

calculate number will run just once per each emp.

D. 11

U. 11

Problem with given difference

6
→ 5 10 3 2 50 80 ✓
i j

Diff → 78 →
↳ Return 1

Approaches

1) Using 2 loops check for every pair
↳ $O(n^2)$

2) Sort + B.S $\Rightarrow O(n \log n)$
↳ $O(n \log n)$

2, 3, 5, 10, 50, 80 78
↑

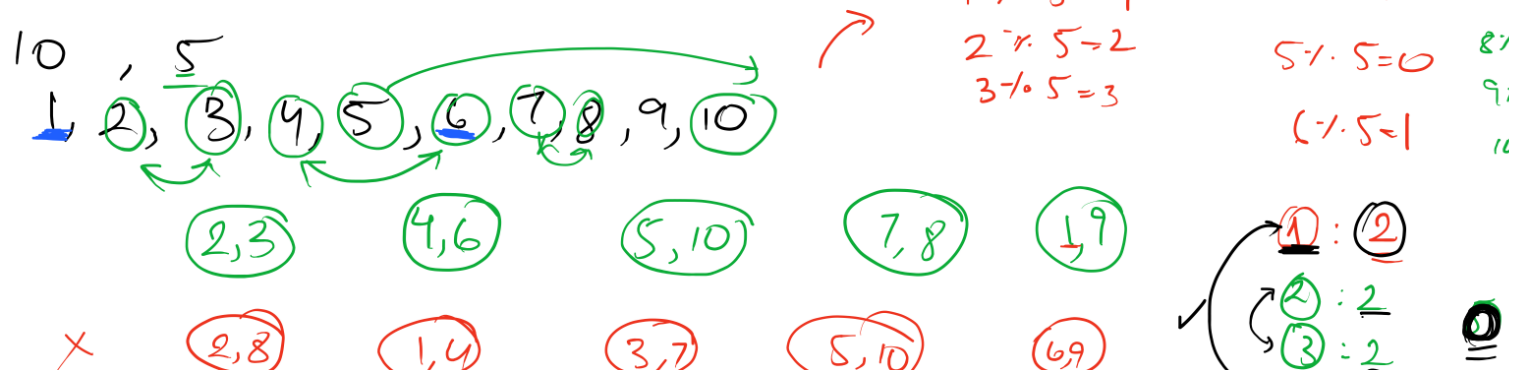
BS to find X = $2 + \text{diff} = 2 + 78 = 80$
↓
 $n \times \log n$ $\hookrightarrow \text{diff} = \underline{\underline{X - 2}}$

✓ 2) Sort $\Rightarrow O(n)$

We can keep a track of visited numbers
 Add to the set & compare
 { 2, 3, 5, 10, 50 }
 Search, add $\rightarrow O(1)$
 $80 \rightarrow 80 + 78 = 158 \times$
 $80 - 78 = 2$
 1st 2nd
 $\downarrow \quad \downarrow$
 $X - Y = \text{Diff}$
 $X = \text{Diff} + Y = \text{Diff} + \text{arr}[i]$
 $Y = X - \text{Diff} = \text{arr}[i] - \text{Diff}$

$3 \rightarrow \begin{cases} 3 + 78 = 81 \\ 3 - 78 = -75 \end{cases}$
 $5 \rightarrow \begin{cases} 5 + 78 = 83 \times \\ 5 - 78 = -73 \times \end{cases}$
 $10 \rightarrow \begin{cases} 10 + 78 = 88 \times \\ 10 - 78 = -68 \times \end{cases}$

Array Pairs divisible by K



Sum of pair should be divisible by K

(K-key)

$$\underline{1} \rightarrow 4, 9$$

$$\underline{5-1} \Rightarrow \underline{4}$$

$$5 \Rightarrow 4 + 1$$

$$\text{Dividend} = \text{Divisor} \times \text{Quotient} + \text{Remainder}$$

$$\underline{9} = \underline{5 \times 1 + 4}$$

$$\underline{4} = 5 \times 0 + 4$$

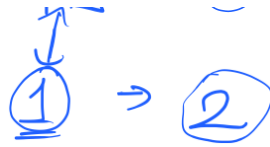
$$\underline{A} \rightarrow \underline{2}$$

$$1 = 5 \times 0 + \underline{1}$$

$$\underline{6} = 5 \times 1 + \underline{1}$$

$$\underline{5-1}$$

$$\underline{(9, 6), (4, 1)}$$



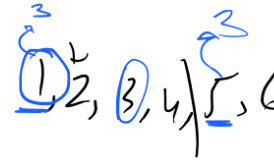
$\rightarrow (9,1), (4,6)$

```

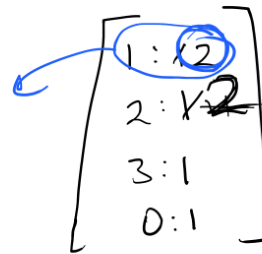
public boolean arrayPairs(int[] arr, int k) {
    HashMap<Integer,Integer> map = new HashMap<>();
    for(int i=0; i<arr.length; i++) {
        int r = arr[i];
        int freq = map.getOrDefault(r,0)+1;
        map.put(r,freq);
    }
    for(Map.Entry<Integer,Integer> entry: map.entrySet()){
        int rem = entry.getKey();
        int count = entry.getValue();

        if(rem==0){
            // remainder=0 need to handle specially
            if(count%2==1){ // count needs to be even
                return false;
            }
        } else {
            int opposite = k - rem;
            int countOfOpposite = map.getOrDefault(opposite,-1);
            if(count != countOfOpposite){
                return false;
            }
        }
    }
    return true;
}

