# A FULLY-AUTOMATED SOLVER FOR MULTIPLE SQUARE JIGSAW PUZZLES USING HIERARCHICAL CLUSTERING

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by

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# The Designated Thesis Committee Approves the Thesis Titled

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#### ABSTRACT

# A Fully-Automated Solver for Multiple Square Jigsaw Puzzles Using Hierarchical Clustering

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Square jigsaw puzzle solving is a variant of a standard jigsaw puzzles, wherein all pieces are equal-sized squares that must be placed adjacent to one another to match an original image. This thesis proposes the first algorithm that can solve multiple jigsaw puzzles simultaneously without any information provided to it beyond the set of pieces. The algorithm has been able to assemble up to 10 images simultaneously with results comparable to the state-of-the-art solver working on each image individually.

This thesis also defines the first set of metrics specifically tailored for multiple puzzle solvers. It also outlines the first visualization standards for viewing solver solutions.

# **DEDICATION**

To my mother.

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#### CHAPTER 1

#### Introduction

Jigsaw puzzles were first introduced in the 1760s when they were made from wood; their name derives from the jigsaws that were used to carve the wooden pieces. The 1930s saw the introduction of the modern jigsaw puzzle where an image was printed on a cardboard sheet that was cut into a set of interlocking pieces [1, 2]. Although jigsaw puzzles had been solved by children for more than two centuries, it was not until 1964 that the first automated jigsaw puzzle solver was proposed by Freeman & Gardner [3]. While an automated jigsaw puzzle solver may seem trivial, the problem has been shown by Altman [4] and Demaine & Demaine [5] to be strongly NP-complete when pairwise compatibility between pieces is not a reliable metric for determining adjacency.

Jig swap puzzles are a specific type of jigsaw puzzle where all pieces are equal sized, non-overlapping squares; an example of these puzzles is shown in Figure 1. Jig swap puzzles<sup>1</sup> are substantially more challenging to solve than traditional jigsaw puzzles since piece shape cannot be considered when determining affinity between pieces. Rather, only the image information on each individual piece is used when solving the puzzle.

Solving a jigsaw puzzle simplifies to reconstructing an object from a set of component pieces. As such, techniques developed for jigsaw puzzles can be generalized to many practical problems. Examples where jigsaw puzzle solving strategies have been used include: reassembly of archaeological artifacts [6, 7],

<sup>&</sup>lt;sup>1</sup>Throughout this thesis, the term "jigsaw puzzle" is used. Almost exclusively, this refers to specifically jig swap puzzles.





(a) Ground-Truth Image

(b) Randomized Jig Swap Puzzle

Figure 1: Jig Swap Puzzle Example

forensic analysis of deleted files [8], image editing [9], reconstruction of shredded documents [10], DNA fragment reassembly [11], and speech descrambling [12]. In most of these practical applications, the original, also known as "ground-truth," input is unknown. This significantly increases the difficulty of the problem as the structure of the complete solution must be determined solely from the bag of component pieces.

This thesis describes the first fully-automated solver for the simultaneous assembly of multiple jigsaw puzzles; unlike all previous solvers, this thesis' algorithm does not require any human input concerning the attributes of the input set of pieces (e.g., the the number of ground-truth (i.e., original) puzzles). What is more, this thesis defines a set of new metrics specifically tailored to quantify the quality of outputs of multi-puzzle solvers. Lastly, it outlines a set of standards for visualizing the accuracy and quality of solver outputs.

#### CHAPTER 2

#### **Previous Work**

Computational jigsaw puzzle solvers have been studied since the 1960s when Freeman & Gardner proposed a solver that relied only on piece shape and could puzzles with up to nine pieces [3]. Since then, the focus of research has gradually shifted from traditional jigsaw puzzles to jig swap puzzles.

Cho et al. [13] proposed in 2010 one of the first modern computational jig swap puzzle solvers; their approach relied on a graphical model built around a set of one or more "anchor piece(s)," which are pieces whose position is fixed in the correct location before the solver began. Cho et al.'s solver also required that the user specify the puzzle's actual dimensions. Future solvers would improve on Cho et al.'s results while simultaneously reducing the amount of information (i.e., beyond the set of pieces) passed to the solver.

A significant contribution of Cho et~al. is that they were first to use the LAB (<u>Lightness</u> and the <u>A/B</u> opponent color dimensions) colorspace to encode image pixels. LAB was selected due to its property of normalizing the lightness and color variation across all three pixel dimensions. Cho et~al. also proposed a measure for quantifying the pairwise distance between two puzzle pieces that became the basis of most future work.

