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```
#include <unistd.h>
const int S = 65536;
inline char RC() {
  static char buf[S], *p = buf, *q = buf;
  return p == q and (q =
        (p = buf) + read(0, buf, S)) == buf ? -1 : *p++;
inline int RI() {
  static char c; int a;
  while (((c = RC
()) < '0' or c > '9') and c != '-' and c != -1);
if (c == '-') { a = 0;
    while ((c = RC()
         ) >= '0' and c <= '9') a *= 10, a -= c ^{\prime}0'; }
  else { a = c ^ '0';
    while ((c = RC()
         ) >= '0' and c <= '9') a *= 10, a += c ^ '0'; }
  return a;
}
```

1.2 Black Magic

```
#include <ext/pb_ds/assoc_container.hpp> // rb_tree
#include <ext/rope> // rope
using namespace __gnu_pbds;
using namespace __gnu_cxx; // rope
typedef __gnu_pbds::priority_queue<int> heap;
int main() {
  heap h1, h2; // max heap
  h1.push(1), h1.push(3), h2.push(2), h2.push(4);
  h1.join(h2); // h1 = {1, 2, 3, 4}, h2 = {};
tree<int, null_type, less<int>, rb_tree_tag
        , tree_order_statistics_node_update > st;
   tree<int, int, less<int>, rb_tree_tag
          tree_order_statistics_node_update > mp;
  for (int x : {0, 2, 3, 4}) st.insert(x);
cout << *st.find_by_order</pre>
  (2) << st.order_of_key(1) << endl; //31 rope <<a href="mailto:rope">char</a> *root[10]; // nsqrt(n)
   root[0] = new rope<char>();
   root[1] = new rope < char > (*root[0]);
  // root[1] -> insert(pos, 'a');
  // root[1]->at(pos); 0-base
  // root[1]->erase(pos, size);
// __int128_t,__float128_t
// for (int i = bs._Find_first
      (); i < bs.size(); i = bs._Find_next(i));
```

2 Graph

2.1 BCC Vertex*

```
vector < int > G[N]; // 1-base
vector < int > nG[N * 2], bcc[N];
int low[N], dfn[N], Time;
int bcc_id[N], bcc_cnt; // 1-base
bool is_cut[N]; // whether is av
bool cir[N * 2];
int st[N], top;
void dfs(int u, int pa = -1) {
  int child = 0;
  low[u] = dfn[u] = ++Time;
  st[top++] = u;
  for (int v : G[u])
    if (!dfn[v]) {
      dfs(v, u), ++child;
      low[u] = min(low[u], low[v]);
      if (dfn[u] <= low[v]) {</pre>
        is_cut[u] = 1;
        bcc[++bcc_cnt].clear();
        int t;
        do {
          bcc_id[t = st[--top]] = bcc_cnt;
           bcc[bcc_cnt].eb(t);
         } while (t != v);
        bcc_id[u] = bcc_cnt;
        bcc[bcc_cnt].eb(u);
    } else if (dfn[v] < dfn[u] && v != pa)</pre>
      low[u] = min(low[u], dfn[v]);
  if (pa == -1 && child < 2) is cut[u] = 0;</pre>
}
void bcc_init(int n) { // TODO: init {nG, cir}[1..2n]
  Time = bcc_cnt = top = 0;
  for (int i = 1; i <= n; ++i)</pre>
    G[i].clear(), dfn[i] = bcc_id[i] = is_cut[i] = 0;
void bcc_solve(int n) {
  for (int i = 1; i <= n; ++i)</pre>
    if (!dfn[i]) dfs(i);
   // block-cut tree
  for (int i = 1; i <= n; ++i)</pre>
    if (is_cut[i])
      bcc_id[i] = ++bcc_cnt, cir[bcc_cnt] = 1;
  for (int i = 1; i <= bcc_cnt && !cir[i]; ++i)
    for (int j : bcc[i])
      if (is_cut[j])
        nG[i].eb(bcc_id[j]), nG[bcc_id[j]].eb(i);
```

2.2 Bridge*

```
int low[N], dfn[N], Time; // 1-base
vector<pii> G[N], edge;
```

```
vector<bool> is bridge:
void init(int n) {
  for (int i = 1; i <= n; ++i)</pre>
    G[i].clear(), low[i] = dfn[i] = 0;
void add_edge(int a, int b) {
 G[a].eb(pii(b, SZ(edge))), G[b].eb(pii(a, SZ(edge)));
  edge.eb(pii(a, b));
void dfs(int u, int f) {
  dfn[u] = low[u] = ++Time;
  for (auto i : G[u])
    if (!dfn[i.X])
     dfs(i.X, i.Y), low[u] = min(low[u], low[i.X]); \\  else if (i.Y != f) low[u] = min(low[u], dfn[i.X]); 
  if (low[u] == dfn[u] && f != -1) is_bridge[f] = 1;
void solve(int n) {
  is_bridge.resize(SZ(edge));
  for (int i = 1; i <= n; ++i)</pre>
    if (!dfn[i]) dfs(i, -1);
```

2.3 2SAT (SCC)*

```
struct SAT { // 0-base
  int low[N], dfn[N], bln[N], n, Time, nScc;
  bool instack[N], istrue[N];
  stack<int> st;
  vector<int> G[N], SCC[N];
  void init(int _n) {
    n = _n; // assert(n * 2 <= N);
for (int i = 0; i < n + n; ++i) G[i].clear();</pre>
  void add_edge(int a, int b) { G[a].eb(b); }
  int rv(int a) {
    if (a >= n) return a - n;
    return a + n;
  void add_clause(int a, int b) {
    add_edge(rv(a), b), add_edge(rv(b), a);
  void dfs(int u) {
    dfn[u] = low[u] = ++Time;
    instack[u] = 1, st.push(u);
    for (int i : G[u])
      if (!dfn[i])
        dfs(i), low[u] = min(low[i], low[u]);
       else if (instack[i] && dfn[i] < dfn[u])</pre>
        low[u] = min(low[u], dfn[i]);
    if (low[u] == dfn[u]) {
       int tmp;
      do {
        tmp = st.top(), st.pop();
         instack[tmp] = 0, bln[tmp] = nScc;
      } while (tmp != u);
       ++nScc;
    }
  bool solve() {
    Time = nScc = 0;
    for (int i = 0; i < n + n; ++i)</pre>
      SCC[i].clear(), low[i] = dfn[i] = bln[i] = 0;
    for (int i = 0; i < n + n; ++i)</pre>
      if (!dfn[i]) dfs(i);
    for (int i = 0; i < n + n; ++i) SCC[bln[i]].eb(i);</pre>
    for (int i = 0; i < n; ++i) {</pre>
      if (bln[i] == bln[i + n]) return false;
      istrue[i] = bln[i] < bln[i + n];</pre>
      istrue[i + n] = !istrue[i];
    return true:
  }
};
```

2.4 MinimumMeanCycle*

```
int road[N][N]; // input here
struct MinimumMeanCycle {
  int dp[N + 5][N], n;
  pii solve() {
```

```
int a = -1, b = -1, L = n + 1;
     for (int i = 2; i <= L; ++i)</pre>
       for (int k = 0; k < n; ++k)
          for (int j = 0; j < n; ++j)
            dp[i][j] =
              min(dp[i - 1][k] + road[k][j], dp[i][j]);
     for (int i = 0; i < n; ++i) {</pre>
       if (dp[L][i] >= INF) continue;
       int ta = 0, tb = 1;
       for (int j = 1; j < n; ++j)</pre>
          \quad \textbf{if} \ (\texttt{dp[j][i]} < \texttt{INF \&\&}
                       j) < (dp[L][i] - dp[j][i]) * tb)</pre>
            ta * (L -
            ta = dp[L][i] - dp[j][i], tb = L - j;
       if (ta == 0) continue;
       if (a == -1 || a * tb > ta * b) a = ta, b = tb;
     if (a != -1) {
       int g = __gcd(a, b);
       return pii(a / g, b / g);
     return pii(-1LL, -1LL);
   void init(int _n) {
     n = _n;
for (int i = 0; i < n; ++i)</pre>
       for (int j = 0; j < n; ++j) dp[i + 2][j] = INF;</pre>
  }
};
```

2.5 Virtual Tree*

