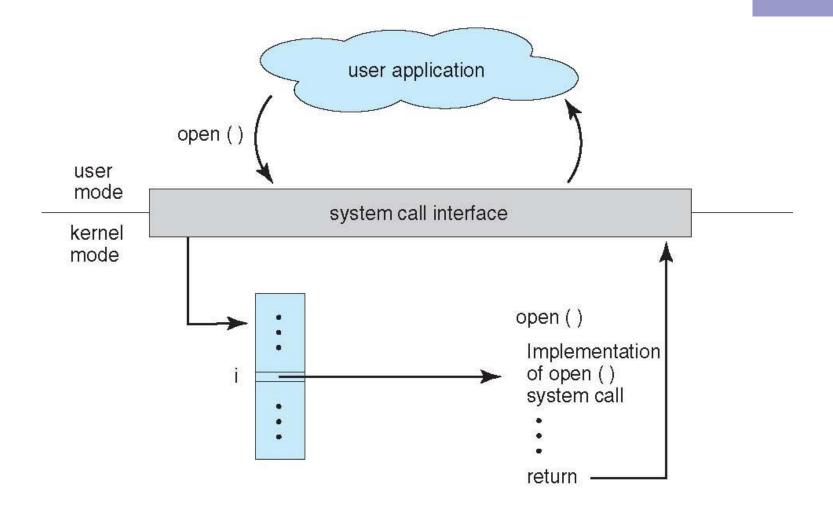
Example

Copy file from A to B Error or Abnormal cases I/O system calls Read file names Display, srcFile, destFile File system calls **Keyboard/mouse** delete file, ... Open srcFile File system calls Create destFile I/O system calls message, ... Read from srcFile File system calls Write into destFile Process system calls Abnormal termination Abort, ... Close srcFile File system calls

and destFile

- How to invoke system calls in high-level language?
 Ex) int open(const char *path, int oflag);
- System-call interface: link between runtime support system of <u>programming language</u> and OS system calls
 - Implementation of I/O functions available in programming language (ex: glibc, MS libc, ...)

- Typically, a number is associated with each system call.
 - System-call interface maintains a table indexed according to these numbers.
- The system call interface invokes intended system call in OS kernel and returns status of the system call and any return values.
- The caller needs to know nothing about how the system call is implemented.
 - Just needs to obey API and understand what OS will do as a result call
 - Most details of OS interface hidden from programmer by API
 - Managed by run-time support library (set of functions built into libraries included with compiler)



Example of system-call interface in Linux

```
User program
int main()
{
    ...
    open();
    ...
}
```

```
System-call Interface (libc)

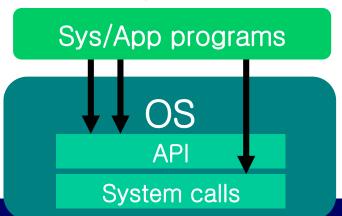
open()
{
...
movl 5, %eax 'system call number
int $0x80 - 'generate interrupt
...
}
```

Interrupt Handling Mechanism **OS** kernel sys_open()

- What does system-call interface do?
 - Passing information to the kernel
 - Switch to kernel mode
 - Any data processing and preparation for execution in kernel mode
 - ETC.
- Cf. System call vs. I/O functions in programming language Ex) read(), vs. fread()
 - read(): provided by OS
 - fread(): standard function defined in C language
 - fread() is implemented using read()

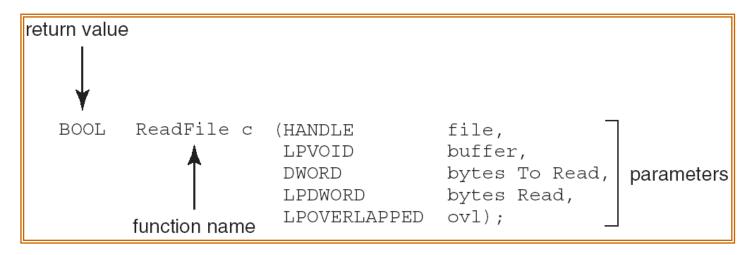
Application Programming Interface

- API: interface that a computer system (OS), library or application provides to allow requests for service
 - A set of functions, parameters, return values available to application programmers.
 - Ex) Win32 API, POSIX API, etc.
 - □ MessageBox(..), CreateWindow(...), ...
 - Can be strongly correlated to system calls
 Ex) POSIX API ≈ UNIX system calls
 - Can provide high-level features implemented with system calls
 - Ex) Win32 API is based on system calls
 - Ex) POSIX thread library API



Example of API

Win32 API function ReadFile() —a function for reading from a file



- A description of the parameters passed to ReadFile()
 - HANDLE file—the file to be read
 - LPVOID buffer—a buffer where the data will be read into and written from
 - DWORD bytesToRead—the number of bytes to be read into the buffer
 - LPDWORD bytesRead—the number of bytes read during the last read
 - LPOVERLAPPED ovl—indicates if overlapped I/O is being used

Application Programming Interface

- Why API rather than system call?
 - Ease to use
 - API provides higher-level interface than system call
 - System portability

Examples of System Calls

	Windows	Unix
Process Control	<pre>CreateProcess() ExitProcess() WaitForSingleObject()</pre>	<pre>fork() exit() wait()</pre>
File Manipulation	<pre>CreateFile() ReadFile() WriteFile() CloseHandle()</pre>	<pre>open() read() write() close()</pre>
Device Manipulation	SetConsoleMode() ReadConsole() WriteConsole()	ioctl() read() write()
Information Maintenance	<pre>GetCurrentProcessID() SetTimer() Sleep()</pre>	<pre>getpid() alarm() sleep()</pre>
Communication	<pre>CreatePipe() CreateFileMapping() MapViewOfFile()</pre>	<pre>pipe() shmget() mmap()</pre>
Protection	<pre>SetFileSecurity() InitlializeSecurityDescriptor() SetSecurityDescriptorGroup()</pre>	<pre>chmod() umask() chown()</pre>