ចំណុចសំខាន់ទាំង ៦ របស់ Algorithms

១. ជាមេតូត is used in computer science to describe a problem solving for implementation as a program

២. Algorithm ភាគច្រើនត្រូវបានប្រើជាចាំបាច់ក្នុងការរៀបចំឱ្យមានដំណើរការ data in the computation ។ ទាំងនេះជាដំណើរនៃការបង្ករឱ្យមានជា Objects ដែលហៅថា Data Structure ហើយវាក៏ជា central objects of study in computer science

៣. Computer ជាជំនួយមួយដ៏ប្រសើរក្នុងការដោះស្រាយបញ្ហា, ករណីមានបញ្ហាតូចក្តី ធំក្តី តំរូវឱ្យមានការឈានឱ្យដល់ នៅការ solves the problem correctly, ដូច្នេះ ក៏តំរូវធ្វើយ៉ាងណាឱ្យមានលក្ខណៈ រហ័ស រស់រវើក ដោយបែង ចែកមេតូត use time or space as efficiently as possible

៤. ក្នុងការធ្វើដំណោះស្រាយបញ្ហាធំៗ ត្រូវមានការប្រុងប្រយ័ត្ន លើប្រសិទ្ធភាពនៃ process of solving => to be developed computer program ដែលត្រូវតែយល់ដឹងឱ្យច្បាស់ នឹងកំណត់អត្ថន័យ problem to be solved => managing its complexity និងបំបែកធាតុជាបំណែកតូចៗ ដែលអាចឱ្យមានការងាយស្រួលក្នុងការអនុវត្តន៏

៥. Computer systems គឺជាអ្នករួមចំណែកមួយក្នុង program => មានការផ្លាស់ប្តូរវិធីសាស្ត្រ computing environments (HW & SW) ដើម្បីសំរេចបានជោគជ័យល្អនៃបញ្ហា (make our solutions more portable and longer lasting)

៦. ការជ្រើសរើស best algorithm សំរាប់ផ្នែកណាមួយនៃការងារ អាចផ្តល់ទង្វើបានជា complicated process, ឬក៏ជា sophisticated mathematical analysis. => លក្ខណៈការសិក្សាបែបនេះ ហៅថា analysis of algorithms (ការសិក្សាវិភាគលើ performance ព្រមនឹង experiences, to comparative performance of the methods)

Algorithm មាន running times សមាមាត្រទៅតាមអនុគមន៏

Algorithms ភាគច្រើនមាន Parameter សំខាន់មួយគឺ (N) ដែលមានឥទ្ធិពលទៅលើ Running time; នឹងដែលមានលក្ខណៈ N degree of polynomial (មានដូចជា ទំហំ files, ចំនួន characters ក្នុង text string)។ ទាំងនេះនាំឱ្យមានជាដំណើរ ដោយពឹងលើ mathematical formulas នឹងបន្ទាប់ផ្តល់បានជា Algorithms មាន running times សមាមាត្រទៅនឹងអនុគមន៏:

1 ជាដំណើរ Instruction ដែលបាន Executes once or a few times => running time is constant.

logN ជាដំណើរនៃ Program slower as N grows -> program នេះត្រូវដោះស្រាយបញ្ហាធំ. N=1000=> logN=3, N អាច double បាន ប៉ុន្តែ logN អាច double ទាល់តែ N កើនឡើងដល់ N^2

N Running time របស់ Program នេះជា Linear, លក្ខណៈការងារប្រភេទនេះ Algorithm វាច្រើនតែ process N inputs ឬ produce N outputs

NlogN ជា Algorithm ដោះស្រាយបញ្ហាដោយបំបែកជាផ្នែកៗនៃ sub problems (ដំណោះស្រាយវា Independently) រួចបន្ទាប់ទើបបញ្ចូលគ្នាវិញ។ ពេល N double => running time មានតំលៃច្រើនជាង N double

N^2 ជា Running time នៃ Algorithm ជា quadratic, អាចកើតមានឡើងនៅពេលវា process all pairs of data items (ករណីវា double -> nested loop) បើ N double => running time កើនដល់ជាបួនដង

2^N ជា algorithm មាន running time ជា exponential, ប្រើក្នុងករណីបង្ខំឱ្យមានដំណោះស្រាយអ្វីមួយ។ បើ N double => running time is squares

ជារួម running time នៃផ្នែកណាមួយរបស់ program មានតំលៃជា **some constant * by one of these terms + some smaller terms** (values of the constant coefficient and the terms វាប់បញ្ចូលលើ results of the analysis on implementation details).