```



$k = 4$



~~2~~

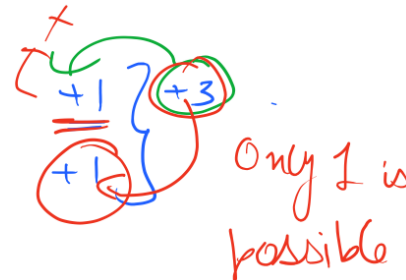
~~4~~ ~~4~~

$4 = 5 \cdot 1 \cdot 4$
 $= 1$

$4 = 6 \cdot 1 \cdot 4$
 $= 2$

$1 = 4 \times 0 + 1 + 3$

$5 = 4 \times 1 + 1 + 3$



Only 1 is possible

5 10
1 2 3 4 5 6

Largest Subarray with sum 0

15, -2, 2, -8, 1, 7, 10, 23

1) To use 2 loops and generate all the possible subarrays.

↳ If any subarray is found with 0 sum

↳ store it's length

↳ return the maximum length

$O(n^2)$

2)

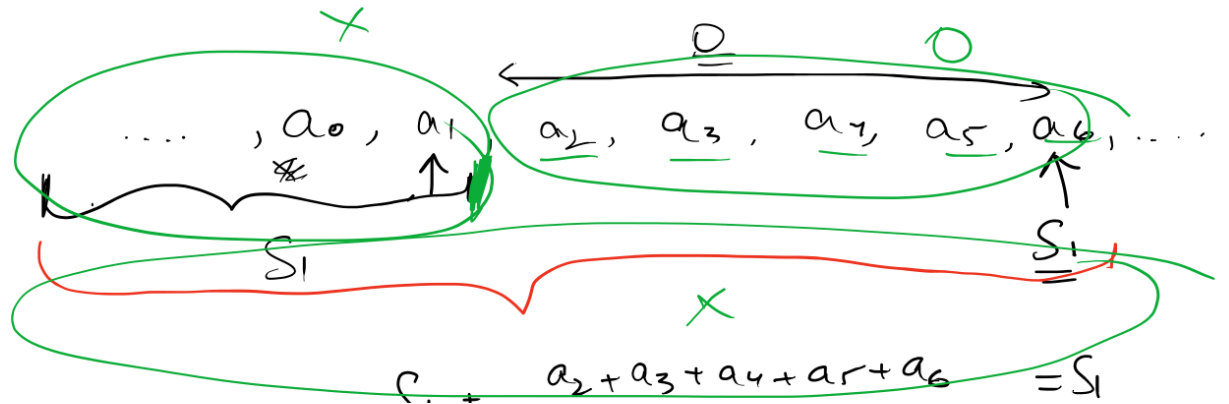
~~Sort~~ + ?

Subarray

-8, -2, 1, 2, 7, 10, 15, 23

$X + 0 = X$

15, -2, 2, -8, 1, 7, 10, 23



21

$\hookrightarrow a_2 + a_3 + a_4 + a_5 + a_6 = 0$

~~15~~ ~~-2~~ ~~2~~ ~~-8~~ ~~1~~ ~~7~~ 10, 23 ✓
 ✓ ✓ ✓ ✓ ✓ ✓
 ✓ ✓ ✓ ✓ ✓ ✓

Sum = 0 15 - 2 = 13 + 2 = 15 - 8 = 7 + 1 = 8 + 7
 len = 0
 2, 5 - 0 = 5
 = 15 + 10
 = 25
 23
 48

Map < Sum, Index >

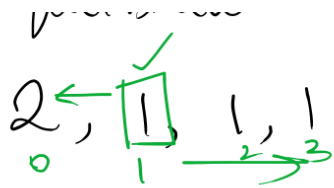
0	-1
15	0
13	1
7	3
8	4
25	6
48	7

[2, -2]
0 1

sum = 2 - 2 = 0
 len = (0, 1 - (-1))
 = (0, 2)

[0, -1]
2: 0

Equilibrium Index



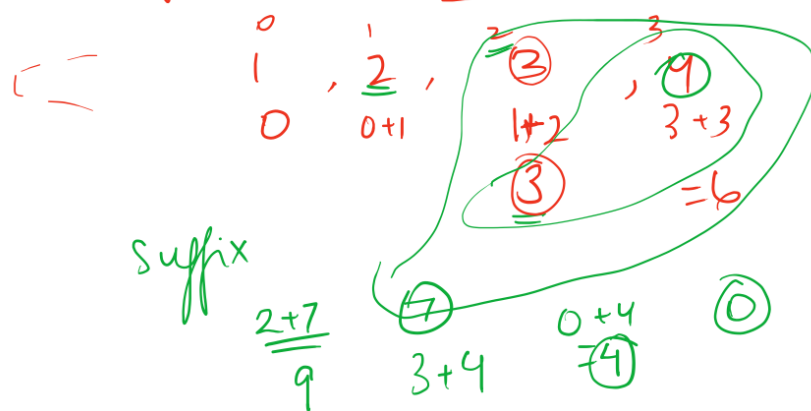
$2 \rightarrow \text{Sum of } [0 \rightarrow i-1] \rightarrow \text{Prefix Sum}$
 $\rightarrow \text{Sum } [i+1, n-1] \rightarrow \text{Suffix Sum}$

1) Using 2 loops

$i \rightarrow$
 $j = 0 \rightarrow i-1$
 $k = i+1 \rightarrow n-1$

2) Using Prefix Sum & Suffix Sum

Prefix Sum $\rightarrow PS[i] = PS[i-1] + arr[i-1]$



\hookrightarrow Calculating the previous sum in $O(1)$