Pomeranz et al. [14] proposed an iterative, greedy jig swap puzzle solver in 2011. Their approach did not rely on anchor pieces, and the only information passed to the solver were the pieces, their orientation, and the size of the puzzle. Pomeranz et al. also generalized and improved on Cho et al.'s inter-piece pairwise distance measure by

proposing a "predictive distance measure." Finally, Pomeranz *et al.* introduced the concept of "best buddies". Equation ?? formally define the best buddy relationship between the side (e.g., top, left, right, bottom),  $s_x$ , of puzzle piece,  $p_i$ , and the side,  $s_y$ , of piece  $p_j$ . Note that  $C(p_i, s_x, p_j, s_y)$  represents the compatibility (i.e., similarity) between the compatibility (i.e., similarity) between the two piece's respective sides.

$$\forall x_k \forall s_c, C(x_i, s_a, x_j, s_b) \ge C(x_i, s_a, x_k, s_c)$$
and
$$(1)$$

$$\forall x_k \forall s_c, C(x_j, s_b, x_i, s_a) \ge C(x_j, s_b, x_k, s_c)$$

Best buddies have served as both an estimation metric for the quality of solver output as well as the foundation of some solvers' assemblers [15].

An additional key contribution of Pomeranz *et al.* is the creation of three image benchmarks. The first benchmark is comprised of twenty 805 piece images; the sizes of the images in the second and third benchmarks are 2,360 and 3,300 pieces respectively.

In 2012, Gallagher [16] formally categorized jig swap puzzles into four primary types. The following is Gallagher's proposed terminology; his nomenclature is used throughout this thesis.

• Type 1 Puzzle: The dimensions of the puzzle (i.e., the width and height of the ground-truth image in number of pixels) is known. The orientation/rotation of each piece is also known, which means that there are exactly four pairwise relationships between any two pieces. A single anchor piece, with a known, correct, location is required with additional anchor pieces being optional. This type of puzzle is the focus of [13, 14].

- Type 2 Puzzle: This is an extension of a Type 1 puzzle, where pieces may be rotated in 90° increments (e.g., 0°, 90°, 180°, or 270°); in comparison to a Type 1 puzzle, this change alone increases the number of possible solutions by a factor of 4<sup>n</sup>, where n is the number of puzzle pieces. What is more, no piece locations are known in advance; this change eliminates the use of anchor piece(s). Lastly, the dimensions of the ground-truth image may be unknown.
- Type 3 Puzzle: All puzzle piece locations are known and only the rotation of the pieces is unknown. This is the least computationally complex of the puzzle variants and is generally considered the least interesting. Type 3 puzzles are not explored as part of this thesis.
- Mixed-Bag Puzzles: The input set of pieces are from multiple puzzles, or there are extra pieces in the input set that belong to no puzzle. The solver may output either a single, merged puzzle, or it may separate the input pieces into disjoint sets that ideally align with the set of ground-truth puzzles. This type of puzzle is the primary focus of this thesis.

Sholomon et al. [17] in 2013 proposed a genetic algorithm based solver for Type 1 puzzles. By moving away from the greedy approach used by Pomeranz et al., Sholomon et al.'s approach is more immune to suboptimal decisions early in the placement process. Sholomon et al.'s algorithm is able to solve puzzles of significantly larger size than previous techniques (e.g., greater than 23,000 pieces). What is more, Sholomon et al. defined three new large image benchmarks; the specific puzzle sizes are 5,015, 10,375, and 22,834 pieces [18].

Paikin & Tal [15] published in 2015 a greedy solver that handles both Type 1 and Type 2 puzzles, even if those puzzles are missing pieces. What is more, their

algorithm is one of the first to support Mixed-Bag Puzzles. While Paikin & Tal's algorithm represents the current state-of-the-art, it has serious limitations:

•

These limitations are addressed by this thesis' Mixed-Bag Solver. Since Paikin & Tal's algorithm represents the current state of the art, it is used as thesis' assembler. What is more, it is used as the baseline of performance.

#### CHAPTER 3

# Mixed-Bag Solver Overview

This thesis' Mixed-Bag Solver supports Type 1, Type 2, and Mixed-Bag Puzzles. It consists of five distinct stages namely: segmentation, stitching, hierarchical clustering of segments, seed piece selection, and final assembly. The flow of the algorithm is shown in Figure 2; the pseudocode for the solver, including the input(s) and output of each stage is shown in Algorithm 1.

The following subsections describe each of Mixed Bag Solver's stages/subfunctions. It also discusses the assembler, which is a separate but associated component of the architecture.

