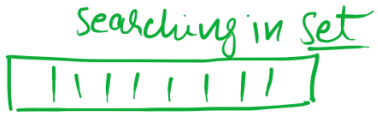
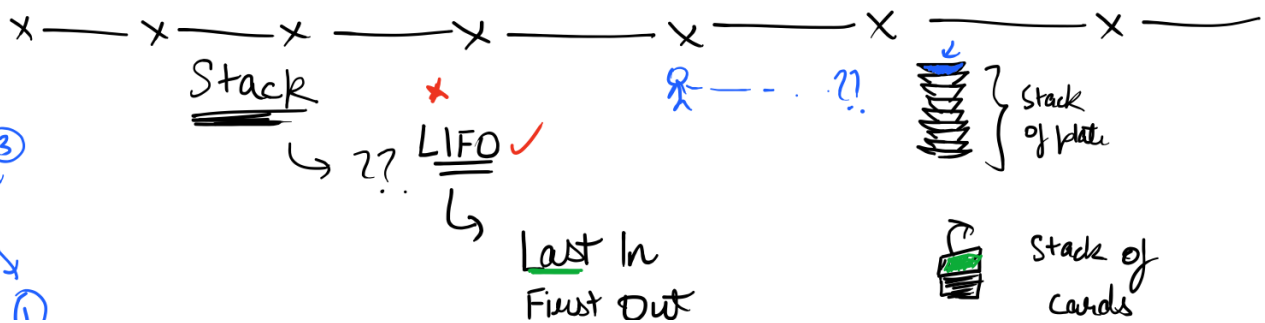
Thursday, 25 July 2024
8:46PM

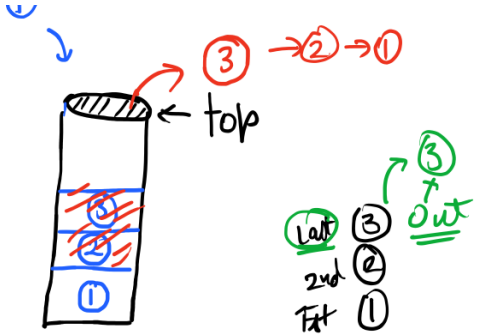
Binary Search
↳ Reduce the search space → Sorted
⇓
O(log n)

↳ Binary Search
Minimize the maximum
Maximize the minimum

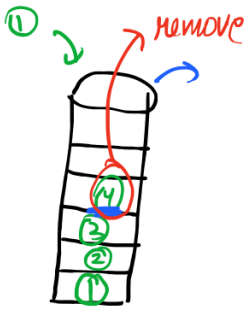
✓Hashing → Searching in O(1) time







→ Stack Memory



* Stacks do not have index

Stack

Adding	Stack: only from 1 end O(1)	Array: Directly at index O(1)
Getting: If top O(1) else O(n)	only get the top element	Directly at the index
Removing: O(1) if top, O(n) else	Only the top most can be removed	Shifting the other elements

Array / Array List

O(n) → When idx is unknown
O(1) → When idx is known

O(n) → Shifting
O(1) → If end is removed

* Stack is like an array without indexes and only 1 end is available.

→ Dynamic

Methods of Stack

Push → Adding an element in the stack — O(1)

Pop → Removes and Returns the top element — O(1)

Peek → Returns the top element but doesn't remove — O(1)

Size → Returns the size of the stack — O(1) * with a variable

Push 1
Push 2
Push 3
Pop ⇒ 3
Peek ⇒ 2
Push 4
Peek ⇒ 4
Push 5
size → 4

* We cannot traverse a stack → Cannot use indexes.

