

Project: Solving proximity constraints

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Version Number	Changes Summary	Author
0.1		Jan-Michael
0.2	added System Model	Jan-Michael
0.3	modified Parser, added Workflow	Jan-Michael
0.4	add ConstraintSimplification	Sophie
0.5	Proximity Relation (as Matrix) added	Jan-Michael
0.6	update system model, update matrix	Sophie
0.7	added UML	Jan-Michael
0.8	added User Interfaces	Jan-Michael
0.9	minor changes in matrix description update constraint simplification	Sophie
0.10	updated UML files	Jan-Michael
0.11	added Input Checker	Jan-Michael
0.12	added Command Line Interface	Jan-Michael

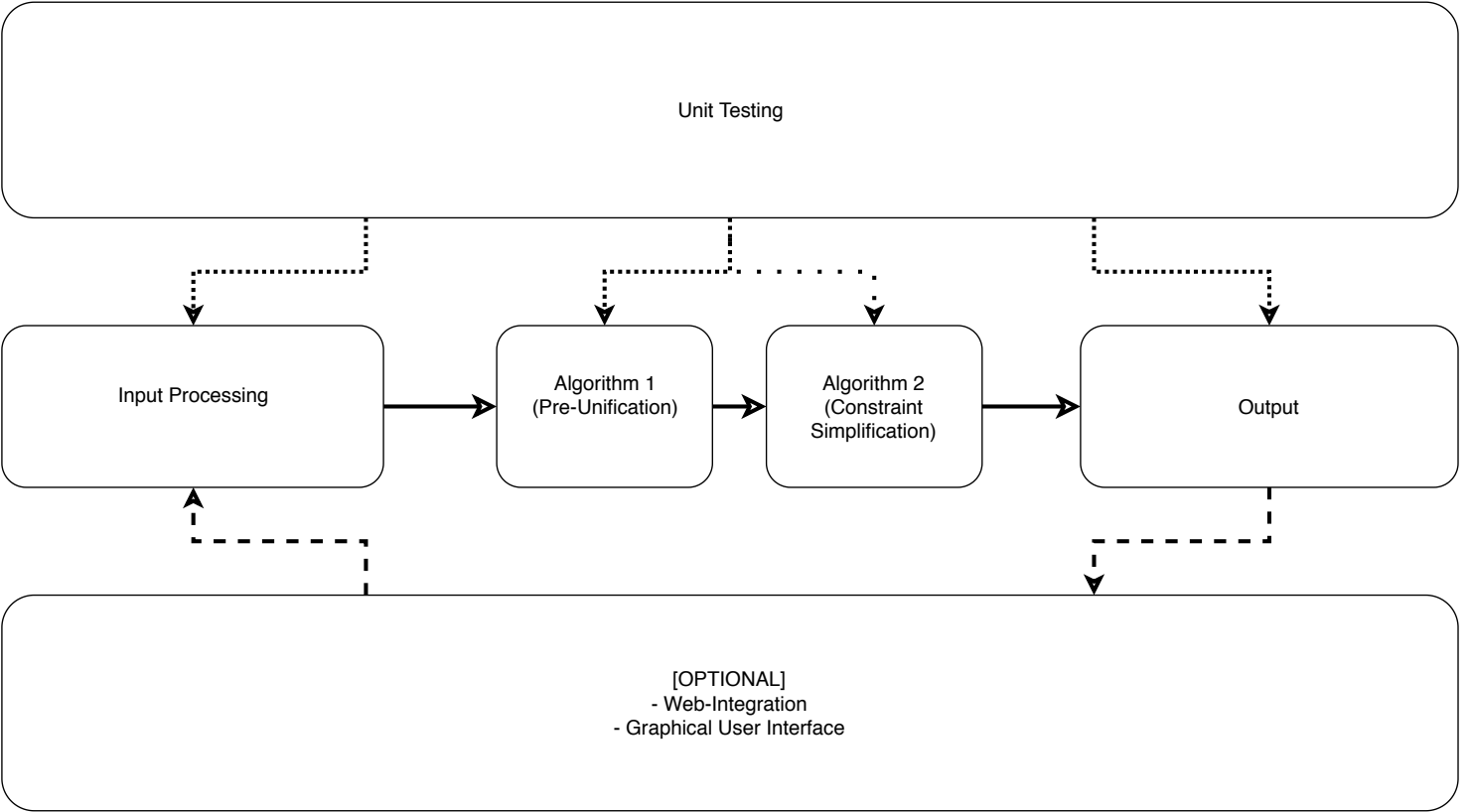
1 System Overview

We split the problem in 4 (5) smaller tasks:

1. Input Processing,
 2. Pre-Unification,
 3. Constraint Simplification,
 4. Input/Output.
- O. Web-Integration.

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1.1 Input Processing

1.1.1 InputParser

The first idea here is to copy, alter and extend the existing code, in the class InputParser.