```
vector<int> vG[N];
int top, st[N];
void insert(int u) {
  if (top == -1) return st[++top] = u, void();
  int p = LCA(st[top], u);
  if (p == st[top]) return st[++top] = u, void();
  while (top >= 1 && dep[st[top - 1]] >= dep[p])
  vG[st[top - 1]].eb(st[top]), --top;
  if (st[top] != p)
    vG[p].eb(st[top]), --top, st[++top] = p;
  st[++top] = u;
}
void reset(int u) {
  for (int i : vG[u]) reset(i);
  vG[u].clear();
void solve(vector<int> &v) {
  top = -1;
  sort(ALL(v),
    [&](int a, int b) { return dfn[a] < dfn[b]; });</pre>
  for (int i : v) insert(i);
  while (top > 0) vG[st[top - 1]].eb(st[top]), --top;
  // do something
  reset(v[0]);
```

2.6 Maximum Clique Dyn*

```
struct MaxClique { // fast when N <= 100</pre>
 bitset < N > G[N], cs[N];
  int ans, sol[N], q, cur[N], d[N], n;
  void init(int _n) {
    n = _n;
    for (int i = 0; i < n; ++i) G[i].reset();</pre>
  void add_edge(int u, int v) {
    G[u][v] = G[v][u] = 1;
  void pre_dfs(vector<int> &r, int l, bitset<N> mask) {
    if (1 < 4) {
      for (int i : r) d[i] = (G[i] & mask).count();
      sort(ALL(r)
          , [&](int x, int y) { return d[x] > d[y]; });
    vector<int> c(SZ(r));
    int lft = max(ans - q + 1, 1), rgt = 1, tp = 0;
    cs[1].reset(), cs[2].reset();
    for (int p : r) {
      while ((cs[k] & G[p]).any()) ++k;
      if (k > rgt) cs[++rgt + 1].reset();
      cs[k][p] = 1;
```

```
if (k < lft) r[tp++] = p;</pre>
    for (int k = lft; k <= rgt; ++k)</pre>
      for (int p = cs[k]._Find_first
           (); p < N; p = cs[k]._Find_next(p))
        r[tp] = p, c[tp] = k, ++tp;
    dfs(r, c, l + 1, mask);
  void dfs(vector<</pre>
      int> &r, vector<int> &c, int l, bitset<N> mask) {
    while (!r.empty()) {
      int p = r.back();
      r.pb(), mask[p] = 0;
      if (q + c.back() <= ans) return;</pre>
      cur[q++] = p;
      vector<int> nr;
      for (int i : r) if (G[p][i]) nr.eb(i);
      if (!nr.empty()) pre_dfs(nr, l, mask & G[p]);
      else if (q > ans) ans = q, copy_n(cur, q, sol);
      c.pb(), --q;
    }
  int solve() {
    vector<int> r(n);
    ans = q = 0, iota(ALL(r), \theta);
    pre_dfs(r, 0, bitset<N>(string(n, '1')));
    return ans;
 }
};
```

2.7 Minimum Steiner Tree*

struct SteinerTree { // 0-base

```
int n, dst[N][N], dp[1 << T][N], tdst[N];</pre>
  int vcst[N]; // the cost of vertexs
  void init(int _n) {
    n = _n;
for (int i = 0; i < n; ++i) {</pre>
       fill_n(dst[i], n, INF);
       dst[i][i] = vcst[i] = 0;
  void chmin(int &x, int val) {
    x = min(x, val);
  void add_edge(int ui, int vi, int wi) {
    chmin(dst[ui][vi], wi);
  void shortest_path() {
    for (int k = 0; k < n; ++k)</pre>
       for (int i = 0; i < n; ++i)</pre>
         for (int j = 0; j < n; ++j)</pre>
           chmin(dst[i][j], dst[i][k] + dst[k][j]);
  int solve(const vector<int>& ter) {
    shortest_path();
    int t = SZ(ter), full = (1 << t) - 1;
for (int i = 0; i <= full; ++i)</pre>
       fill_n(dp[i], n, INF);
    copy_n(vcst, n, dp[0]);
    for (int msk = 1; msk <= full; ++msk) {</pre>
       if (!(msk & (msk - 1))) {
         int who =
                    __lg(msk);
         for (int i = 0; i < n; ++i)</pre>
           dp[msk
               ][i] = vcst[ter[who]] + dst[ter[who]][i];
      for (int i = 0; i < n; ++i)</pre>
         for (int sub = (
             msk - 1) & msk; sub; sub = (sub - 1) & msk)
           chmin(dp[msk][i],
               dp[sub][i] + dp[msk ^ sub][i] - vcst[i]);
       for (int i = 0; i < n; ++i) {</pre>
         tdst[i] = INF;
         for (int j = 0; j < n; ++j)</pre>
           chmin(tdst[i], dp[msk][j] + dst[j][i]);
      copy_n(tdst, n, dp[msk]);
    return *min_element(dp[full], dp[full] + n);
\}; // O(V 3^T + V^2 2^T)
```

2.8 Dominator Tree*

```
| struct dominator_tree { // 1-base
```

```
vector < int > G[N], rG[N];
int n, pa[N], dfn[N], id[N], Time;
   int semi[N], idom[N], best[N];
vector < int > tree[N]; // dominator_tree
   void init(int _n) {
     n = _n;
for (int i = 1; i <= n; ++i)</pre>
       G[i].clear(), rG[i].clear();
   void add_edge(int u, int v) {
     G[u].eb(v), rG[v].eb(u);
   void dfs(int u) {
     id[dfn[u] = ++Time] = u;
     for (auto v : G[u])
       if (!dfn[v]) dfs(v), pa[dfn[v]] = dfn[u];
   int find(int y, int x) {
     if (y <= x) return y;</pre>
     int tmp = find(pa[y], x);
     if (semi[best[y]] > semi[best[pa[y]]])
       best[y] = best[pa[y]];
     return pa[y] = tmp;
   void tarjan(int root) {
     Time = 0;
     for (int i = 1; i <= n; ++i) {</pre>
       dfn[i] = idom[i] = 0;
       tree[i].clear();
       best[i] = semi[i] = i;
     dfs(root);
     for (int i = Time; i > 1; --i) {
       int u = id[i];
       for (auto v : rG[u])
   if (v = dfn[v]) {
            find(v, i);
            semi[i] = min(semi[i], semi[best[v]]);
       tree[semi[i]].eb(i);
       for (auto v : tree[pa[i]]) {
          find(v, pa[i]);
            semi[best[v]] == pa[i] ? pa[i] : best[v];
       tree[pa[i]].clear();
     for (int i = 2; i <= Time; ++i) {</pre>
       if (idom[i] != semi[i]) idom[i] = idom[idom[i]];
       tree[id[idom[i]]].eb(id[i]);
  }
};
```

2.9 Minimum Arborescence*

```
/* TODO
DSU: disjoint set
- DSU(n), .boss(x), .Union(x, y)
min heap <
    T, Info>: min heap for type {T, Info} with lazy tag
 .push({w, i}),
    .top(), .join(heap), .pop(), .empty(), .add_lazy(v)
struct E { int s, t; int w; }; // O-base
vector<int> dmst(const vector<E> &e, int n, int root) {
  vector<min_heap<int, int>> h(n * 2);
  for (int i = 0; i < SZ(e); ++i)</pre>
    h[e[i].t].push({e[i].w, i});
  DSU dsu(n * 2);
  vector<int> v(n * 2, -1), pa(n * 2, -1), r(n * 2);
  v[root] = n + 1;
  int pc = n;
  for (int i = 0; i < n; ++i) if (v[i] == -1) {</pre>
    for (int p = i; v[p]
          == -1 \mid \mid v[p] == i; p = dsu.boss(e[r[p]].s)) {
      if (v[p] == i) {
         int q = p; p = pc++;
           h[q].add_lazy(-h[q].top().X);
           pa[q] = p, dsu.Union(p, q), h[p].join(h[q]);
        } while ((q = dsu.boss(e[r[q]].s)) != p);
      v[p] = i;
      while (!h[p].
           empty() \&\& dsu.boss(e[h[p].top().Y].s) == p)
```

2.10 Vizing's theorem*

```
namespace vizing { // returns
  edge coloring in adjacent matrix G. 1 - based
const int N = 105
int C[N][N], G[N][N], X[N], vst[N], n;
void init(int _n) { n = _n;
  for (int i = 0; i <= n; ++i)
  for (int j = 0; j <= n; ++j)</pre>
      C[i][j] = G[i][j] = 0;
void solve(vector<pii> &E) {
  auto update = [&](int u)
  { for (X[u] = 1; C[u][X[u]]; ++X[u]); };
auto color = [&](int u, int v, int c) {
    int p = G[u][v];
    G[u][v] = G[v][u] = c;
    C[u][c] = v, C[v][c] = u;
    C[u][p] = C[v][p] = 0;
    if (p) X[u] = X[v] = p;
    else update(u), update(v);
    return p;
  };
  auto flip = [&](int u, int c1, int c2) {
    int p = C[u][c1];
    swap(C[u][c1], C[u][c2]);
    if (p) G[u][p] = G[p][u] = c2;
if (!C[u][c1]) X[u] = c1;
    if (!C[u][c2]) X[u] = c2;
    return p;
  fill_n(X + 1, n, 1);
for (int t = 0; t < SZ(E); ++t) {
    int u = E[t
         ].X, v0 = E[t].Y, v = v0, c0 = X[u], c = c0, d;
    vector<pii> L;
    fill_n(vst + 1, n, 0);
    while (!G[u][v0]) {
       L.emplace_back(v, d = X[v]);
       if (!C[v][c]) for (int a = SZ(
           L) - 1; a >= 0; --a) c = color(u, L[a].X, c);
       else if (!C[u][d]) for (int a = SZ(L
           ) - 1; a >= 0; --a) color(u, L[a].X, L[a].Y);
       else if (vst[d]) break;
       else vst[d] = 1, v = C[u][d];
    if (!G[u][v0]) {
  for (; v; v = flip(v, c, d), swap(c, d));
       if (int a; C[u][c0]) {
              a = SZ(L) - 2; a >= 0 && L[a].Y != c; --a);
         for (; a >= 0; --a) color(u, L[a].X, L[a].Y);
       else --t;
    }
  }
} // namespace vizing
```

2.11 Minimum Clique Cover*

```
struct Clique_Cover { // 0-base, O(n2^n)
  int co[1 << N], n, E[N];
  int dp[1 << N];
  void init(int _n) {
    n = _n, fill_n(dp, 1 << n, 0);
    fill_n(E, n, 0), fill_n(co, 1 << n, 0);
  }
  void add_edge(int u, int v) {
    E[u] |= 1 << v, E[v] |= 1 << u;
  }</pre>
```

```
int solve() {
      for (int i = 0; i < n; ++i)</pre>
        co[1 << i] = E[i] | (1 << i);
      co[0] = (1 << n) - 1;

dp[0] = (n & 1) * 2 - 1;
      for (int i = 1; i < (1 << n); ++i) {</pre>
        int t = i & -i;
dp[i] = -dp[i ^ t];
         co[i] = co[i ^ t] & co[t];
      for (int i = 0; i < (1 << n); ++i)</pre>
        co[i] = (co[i] \& i) == i;
      fwt(co, 1 << n, 1);
      for (int ans = 1; ans < n; ++ans) {
  int sum = 0; // probabilistic</pre>
        for (int i = 0; i < (1 << n); ++i)
  sum += (dp[i] *= co[i]);</pre>
        if (sum) return ans;
      return n;
   }
};
```

2.12 NumberofMaximalClique*

```
struct BronKerbosch { // 1-base
  int n, a[N], g[N][N];
  int S, all[N][N], some[N][N], none[N][N];
  void init(int _n) {
    for (int j = 1; j <= n; ++j) g[i][j] = 0;</pre>
  void add_edge(int u, int v) {
    g[u][v] = g[v][u] = 1;
  void dfs(int d, int an, int sn, int nn) {
    if (S > 1000) return; // pruning
    if (sn == 0 && nn == 0) ++S;
    int u = some[d][0];
    for (int i = 0; i < sn; ++i) {</pre>
      int v = some[d][i];
      if (g[u][v]) continue;
int tsn = 0, tnn = 0;
      copy_n(all[d], an, all[d + 1]);
       all[d + 1][an] = v;
      for (int j = 0; j < sn; ++j)</pre>
         if (g[v][some[d][j]])
      some[d + 1][tsn++] = some[d][j];
for (int j = 0; j < nn; ++j)
         if (g[v][none[d][j]])
      none[d + 1][tnn++] = none[d][j];
dfs(d + 1, an + 1, tsn, tnn);
      some[d][i] = 0, none[d][nn++] = v;
    }
  int solve() {
    iota(some[0], some[0] + n, 1);
    S = 0, dfs(0, 0, n, 0);
    return S;
  }
};
```

3 Data Structure

3.1 Interval Container*

```
/* Add and
      remove intervals from a set of disjoint intervals.
   Will merge the added interval with
      any overlapping intervals in the set when adding.
 * Intervals are [inclusive, exclusive). */
set<pii>::
    iterator addInterval(set<pii>& is, int L, int R) {
  if (L == R) return is.end();
auto it = is.lower_bound({L, R}), before = it;
  while (it != is.end() && it->X <= R) {</pre>
    R = max(R, it->Y);
    before = it = is.erase(it);
  if (it != is.begin() && (--it)->Y >= L) {
    L = min(L, it->X);
R = max(R, it->Y);
    is.erase(it);
  }
  return is.insert(before, pii(L, R));
```

```
}
void removeInterval(set<pii>& is, int L, int R) {
   if (L == R) return;
   auto it = addInterval(is, L, R);
   auto r2 = it->Y;
   if (it->X == L) is.erase(it);
   else (int&)it->Y = L;
   if (R != r2) is.emplace(R, r2);
}
```

3.2 Heavy light Decomposition

```
struct Heavy_light_Decomposition { // 1-base
  int n, ulink[N], deep[N], mxson[N], w[N], pa[N];
  int t, pl[N], data[N], dt[N], bln[N], edge[N], et;
  vector<pii> G[N];
 G[i].clear(), mxson[i] = 0;
  void add_edge(int a, int b, int w) {
   G[a].eb(pii(b, et));
    G[b].eb(pii(a, et));
    edge[et++] = w;
  void dfs(int u, int f, int d) {
  w[u] = 1, pa[u] = f, deep[u] = d++;
    for (auto &i : G[u])
      if (i.X != f) {
        dfs(i.X, u, d), w[u] += w[i.X];
        if (w[mxson[u]] < w[i.X]) mxson[u] = i.X;
      } else bln[i.Y] = u, dt[u] = edge[i.Y];
  void cut(int u, int link) {
    data[pl[u] = t++] = dt[u], ulink[u] = link;
    if (!mxson[u]) return;
    cut(mxson[u], link);
    for (auto i : G[u])
      if (i.X != pa[u] && i.X != mxson[u])
        cut(i.X, i.X);
  void build() { dfs(1, 1, 1), cut(1, 1), /*build*/; }
  int query(int a, int b) {
    int ta = ulink[a], tb = ulink[b], re = 0;
    while (ta != tb)
      if (deep[ta] < deep[tb])
  /*query*/, tb = ulink[b = pa[tb]];</pre>
      else /*query*/, ta = ulink[a = pa[ta]];
    if (a == b) return re;
    if (pl[a] > pl[b]) swap(a, b);
    /*query*/
    return re;
};
```

3.3 Centroid Decomposition*

```
struct Cent_Dec { // 1-base
  vector<pii> G[N];
 pii info[N]; // store info. of itself
pii upinfo[N]; // store info. of climbing up
  int n, pa[N], layer[N], sz[N], done[N];
  int dis[__lg(N) + 1][N];
  void init(int _n) {
    n = _n, layer[0] = -1;
    fill_n(pa + 1, n, 0), fill_n(done + 1, n, 0);
    for (int i = 1; i <= n; ++i) G[i].clear();</pre>
  void add_edge(int a, int b, int w) {
    G[a].eb(pii(b, w)), G[b].eb(pii(a, w));
  void get_cent(
    int u, int f, int &mx, int &c, int num) {
    int mxsz = 0;
    sz[u] = 1;
    for (pii e : G[u])
      if (!done[e.X] && e.X != f) {
        get_cent(e.X, u, mx, c, num);
        sz[u] += sz[e.X], mxsz = max(mxsz, sz[e.X]);
    if (mx > max(mxsz, num - sz[u]))
      mx = max(mxsz, num - sz[u]), c = u;
  void dfs(int u, int f, int d, int org) {
    // if required, add self info or climbing info
```

```
dis[layer[org]][u] = d;
    for (pii e : G[u])
      if (!done[e.X] && e.X != f)
        dfs(e.X, u, d + e.Y, org);
  int cut(int u, int f, int num) {
    int mx = 1e9, c = 0, lc;
    get_cent(u, f, mx, c, num);
    done[c] = 1, pa[c] = f, layer[c] = layer[f] + 1;
    for (pii e : G[c])
      if (!done[e.X]) {
        if (sz[e.X] > sz[c])
          lc = cut(e.X, c, num - sz[c]);
        else lc = cut(e.X, c, sz[e.X]);
        upinfo[lc] = pii(), dfs(e.X, c, e.Y, c);
    return done[c] = 0, c;
  void build() { cut(1, 0, n); }
  void modify(int u) {
    for (int a = u, ly = layer[a]; a;
    a = pa[a], --ly) {
      info[a].X += dis[ly][u], ++info[a].Y;
      if (pa[a])
        upinfo[a].X += dis[ly - 1][u], ++upinfo[a].Y;
    }
  int query(int u) {
    int rt = 0;
    for (int a = u, ly = layer[a]; a;
         a = pa[a], --ly) {
      rt += info[a].X + info[a].Y * dis[ly][u];
      if (pa[a])
        rt -:
          upinfo[a].X + upinfo[a].Y * dis[ly - 1][u];
    return rt;
};
```

3.4 Link cut tree*

```
struct Splay { // xor-sum
  static Splay nil;
  Splay *ch[2], *f;
  int val, sum, rev, size;
  Splay(int _val = 0)
     val(_val), sum(_val), rev(0), size(1) {
    f = ch[0] = ch[1] = &nil;
  bool isr() {
    return f->ch[0] != this && f->ch[1] != this;
  int dir() { return f->ch[0] == this ? 0 : 1; }
  void setCh(Splay *c, int d) {
    ch[d] = c;
    if (c != &nil) c->f = this;
    pull();
 void give_tag(int r)
{ if (r) swap(ch[0], ch[1]), rev ^= 1; }
  void push() {
    if (ch[0] != &nil) ch[0]->give_tag(rev);
    if (ch[1] != &nil) ch[1]->give_tag(rev);
    rev = 0:
  void pull() {
    // take care of the nil!
    size = ch[0] -> size + ch[1] -> size + 1;
    sum = ch[0] -> sum ^ ch[1] -> sum ^ val;
    if (ch[0] != &nil) ch[0]->f = this;
    if (ch[1] != &nil) ch[1]->f = this;
} Splay::nil;
Splay *nil = &Splay::nil;
void rotate(Splay *x) {
  Splay *p = x -> f;
  int d = x->dir();
  if (!p->isr()) p->f->setCh(x, p->dir());
  else x->f = p->f;
  p->setCh(x->ch[!d], d);
  x->setCh(p, !d);
 p->pull(), x->pull();
void splay(Splay *x) {
 vector<Splay *> splayVec;
```

```
for (Splay *q = x;; q = q->f) {
    splayVec.eb(q);
    if (q->isr()) break;
  reverse(ALL(splayVec));
  for (auto it : splayVec) it->push();
  while (!x->isr()) {
    if (x->f->isr()) rotate(x);
    else if (x->dir() == x->f->dir())
     rotate(x->f), rotate(x);
    else rotate(x), rotate(x);
 }
Splay *access(Splay *x) {
 Splay *q = nil;
 for (; x != nil; x = x->f)
   splay(x), x->setCh(q, 1), q = x;
void root_path(Splay *x) { access(x), splay(x); }
void chroot(Splay *x) {
 root_path(x), x->rev ^= 1;
  x->push(), x->pull();
void split(Splay *x, Splay *y) {
 chroot(x), root_path(y);
void link(Splay *x, Splay *y) {
 root_path(x), chroot(y);
 x->setCh(y, 1);
void cut(Splay *x, Splay *y) {
 split(x, y);
  if (y->size != 5) return;
 y->push();
 y - ch[0] = y - ch[0] - f = nil;
Splay *get_root(Splay *x) {
 for (root_path(x); x->ch[0] != nil; x = x->ch[0])
   x->push();
  splay(x);
  return x;
bool conn(Splay *x, Splay *y) {
 return get_root(x) == get_root(y);
Splay *lca(Splay *x, Splay *y) {
 access(x), root_path(y);
  if (y->f == nil) return y;
  return v->f;
void change(Splay *x, int val) {
 splay(x), x->val = val, x->pull();
int query(Splay *x, Splay *y) {
 split(x, y);
  return y->sum;
```

3.5 KDTree

```
namespace kdt {
int root, lc[maxn], rc[maxn], xl[maxn], xr[maxn],
 yl[maxn], yr[maxn];
point p[maxn];
int build(int l, int r, int dep = 0) {
  if (l == r) return -1;
  function < bool(const point &, const point &) > f =
    [dep](const point &a, const point &b) {
      if (dep & 1) return a.x < b.x;</pre>
      else return a.y < b.y;</pre>
    };
  int m = (l + r) >> 1;
  nth_element(p + l, p + m, p + r, f);
  xl[m] = xr[m] = p[m].x;
  yl[m] = yr[m] = p[m].y;
  lc[m] = build(l, m, dep + 1);
  if (~lc[m]) {
    xl[m] = min(xl[m], xl[lc[m]]);
    xr[m] = max(xr[m], xr[lc[m]]);
yl[m] = min(yl[m], yl[lc[m]]);
    yr[m] = max(yr[m], yr[lc[m]]);
  rc[m] = build(m + 1, r, dep + 1);
  if (~rc[m]) {
    xl[m] = min(xl[m], xl[rc[m]]);
```

```
xr[m] = max(xr[m], xr[rc[m]]);
    yl[m] = min(yl[m], yl[rc[m]]);
    yr[m] = max(yr[m], yr[rc[m]]);
  return m;
bool bound(const point &q, int o, long long d) {
  double ds = sqrt(d + 1.0);
  if (q.x < xl[o] - ds || q.x > xr[o] + ds ||
    q.y < yl[o] - ds || q.y > yr[o] + ds)
    return false:
  return true:
long long dist(const point &a, const point &b) {
  return (a.x - b.x) * 111 * (a.x - b.x) +
    (a.y - b.y) * 1ll * (a.y - b.y);
void dfs(
  const point &q, long long &d, int o, int dep = 0) {
  if (!bound(q, o, d)) return;
  long long cd = dist(p[o], q);
if (cd != 0) d = min(d, cd);
  if ((dep & 1) && q.x < p[o].x ||</pre>
    !(dep & 1) && q.y < p[o].y) {
     if (~lc[o]) dfs(q, d, lc[o], dep + 1);
    if (~rc[o]) dfs(q, d, rc[o], dep + 1);
  } else {
    if (~rc[o]) dfs(q, d, rc[o], dep + 1);
if (~lc[o]) dfs(q, d, lc[o], dep + 1);
  }
void init(const vector<point> &v) {
  for (int i = 0; i < v.size(); ++i) p[i] = v[i];</pre>
  root = build(0, v.size());
long long nearest(const point &q) {
  long long res = 1e18;
  dfs(q, res, root);
  return res;
} // namespace kdt
```

4 Flow/Matching

4.1 Kuhn Munkres*

```
struct KM { // 0-base, maximum matching
int w[N][N], hl[N], hr[N], slk[N];
  int fl[N], fr[N], pre[N], qu[N], ql, qr, n;
bool vl[N], vr[N];
  void init(int _n) {
    n = _n;
for (int i = 0; i < n; ++i)</pre>
      fill_n(w[i], n, -INF);
  void add_edge(int a, int b, int wei) {
    w[a][b] = wei;
  bool Check(int x) {
    if (vl[x] = 1, \sim fl[x])
      return vr[qu[qr++] = fl[x]] = 1;
    while (\sim x) swap(x, fr[fl[x] = pre[x]]);
    return 0;
  void bfs(int s) {
    fill_n(slk
         , n, INF), fill_n(vl, n, 0), fill_n(vr, n, 0);
    ql = qr = 0, qu[qr++] = s, vr[s] = 1;
    for (int d; ; ) {
      while (ql < qr)
for (int x = 0, y = qu[ql++]; x < n; ++x)</pre>
           if (!vl[x] && slk
                [x] >= (d = hl[x] + hr[y] - w[x][y])) {
              if (pre[x] = y, d) slk[x] = d;
             else if (!Check(x)) return;
      d = INF;
      for (int x = 0; x < n; ++x)
         if (!vl[x] && d > slk[x]) d = slk[x];
      for (int x = 0; x < n; ++x) {
         if (vl[x]) hl[x] += d;
         else slk[x] -= d;
         if (vr[x]) hr[x] -= d;
       for (int x = 0; x < n; ++x)
         if (!vl[x] && !slk[x] && !Check(x)) return;
```

4.2 MincostMaxflow*