Heap → Stack object will be present in heap memory

Internally, Stack uses arrays.

int vs Integer

int → data type

objects ✗

Integer → Class → objects

Class on Sat → Stack
11-3, 12-4 ✓

6 ⇒ 100% ✓
↳ contest on
1.5 BS + Hashing

Sunday 12-3-3.30 ✓
↳ Doubt + Question
BS + Hashing

()(()) ✓

))((✗
)()(✗
(()) ✗

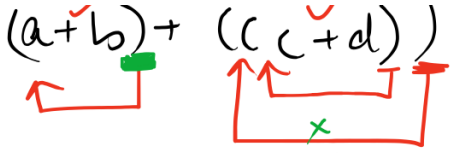
Q1 Extra Brackets ✓

(a+b) + ((c+d)) ✓

((a+b)+(c+d))

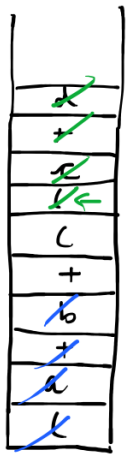
(a+b) + (((c+d)+e))

No extra brackets



↳ we are not doing any calculation

T(n) : O(n)
S(n) : O(n)



⇐ top of stack

(a+b) + ((c+d))

'(' ↳ x Invalid

⇒ return true

⇒ extra bracket

Next Greater Right

6 5 8 0 2 3 1

8 8 -1 2 3 -1 -1

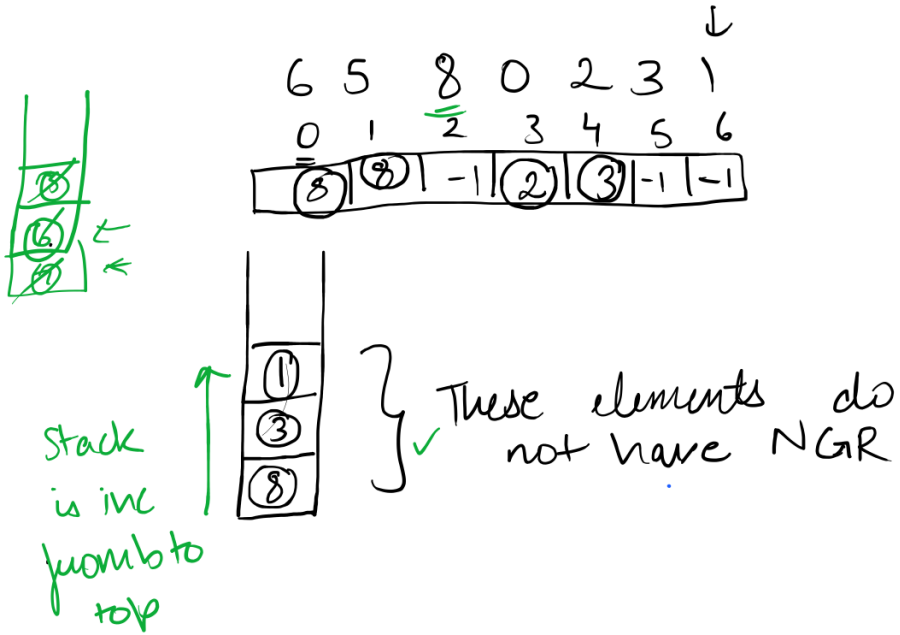
1) Using 2 loops find the greater element on right side.

for (i=0 → n)

TC : O(n²)
SC : O(1)

```
for(j=i+1 → n){
  if(a[j] > a[i]){
    ans[i] = a[j]
    break
}}}
```

2) Using Stack



Traverse from left to right
↳ while the current element is greater than the top of stac
↳ current el is the NG from to

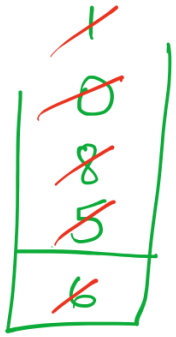
TC : O(n)
SC : O(n)

-1 -1 -1 -1 -1

5, 4, 3, 2, 1

0 1 2 3 4

$n \Rightarrow 1+2+3+4+\dots(n-1) = \frac{(n-1)(n-1+1)}{2} = \frac{n(n-}{2}$



6, 5, 8, 0, 1

$n = n \times = 2 \times n \approx O(n)$