

Q:

For what values of  $b$  does

$$2x_1^2 + 2x_2^2 + 2bx_1x_2 \leq 3$$

define an ellipse?

A:

$$\boxed{\phantom{2x_1^2 + 2x_2^2 + 2bx_1x_2 \leq 3}} = x^T Q x \text{ for } Q = \begin{bmatrix} 2 & b \\ b & 2 \end{bmatrix}.$$

$$Q \succ 0 \text{ when } 4 - b^2 > 0.$$

Thus  $2x_1^2 + 2x_2^2 + 2bx_1x_2 \leq 3$  defines an ellipse when  $b^2 < 4$  or  $b \in (-2, 2)$ .

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Q: The constraint

$$(*) (2x_1 - x_2 + 1)^2 + (-x_1 + x_2 + 1)^2 \leq (4x_1 - x_2 + 2)^2$$

A: is NOT a SOC constraint.

$$\text{let } y_1 = 2x_1 - x_2 + 1$$

$$y_2 = -x_1 + x_2 + 1$$

$$z = 4x_1 - x_2 + 2$$

Add constraint  $z \geq 0$ .

Then  $(*)$  is equivalent to



$$(\pm, y) \in \mathcal{Z} \quad \text{or} \quad \pm \geq \|y\|_2,$$

$$\pm \geq \sqrt{y_1^2 + y_2^2}$$


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Q:  $\|x-a\| + \|y-b\|^2 \leq 0$

A: only feasible solution (since  $\|\cdot\|$  always  $\geq 0$ )  
is  $x=a, y=b$ . [Linear constraint]

Q:  $\|x-a\|_2^2 - r^2 \leq 0, \quad r \geq 0$

A: This is a SOC:  
 $r \geq \|x-a\|_2$

Note that  
 $r \geq 0$  is  
important!

Q:  $\|x-a\|_2 \leq \|x-b\|_2$

A:  $\|x-a\|_2 \leq \|x-b\|_2 \Leftrightarrow \|x-a\|^2 \leq \|x-b\|^2$

$\Leftrightarrow (x-a)^T(x-a) \leq (x-b)^T(x-b)$

$x^T x - 2a^T x + a^T a \leq x^T x - 2b^T x + b^T b$

$2(b-a)^T x \leq b^T b - a^T a$



## Linear Constraint

Q: prove for  $\lambda > 0$  it is impossible to find  $x$  such that  $J_1(x) < J_1(x_2)$  AND  $J_2(x) < J_2(x_2)$

A: Suppose for a contradiction that such an  $x$  existed, then we have

$$J_1(x) + \lambda J_2(x) < J_1(x_2) + \lambda J_2(x_2)$$

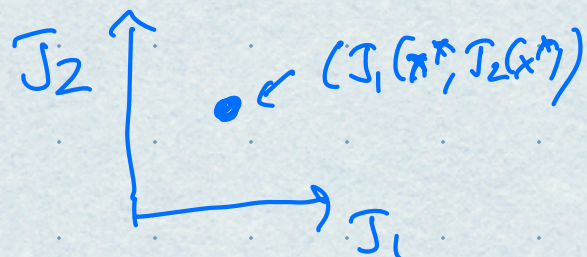
when  $\lambda > 0$ , but this means that  $x_2$  was not the minimizer  $\square$

Q: Suppose  $\arg \min_x J_1(x) = \arg \min_x J_2(x) = x^*$

what does pareto curve look like?

A: whole pareto curve is a single point.

Suppose  $\exists x = x^*$  with





$$J_1(x) + \lambda J_2(x) < J_1(x^*) + \lambda J_2(x^*)$$

we know that  $J_1(x^*) \leq J_1(x) \quad \forall x$   
 $J_2(x^*) \leq J_2(x) \quad \forall x,$   
 so such an  $x$  cannot exist.

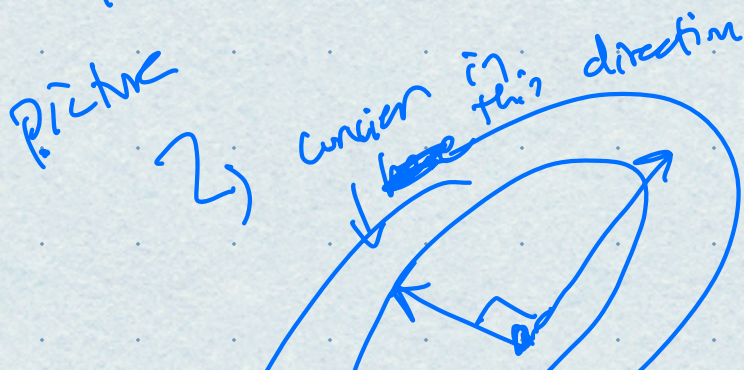
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Q:  $(x \ y) \begin{pmatrix} a & b \\ c & d \end{pmatrix} \begin{pmatrix} x \\ y \end{pmatrix} \leq 1$  (\*)

$$Q = \begin{bmatrix} 3/5 & -4/5 \\ 4/5 & 3/5 \end{bmatrix} \begin{bmatrix} 4/25 & 0 \\ 0 & 4/25 \end{bmatrix} \begin{bmatrix} 3/5 & -4/5 \\ 4/5 & 3/5 \end{bmatrix}^T$$

A: This is an ellipsoid  
 which is a convex set

[ Ellipsoid since eigenvalues all  $> 0$  ]