Collection of simple algorithms in C++

Recursive Algorithm of factorial

```
long fact (int n)
{ long y;
  if (n==0)
    return 1;
  else
       { y=fact(n-1);
         return n*y;
       }
}
```

or return n*fact(n-1);

Algorithm Iteration of factorial

```
long fact (int n)
{ long f;
  int i;
  f=1;
  for (i=1; i<=n; i++)
      f=f*i;
      return f;
}
```

Algorithm of Euclid's (GCD)

```
int gcd (int u; int v)
{ int t;
  while (u>0)
  { if (u<v)
      { t=u;
        u=v;
        v=t;
      }
      u=u-v;
   }
    return v;
}
```

Algorithm Tower of Hanoi

```
typedef char column;
int n;
void change(column x, column y)
{ cout <<"Change Disk " <<x <<" to "<<y <<"\n"; }

void change_column(int n, column a, column b, column c)
{ if (n>0)
  { change_column(n-1, a, c, b);
    change(a, c);
    change_column(n-1, b, a, c);
   }
}
```

Algorithm of Fibonacci

```
fib(n)=  if (n=0) then 0
         if (n=1) then 1
         else fib (n-1) + fib(n-2)
```

$f_0=0$; $f_1=1$; $f_n=f_{n-2}+f_{n-1}$

n	0	1	2	3	4	5	6	7	8	9
f_0	0	1	1	2	3	5	8	13	21	33

Recursive Algorithm of Fibonacci

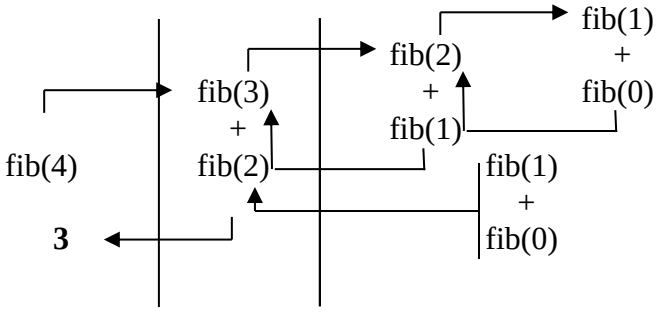
```
long fib(int n)
{ long x,y;
  if (n<=1)
    return n;
  else
      { x=fib(n-1);
        y=fib(n-2);
        return x+y;
      }
}
```

or return fib(n-1) + fib(n-2);

Algorithm Iteration of Fibonacci

```
long fib(int n)
{ long fib2=0;
  long fib1=1;
  long fibn; int i;
  if (n<=1)
    fibn=n;
  else
       { for (i=2; i<=n; i++)
         { fibn=fib1+fib2;
           fib2=fib1;
           fib1=fibn;
         }
        }
   return fibn;
}
```

Example solution of Fibonacci showed as graphic

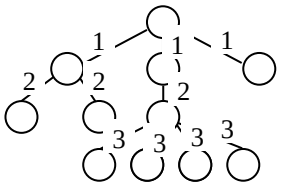


Algorithm of Prime number

```
const int N=100;
void main()
{ int i, j, a[N+1];
  for (a[1]=0; i=2; i<=N; i++)
       a[i]=1;
      for (i=2; i<=N/2; i++)
         for (j=2; j<=N/i; j++)
              a[i*j]=0;
              for (i=1; i<=N; i++)
   if (a[i])
      cout<<i <<' ' ;
      cout<<'\n';
}
```

How to calculate path length in the tree

PL=∑level nodes = ∑lengths
= (1+1+1)+(2+2+2)+(3+3+3+3)=21
= (1+2+2)+(1+2+(3+3+3+3)+1=21



Big-Oh Notation

The mathematical artifact that allows us to suppress detail when we are analyzing algorithms is called the *O-notation, or "big-Oh notation,"* which is defined as follows.

Definition 3.1 *A function g(N) is said to be O(f (N)) if there exist constants c_o and N_o such that $g(N) < c_o f(N)$ for all $N > N_o$.*

We use the O-notation for three distinct purposes:

- To bound the error that we make when we ignore small terms in mathematical formulas
- To bound the error that we make when we ignore parts of a program that contribute a small amount to the total being analyzed
- To allow us to classify algorithms according to upper bounds on their total running times