1.1.2 Input Checker

We want to be able to check the input for example for missing parenthesis and so on. Therefore we will allow the user to specify an InputChecker. The idea is, to create an Interface, with just a method `check(String):boolean`. This then can be implemented by the implementation of the checkers.

1.1.3 Proximity Relations

For the Proximity Relations \mathcal{R} and the λ -cut we have the following idea:

1. We try to get the number of function symbols with the same arity.
2. We let the user input the values to construct a symmetric matrix that consists of values in $[0, 1]$. This matrix must have a 1 in the main diagonal. All values below are 0. Therefore they will not be stored in the implementation.
3. We let the user (later) input $\lambda \in [0, 1]$ and calculate the set \mathcal{R}_λ .

The current implementation of InputParser creates a list of all Functions. This list is then sorted by arity, i.e. let f, g be functions with arity $a, b \in \mathbb{N}$ respectively. Then $a < b \Rightarrow f \prec g$.

Let now $n \in \mathbb{N}$ be the size of the list, the list then represents the caption of a $n \times n$ matrix. Assume the problem contains functions with arity $0, 1, \dots, m$. Let k_l be the number of Functions with arity l . Then

$$\sum_{l=0}^m k_l = n.$$

The matrix then looks like

$$\begin{matrix} & f_{0_1} & f_{0_2} & \dots & f_{0_{k_0}} & f_{1_1} & f_{1_2} & \dots & f_{1_{k_1}} & \dots & f_{m_1} & f_{m_2} & \dots & f_{m_{k_m}} \\ \begin{matrix} f_{0_1} \\ f_{0_2} \\ \vdots \\ f_{0_{k_0}} \\ f_{1_1} \\ f_{1_2} \\ \vdots \\ f_{1_{k_1}} \\ \vdots \\ f_{m_1} \\ f_{m_2} \\ \vdots \\ f_{m_{k_m}} \end{matrix} & \left(\begin{array}{cccccccccccc} 1 & (*) & (*) & (*) & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ * & 1 & (*) & (*) & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ * & * & \ddots & (*) & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ * & * & * & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & (*) & (*) & (*) & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & * & 1 & (*) & (*) & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & * & * & \ddots & (*) & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & * & * & * & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & \ddots & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & (*) & (*) & (*) & (*) \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & * & 1 & (*) & (*) & (*) \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & * & * & \ddots & (*) & (*) \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & * & * & * & 1 & \end{array} \right) \end{matrix},$$

where:

- the 1's in the diagonal are fixed, as a function must have proximity 1 to itself,
- the 0's are fixed, as functions with different arities can't be close,
- the values where the entry is * must be of the same value as their (*) counterpart, as the matrix is symmetric.

Hence, we only need to get the values for the positions marked with (*). We call such a position “open case”, and we make a list of open cases. Before asking the user to enter relation values, we will generate the list of open cases inside of the matrix. After generation of the list, we will wait for user input. With the valid user input we will build the proximity relation matrix.

If we want to retrieve a value, i.e. if we want to know $\mathcal{R}(s_1, s_2)$ for given s_1, s_2 , we can call *getRelation*(s_1, s_2). To get all relation pairs of s_1 , that have a relation value greater or equal to λ , the method *getRelations*(s_1, λ) can be used. Moreover it is possible to get all relation elements of the matrix by calling *getListOfFunctions*(). This will return all $f_{i_{k,j}}$, where $0 \leq i \leq m$ and $0 \leq j \leq m$.

1.2 Algorithms

We implement the Algorithms in an own class, that has two public static functions, preUnification and constraintSimplification. The other methods are only used for the construction of the two simplification algorithms and therefore they are private and static. The methods should take two inputs, mandatory an unification problem, and optional a StringBuffer, to log certain events.

1.2.1 Pre-Unification Algorithm

The preUnification method consists of a loop, that runs until either $P = \emptyset$ or it is detected, that there is no solution to the problem.

Inside the loop body, the 7 pre-unification rules are iteratively applied to the first element (which gets popped by doing so). The method returns true, iff the pre-unification was successful.

The method changes the problems constraints and pre-unifier accordingly.