```
struct MinCostMaxFlow { // 0-base
  struct Edge {
    int from, to, cap, flow, cost, rev;
  } *past[N];
  vector<Edge> G[N];
  int inq[N], n, s, t;
  int dis[N], up[N], pot[N];
  bool BellmanFord() {
    fill_n(dis, n, INF), fill_n(inq, n, 0);
    queue<int> q;
    auto relax = [&](int u, int d, int cap, Edge *e) {
      if (cap > 0 && dis[u] > d) {
        dis[u] = d, up[u] = cap, past[u] = e;
        if (!inq[u]) inq[u] = 1, q.push(u);
      }
    };
    relax(s, 0, INF, 0);
    while (!q.empty()) {
      int u = q.front();
      q.pop(), inq[u] = 0;
      for (auto &e : G[u]) {
        int d2 = dis[u] + e.cost + pot[u] - pot[e.to];
        relax
            (e.to, d2, min(up[u], e.cap - e.flow), &e);
     }
    }
    return dis[t] != INF;
  }
  void solve(int _s,
       int _t, int &flow, int &cost, bool neg = true) {
    s = _s, t = _t, flow = 0, cost = 0;
    if (neg) BellmanFord(), copy_n(dis, n, pot);
    for (; BellmanFord(); copy_n(dis, n, pot)) {
      for (int
          i = 0; i < n; ++i) dis[i] += pot[i] - pot[s];
      flow += up[t], cost += up[t] * dis[t];
      for (int i = t; past[i]; i = past[i]->from) {
        auto &e = *past[i];
        e.flow += up[t], G[e.to][e.rev].flow -= up[t];
   }
  void init(int _n) {
    n = _n, fill_n(pot, n, \theta);
    for (int i = 0; i < n; ++i) G[i].clear();</pre>
  void add_edge(int a, int b, int cap, int cost) {
    G[a].eb(Edge{a, b, cap, 0, cost, SZ(G[b])});
    G[b].eb(Edge{b, a, 0, 0, -cost, SZ(G[a]) - 1});
};
```

4.3 Maximum Simple Graph Matching

```
struct GenMatch { // 1-base
  int V, pr[N];
  bool el[N][N], inq[N], inp[N], inb[N];
  int st, ed, nb, bk[N], djs[N], ans;
  void init(int _V) {
    V = _V;
    for (int i = 0; i <= V; ++i) {
        for (int j = 0; j <= V; ++j) el[i][j] = 0;
        pr[i] = bk[i] = djs[i] = 0;
        inq[i] = inp[i] = inb[i] = 0;
    }
}
void add_edge(int u, int v) {
    el[u][v] = el[v][u] = 1;
}
int lca(int u, int v) {</pre>
```

```
while (1)
       if (u = djs[u], inp[u] = true, u == st) break;
       else u = bk[pr[u]];
     while (1)
       if (v = djs[v], inp[v]) return v;
       else v = bk[pr[v]];
     return v;
   void upd(int u) {
     for (int v; djs[u] != nb;) {
      v = pr[u], inb[djs[u]] = inb[djs[v]] = true;
       u = bk[v];
       if (djs[u] != nb) bk[u] = v;
     }
   }
   void blo(int u, int v, queue<int> &qe) {
     nb = lca(u, v), fill_n(inb, V + 1, 0);
     upd(u), upd(v);
     if (djs[u] != nb) bk[u] = v;
     if (djs[v] != nb) bk[v] = u;
     for (int tu = 1; tu <= V; ++tu)</pre>
       if (inb[djs[tu]])
         if (djs[tu] = nb, !inq[tu])
           qe.push(tu), inq[tu] = 1;
   void flow() {
     fill_n(inq + 1, V, 0), fill_n(bk + 1, V, 0);
     iota(djs + 1, djs + V + 1, 1);
     queue<int> qe;
     \dot{q}e.push(st), inq[st] = 1, ed = 0;
     while (!qe.empty()) {
       int u = qe.front();
       qe.pop();
       for (int v = 1; v <= V; ++v)</pre>
         if (el[u][v] && djs[u] != djs[v] &&
           pr[u] != v) {
           if ((v == st) ||
             (pr[v] > 0 \&\& bk[pr[v]] > 0)) {
             blo(u, v, qe);
           } else if (!bk[v]) {
             if (bk[v] = u, pr[v] > 0) {
               if (!inq[pr[v]]) qe.push(pr[v]);
             } else {
               return ed = v, void();
             }
           }
         }
    }
   void aug() {
     for (int u = ed, v, w; u > 0;)
       v = bk[u], w = pr[v], pr[v] = u, pr[u] = v,
   int solve() {
     fill_n(pr, V + 1, \theta), ans = \theta;
     for (int u = 1; u <= V; ++u)</pre>
       if (!pr[u])
         if (st = u, flow(), ed > \theta) aug(), ++ans;
     return ans;
};
```

 $fill_n(inp, V + 1, 0);$

4.4 Minimum Weight Matching (Clique version)*

```
if (onstk[m] || SPFA(m)) return 1;
            --tp, onstk[v] = 0;
         }
     onstk[u] = 0, --tp;
     return 0;
   int solve() { // find a match
     for (int i = 0; i < n; ++i) match[i] = i ^ 1;</pre>
     while (1) {
       int found = 0;
       fill_n(dis, n, 0);
       fill_n(onstk, n, 0);
for (int i = 0; i < n; ++i)
         if (tp = 0, !onstk[i] && SPFA(i))
           for (found = 1; tp >= 2;) {
  int u = stk[--tp];
              int v = stk[--tp];
              match[u] = v, match[v] = u;
       if (!found) break;
     int ret = 0;
     for (int i = 0; i < n; ++i)</pre>
       ret += edge[i][match[i]];
     return ret >> 1;
};
```

4.5 SW-mincut

```
struct SW{ // global min cut, O(V^3)
   #define REP for (int i = 0; i < n; ++i)
   static const int MXN = 514, INF = 2147483647;</pre>
  int vst[MXN], edge[MXN][MXN], wei[MXN];
  void init(int n) {
    REP fill_n(edge[i], n, 0);
  void addEdge(int u, int v, int w){
    edge[u][v] += w; edge[v][u] += w;
  int search(int &s, int &t, int n){
    fill_n(vst, n, 0), fill_n(wei, n, 0);
    s = t = -1;
     int mx, cur;
     for (int j = 0; j < n; ++j) {</pre>
       mx = -1, cur = 0;
       REP if (wei[i] > mx) cur = i, mx = wei[i];
       vst[cur] = 1, wei[cur] = -1;
       s = t; t = cur;
       REP if (!vst[i]) wei[i] += edge[cur][i];
    }
    return mx;
  int solve(int n) {
    int res = INF;
    for (int x, y; n > 1; n--){
  res = min(res, search(x, y, n));
       REP edge[i][x] = (edge[x][i] += edge[y][i]);
       REP {
         edge[y][i] = edge[n - 1][i];
         edge[i][y] = edge[i][n - 1];
       }
    return res;
} sw;
```

4.6 BoundedFlow*(Dinic*)

```
struct BoundedFlow { // O-base
    struct edge {
        int to, cap, flow, rev;
    };
    vector<edge> G[N];
    int n, s, t, dis[N], cur[N], cnt[N];
    void init(int _n) {
        n = _n;
        for (int i = 0; i < n + 2; ++i)
            G[i].clear(), cnt[i] = 0;
    }
    void add_edge(int u, int v, int lcap, int rcap) {
        cnt[u] -= lcap, cnt[v] += lcap;
        G[u].eb(edge{v, rcap, lcap, SZ(G[v])});
        G[v].eb(edge{u, 0, 0, SZ(G[u]) - 1});
    }</pre>
```

```
void add_edge(int u, int v, int cap) {
  G[u].eb(edge{v, cap, 0, SZ(G[v])});
     G[v].eb(edge{u, 0, 0, SZ(G[u]) - 1});
   int dfs(int u, int cap) {
     if (u == t || !cap) return cap;
     for (int &i = cur[u]; i < SZ(G[u]); ++i) {</pre>
        edge &e = G[u][i];
        if (dis[e.to] == dis[u] + 1 && e.cap != e.flow) {
          int df = dfs(e.to, min(e.cap - e.flow, cap));
          if (df) {
            e.flow += df, G[e.to][e.rev].flow -= df;
            return df;
       }
     dis[u] = -1;
     return 0;
   bool bfs() {
     fill_n(dis, n + 3, -1);
     queue<int> q;
     q.push(s), dis[s] = 0;
     while (!q.empty()) {
       int u = q.front();
       q.pop();
        for (edge &e : G[u])
          if (!~dis[e.to] && e.flow != e.cap)
            q.push(e.to), dis[e.to] = dis[u] + 1;
     return dis[t] != -1;
   int maxflow(int _s, int _t) {
     s = _s, t = _t;
int flow = 0, df;
     while (bfs()) {
       fill_n(cur, n + 3, 0);
        while ((df = dfs(s, INF))) flow += df;
     return flow;
   bool solve() {
     int sum = 0;
     for (int i = 0; i < n; ++i)</pre>
       if (cnt[i] > 0)
        add_edge(n + 1, i, cnt[i]), sum += cnt[i];
else if (cnt[i] < 0) add_edge(i, n + 2, -cnt[i]);
     if (sum != maxflow(n + 1, n + 2)) sum = -1;
     for (int i = 0; i < n; ++i)</pre>
       if (cnt[i] > 0)
       G[n + 1].pb(), G[i].pb();
else if (cnt[i] < 0)</pre>
     G[i].pb(), G[n + 2].pb();
return sum != -1;
   int solve(int _s, int _t) {
     add_edge(_t, _s, INF);
     if (!solve()) return -1; // invalid flow
     int x = G[_t].back().flow;
     return G[_t].pb(), G[_s].pb(), x;
  }
};
```

4.7 Gomory Hu tree*

```
MaxFlow Dinic;
int g[maxn];
void GomoryHu(int n) { // 0-base
  fill_n(g, n, 0);
  for (int i = 1; i < n; ++i) {
    Dinic.reset();
    add_edge(i, g[i], Dinic.maxflow(i, g[i]));
    for (int j = i + 1; j <= n; ++j)
        if (g[j] == g[i] && ~Dinic.dis[j])
        g[j] = i;
  }
}</pre>
```

4.8 Minimum Cost Circulation*

```
struct MinCostCirculation { // 0-base
   struct Edge {
     int from, to, cap, fcap, flow, cost, rev;
   } *past[N];
   vector<Edge> G[N];
   int dis[N], inq[N], n;
```

```
void BellmanFord(int s) {
    fill_n(dis, n, INF), fill_n(inq, n, 0);
    queue<int> q;
    auto relax = [&](int u, int d, Edge *e) {
       if (dis[u] > d) {
         dis[u] = d, past[u] = e;
if (!inq[u]) inq[u] = 1, q.push(u);
      }
    };
    relax(s, 0, 0);
    while (!q.empty()) {
  int u = q.front();
       q.pop(), inq[u] = 0;
       for (auto &e : G[u])
         if (e.cap > e.flow)
           relax(e.to, dis[u] + e.cost, &e);
    }
  void try_edge(Edge &cur) {
    if (cur.cap > cur.flow) return ++cur.cap, void();
    BellmanFord(cur.to);
    if (dis[cur.from] + cur.cost < 0) {</pre>
       ++cur.flow, --G[cur.to][cur.rev].flow;
       for (int
             i = cur.from; past[i]; i = past[i]->from) {
         auto &e = *past[i];
++e.flow, --G[e.to][e.rev].flow;
      }
    ++cur.cap;
  }
  void solve(int mxlg) {
    for (int b = mxlg; b >= 0; --b) {
       for (int i = 0; i < n; ++i)</pre>
         for (auto &e : G[i])
       e.cap *= 2, e.flow *= 2;
for (int i = 0; i < n; ++i)
         for (auto &e : G[i])
           if (e.fcap >> b & 1)
             try_edge(e);
    }
  void init(int _n) { n = _n;
    for (int i = 0; i < n; ++i) G[i].clear();</pre>
  void add_edge(int a, int b, int cap, int cost) {
    G[a].eb(Edge
         {a, b, 0, cap, 0, cost, SZ(G[b]) + (a == b)});
    G[b].eb(Edge{b, a, 0, 0, 0, -cost, SZ(G[a]) - 1});
} mcmf; // O(VE * ElogC)
```

4.9 Flow Models

- Maximum/Minimum flow with lower bound / Circulation problem
 - 1. Construct super source ${\cal S}$ and sink ${\cal T}$.
 - 2. For each edge (x,y,l,u), connect $x \rightarrow y$ with capacity u-l.
 - 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
 - 4. If in(v)>0, connect $S\to v$ with capacity in(v), otherwise, connect $v\to T$ with capacity -in(v).
 - To maximize, connect $t \to s$ with capacity ∞ (skip this in circulation problem), and let f be the maximum flow from S to T. If $f \neq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, the maximum flow from s to t is the answer.
 - To minimize, let f be the maximum flow from S to T. Connect $t \to s$ with capacity ∞ and let the flow from S to T be f'. If $f+f' \neq \sum_{v \in V, in(v)>0} in(v)$, there's no solution. Otherwise, f' is the answer.
 - 5. The solution of each edge e is l_e+f_e , where f_e corresponds to the flow of edge e on the graph.
- Construct minimum vertex cover from maximum matching M on bipartite graph (X,Y)
 - 1. Redirect every edge: $y \rightarrow x$ if $(x,y) \in M$, $x \rightarrow y$ otherwise.
 - 2. DFS from unmatched vertices in \hat{X} .
 - 3. $x \in X$ is chosen iff x is unvisited.
 - 4. $y \in Y$ is chosen iff y is visited.
- Minimum cost cyclic flow
 - 1. Consruct super source ${\cal S}$ and sink ${\cal T}$
- 2. For each edge (x,y,c), connect $x \to y$ with (cost,cap)=(c,1) if c>0, otherwise connect $y \to x$ with (cost,cap)=(-c,1)
- 3. For each edge with c < 0, sum these cost as K, then increase d(y) by 1, decrease d(x) by 1
- 4. For each vertex v with d(v)>0, connect $S\to v$ with (cost, cap)=(0, d(v))

- 5. For each vertex v with d(v)<0 , connect $v\to T$ with (cost, cap) = (0, -d(v))
- 6. Flow from S to T, the answer is the cost of the flow C+K
- · Maximum density induced subgraph
 - 1. Binary search on answer, suppose we're checking answer ${\cal T}$
 - 2. Construct a max flow model, let K be the sum of all weights
 - 3. Connect source $s \rightarrow v$, $v \in G$ with capacity K
 - 4. For each edge (u,v,w) in G, connect $u \to v$ and $v \to u$ with capacity w
 - 5. For $v\in G$, connect it with sink $v\to t$ with capacity $K+2T-(\sum_{e\in E(v)}w(e))-2w(v)$
 - 6. T is a valid answer if the maximum flow f < K|V|
- · Minimum weight edge cover
 - 1. For each $v \in V$ create a copy v', and connect $u' \to v'$ with weight w(u,v).
 - 2. Connect $v \to v'$ with weight $2\mu(v)$, where $\mu(v)$ is the cost of the cheapest edge incident to v.
 - 3. Find the minimum weight perfect matching on G'.
- Project selection problem
 - 1. If $p_v>0$, create edge (s,v) with capacity p_v ; otherwise, create edge (v,t) with capacity $-p_v$.
 - 2. Create edge (u,v) with capacity w with w being the cost of choosing u without choosing v.
 - 3. The mincut is equivalent to the maximum profit of a subset of projects.
- Dual of minimum cost maximum flow
 - 1. Capacity c_{uv} , Flow f_{uv} , Cost w_{uv} , Required Flow difference for vertex b_{uv} .
 - 2. If all w_{uv} are integers, then optimal solution can happen when all p_u are integers.

$$\begin{aligned} \min \sum_{uv} w_{uv} f_{uv} \\ -f_{uv} \geq -c_{uv} &\Leftrightarrow \min \sum_{u} b_{u} p_{u} + \sum_{uv} c_{uv} \max(0, p_{v} - p_{u} - w_{uv}) \\ \sum_{v} f_{vu} - \sum_{v} f_{uv} = -b_{u} \end{aligned}$$

4.10 MCMF HLPP

// Push-Relabel