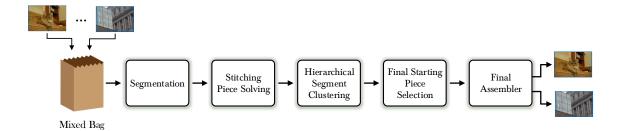


Figure 2: Components of the Mixed-Bag Puzzle Solver

# **Algorithm 1** Pseudocode for the Mixed Bag Solver

- 1: **function** MIXEDBAGSOLVER(all pieces)
- 2:  $solved segments \leftarrow Segmentation(all pieces)$
- 3:  $overlap \ matrix \leftarrow Stitching(solved \ segments, \ all \ pieces)$
- 4:  $segment\ clusters \leftarrow PerformHierarchicalClustering(solved\ segments,\ overlap$
- 5:  $puzzle \ \overline{start} \ pieces \leftarrow FINDSTARTINGPIECES(segment_clusters)$
- 6:  $solved puzzles \leftarrow RunFinalAssembly(puzzle start pieces, all pieces)$
- 7: **return** solved puzzles

#### 3.1 Assembler

The assembler assigns the placement (and optionally rotation) of the puzzle pieces in the solved puzzle. This solver architecture is largely independent of the particular assembler used. Hence, any improvements or modifications to the assembler can be directly incorporated into the Mixed-Bag Solver to improve its performance. What is more, if a particular solver performs better than another for a particular application, the solvers can be interchanged. This provides the Mixed-Bag Solver with significant flexibility and upgradability to maximize performance across a wide range of applications.

For all experiments in this thesis, we used the solver proposed by Paikin & Tal [15]. It was selected because it the current state-of-the-art. What is more, since it supports mixed bag puzzles, it can be used for direct comparison of performance.

#### 3.2 Segmentation

As shown in Algorithm 1, segmentation stages takes as input only the bag of puzzle pieces created from the original images; unlike all other solvers to date, this algorithm takes no other inputs. The role of stage is to provide structure to the unordered input. This is done by partitioning the pieces into disjoint sets, referred to here as segments. These segments are groups of puzzle pieces where there is a high degree of confidence that the pieces are assembled correctly.

Algorithm 2 outlines the basic segmentation framework. The algorithm is iterative and will have one or more rounds. In each round, all pieces that have not yet been assigned to a segment are assembled as if they all belong to a single ground truth image. This is done as it eliminates the need to make any assumptions regarding the input at this early stage of the solver. While it may lead to a single

puzzle being divided into multiple segments, these segments can be merged together later in a later stage of the solver.

After the single puzzle is assembled in each round, the solved puzzle is then divided into segments; the procedure for this is described in Section . Assuming the largest segment exceeds the minimum allowed size<sup>1</sup>, it is passed to the next stage of the Mixed-Bag Solver.

The term " $\alpha$ " in Algorithm 2 defines which segments other than the largest one constructed puzzle are immediately passed to the next solver stage. In this thesis,  $\alpha$  was set to 0.5, meaning any segment that was at least half the size of the largest segment in that round (and larger than the minimum segment size) is saved. This scalar value provides sufficient balance between ensuring the largest segments for analysis with limiting the execution time of this stage.

Once a piece is assigned to a saved segment, it is removed from the set of unassigned pieces. Hence, those pieces will not be placed in the next segmentation round. Segmentation continues until all pieces have been assigned to sufficiently large segments, or no segment exceeds the minimum allowed segment size.

## 3.2.1 The SegmentPuzzle Function

The segment function shown in Algorithm 3 is adapted from the kernel growing segmentation procedure proposed by Pomeranz *et al.*, where it was shown to have greater than 99.7% accuracy identifying genuine neighbors [14]. The kernel of each segment is a single seed piece.

Whenever a piece is added to a segment, it is removed from the set of

<sup>&</sup>lt;sup>1</sup>For this thesis, it was observed that a minimum segment size of 7 provided the best balance between solution quality and algorithm execution time.

# Algorithm 2 Pseudocode for the Segmentation Algorithm

```
1: function Segmentation(all pieces)
       solved\ segments \leftarrow \{\}
 2:
       unassigned pieces \leftarrow \{all \ pieces\}
 3:
       repeat
 4:
          solved puzzle \leftarrow RunSinglePuzzleAssembly(unassigned pieces)
 5:
 6:
          puzzle \ segments \leftarrow SegmentPuzzle(solved \ puzzle)
 7:
          max segment size \leftarrow maximum size of segment in solved segments
          for each segment \in puzzle segments do
 8:
              if |segment| > \alpha \times max segment size then
9:
                 add segment to solved segments
10:
                 remove pieces in segment from unassigned pieces
11:
       until max segment size < smallest allowed or |unplaced\ pieces| < 0
12:
       return solved segments
13:
```

unassigned pieces. What is more, the algorithm check's all pieces directly adjacent to the added piece. If the adjacent piece and the added piece are "best buddies" (i.e., each is more similar to the other on their respective sides than they are to any other piece as defined by [14] and [15]), then the adjacent piece is also added to the segment. This process continues until no there are no pieces adjacent to a segment member that fulfill the best buddy criteria.

If Pomeranz et al.'s original segmentation algorithm is used for Mixed-Bag puzzles, two correctly assembled segments from different input puzzles can be merged into a single segment. This is usually in the form of narrow bridges no wider than a single piece. This necessitates that each segment goes through post-processing to identify and remove these single point bridges; the procedure for doing this is described in Section 3.2.2.

Once these single piece bridges have been broken, one or more pieces become disconnected from the segment. These pieces are then returned to the set of unassigned pieces to be assigned to a different segment. At the end of segment, each

# **Algorithm 3** Pseudocode for the segment Function

```
1: function SegmentPuzzle(solved puzzle)
       puzzle \ segments \leftarrow \{\}
 2:
       unassigned pieces \leftarrow \{all pieces in solved puzzle\}
 3:
       while |unassigned pieces| > 0 do
 4:
           segment \leftarrow \text{new empty segment}
 5:
 6:
           seed piece \leftarrow next piece in unassigned pieces
           queue \leftarrow [seed\_piece]
 7:
           while |queue| > 0 do
 8:
              piece \leftarrow next piece in queue
 9:
              add piece to segment
10:
              for each neighbor piece of piece do
11:
                  if IsBestBuddies(neighbor piece, piece) then
12:
                     add neighbor piece to queue
13:
                     remove neighbor piece from unassigned pieces
14:
15:
           articulation points \leftarrow FINDARTICULATIONPOINTS(segment)
           remove articulation points from segment
16:
           disconnected\ pieces \leftarrow FINDDISCONNECTEDPIECES(segment, seed\ piece)
17:
           add segment to puzzle segments
18:
19:
           remove disconnected points from segment
           add articulation points and disconnected pieces to unassigned pieces
20:
21:
       return puzzle segments
```

piece in the puzzle will be assigned to exactly one segment.

#### 3.2.2 Articulation Points

A segment can be modeled as a graph with a single connected component. The individual puzzle pieces represent the vertices while the edges are the best buddy relationships as defined in Chapter ?? and determined in Algorithm ?? by the function "ISBESTBUDDIES". An articulation point is any vertex (i.e., puzzle piece)

whose removal increases the number of connected components. The Mixed-Bag Solver uses the algorithm proposed by ?? for identifying articulation points. While most implementations of this algorithm are recursive, this thesis instead uses an iterative approach as segment can be several thousand pieces in sizes. As such, a recursive implementation is prone to stack overflows.

# 3.3 Stiching

As defined previously, a segment represents a partial assembly where there is a particularly high degree of confidence that the placement is correct. During the segmentation stage, puzzle pieces from a single ground truth image segment may be partitioned into multiple segments. If two segments are similar, it is likely that if one segment is allowed to grow that it would eventually merge with its adjacent segment. However, uncontrolled segment expansion can cause a segment to grow beyond the boundaries of its ground truth image and merge with a segment from a different input image. This thesis controls segmentation growth through the use of multiple "mini solvers," which each have a "stitching piece" as their seed. The output from these mini solvers are used to determine inter-segment similarity as described in the following subsection.

## 3.3.1 Definition of a Stitching Piece

As the name indicates, the set of stitching pieces assist in the "stitching" together of associated segments. Since segments are disjoint, a segment will need to grow in order to join with another segment.

#### 3.3.2 defi

An "open location" with respect to a segment is any puzzle location that is not populated with a member of that segment.

#### 3.3.3 Determine the

Stitching pieces should be near the edge of a segment to increase the likelihood of growth towards a neighboring segment. However, if a stitching piece is too close to the edge of the segment, the algorithm may make erroneous segmentation associations.

As outlined in Algorithm ??, the Mixed-Bag Solver uses iterative boundary tracing to determine each stitching piece's distance to the nearest puzzle location that is not populated with a piece that is a member of the stitching piece's segment..

The algorithm begins by finding set of puzzle locations adjacent to pieces in the segment that are not filled by segment members. These locations represent the .

Algorithm 4 is run for each of the segments. A single iteration of the algorithm has time complexity of O(n), where n is the number of pieces in the segment. What is more, the algorithm is robust enough to handle voids within a segment as well as potential necking within the segment where two large segment components are joined by a narrower bridge.

#### 3.3.4 Stitching Pieces

If stitching pieces are placed to close together, the solver will perform many redundant solvings. While this may not have a deleterious effect on the solver output, it can significantly increase the execution time of this stage. In contrast, if stitching **Algorithm 4** Pseudocode for Determining a Segment Point's Manhattan Distance to the Nearest Open Location

```
1: procedure FINDPIECEDISTANCETOOPEN(segment pieces)
 2:
       explored pieces \leftarrow \{\}
       locations at prev dist \leftarrow \{segment\ pieces\}
 3:
       distance to open \leftarrow 1
 4:
       while |explored pieces| > 0 do
5:
           locations at\_current\_dist \leftarrow \{\}
 6:
 7:
           for each prev dist loc \in locations at prev dist do
              for each adjacent loc of prev dist loc do
 8:
                  if \exists piece at adjacent loc and piece \notin explored pieces then
9:
                      set distance to open for piece
10:
11:
                      remove piece from explored pieces
12:
                      add adjacent loc to locations at current dist
           locations at prev dist \leftarrow locations at current dist
13:
           distance to open \leftarrow distance to open + 1
14:
```

pieces are placed too far apart, subtle segment pairings may be missed. To achieve

Figure ?? shows an original image along with an segment found during the segmentation. This figure also shows the location of the stitching pieces (denoted with a white cross).

## 3.3.5 Quantifying Inter-Segment Similarity

## 3.4 Hierarchical Clustering of Segments

Agglomerative hierarchical clustering is a bottom-up clustering algorithm where in each clustering round, two clusters are merged. Algorithm 5 shows the basic flow of the hierarchical clustering algorithm of the Mixed-Bag Solver; it is adapted from [19].

The only inputs to the hierarchical clustering algorithm are the segments found in the segmentation stage and the Segment Overlap Matrix from the Stitching stage.

# Algorithm 5 Pseudocode for the Hierarchical Clustering Algorithm

- 1: **function** PerformHierarchicalClustering(solved\_segments, overlap\_matrix)
- 2:  $segment\ clusters = \{\}$
- 3: for each  $segment_i \in solved$  segments do
- 4: add new segment cluster  $\Phi_i$  containing segment<sub>i</sub> to segment\_clusters
- 5: Compute the similarity matrix,  $\Gamma$
- 6: while maximum similarity in  $\Gamma > min\ cluster\ similarity\ do$
- 7: Merge the two most similar clusters  $\Phi_i$  and  $\Phi_j$  in segment\_clusters
- 8: Update the similarity matrix,  $\Gamma$  for the merged clusters
- 9: **return** cluster segments

## 3.4.1 Calculating the Initial Similarity Matrix

The Segment Overlap Matrix is a form of hollow matrix, where all elements in the matrix, except those along the diagonal, are populated with meaningful values. In contrast, hierarchical clustering merges segments using a triangular, similarity matrix. Equation (2) defines the similarity,  $\omega(i,j)$  between any two clusters  $\Phi_i$  to  $\Phi_j$ .

$$\omega_{i,j} = \frac{Overlap(\Phi_i, \Phi_j) + Overlap(\Phi_j, \Phi_i)}{2}$$
 (2)

If there are n solved segments found during segmentation, then the initial similarity matrix  $\Gamma$  is size n by n. Each element in  $\Gamma$  is defined by Equation (3). Both i and j are bounded between 1 and n inclusive. What is more, all elements in  $\Gamma$  are normalized between 0 and 1, also inclusive.

$$\Gamma = \begin{cases} 0 & j >= i \\ \omega_{i,j} & i < j \end{cases}$$
 (3)

## 3.4.2 Updating the Similarity Matrix via Single Linking

The Mixed-Bag Solver uses the Single Link version of hierarchical clustering. Hence, the similarity between any two cluster segments is defined as the similarity between the two most similar segments in either cluster. This approach is required because two segments clusters may only be adjacent along the border of two of the composite segments.

Equation (4) defines the similarity between any a merged cluster containing segment clusters,  $\Psi_x$  and  $\Psi_y$ , and any other segment cluster  $\Psi_z$ . Note that segment  $\Phi_i$  is a member of the union of  $\Psi_x$  and  $Psi_y$  while segment  $\Phi_j$  is a member of segment cluster  $\Psi_z$ .

$$\omega_{x \cup y, z} = \underset{\Phi_i \in \Psi_x \cup \Psi_y}{\operatorname{arg max}} \left( \underset{\Phi_j \in \Psi_z}{\operatorname{arg max}} \omega_{i, j} \right) \tag{4}$$

## 3.4.3 Terminating Hierarchical Clustering

Unlike traditional hierarchical clustering, the Mixed-Bag Solver does not always merge all segment clusters until only a cluster remains. Instead clustering continues until the similarity between any of the remaining clusters drops below a predefined threshold. In this thesis, a minimum similarity of 0.1 provided sufficient clustering accuracy without merging unrelated segments.

