

1.

(a)

A subnet with pre x 128.119.40.192/26 means $32-26=6$ bits to use $\rightarrow 2^6=64$ addresses.

In practice, I get $64-2=62$ addresses that can be assigned to hosts (all 0s and all 1s cases are special).

So, in the range of 128.119.40.193---128.119.40.254, we can assign 128.119.40.222 to a device in this network.

(b)

We want to create four subnets from this block, so we need 2 bits for that. So the prefix should be a.b.c.d/28 (because $26+2=28$). For each subnet, it can include $2^4=16$ hosts.

1. 128.119.40.192/28

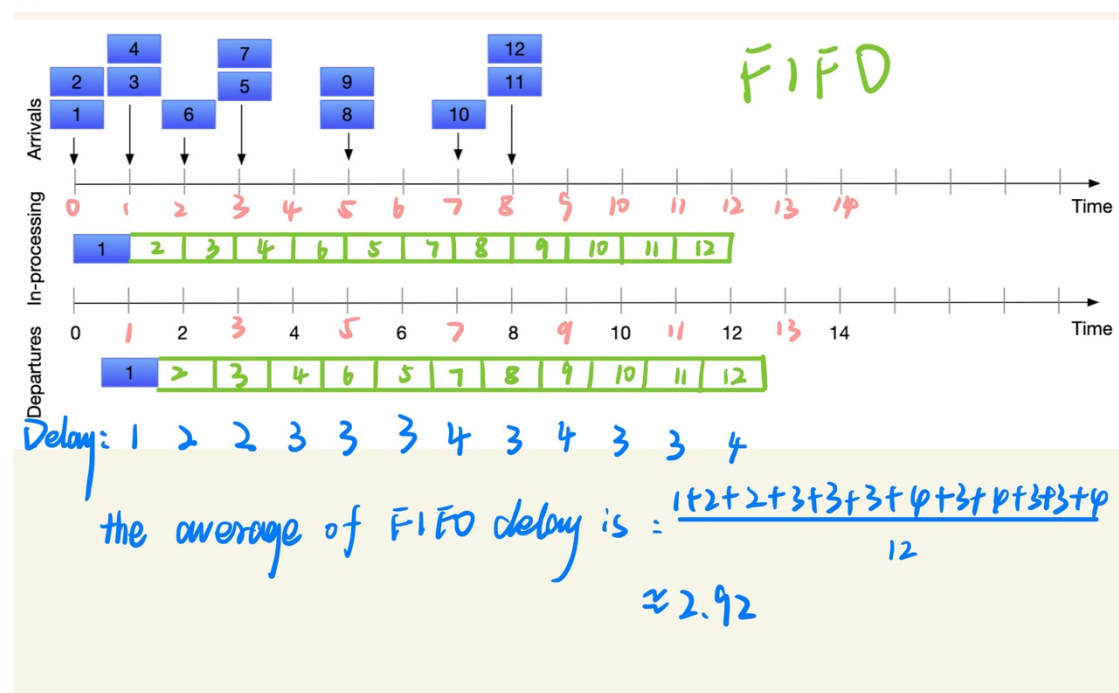
2. 128.119.40.208/28

3. 128.119.40.224/28

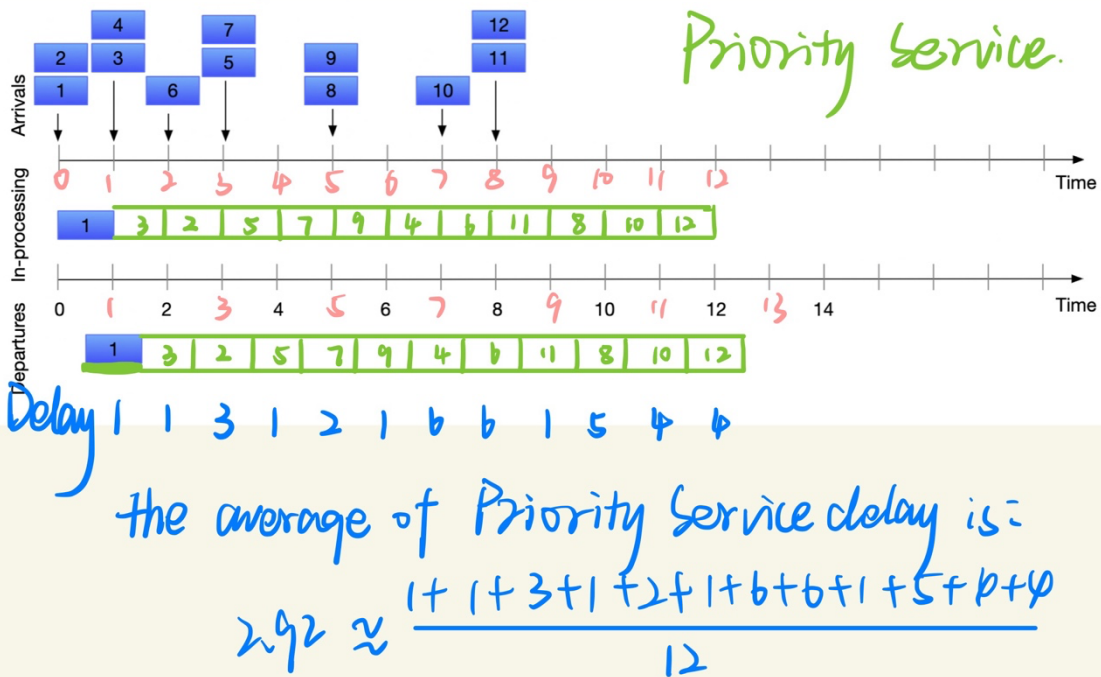
4. 128.119.40.240/28

2.

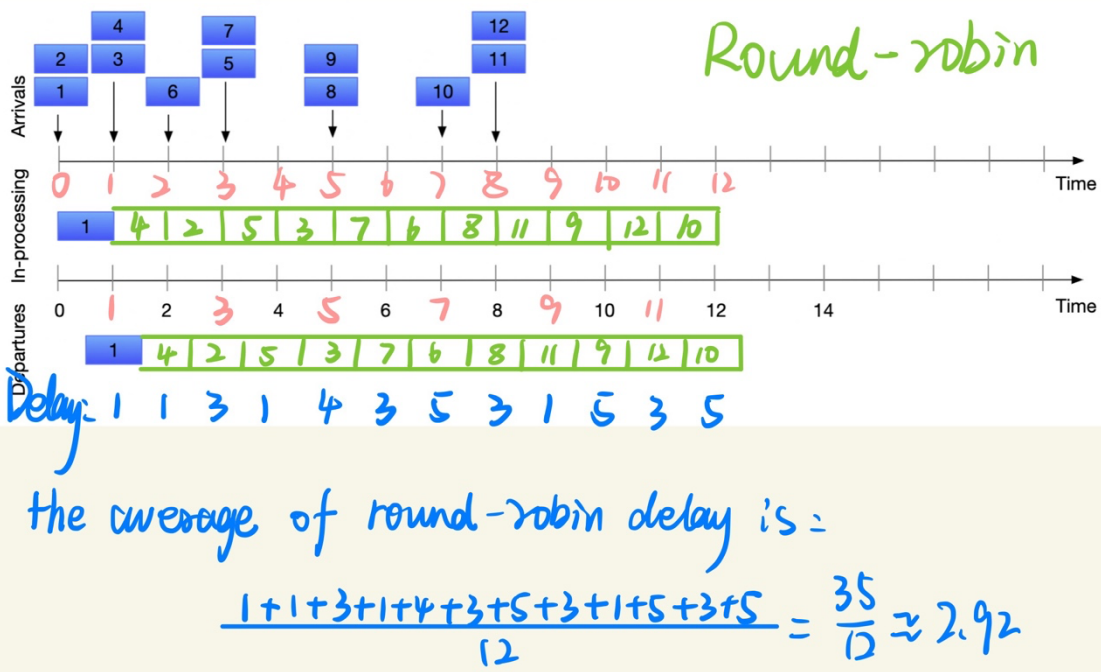
(a) FIFO service:



(b) priority service:

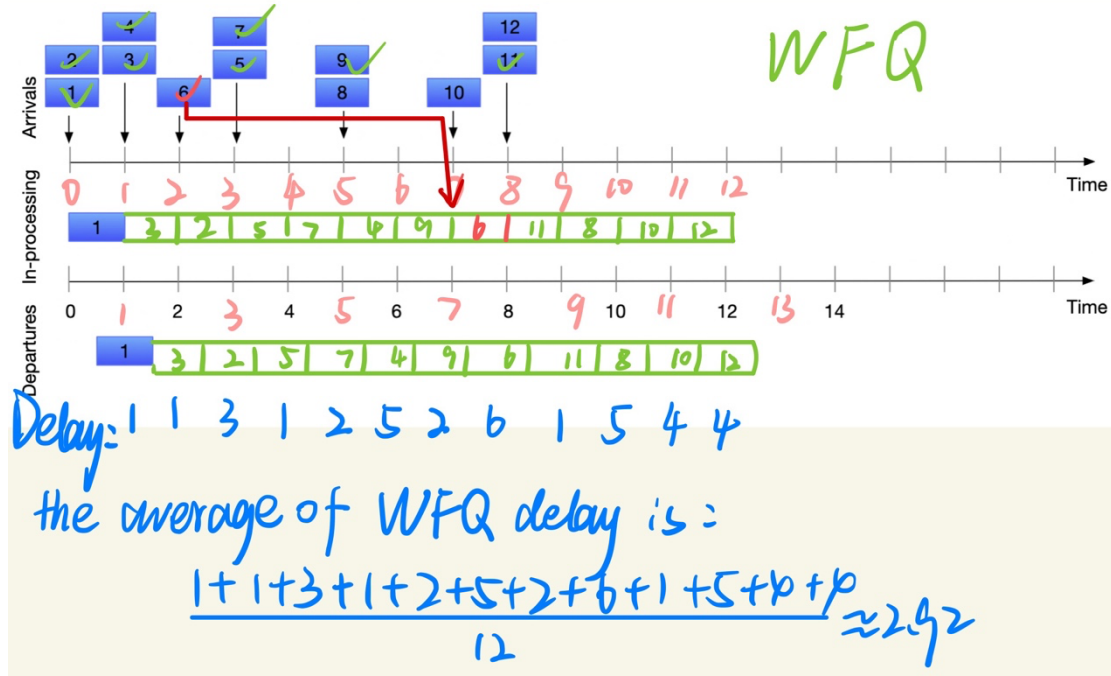


(c) round-robin service:



(d) WFQ:

In this situation, the position of packet 6, it's supposed to be packet 11 (but it hasn't arrived which means the queue for class 1 is empty), so for the next processing we should choose it from class 2. At that time, the first packet in class 2 is packet 6, so we pick packet 6 to go into service.



(e) comparison:

As we can see, the average delay for all these four services is the same.

FIFO does not classify messages; FIFO lets messages enter the queue in the order they arrive at the interface, and lets messages exit the queue in the order they entered, with first-in messages going out first and last-in messages going out second.

P (Priority) scheduling is the scheduling of queues in order of priority.

RR scheduling uses polling for multiple queues. RR polls multiple queues in a circular fashion. If the polled queue is not empty, a message is fetched from that queue; if the queue is empty, the queue is skipped directly, and the scheduler does not wait.

Advantages of WFQ.

1 Different queues get a fair chance of scheduling, which balances the delay of each flow in general.

2、 Short and long messages get fair scheduling: If multiple long and short messages are waiting to be sent simultaneously between different queues, let the short messages get priority scheduling, thus reducing the inter-message jitter of each stream in general.

3. Statistically speaking, the smaller the weight, the less bandwidth is allocated. The larger the weight, the more bandwidth is allocated.