

Solving Proximity Constraints

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1 Introduction

2 System Model



Motivation

For proving theorems, a frequently occurring problem is to find common instances of formulae.

Example 1

Let f be a function, a, b constants and x a variable. The two expressions

$$f(a,x)$$
 and $f(a,b)$

can be unified with $\{x \mapsto b\}$.



Motivation

For proving theorems, a frequently occurring problem is to find common instances of formulae.

Example 2

Let f, g be functions, a, b constants and x a variable. The two expressions

$$f(a,x)$$
 and $g(a,b)$

cannot be unified as $f \neq g$.



Motivation

In 1965 Robinson presented his unification algorithm and solved this problem, his algorithm was improved for better(=faster) performance since.

If we consider now the unification problem

$$f(a,x) \simeq^{?} g(a,b)$$

again, we might wonder, if we could not ignore $f \neq g$, if they are "close" to each other, i.e. if they are equal in a fuzzy logic sense. Being close is represented as a proximity relation, which are symmetric and reflexive, but not necessarily transitive. C. Pau and T. Kutsia solved this problem, presenting an algorithm, which we implemented.



Introduction

4 sets:

- P: unification problem to be solved
- C: neighborhood constraint
- lacksquare σ : set of pre-unifier
- Φ: name-class mapping





Pre-Unification rules

(Tri)
$$\{x \simeq^{?} x\} \uplus P; C; \sigma \Rightarrow P; C; \sigma$$

(Dec) $\{F(\overline{s_n}) \simeq^{?} G(\overline{t_n})\} \uplus P; C; \sigma \Rightarrow \{\overline{s_n} \simeq^{?} t_n\} \cup P; \{F \approx^{?} G\} \cup C; \sigma$
(VE) $\{x \simeq^{?} t\} \uplus P; C; \sigma \Rightarrow \{t' \simeq^{?} t\} \cup Px \mapsto t'; C; \sigma\{x \mapsto t'\}$
(Ori) $\{t \simeq^{?} x\} \uplus P; C; \sigma \Rightarrow \{x \simeq^{?} t\} \cup P; C; \sigma$
(Cla) $\{F(\overline{s_n}) \simeq^{?} G(\overline{t_n})\} \uplus P; C; \sigma \Rightarrow \bot$ if $m \neq n$
(Occ) $\{x \simeq^{?} t\} \uplus P; C; \sigma \Rightarrow \bot$ if there is an occurence cycle of x in t
(VO) $\{x \simeq^{?} y, \overline{x_n} \simeq^{?} y_n\}; C; \sigma \Rightarrow \{\overline{x_n} \simeq^{?} y_n\}\{x \mapsto y\}; C; \sigma\{x \mapsto y\}$



Rules for Neigborhood Constraints

```
(FFS) \{f \approx^? g\} \uplus C; \Phi \Rightarrow C; \Phi; \text{ if } \mathcal{R}(f,g) \geq \lambda
(NFS) \{N \approx^? g\} \uplus C; \Phi \Rightarrow C; update(\Phi, N \rightarrow pc(g, \mathcal{R}, \lambda))
(FSN) \{g \approx^? N\} \uplus C; \Phi \Rightarrow \{N \approx^? g\} \cup C; \Phi
(NN1)
\{N \approx^{?} M\} \uplus C; \Phi \Rightarrow C; update(\Phi, N \rightarrow f, M \rightarrow pc(f, \mathcal{R}, \lambda)),
where N \in dom(\Phi), f \in \Phi(N)
(NN2) \{M \approx^? N\} \uplus C; \Phi \Rightarrow \{N \approx^? M\} \cup C; \Phi, \text{ where }
M \notin dom(\Phi), N \in dom(\Phi)
(Fail1) \{f \approx^? g\} \uplus C; \Phi \Rightarrow \bot, if \mathcal{R}(f,g) < \lambda
```

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(Fail2) $C; \Phi \Rightarrow \bot$, if there exists $N \in dom(\Phi)$ such that $\Phi(N) = \emptyset$



Simple example

Pre - fail: p(x, a) = ?q(f(a), g(b))

Test 4 in Test.java

Pre - ok, CS - fail: - des ging bei mir, i hätt 5 Schritte und 2 Branches, die aber failen-

$$p(x,y,x) = ?q(f(a),g(b),y); R = \{(b,c),(c,d),(f,g),(p,q)\}$$

Test 3 in Test.java

Pre - ok, CS - ok: - des ging bei mir, i hätt 4 Schritte und 1 Branch-

$$p(x,z) = q(f(b), f(x)); R = \{(a,a'), (a',b), (b,c'), (c',c), (p,q)\}$$

Solution: ...



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System Model

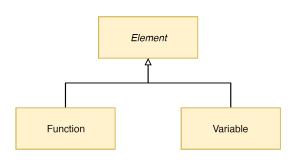
Project consists of 4 packages:

- elements : deals with elements
- tool : offers important tools (e.g. read input)
- unificationProblem : treats the unification problem
- userInterfaces : allow user interfaces





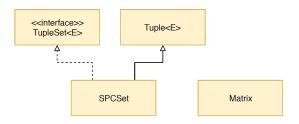
Package elements

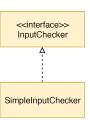






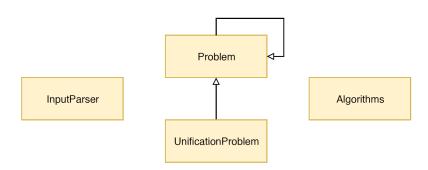
Package tool







Package unificationProblem





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Package userInterfaces





2 System Model