#### 3.5 Selecting Seed Pieces

#### 3.6 Final Assembly

#### CHAPTER 4

# Quantifying and Visualizing the Quality of a Mixed-Bag Solver Output

Modern jig swap puzzle solvers are not able to perfectly reconstruct the ground-truth input in most cases. As such, quantifiable metrics are required to objectively compare the quality of outputs from different solvers. Cho et al. [13] defined two such metrics namely: direct accuracy and neighbor accuracy. These metrics have been used by others including [17, 14, 15, 20, 16]. This section describes the existing quality metrics, their weaknesses, and proposes enhancements to these metrics to make them more meaningful for Type 2 and Mixed-Bag puzzles.

In the final subsection, tools developed as part of this thesis to visualize the solver output quality are discussed.

## 4.1 Direct Accuracy

Direct accuracy is a relatively naïve quality; it is defined as the fraction of pieces placed in the same location in the ground-truth (i.e., original) and solved image with respect to the total number of pieces. Equation (5) shows the formal definition of direct accuracy (DA), where n is the total number of pieces and c is the number of pieces placed in their original (i.e., correct) location.

$$DA = \frac{c}{n} \tag{5}$$

Direct accuracy is vulnerable to shifts in the solved image where even a few misplaced pieces can cause a significant decrease in accuracy. As shown in Section 4.1.2, this can be particularly true when the ground-truth image's dimensions

are not known by the solver.

This thesis proposes two new direct accuracy metrics namely: Enhanced Direct Accuracy Score (EDAS) and Shiftable Enhanced Direct Accuracy Score (SEDAS).

They are described in the following two subsections; the complementary relationship between EDAS and SEDAS is described in the third subsection.

# 4.1.1 Enhanced Direct Accuracy Score

The standard direct accuracy metric does not account for the possibility that there may be pieces from multiple puzzles in the same solver output. For a given a puzzle  $P_i$  in the set of input puzzles P (where  $P_i \in P$ ) and a set of solved puzzles S where  $S_j$  is in S, Enhanced Direct Accuracy Score (EDAS) is defined as shown in Equation (6).

$$EDAS_{P_i} = \underset{S_j \in S}{\arg\max} \frac{c_{i,j}}{n_i + \sum_{k \neq i} (m_{k,j})}$$

$$\tag{6}$$

 $c_{ij}$  is the number of pieces from input puzzle  $P_i$  correctly placed (with no rotation for Type 2 puzzles) in solved puzzle  $S_j$  while  $n_i$  is the number of pieces in puzzle  $P_i$ .  $m_{k,j}$  is the number of pieces from an input puzzle  $P_k$  (where  $k \neq i$ ) that are also in  $S_j$ .

When solving only a single puzzle, EDAS and standard direct accuracy as defined Equation (5) are equivalent. When solving multiple puzzles simultaneously, EDAS necessarily marks as incorrect any pieces from  $P_i$  that are not in  $S_j$  by dividing by  $n_i$ . What is more, the summation of  $m_{k,j}$  in EDAS is used to penalize for any puzzle pieces not from  $P_i$ . Combined, these two factors enable EDAS to penalize for both extra and misplaced pieces.

It is important to note that EDAS is a score and not a measure of accuracy.





(a) Ground-Truth Image

(b) Solver Output

Figure 3: Solver Output where a Single Misplaced Piece Catastrophically Affects the Direct Accuracy

While its value is bounded between 0 and 1 (inclusive), it is not specifically defined as the number of correct placements divided by the total number of placements since the denominator of Equation (6) is greater than or equal to the number of pieces in both  $P_i$  and  $S_j$ .

# 4.1.2 Shiftable Enhanced Direct Accuracy Score

As mentioned previously, the direct accuracy decreases if there are shifts in the solved image. In many cases such, direct accuracy is overly punitive.

Figure 3 shows a ground-truth image and an actual solver output when the puzzle boundaries were not fixed. Note that only a single piece is misplaced; this shifted all other pieces to the right one location causing the direct accuracy to drop to zero. Had this same piece been misplaced along either the right or bottom side of the image, the direct accuracy would have been largely unaffected. The fact that direct accuracy can give such vastly differing results for essentially the same error shows that direct accuracy has a serious flaw. This thesis proposes Shiftable Enhanced Direct Accuracy Score (SEDAS) to address the often misleadingly punitive nature of direct accuracy.

