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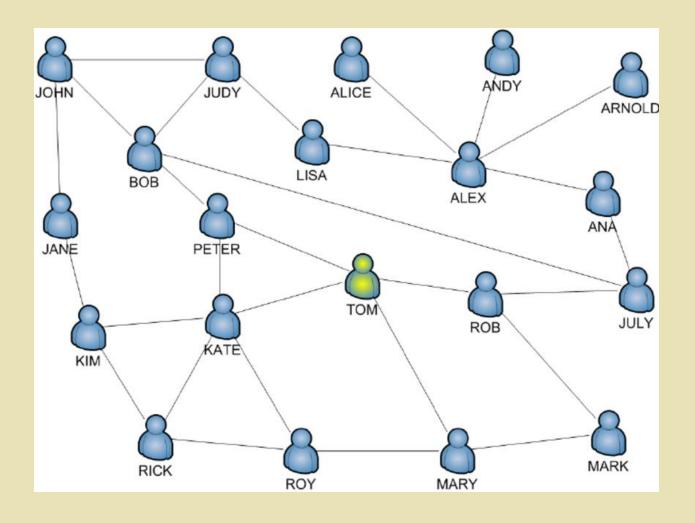
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 - Adjacency list
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 - Breadth-first search (BFS)
 - Depth-first search (DFS)
- Directed Graphs
- Weighted Graphs



Learning Objectives

- 1. Apply graph terminology to describe a graph
- 2. Given a graph, represent it using the adjacency list and adjacency matrix
- 3. Describe the advantages and disadvantages of adjacency list and adjacency matrix representation
- 4. Name different types of graphs introduced
- 5. Understand and implement the breadth-first search and depth-first search
- 6. Calculate the path lengths in a graph

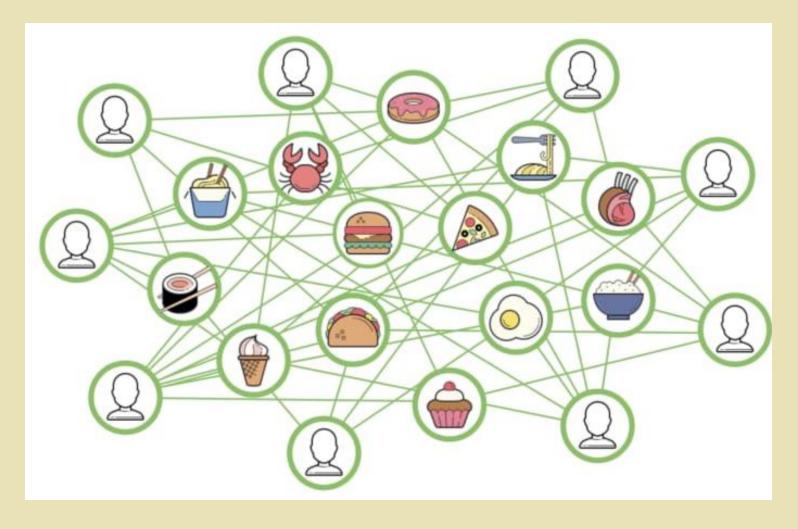




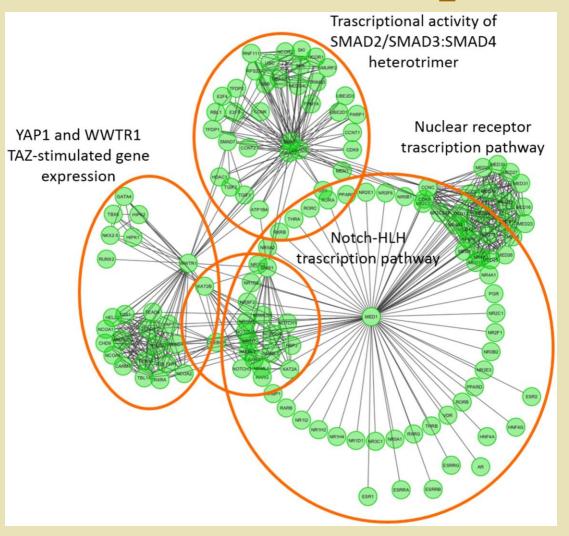


- The small-world experiment comprised several experiments conducted by Stanley Milgram and other researchers examining the average path length for social networks of people in the United States.
- The research was groundbreaking in that it suggested that human society is a small-worldtype network characterized by short path-lengths