Sorting Methods

```
BubbleSort(size)
{   for (i=size-1; i > 0; i--)
  {      for (j=0; j < i; j++)
    {         if (a[j] > a[j+1])
                swap a[j] and a[j+1];
    }
  }
} Array size:   10
Last time:       43.182 sec
```

```
InsertionSort(size)
{   for (i = 1; i < size; i++)
  {      if (a[i] < a[i-1])
    {         temp = a[i];
              for (j = i; j>0 && a[j-1]>temp; j--)
                   a[j] = a[j-1];   a[j]=temp;
    }
  }
} Array size:   10
Last time:       33.365 sec
```

```
BubbleSort(size)
{   for (i=size-1; i > 0; i--)
  {      swaps = false;
         for (j=0; j < i; j++)
    {         if (a[j] > a[j+1])
      {          swap a[j] and a[j+1];
                 swaps = true;
      }
    }
     if (!swaps) break;
  }
} Array size:   10
Last time:       51.258 sec
```

```
ShellSort(size)
{   for (inc = size/2; inc>0; inc /= 2)
  {      for (i = inc; i < size; i++)
    {         j = i - inc;
              while (j >= 0)
      {       if (a[j] > a[j+inc])
        {       swap a[j] and a[j+inc];   j -= inc;
        }  else  {      j = -1;       }
      }
    }
  }
} Array size:   10
Last time:       40.66 sec
```

```
QuickSort(low, high)
{   if (high-low <= 1) return;
       pivot = a[high-1];
      split = low;
   for (i=low; i<high-1; i++)
        if (a[i] <pivot)
    {      swap a[i] and a[split];
           split++;
    }
       swap a[high-1] and a[split];
      QuickSort(low, split);
      QuickSort(split+1, high);
   return;
} Array size:   10
Last time:       48.819 sec
```

```
QuickSort(low, high)
{   if (high-low <= 1) return;
       pivot = MedianOf3(low, high);
       split = low;
   for (i=low; i<high-1; i++)
        if (a[i] <pivot)
    {     swap a[i] and a[split];    split++;     }
       swap a[high-1] and a[split];
      QuickSort(low, split);
      QuickSort(split+1, high);
    return;
}
MedianOf3(low, high)
{   middle = (low + high) / 2;
        if (a[low] > a[middle])
        swap a[low] and a[middle];
        if (a[low] > a[high-1])
        swap a[low] and a[high-1];
       if (a[middle] < a[high-1])
      swap a[middle] and a[high-1];
     return a[high - 1];
} Array size:   10
Last time:       38.075 sec
```

```
ShellSort(size)
{   for (inc = size/2; inc>0; inc /= 2)
  {      for (i = inc; i < size; i++)
    {         j = i - inc;
```

```
MergeSort(low, high)
{   if (high-low <= 1) return;
       split = (low+high) / 2;
       MergeSort(low, split);
```

```
            while (j >= 0)
        {           if (a[j] > a[j+inc])
          {             swap a[j] and a[j+inc];
                        j -= inc;
          } else  {   j = -1;     }
        }
      }
    }
} Array size:   10
Last time:      40.66 sec
```

```
       MergeSort(split, high);
      copy a[low ... split-1] to scratch array;
            m1 = 0;   m2 = split;   i = low;
      while (i < m2 && m2 < high)
        if (scratch[m1] <= a[m2])
           a[i++]=scratch[m1++];
        else    a[i++]=a[m2++];
      while (i < m2)   a[i++]=scratch[m1++];
} Array size:   10
Last time:      46.75 sec
```

```
HeapSort(size)
{   for (i = size/2; i >= 0; i--)
     ReHeap(size, i);
        for (i = size-1; i > 0; i--)
  {         swap a[i] and a[0];
           ReHeap(i, 0);
  }
}
ReHeap(len, parent)
{   temp = a[parent];
   child = 2*parent + 1;
   while (child < len)
  {      if (child<len-1 && a[child]<a[child+1])
           child++;
       if (temp >= a[child]) break;
          a[parent] = a[child];
         parent = child;
         child = 2*parent + 1;
  }
   a[parent] = temp;
  return;
} Array size:   10
Last time:      60.302 sec
```

```
MergeSort(low, high)
{   if (high-low <= 1) return;
       split = (low+high) / 2;
       MergeSort(low, split);
       MergeSort(split, high);
      copy a[low ... split-1] to scratch array;
            m1 = 0;
            m2 = split;
            i = low;
     while (i < m2 && m2 < high)
     if (scratch[m1] <= a[m2])
        a[i++]=scratch[m1++];
     else
         a[i++]=a[m2++];
     while (i < m2)
        a[i++]=scratch[m1++];
} Array size:   10
Last time:      46.75 sec
```

```
Radix(size)
{   allZero = false;
   for (i=0; !allZero; i++)
  {  allZero = true;
      for (j=m=n=0; j<size; j++)
     {   tmp = a[j] >> i;
              if (tmp > 1) allZero = false;
              if ((tmp & 1) == 1)
              bucket[m++] = a[j];
       else  a[n++] = a[j];
     }
       m = 0;
      while (n<size)
       a[n++] = bucket[m++];
  }
} Array size:   10
Last time:      70.517 sec
```

```
void SelectionSort(apvector <int> &num)
{    int i, j, first, temp;
     int numLength = num.length( );
     for (i= numLength - 1; i > 0; i--)
    {    first = 0;
        // initialize to subscript of first element
         for (j=1; j<=i; j++)
       // locate smallest between positions 1 and i.
        {     if (num[j] < num[first])
               first = j;
        }    temp = num[first];
 // Swap smallest found with element in position i.
              num[first] = num[i];
              num[i] = temp;
    }  return;
}
```