Let  $d_{min}$  be the Manhattan distance between the upper left corner of the solved image and the nearest placed puzzle piece. Also let L be the set of all puzzle piece locations within distance  $d_{min}$  of the upper left corner of the image. Given that l is a location in L that is used as the reference point for determining the absolute location of all pieces, then SEDAS is defined as shown in Equation (7).

$$SEDAS_{P_i} = \underset{l \in L}{\operatorname{arg max}} \left( \underset{S_j \in S}{\operatorname{arg max}} \frac{c_{i,j,l}}{n_i + \sum_{k \neq i} (m_{k,j})} \right)$$
(7)

In the standard definition of direct accuracy proposed by Cho  $et\ al.$ , l is fixed at the upper left corner of the image. In contrast, SEDAS shifts this reference point within a radius of the upper left corner of the image in order to find a more meaningful value for direct accuracy.

Rather than defining SEDAS based off the distance  $d_{min}$ , an alternative approach is to use the point anywhere in the image that maximizes Equation (7). However, that approach can be significantly more computationally complex in particular in large puzzles with several thousand pieces. Hence, this thesis' approach balances finding a meaningful direct accuracy score with computational efficiency.

## 4.1.3 The Necessity to Use Both EDAS and SEDAS

While EDAS can be misleadingly punitive, it cannot be wholly replaced by SEDAS. Rather, EDAS and SEDAS serve complementary roles. First, EDAS must necessarily be calculated as part of SEDAS since the upper left corner location is inherently a member of the set L. Hence, there is no additional time required to calculate EDAS. What is more, by continuing to use EDAS along with SEDAS, some shifts in the solved image may be quantified; this would not be possible if SEDAS

was used alone.

## 4.2 Neighbor Accuracy

Cho et al. [13] defined neighbor accuracy as the ratio of the number of puzzle piece sides adjacent to the same piece's side in both the ground-truth and solved image versus the total number of puzzle piece sides. Formally, let q be the number of sides each piece has (i.e., four in a jig swap puzzle) and n be the number of pieces. If a is the number of puzzle piece sides adjacent in both the ground-truth and solved images, then the neighbor accuracy, NA, is defined as shown in Equation (8).

$$NA = \frac{a}{n \ q} \tag{8}$$

Unlike direct accuracy, neighbor accuracy is largely unaffected by shifts in the solved image since it considers only a piece's neighbors and not its absolute location. However, the standard definition of neighbor accuracy cannot encompass the case where pieces from multiple input puzzles may be present in the same solver output.

#### 4.2.1 Enhanced Neighbor Accuracy Score

Enhanced Neighbor Accuracy Score (ENAS) improves the neighbor accuracy metric by providing a framework to quantify the quality of Mixed-Bag solver outputs.

Let  $n_i$  be the number of puzzles pieces in the input puzzle  $P_i$  and  $a_{i,j}$  be the number of puzzle pieces sides adjacent in  $P_i$  and  $S_j$ . If  $m_{k,j}$  is the number of puzzle pieces from an input puzzle  $P_k$  (where  $k \neq i$ ) in  $S_j$ , then the ENAS for  $P_i$  is defined as shown in Equation (9).

$$ENAS_{P_i} = \underset{S_j \in S}{\operatorname{arg max}} \frac{a_{i,j}}{q \left( n_i + \sum_{k \neq i} (m_{k,j}) \right)}$$
(9)

In the same fashion as the technique described for EDAS in Section 4.1.1, ENAS divides by the number of pieces  $n_i$  in input puzzle  $P_i$ . By doing so, it effectively marks as incorrect any pieces from  $P_i$  that are not in  $S_j$ . What is more, by including a summation of all  $m_{k,j}$  in the denominator of (9), ENAS marks as incorrect any pieces not from  $P_i$  that are in  $S_j$ . The combination of these two factors allows ENAS to account for extra and misplaced pieces.

# 4.3 Visualizing Solver Output Quality

In images with thousands of pieces, it is often difficult to visually determine the location of individual pieces that are incorrectly placed. What is more, visual tools help developers quickly detect and fix latent bugs.

The following two subsections describe the tools developed as part of this thesis for visualizing direct accuracy and neighbor accuracy.

#### 4.3.1 Visualizing EDAS and SEDAS

In standard direct accuracy, EDAS, and SEDAS, each puzzle piece is assigned a single value (i.e., correct or incorrect). Due to that, the direct accuracy visualization represents each puzzle by a square filled with a solid color. One additional refinement used in this thesis is to subdivide the "incorrect" placements into a set of subcategories; they are (in order of precedence): wrong puzzle, wrong location, and wrong rotation. Table 1 shows the colors assigned to puzzle pieces depending on their direct accuracy classification.