1.2.2 Constraint Simplification Algorithm

The constraintSimplification method will be called, if the preUnification method returns true. As the preUnification method, the constraintSimplification method consists of a loop. This loop runs until the set of constraints is empty. As input the method needs the problem, more precisely the set of constraints, and the relation matrix \mathcal{R} . If the set of constraints could be simplified, the method returns true and otherwise false.

Inside the loop, the seven rules to simplify the constraint set, will be applied. Three of the seven rules are hidden, because the NN2 rule is integrated in the NN1 rule, the Fail1 is implemented in FFS and the Fail2 is integrated in the NN1 rule. Moreover it is possible to run the constraint simplification algorithm with and without logging.

1.3 Input/Output

We provide a command line user interface for the program. This user interface is implemented in the “SPC_CL” class(SolvingProximityConstraints_via_CommandLine). The program can be accessed with no or some of this optional arguments: Note: the first argument if any are given is always the equation in the format “lhs =? rhs” where the blanks in between are mandatory.

the equation in the format “lhs =? rhs” (left hand side and right hand side of the equation.

-f <PATH> a path given to a file where (parts) of the proximity relations are stored the format must follow the convention, that is mentioned below.

-s silent mode on/off, default is off, -s with no argument switches on, optional argument “n” switches off.

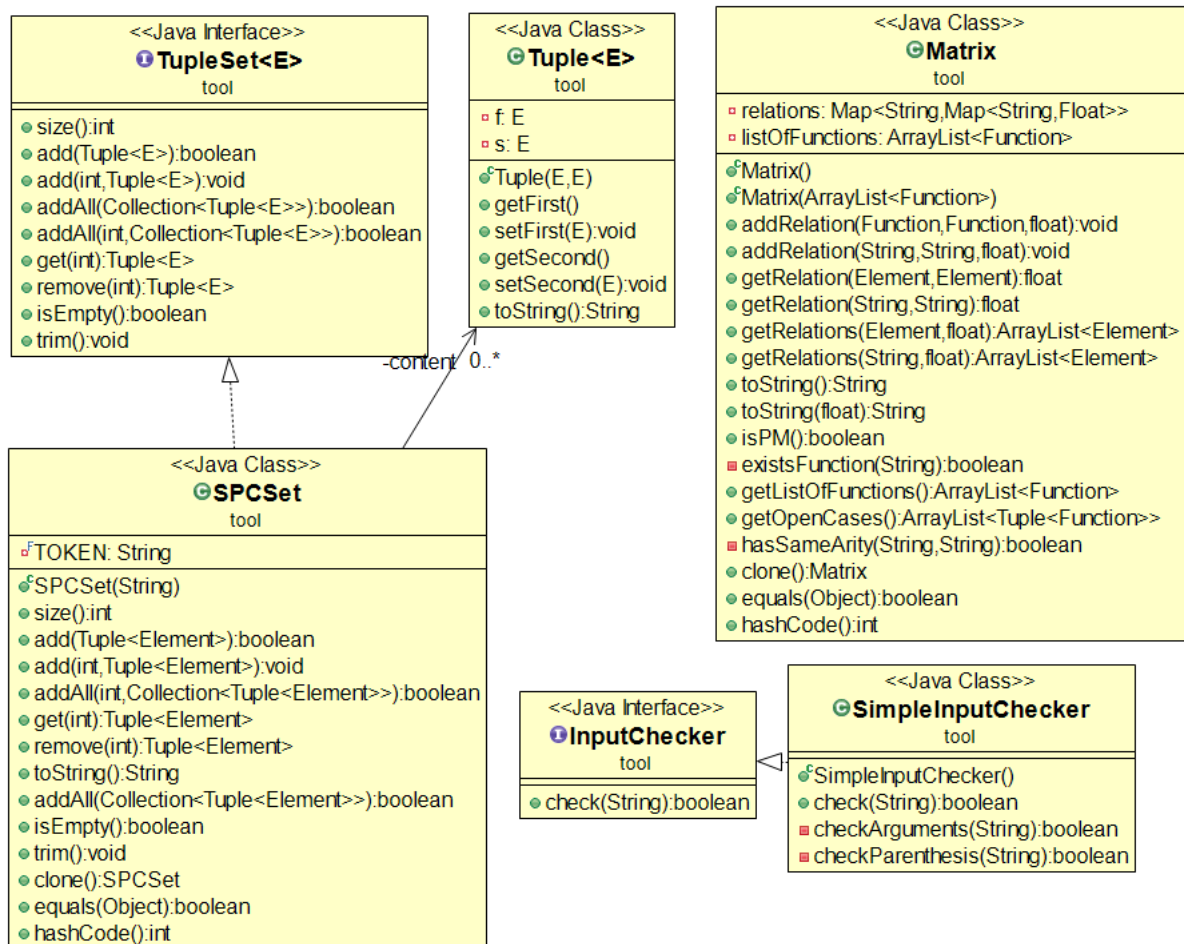
-l followed by the lambda value.