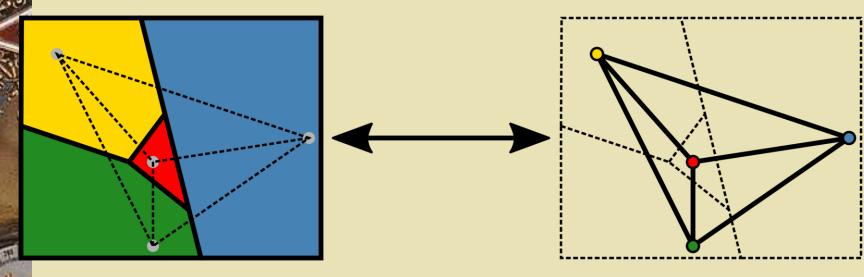








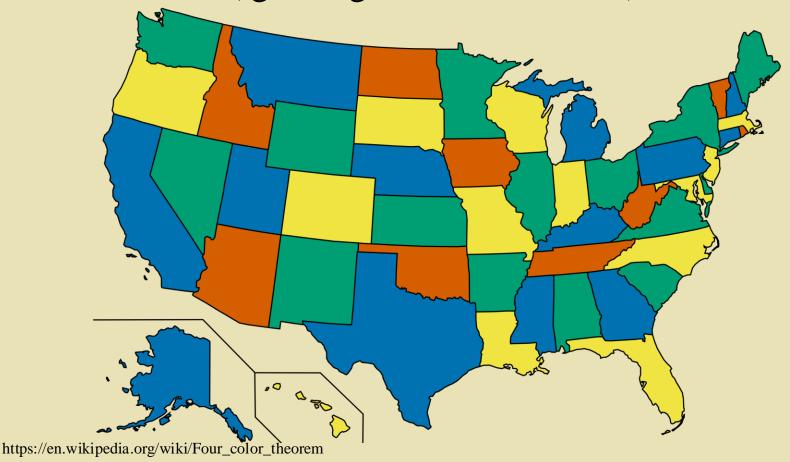
• Four color map theorem, states that no more than four colors are required to color the regions of any map so that no two adjacent regions have the same color.



a **coloring** of a graph almost always refers to a *proper vertex coloring*.



• Example: A four-colored map of the states of the United States (ignoring lakes and oceans)





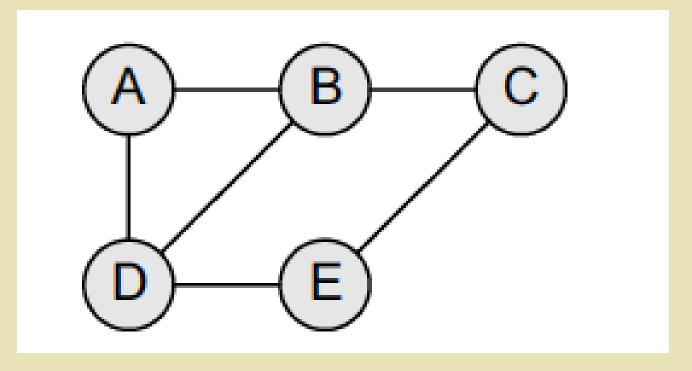
Definition

- A graph is a non-linear data structure, a collection of nodes (vertices) connected to each other via edges.
- A graph can be defined as G = (V, E)
 - V is a set of vertices $\{v_1, v_2, \dots v_n\}$
 - E is a set of edges, $e_1 = (v_1, v_2), e_2 = (v_1, v_3), ...$
 - Edges can be directed or undirected (more on this later)
- A graph is often viewed as a generalization of the tree structure



Definition -- Example

V(G) = {A, B, C, D and E} and E(G) = {(A, B),
 (B, C), (A, D), (B, D), (D, E), (C, E)}.



Is it a directed graph or undirected graph?



