Coding Conventions

Formatting

• Run AStyle script

RunAstyle.pl

I'm not even slightly happy about the way this looks but I've found no better alternative.

- Spacing in header files
 - o No whatspace before or after text of file
 - o One space between leading include #define and includes
 - One space between releated sections of #includes
 - o 3 spaces between major file sections (before \file comment and after)
 - Inside namespaces (where code declared) two spaces between major related declarations.

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EXAMPLE

```
* Copyright(c) Sophist Solutions, Inc. 1990-2013. All rights reserved
#ifndef Stroika Foundation Configuration Enumeration h
#define Stroika Foundation Configuration Enumeration h
#include
          "../StroikaPreComp.h"
#include
           "Common.h"
* \file
 * TODO:
       @todo - maybe stuff like Add(ENUM, ENUM), and DIFF (ENum, ENUM) to
workouarnd
               issues with too-strong typing with enum class?? (avoid so many
casts)
 */
namespace Stroika {
   namespace Foundation {
       namespace Configuration {
               \brief Increment the given enumeration safely, without a bunch
of casts.
                   \req ENUM uses Define Start End Count() to define eSTART,
eEND
                   \req e >= typename ENUM::eSTART and e < typename
ENUM::eEND
```

Begin/End versus start/length

STL is reasonably consistent, with most APIs using T* start, T* end, but some APIs use length instead of end. The Stroika convention is to always use T* start, T* end.

Rationale

One, this gives more consistent expectations. That's especially important for APIs that use offsets (like String) – so that it's obvious the meaning of integer parameters.

And it avoids problems with overflow. For example, if you had an API like:

```
basic_string substr(
   size_type _Off = 0,
   size_type _Count = npos
) const
```

To map this to an internal representation you have todo:

```
char* s = m_bufPtr + _Off;
char* e = m_bufPtr + _Off + _Count;
```

but if count was numeric_limits<size_t>::max(), then the e pointer computation would overflow. There are ways around this, but mixing the two styles creates a number of problems - but for implementations - and for use.

mk Factories

Stroika doesn't make much use of the factory pattern, but occasionally – it is useful. If the type provided by the factory is exactly the type of a given class, then we generally use

```
struct T {
    static T mk();
};
```

Of course in this case, there was little obvious motivation to use a factory instead of regular constructor. However, if the class T is effectively a smart-pointer wrapper on some underlying dynamic 'rep' – this pattern may make sense.

But – for shared_ptr types, and typedefs, we generally use

```
struct X;
typedef shared_ptr<X> XPtr;
XPtr mkXPtr ();
```

Compare () and operator<, operator>, etc...

For types Stroika defines, it generally uses the convention of providing a compare function:

bool operator<, operator>=, operator>=, operator==, operator!= which inline trivially maps to this.

Stroika code which COUNTS on comparison doesn't directly call Compare(), but instead uses 'a < b', etc. This applies to things like Stroika containers. The reason for this later choice include:

- Working with builtin types (e.g. in)
- ➤ Working with STL types, and 3rd-party libraries
- Probably more likely to seamlessly fit with user code

Note that we choose to use member function operators for comparison – instead of global (namespace) functions (with two arguments) – because

- The namespace based 'global' operators only get overloaded if you import the entire namespace (or at least import those functions)
 - This is either very awkward to use or encourages namespace conflicts
- The namespace/global approach CAN lead to confusing conflicts of inappropriately colliding chained conversions

The downside of this approach is that stuff like:

```
if (L"aa" < String (L"ss")) {
}</pre>
```

Fails. You must have the left most object be already a String (or other stroika) object. Sigh. Seems like the best compromise?