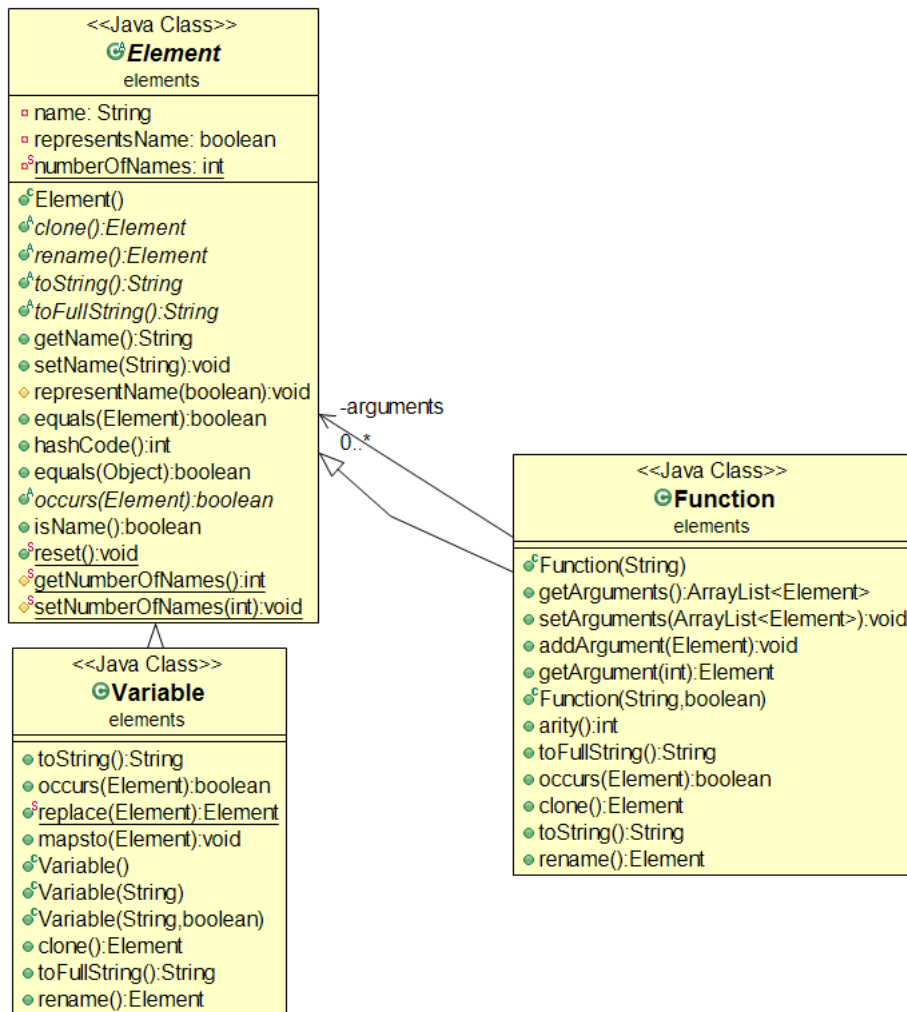
The proximity relations must be stored in a text file, one line for each relation, where the two functions (or constants) are separated with “,” (comma) and the lambda value, also separated with “,”.