Definition -- Exercise

Given V(G) = {W, X, Y, Z} and E(G) = {(W, X), (W, Y), (W, Z), (Y, Z)}. Draw the undirected graph.





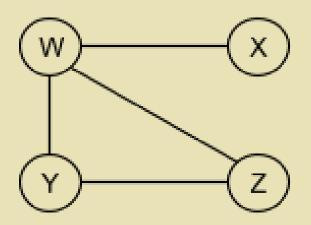
Terminology

- Two nodes are **adjacent** if there is an edge connecting them.
- The **degree** of a node is the number of edges connected to it.
- A path is a sequence of edges leading from a source (starting) vertex to a destination (ending) vertex. The path length is the number of edges in the path.



Terminology

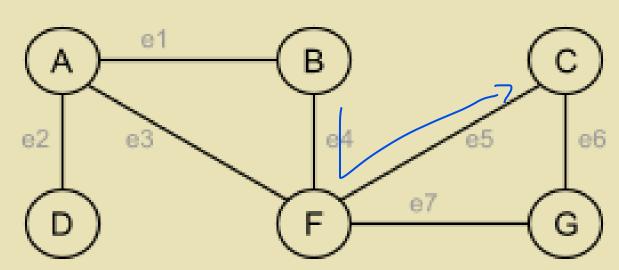
- A cycle is a path that starts and ends at the same vertex, forming a closed loop.
- The **distance** between two vertices is the number of edges on the shortest path between those vertices.





Definition -- Exercise

- 1. Are A and B adjacent?
- 2. Is there a cycle in this graph? \square
- 3. Name a path (vertices visited) from B to C.
- 4. What is the distance between B and C?





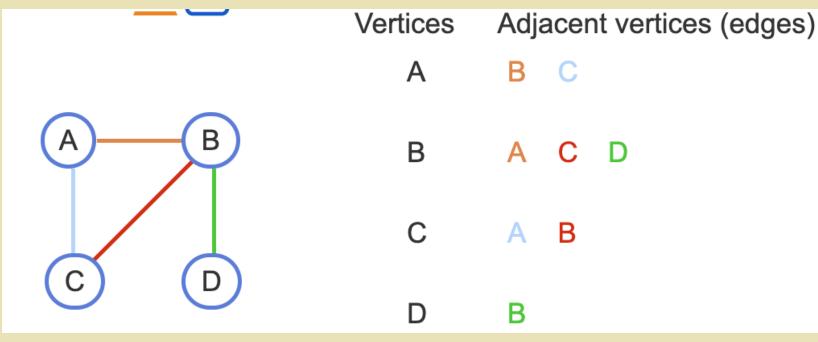
Terminology

- Types of Graph
 - Directed Graph (Digraph): Edges have a direction, indicating a one-way relationship.
 - Undirected Graph: Edges have no direction, representing mutual relationships.
 - Weighted Graph: Each edge has a weight or cost associated with it, often used in optimization problems.
 - Cyclic Graph: Contains at least one cycle (a path that starts and ends at the same node).
 - Acyclic Graph: Does not contain any cycles.
 - Connected Graph: There is a path between every pair of nodes.



Graph representations: Adjacency lists

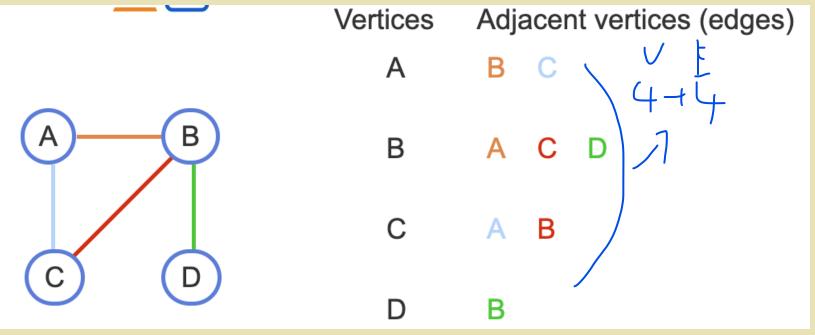
• In an **adjacency list** graph representation, each vertex has a list of adjacent vertices, each list item representing an edge.





Graph representations: Adjacency lists

- A key advantage of an adjacency list graph representation is a size of O(V + E)
 - each vertex appears once
 - each edge appears twice





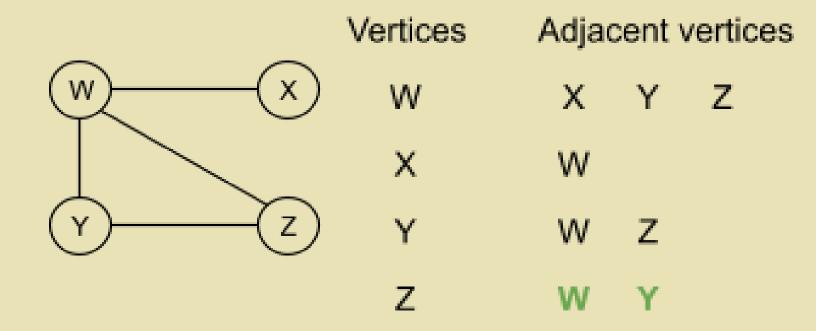
Graph representations: Adjacency lists

- A disadvantage is that determining whether two vertices are adjacent is O(V)
 - must be traversed all other vertices to check if they are connected to this vertex, and that list could have V items.
 - However, in most applications, a vertex is only adjacent to a small fraction of the other vertices, yielding a sparse graph.
- A sparse graph has far fewer edges than the maximum possible.



Adjacency lists Exercise

 Write down the adjacency lists of the following graph

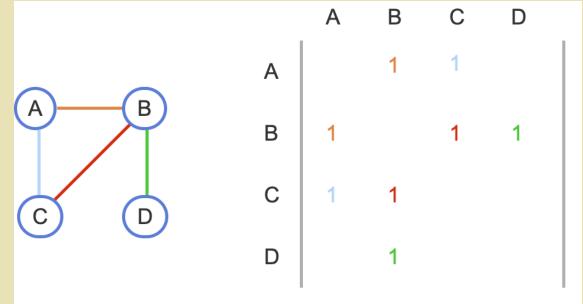




Graph representations: Adjacency matrix

• In an adjacency matrix graph representation, each vertex is assigned to a matrix row and column, and a matrix element is 1 if the corresponding two vertices have an edge or is 0

otherwise.





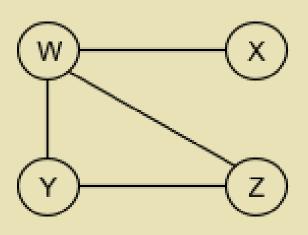
Graph representations: Adjacency matrix

- Assuming the common implementation as a two-dimensional array whose elements are accessible in O(1), then an adjacency matrix's key benefit is O(1) determination of whether two vertices are adjacent: The corresponding element is just checked for 0 or 1.
- A key drawback is $O(V^2)$ size.
 - Ex: A graph with 1000 vertices would require a 1000 x 1000 matrix, meaning 1,000,000 elements. An adjacency matrix's large size is inefficient for a sparse graph, in which most elements would be 0's.



Adjacency Matrix Exercise

 Write down the adjacency matrix of the following graph



	W	X	Y	Z
W	0	1	1	1
X	1	0	0	0
Y	1	0	0	1
Z	1	0	1	0



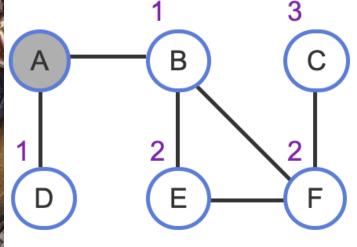
Adjacency Matrix Exercise

- Given the adjacency matrix,
 - 1. How many edges does D have?
 - 2. Is there a path from A to E? \nearrow

	Α	В	С	D	E
Α	0	1	1	1	0
В	1	0	1	1	1
С	1	1	0	1	1
D	1	1	1	0	0
Е	0	1	1	0	0



• A breadth-first search (BFS) is a traversal that visits a starting vertex, then all vertices of distance 1 from that vertex, then of distance 2, and so on, without revisiting a vertex.



