## Operating Systems: Processes CIS\*3110: Operating Systems

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# Mind-bending POSIX tools to create and manage processes

wait(2): wait for a child process to complete, and find out what its exit status was

exec(3): take a running process image, and load up a new program's text and data into it, resetting the stack.

#### fork(2)

- - "returns twice" return status indicates whether the program is still in the parent, or now in the child

```
int pid;
int childstatus;
pid = fork();
if (pid < 0) {
     ... something bad happened ...
} else if (pid == 0) {
     ...do whatever child is to do ...
     exit(someStatusReportedFromChild);
} else {
     ...do whatever parent is to do, ...and eventually call:
     pid = wait(&childstatus);
```

#### wait(2)

wait(2): wait for a child process to complete, and find out what its exit status was

- pid\_t wait(int \*status) waits for child to return; once this happens, returns the child PID, and status will be loaded with child's exit status.
- interpret \*status via:
  - WIFEXITED(status) did child crash?

#### wait(2)

#### wait(2): has additional friends:

- pid\_t waitpid(pid\_t wpid, int \*status, int options) provides greater control than simple wait(int \*status);
  - the pid\_t wpid argument restricts what to wait for:
    - -1 wait for anyone (just like wait())
    - wait for anyone in our process group (more on this later)
    - > 0 wait specifically for the process with this PID
  - the options argument is formed of a bitwise OR of 0 and several choices including:

NOHANG do not **block** when there are no processes ready to report status

If NOHANG is used, and there is nothing to wait for, waitpid() returns a 0.

## exec(3)

- exec(3): take a running process image, and load up a new program's text and data into it, resetting the stack.
  - if it succeeds, our program is gone, so only returns an error

```
int pid;
int childstatus;

pid = fork();
if (pid < 0) {
    ...something bad happened ...
} else if (pid == 0) {
    exec << ? >>(...)
    exit(somestatus);
}
```

#### The exec(3) family: exec1\*

```
execl(3): take a list of arguments to turn into argv,
           terminated by a NULL
           "l" ("ell") \rightarrow "list of arguments"
execl(char *fullpathtoprogram, char *arg, ...)
  this runs the indicated program
execlp(char *nameofprogram, char *arg, ...)
  this will search the path for the program by name
execle(char *path, char *arg, ..., char *envp[])
  this runs the indicated program, with the indicated en-
  vironment variable list
```

#### **Example:**

```
execl("/bin/ls", "ls", "-l", "myprog.c", NULL);
execlp("ls", "ls", "-l", "myprog.c", NULL);
```

## The exec(3) family: execv\*

```
execv(3): take as a single argument a vector (array) of
           strings, which will be terminated by a NULL, and
           use it as argv
           "v" \rightarrow "vector (i.e.; an array) of arguments"
execv(char *fullpathtoprogram, char *argv[])
  this runs the indicated program
execvp(char *nameofprogram, char *argv[])
  this will search the path for the program by name
execve(char *path, char *argv[], char *envp[])
  this runs the indicated program, with the indicated
  environment variable list (and all the others call
  execve(2) internally)
```

#### **Example:**

```
char *arglist[] = { "ls", "-l", "myprog.c", NULL };
execv("/bin/ls", arglist);
```

## Threads versus processes

Process	Thread	Notes
fork()	<pre>pthread_create()</pre>	Creates a new <b>context</b> structure with new PC & other registers
exec()	No meaningful (or possible) equivalent.	
exit()	<pre>pthread_exit()</pre>	Cleans up the thread <b>context</b> .  No exit status provided (or needed).
wait()	<pre>pthread_join()</pre>	Synchronizes parent thread. Waits until the indicated child calls pthread_exit()