Wrong	Wrong	Wrong	Correct	No Piece
Puzzle	Location	Rotation	Location	Present

Table 1: Color Scheme for Puzzles Pieces in Direct Accuracy Visualizations

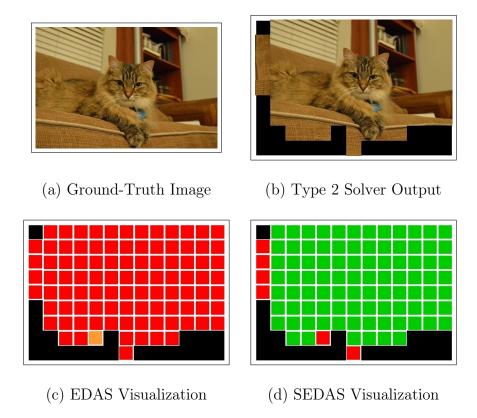


Figure 4: Example Solver Output Visualizations for EDAS and SEDAS

Figure 4 shows a Type 2 solver output along with the associated EDAS and SEDAS visualizations. Since four puzzle pieces were erroneously placed on the left of the image, almost all pieces had the wrong location according to EDAS; the only exception is a single piece that had the right location but wrong rotation. In contrast, almost all pieces have the correct location in the SEDAS representation; note that the piece in the correct location but with wrong rotation in EDAS has the wrong location in SEDAS.

## 4.3.2 Visualizing ENAS

The visualization for neighbor accuracy is very similar to the techniques described in Section 4.3.3 for visualizing best buddies where each puzzle piece is divided into four equal-sized isosceles triangles (one for each side). The triangles are assigned colors depending on whether their neighbors in the solver output and ground-truth image match. The visualization includes a subcategory known as "wrong puzzle" which is a special case that occurs when solving Mixed-Big puzzles and some of the pieces in the solved puzzle are not from the puzzle of interest,  $P_i$ . Table 2 defines the colors used to represent the different classifications of puzzle piece sides in neighbor accuracy visualizations.

Wrong	Wrong	Correct	No Piece
Puzzle	Neighbor	Neighbor	Present

Table 2: Color Scheme for Puzzles Piece Sides in Neighbor Accuracy Visualizations

Figure 5 shows an actual output when solving a Mixed-Bag puzzle with two images. Note that that the puzzle of interest  $P_i$  is the glass and stone building while the other puzzle  $P_k$  is a rainforest house.

All pieces that came from the rainforest house image are shown as blue despite being assembled correctly; this is because they are not from the puzzle of interest. All neighbors from the puzzle of interest (i.e., the glass and stone building) that are placed next to their original neighbor are represented by green triangles while all incorrect neighbors, such as those bordering the rainforest house image, are shown with red triangles.



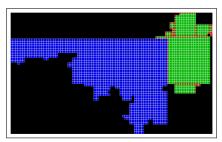
(a) Input Image # 1 - Rainforest House [21]



(b) Input Image # 2 - Building Exterior [22]



(c) Solver Output



(d) ENAS Visualization

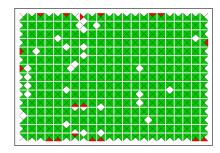
Figure 5: Example Solver Output Visualization for ENAS

# 4.3.3 Visualizing Best Buddies

There is currently no standard for visualizing an image's best buddies. This thesis proposes such a standard for the first time. As a nomenclature, any best buddies that are neighbors in the ground-truth image are referred to as "adjacent best buddies" while any best buddies that are not neighbors are referred to as "non-adjacent best buddies."

In a jig swap puzzle, a piece may have best buddies on up to four sides (since the pieces are square). As such, each piece in the best buddy visualization is divided into four isosceles triangles; the base of each triangle is along the side of the puzzle piece whose best buddy is being represented. A puzzle piece's four isosceles triangles all share a common, non-base vertex at the piece's center.





- (a) Original Image
- (b) Best Buddy Visualization

Figure 6: Visualization of Best Buddies in an Image

No Best	Non-Adjacent	Adjacent Best	No Piece
Buddy	Best Buddy	Buddy	Present

Table 3: Color Scheme for Puzzles Piece Sides in Best Buddy Visualizations

Figure 6 shows an image and its best buddy visualization. The color scheme used to denote the different best buddy relationships is shown in Table 3.

In previous research, all best buddies (in particular best buddy errors) were treated the same. However, non-adjacent best buddies are more likely to occur on piece sides that are adjacent to an open space (i.e., where there is no neighbor piece). Hence, such piece sides are more like to couple with unrelated pieces at a higher rate. For example, the image [?] in Figure 6 has 4 interior and 14 exterior non-adjacent best buddies despite there being 16-times more interior sides.

When using best buddy accuracy as an estimation metric, this thesis differentiates between interior and exterior best buddy errors by giving interior best buddy errors a higher weight.

# CHAPTER 5

# **Experimental Results**

# CHAPTER 6

# Conclusions

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# **APPENDIX**

Example Output of a Single Segmentation Round