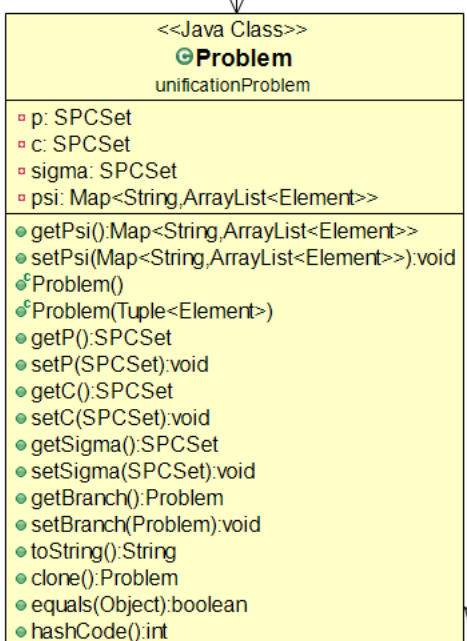
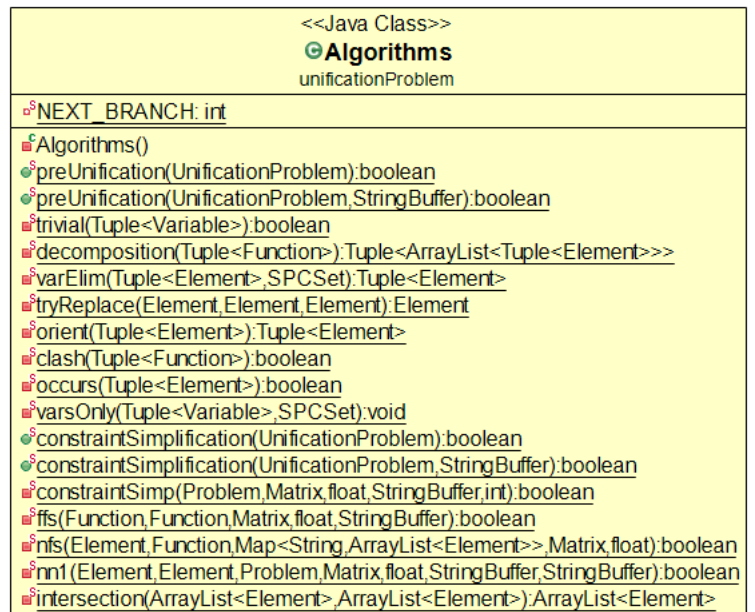
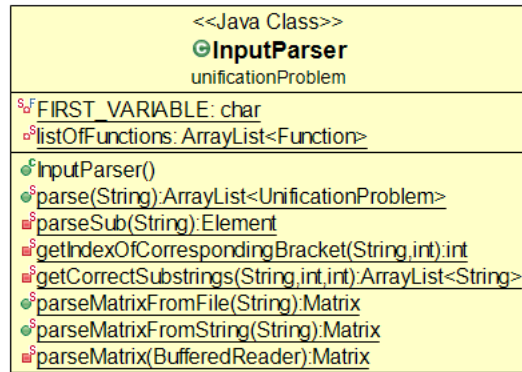
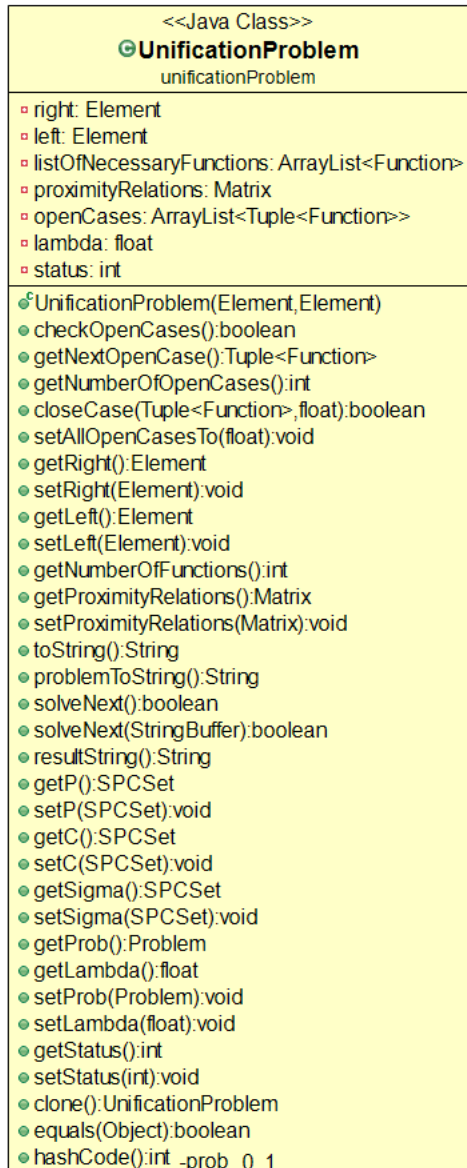
2 System Model

The program consists of 4 packages,

- tool
- elements
- unificationProblem
- userInterfaces














-branch

0..1

<<Java Class>>  SPC_CL userInterfaces
 br: <u>BufferedReader</u>  intermediateResult: <u>UnificationProblem</u>
 SPC_CL()  <u>main(String[]):int</u>  <u>getLambda():float</u>  <u>checkInput(String):boolean</u>

3 Work Flow

The typical workflow looks like this:

