

Solving Proximity Constraints

Sophie Hofmanninger & Jan-Michael Holzinger

26.06.2019

1 Introduction

2 System Model

3 Workflow

Motivation

For proving theorems, a frequently occurring problem is to find common instances of formulae.

Example 1

Let f be a function, a, b constants and x a variable. The two expressions

$$f(a, x) \quad \text{and} \quad f(a, b)$$

can be unified with $\{x \mapsto b\}$.

Motivation

For proving theorems, a frequently occurring problem is to find common instances of formulae.

Example 2

Let f, g be functions, a, b constants and x a variable. The two expressions

$$f(a, x) \quad \text{and} \quad g(a, b)$$

cannot be unified as $f \neq g$.

Motivation

In 1965 Robinson presented his unification algorithm and solved this problem, his algorithm was improved for better(=faster) performance since.

If we consider now the unification problem

$$f(a, x) \simeq^? g(a, b)$$

again, we might wonder, if we could not ignore $f \neq g$, if they are “close” to each other, i.e. if they are equal in a fuzzy logic sense. Being close is represented as a proximity relation, which are symmetric and reflexive, but not necessarily transitive. C. Pau and T. Kutsia solved this problem, presenting an algorithm, which we implemented.

Introduction

4 sets:

- P : unification problem to be solved
- C : neighborhood constraint
- σ : set of pre-unifier
- Φ : name-class mapping

Pre-Unification rules

(Tri) Trivial: ...

(Dec) Decomposition: ...

...

Rules for Neighborhood Constraints

(FFS) $\{f \approx^? g\} \uplus C; \Phi \Rightarrow C; \Phi$; if $\mathcal{R}(f, g) \geq \lambda$

(NFS) $\{N \approx^? g\} \uplus C; \Phi \Rightarrow C; \text{update}(\Phi, N \rightarrow \mathbf{pc}(g, \mathcal{R}, \lambda))$

(FSN) $\{g \approx^? N\} \uplus C; \Phi \Rightarrow \{N \approx^? g\} \cup C; \Phi$

(NN1)

$\{N \approx^? M\} \uplus C; \Phi \Rightarrow C; \text{update}(\Phi, N \rightarrow f, M \rightarrow \mathbf{pc}(f, \mathcal{R}, \lambda))$,
where $N \in \text{dom}(\Phi)$, $f \in \Phi(N)$

(NN2) $\{M \approx^? N\} \uplus C; \Phi \Rightarrow \{N \approx^? M\} \cup C; \Phi$, where
 $M \notin \text{dom}(\Phi)$, $N \in \text{dom}(\Phi)$

(Fail1) $\{f \approx^? g\} \uplus C; \Phi \Rightarrow \perp$, if $\mathcal{R}(f, g) < \lambda$

(Fail2) $C; \Phi \Rightarrow \perp$, if there exists $N \in \text{dom}(\Phi)$ such that $\Phi(N) = \emptyset$

Simple example

Problem to solve: $p(x, y) = ? q(f(a), g(b))$

Solution: ...

1 Introduction

2 System Model

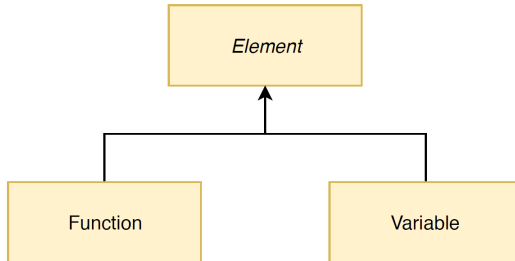
3 Workflow

System Model

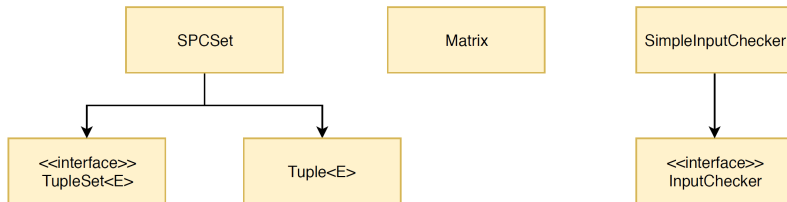
Project consists of 4 packages:

- elements : deals with elements
- tool : offers important tools (e.g. read input)
- unificationProblem : treats the unification problem
- userInterfaces : allow user interfaces

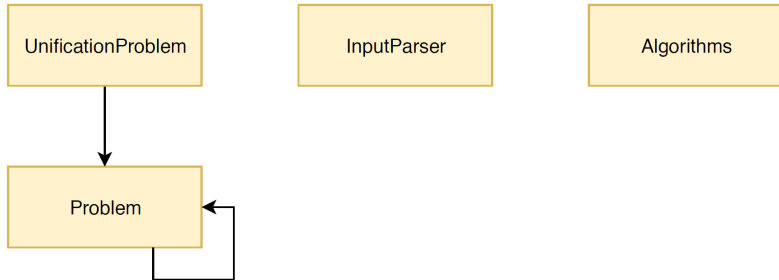
Package elements



Package tool



Package unificationProblem



Package userInterfaces

1 Introduction

2 System Model

3 Workflow

Workflow