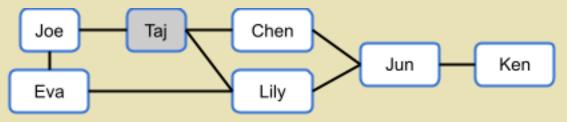
Starting at A: A B D E F C



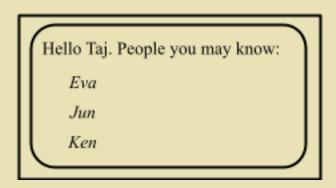
- Example: Social networking connection recommender
 - Social networking sites like Facebook and LinkedIn use graphs to represent connections among people. For a particular user, a site may wish to recommend new connections. One approach does a breadth-first search starting from the user, recommending new connections starting at distance 2 (distance 1 people are already connected with the user).



 Example: Social networking connection recommender



Breadth-first	Taj	Joe	Chen	Lily	Eva	Jun	Ken
traversal:	0	1	1	1	2	2	3



```
BFS(startVertex) {
  initialize queue with startVertex
  initialize discoveredSet with startVertex
```

Graphs: Breadthfirst search algorithm

```
while (queue is not empty) {
   currentVertex = dequeue from queue
   // Visit or process the current vertex
   process(currentVertex) // can be print or others
   // Explore each adjacent vertex of the current vertex
   for each neighbor of currentVertex {
     if (neighbor is not in discoveredSet) {
       // Enqueue undiscovered neighbor and mark it as
discovered
       enqueue neighbor in queue
       add neighbor to discoveredSet
```



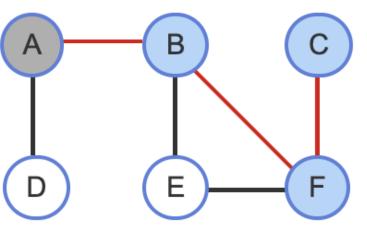
- When the BFS algorithm first encounters a vertex, that vertex is said to have been **discovered**.
- In the BFS algorithm, the vertices in the queue are called the **frontier**, being vertices thus far discovered but not yet visited.
- Because each vertex is visited at most once, an already-discovered vertex is not enqueued again.
- A "visit" may mean to print the vertex, append the vertex to a list, compare vertex data to a value and return the vertex if found, etc.



• A depth-first search (DFS) is a traversal that visits a starting vertex, then visits every vertex along each path starting from that vertex to the path's end before backtracking.



Example

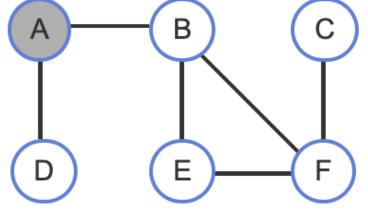


Starting at A: A B F C

Starting at A descends along a path to the path's end before backtracking.



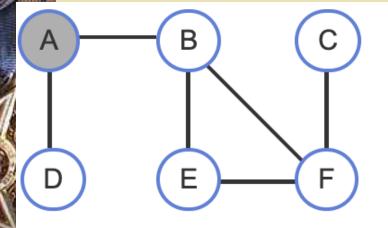
Example



Starting at A: A B F C E

Reach path's end: Backtrack to F, visit F's other adjacent vertex (E). B already visited, backtrack again.

Example



Starting at A: A B F C E D

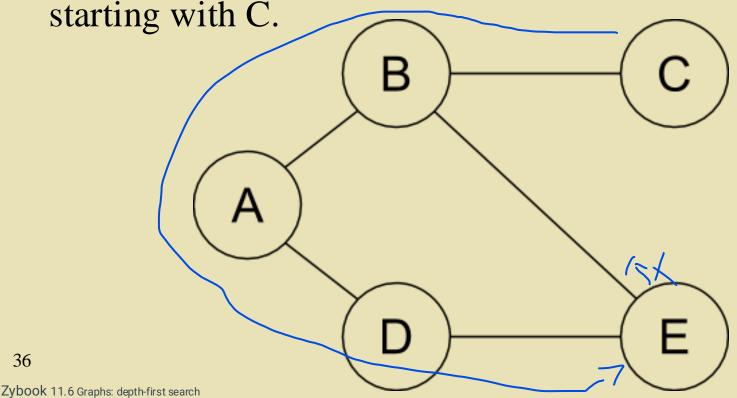
Backtracked all the way to A. Visit A's other adjacent vertex (D). No other adjacent vertices: Done.



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Graphs: Depth-first search

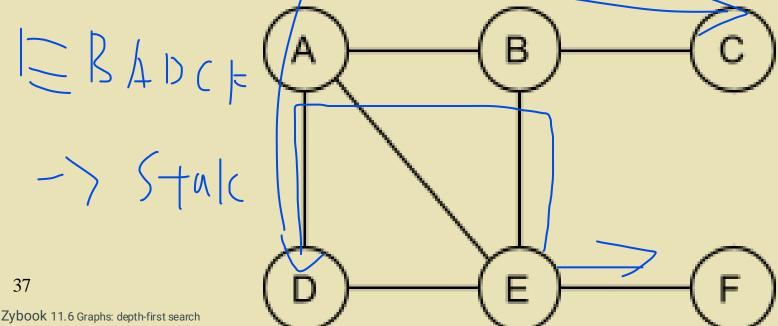
• Exercise: Perform a valid depth-first search of the graph below. Assume the starting vertex is C. Write down the order of the vertices visited,





• Exercise: Perform a depth-first search of the graph below. Assume the starting vertex is E. Write down the order of vertices visited (start with E).

- When a vertex has multiple connected vertices, the algorithm visit them using alphabet order.



```
DFS(startVertex) {
                      initialize stack with start Vertex
                      while (stack is not empty) {
                        currentVertex = pop from stack
                        if (currentVertex is not in visitedSet) {
                          // Visit or process the current vertex
                          process(currentVertex)
Depth-first
                          // Mark the current vertex as visited
                          add currentVertex to visitedSet
                          // Push each unvisited adjacent vertex to the stack
                          for each neighbor of currentVertex {
                            if (neighbor is not in visitedSet) {
                              push neighbor to stack
```

Graphs:

search

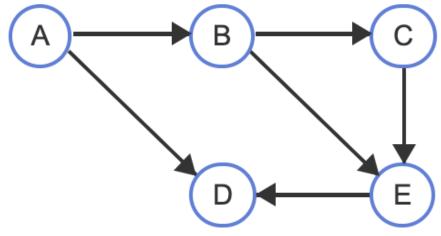
```
RecursiveDFS(vertex) {
                    // If the vertex has not been visited yet
                    if (vertex is not in visitedSet) {
                      // Mark the vertex as visited
                      add vertex to visitedSet
  Graphs:
                      // Visit or process the current vertex
Depth-first
                      process(vertex)
    search
                      // Recursively visit each adjacent vertex
                      for each neighbor of vertex {
recuersive
                       RecursiveDFS(neighbor)
 algorithm }
```



Directed Graphs

Example

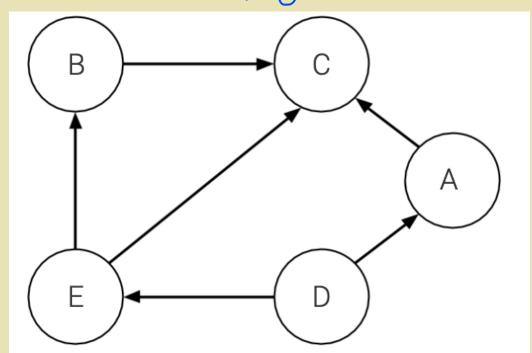






Directed Graphs

- Exercise
 - 1. Which vertices are adjacent to D?
 - 2. B is adjacent to <u>C</u>.
 - 3. Are there circles in this graph? $\sqrt{}$

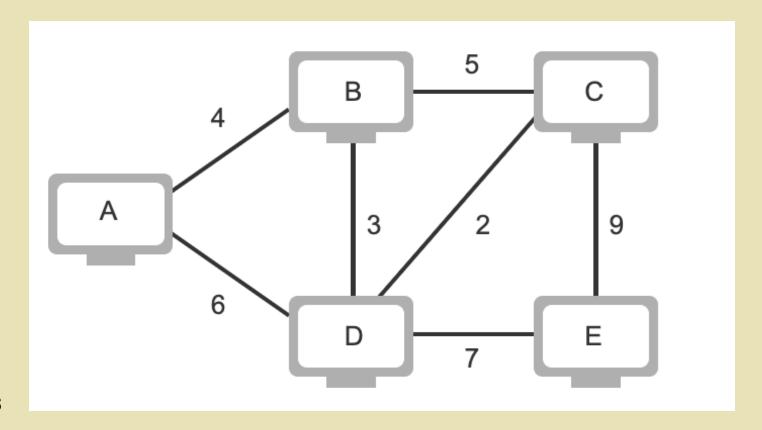




- A weighted graph associates a weight (w) with each edge.
 - E.g. e=(A,B), w_e=5
- A graph edge's weight, or cost, represents some numerical value between vertex items, such as flight cost between airports, connection speed between computers, or travel time between cities.
- A weighted graph may be directed or undirected.
- In a weighted graph, the **path length** is the sum of the edge weights in the path.
 - Find shortest path problem
- Unweighted graph can be viewed as w=1 for all edges.

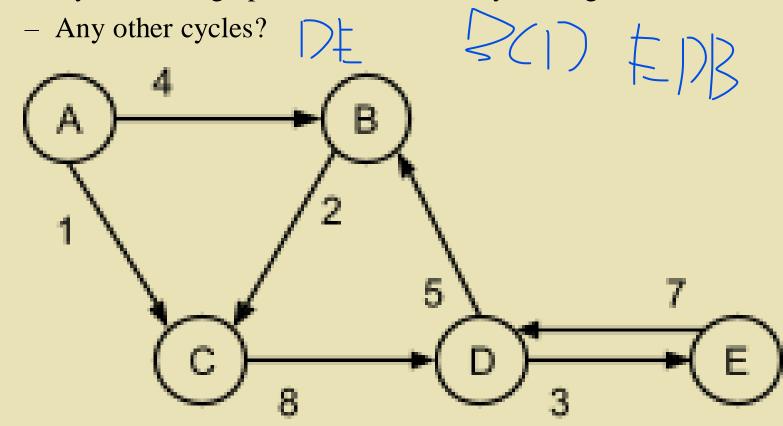


- Examples
 - Path A-B-C, the path length is 9



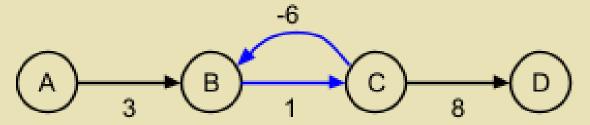


- Directed weighted graph example
 - Cycle in this graph $B \rightarrow C \rightarrow D \rightarrow B$, cycle length is 15.





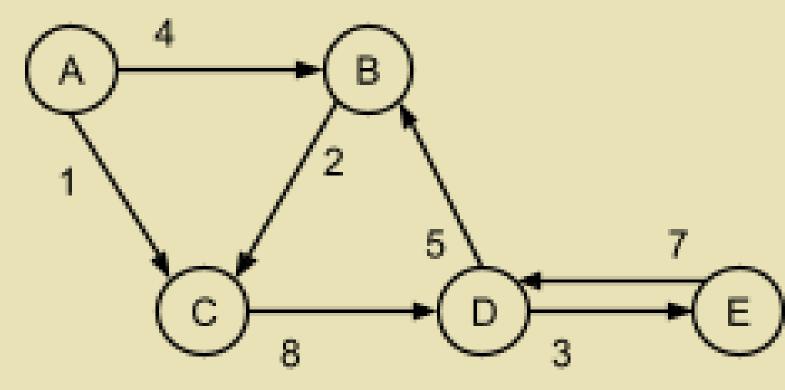
- The **cycle length** is the sum of the edge weights in a cycle.
- A negative edge weight cycle has a cycle length less than 0.
- A shortest path does not exist in a graph with a negative edge weight cycle
 - each loop around the negative edge weight cycle further decreases the cycle length, so no minimum exists.





Weighted Graphs -- Exercise

- Find a path from A to D \(\) \(\) = 9
- Fina the shortest path from A to D.





References and Useful Resources