Q: Find  $A, b$  such that (\*)  
can be written like  $\|A \begin{pmatrix} x \\ y \end{pmatrix}\| \leq b$

A:  $x^T Q x \leq 1 \Leftrightarrow \|Q^{\frac{1}{2}} x\| \leq 1$

$$Q = U \Lambda U^T, \text{ so } Q^{\frac{1}{2}} = U \Lambda^{\frac{1}{2}}$$

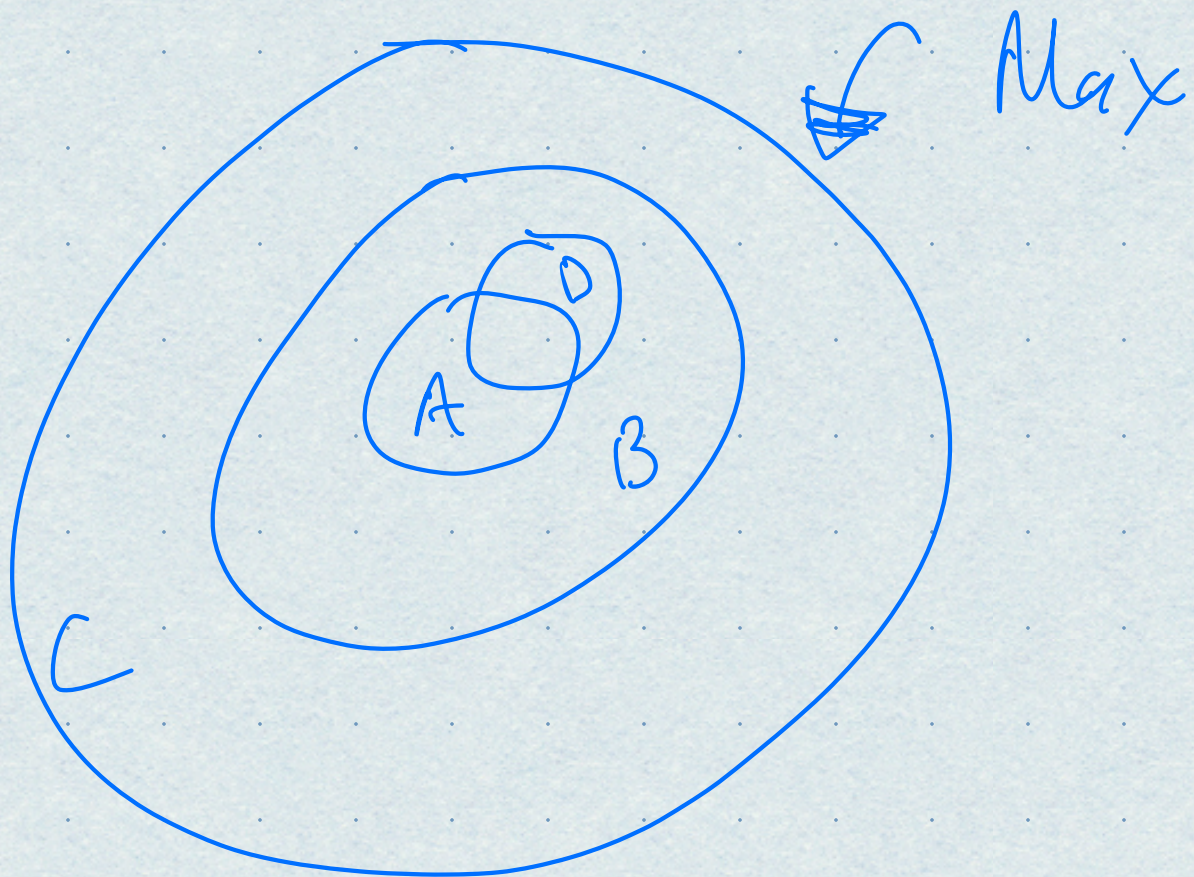
$$\begin{bmatrix} 3 & 4 \\ 5 & 5 \end{bmatrix} \begin{bmatrix} 4/5 & 0 \\ 0 & 2/5 \end{bmatrix} = \begin{bmatrix} 3 & -8 \\ 4 & 6 \\ 25 & 25 \end{bmatrix}$$

$$A = U \Lambda^{\frac{1}{2}} = \frac{1}{25} \begin{bmatrix} 3 & -8 \\ 4 & 6 \end{bmatrix}$$

$$b = 1.$$



Q: Relaxations:



E: Same feas region as A,  
except 'min'

F: Dual LP to ~~the~~ C.



A:

$$z_A \leq z_B$$

$$z_A \times z_D$$

$$z_A \geq z_E$$

$$z_F = z_C$$

$$z_D \leq z_B$$

Beer location:

$$2.1 \quad \min \sum_{i \in I} \sum_{j \in J} c_{ij} x_{ij} + \sum_{i \in I} h_i z_i$$



2.2

$$\sum_{i \in I} x_{ij} \geq b_j \quad \forall j \in J$$

Note this implements  
"fixed cost"

2.3

$$\sum_{j \in J} x_{ij} \leq M_i z_i \quad \forall i \in I$$

where  $M_i, i \in I$  is upper bound

on  $\sum_{j \in J} x_{ij}$ , say

$$M_i = \sum_{j \in J} b_j \quad \forall i \in I.$$

2.4

$$x_{ij} \geq 0 \quad \forall i \in I, j \in J$$

$$z_i \in \{0, 1\} \quad \forall i \in I$$