```
implementation of the cost-scaling algorithm
   Runs in O( <max_flow
> * log(V * max_edge_cost)) = O( V^3 * log(V * C))
// Operates on integers
#include < bits / stdc++.h>
using namespace std;
template < typename flow_t = int, typename cost_t = int>
struct mcSFlow{
  struct Edge{
    cost t c;
    flow t f;
    int to, rev;
    Edge(int _to, cost_t _c, flow_t _f
         , int _rev):c(_c), f(_f), to(_to), rev(_rev){}
  }:
  const
      cost t INFCOST = numeric limits < cost t >:: max()/2;
  const
      cost_t INFFLOW = numeric_limits<flow_t>::max()/2;
  cost_t epsilon;
  int N, S, T;
  vector<vector<Edge> > G:
  vector<unsigned int> isEnqueued, state;
  mcSFlow(int _N, int _S,
        int _T):epsilon(0), N(_N), S(_S), T(_T), G(_N){}
  void add_edge(int a, int b, cost_t cost, flow_t cap){
  if(a==b){assert(cost>=0); return;}
    cost*=N;// to preserve integer-values
    epsilon = max(epsilon, abs(cost));
    assert(a>=0&&a<N&&b>=0&&b<N);
    G[a].emplace_back(b, cost, cap, G[b].size());
G[b].emplace_back(a, -cost, 0, G[a].size()-1);
  flow_t calc_max_flow(){ // Dinic max-flow
    vector<flow_t> dist(N), state(N);
    vector < Edge* > path(N);
    auto cmp = [](Edge*a, Edge*b){return a->f < b->f;};
    flow_t addFlow, retflow=0;;
      fill(dist.begin(), dist.end(), -1);
      dist[S]=0;
      auto head = state.begin(), tail = state.begin();
       for(*tail++ = S;head!=tail;++head){
         for(Edge const&e:G[*head]){
           if(e.f && dist[e.to]==-1){
             dist[e.to] = dist[*head]+1;
             *tail++=e.to;
```

```
}
    addFlow = 0;
    fill(state.begin(), state.end(), 0);
    auto top = path.begin();
    Edge dummy(S, 0, INFFLOW, -1);
    *top++ = &dummy;
    while(top != path.begin()){
      int n = (*prev(top))->to;
      if(n==T){
        auto next_top
             = min_element(path.begin(), top, cmp);
        flow_t flow = (*next_top)->f;
        while(--top!=path.begin()){
          Edge &e=**top, &f=G[e.to][e.rev];
          e.f-=flow:
          f.f+=flow;
        addFlow=1;
        retflow+=flow;
        top = next_top;
        continue;
      for(int &i=state[n], i_max
           = G[n].size(), need = dist[n]+1;;++i){
        if(i==i_max){
          dist[n]=-1;
          --top;
          break;
        if(dist[G[n][i].to] == need && G[n][i].f){
          *top++ = &G[n][i];
          break:
        }
     }
  }while(addFlow);
  return retflow;
vector<flow_t> excess;
vector<cost_t> h;
void push(Edge &e, flow_t amt){
                                                           };
  //cerr << "push:</pre>
       << G[e.to][e.rev].to << " -> " << e.to << " ("
       << e.f << "/" << e.c << ") : " << amt << "\n";
  if(e.f < amt) amt=e.f;</pre>
  e.f-=amt;
  excess[e.to]+=amt;
  G[e.to][e.rev].f+=amt;
  excess[G[e.to][e.rev].to]-=amt;
void relabel(int vertex){
  cost_t newHeight = -INFCOST;
  for(unsigned int i=0;i<G[vertex].size();++i){</pre>
    Edge const&e = G[vertex][i];
    if(e.f && newHeight < h[e.to]-e.c){</pre>
      newHeight = h[e.to] - e.c;
      state[vertex] = i;
    }
  h[vertex] = newHeight - epsilon;
                                                           }
const int scale=2;
pair<flow_t, cost_t> minCostFlow(){
  cost_t retCost = 0;
  for(int i=0;i<N;++i){</pre>
    for(Edge &e:G[i]){
      retCost += e.c*(e.f);
    }
  }
  //find feasible flow
  flow_t retFlow = calc_max_flow();
  excess.resize(N);h.resize(N);
  queue < int > q;
  isEnqueued.assign(N, 0); state.assign(N,0);
                                                           }
  for(;epsilon;epsilon>>=scale){
    //refine
    fill(state.begin(), state.end(), 0);
for(int i=0;i<N;++i)</pre>
      for(auto &e:G[i])
        if(h[i]
            + e.c - h[e.to] < 0 && e.f) push(e, e.f);
    for(int i=0;i<N;++i){</pre>
      if(excess[i]>0){
        q.push(i);
```

```
isEngueued[i]=1:
     }
    while(!q.empty()){
      int cur=q.front();q.pop();
      isEnqueued[cur]=0;
      // discharge
      while(excess[cur]>0){
        if(state[cur] == G[cur].size()){
         relabel(cur);
       for(unsigned int &i=state[
            cur], max_i = G[cur].size();i<max_i;++i){</pre>
         Edge &e=G[cur][i];
         if(h[cur] + e.c - h[e.to] < 0){
           push(e, excess[cur]);
            if(excess
                [e.to]>0 && isEnqueued[e.to]==0){
             q.push(e.to);
             isEnqueued[e.to]=1;
           if(excess[cur]==0) break;
         }
       }
     }
    if(epsilon>1 && epsilon>>scale==0){
     epsilon = 1<<scale;</pre>
  for(int i=0;i<N;++i){</pre>
   for(Edge &e:G[i]){
     retCost -= e.c*(e.f);
  return make_pair(retFlow, retCost/2/N);
flow_t getFlow(Edge const &e){
  return G[e.to][e.rev].f;
}
  String
```

5

5.1 KMP

```
int F[maxn];
vector<int> match(string A, string B) {
 vector<int> ans;
 F[0] = -1, F[1] = 0;
 for (int i = 1, j = 0; i < SZ(B); F[++i] = ++j) {</pre>
    if (B[i] == B[j]) F[i] = F[j]; // optimize
    while (j != -1 && B[i] != B[j]) j = F[j];
 for (int i = 0, j = 0; i < SZ(A); ++i) {
    while (j != -1 && A[i] != B[j]) j = F[j];
    if (++j == SZ(B)) ans.eb(i + 1 - j), j = F[j];
 return ans;
```

5.2 Z-value*

```
int z[maxn];
void make_z(const string &s) {
  int l = 0, r = 0;
  for (int i = 1; i < SZ(s); ++i) {</pre>
    for (z[i] = max(0, min(r - i + 1, z[i - l]));
         i + z[i] < SZ(s) && s[i + z[i]] == s[z[i]];
         ++z[i])
    if (i + z[i] - 1 > r) l = i, r = i + z[i] - 1;
  }
```

5.3 Manacher*

```
int z[maxn]; // 0-base
/* center i: radius z[i * 2 + 1] / 2
center i, i + 1: radius z[i * 2 + 2] / 2
    both aba, abba have radius 2 */
void Manacher(string tmp) {
  string s = "%";
  int l = 0, r = 0;
```

```
namespace sfx {
bool _t[N * 2];
int SA[N * 2], H[N], RA[N];
int _s[N * 2], _c[N * 2], x[N],
                                    _p[N], _q[N * 2];
// zero based, string content MUST > 0
// SA[i]: SA[i]-th
    suffix is the i-th lexigraphically smallest suffix.
// H[i]: longest
    common prefix of suffix SA[i] and suffix SA[i-1].
void pre(int *sa, int *c, int n, int z)
{ fill_n(sa, n, 0), copy_n(c, z, x); }
void induce
    (int *sa, int *c, int *s, bool *t, int n, int z) {
  copy_n(c, z - 1, x + 1);
  for (int i = 0; i < n; ++i)</pre>
    if (sa[i] && !t[sa[i] - 1])
      sa[x[s[sa[i] - 1]]++] = sa[i] - 1;
  copy_n(c, z, x);

for (int i = n - 1; i >= 0; --i)
    if (sa[i] && t[sa[i] - 1])
      sa[--x[s[sa[i] - 1]]] = sa[i] - 1;
void sais(int *s, int *sa
      int *p, int *q, bool *t, int *c, int n, int z) {
  bool uniq = t[n - 1] = true;
  int nn = 0,
      nmxz = -1, *nsa = sa + n, *ns = s + n, last = -1;
  fill_n(c, z, 0);
  for (int i = 0; i < n; ++i) uniq &= ++c[s[i]] < 2;</pre>
  partial_sum(c, c + z, c);
  if (uniq) {
    for (int i = 0; i < n; ++i) sa[--c[s[i]]] = i;</pre>
    return:
  for (int i = n - 2; i >= 0; --i)
    t[i] = (
        s[i] == s[i + 1] ? t[i + 1] : s[i] < s[i + 1]);
  pre(sa, c, n, z);
  for (int i = 1; i <= n - 1; ++i)</pre>
    if (t[i] && !t[i - 1])
      sa[--x[s[i]]] = p[q[i] = nn++] = i;
  induce(sa, c, s, t, n, z);
for (int i = 0; i < n; ++i)</pre>
    if (sa[i] && t[sa[i]] && !t[sa[i] - 1]) {
      bool neq = last < 0 || !equal
           (s + sa[i], s + p[q[sa[i]] + 1], s + last);
      ns[q[last = sa[i]]] = nmxz += neq;
  sais(ns,
       nsa, p + nn, q + n, t + n, c + z, nn, nmxz + 1);
  pre(sa, c, n, z);
for (int i = nn - 1; i >= 0; --i)
    sa[--x[s[p[nsa[i]]]]] = p[nsa[i]];
  induce(sa, c, s, t, n, z);
void mkhei(int n) {
  for (int i = 0, j = 0; i < n; ++i) {
    if (RA[i])
    for (; _s[i + j] == _s[SA[RA[i] - 1] + j]; ++j);
H[RA[i]] = j, j = max(0, j - 1);
  }
void build(int *s, int n) {
 copy_n(s, n, _s), _s[n] = 0;
sais(_s, SA, _p, _q, _t, _c, n + 1, 256);
copy_n(SA + 1, n, SA);
  for (int i = 0; i < n; ++i) RA[SA[i]] = i;</pre>
  mkhei(n);
```

5.5 Aho-Corasick Automatan

```
const int len = 400000, sigma = 26;
struct AC_Automatan {
  int nx[len][sigma], fl[len], cnt[len], pri[len], top;
```

```
int newnode() {
  fill(nx[top], nx[top] + sigma, -1);
  return top++;
void init() { top = 1, newnode(); }
int input(
  string &s) { // return the end_node of string
  int X = 1;
  for (char c : s) {
  if (!~nx[X][c - 'a']) nx[X][c - 'a'] = newnode();
  X = nx[X][c - 'a'];
  return X;
void make_fl() {
  queue < int > q;
  q.push(1), fl[1] = 0;
  for (int t = 0; !q.empty();) {
    int R = q.front();
    q.pop(), pri[t++] = R;
for (int i = 0; i < sigma; ++i)</pre>
       if (~nx[R][i]) {
         int X = nx[R][i], Z = fl[R];
         for (; Z && !~nx[Z][i];) Z = fl[Z];
         fl[X] = Z ? nx[Z][i] : 1, q.push(X);
  }
void get_v(string &s) {
  int X = 1;
  fill(cnt, cnt + top, 0);
  for (char c : s) {
    while (X \&\& !~nx[X][c - 'a']) X = fl[X];
    X = X ? nx[X][c -
                          'a'] : 1, ++cnt[X];
  for (int i = top - 2; i > 0; --i)
  cnt[fl[pri[i]]] += cnt[pri[i]];
```

5.6 De Bruijn sequence*

```
constexpr int maxc = 10, maxn = 1e5 + 10;
struct DBSeq {
  int C, N, K, L, buf[maxc * maxn]; // K <= C^N</pre>
   void dfs(int *out, int t, int p, int &ptr) {
     if (ptr >= L) return;
     if (t > N) {
       if (N % p) return;
       for (int i = 1; i <= p && ptr < L; ++i)</pre>
         out[ptr++] = buf[i];
     } else {
       buf[t] = buf[t - p], dfs(out, t + 1, p, ptr);
       for (int j = buf[t - p] + 1; j < C; ++j)</pre>
         buf[t] = j, dfs(out, t + 1, t, ptr);
  void solve(int _c, int _n, int _k, int *out) {
    int p = 0;
    C = _c, N = _n, K = _k, L = N + K - 1;
dfs(out, 1, 1, p);
    if (p < L) fill(out + p, out + L, 0);</pre>
} dbs;
```

5.7 Extended SAM*

```
struct exSAM {
  int len[N * 2], link[N * 2]; // maxlength, suflink
  int next[N * 2][CNUM], tot; // [0, tot), root = 0
  int lenSorted[N * 2]; // topo. order
  int cnt[N * 2]; // occurence
  int newnode() {
    fill_n(next[tot], CNUM, 0);
    len[tot] = cnt[tot] = link[tot] = 0;
    return tot++;
  }
  void init() { tot = 0, newnode(), link[0] = -1; }
  int insertSAM(int last, int c) {
    int cur = next[last][c];
    len[cur] = len[last] + 1;
    int p = link[last];
  while (p!= -1 && !next[p][c])
    next[p][c] = cur, p = link[p];
  if (p == -1) return link[cur] = 0, cur;
  int q = next[p][c];
```

```
if (len
        [p] + 1 == len[q]) return link[cur] = q, cur;
    int clone = newnode();
    for (int i = 0; i < CNUM; ++i)</pre>
      next[
          clone][i] = len[next[q][i]] ? next[q][i] : 0;
    len[clone] = len[p] + 1;
    while (p != -1 && next[p][c] == q)
      next[p][c] = clone, p = link[p];
    link[link[cur] = clone] = link[q];
    link[q] = clone;
    return cur;
  void insert(const string &s) {
    int cur = 0;
    for (auto ch : s) {
      int &nxt = next[cur][int(ch - 'a')];
      if (!nxt) nxt = newnode();
      cnt[cur = nxt] += 1;
    }
  }
  void build() {
    queue < int > q;
    q.push(0);
    while (!q.empty()) {
      int cur = q.front();
      q.pop();
      for (int i = 0; i < CNUM; ++i)</pre>
        if (next[cur][i])
          q.push(insertSAM(cur, i));
    vector<int> lc(tot);
    for (int i = 1; i < tot; ++i) ++lc[len[i]];</pre>
    partial_sum(ALL(lc), lc.begin());
    for (int i
        = 1; i < tot; ++i) lenSorted[--lc[len[i]]] = i;
  void solve() {
    for (int i = tot - 2; i >= 0; --i)
      cnt[link[lenSorted[i]]] += cnt[lenSorted[i]];
};
```

5.8 PalTree*

```
struct palindromic_tree {
  struct node {
    int next[26], fail, len;
    int cnt, num; // cnt: appear times, num: number of
    // pal. suf.
node(int l = 0) : fail(0), len(l), cnt(0), num(0) {
      for (int i = 0; i < 26; ++i) next[i] = 0;</pre>
 }:
  vector<node> St:
 vector<char> s;
  int last, n;
 palindromic_tree() : St(2), last(1), n(0) {
   St[0].fail = 1, St[1].len = -1, s.eb(-1);
 inline void clear() {
   St.clear(), s.clear(), last = 1, n = 0;
    St.eb(0), St.eb(-1);
    St[0].fail = 1, s.eb(-1);
  inline int get_fail(int x) {
    while (s[n - St[x].len - 1] != s[n])
     x = St[x].fail;
    return x;
  inline void add(int c) {
    s.eb(c -= 'a'), ++n;
    int cur = get_fail(last);
    if (!St[cur].next[c]) {
      int now = SZ(St):
      St.eb(St[cur].len + 2);
      St[now].fail =
        St[get_fail(St[cur].fail)].next[c];
      St[cur].next[c] = now;
      St[now].num = St[St[now].fail].num + 1;
    last = St[cur].next[c], ++St[last].cnt;
  inline void count() { // counting cnt
    auto i = St.rbegin();
    for (; i != St.rend(); ++i) {
```

```
St[i->fail].cnt += i->cnt;
}
}
inline int size() { // The number of diff. pal.
    return SZ(St) - 2;
}
};
```

6 Math

6.1 ax+by=gcd(only exgcd *)

```
pii exgcd(int a, int b) {
   if (b == 0) return pii(1, 0);
   int p = a / b;
   pii q = exgcd(b, a % b);
   return pii(q.Y, q.X - q.Y * p);
}
/* ax+by=res, let x be minimum non-negative
g, p = gcd(a, b), exgcd(a, b) * res / g
   if p.X < 0: t = (abs(p.X) + b / g - 1) / (b / g)
   else: t = -(p.X / (b / g))
   p += (b / g, -a / g) * t */</pre>
```

6.2 Floor and Ceil

```
int floor(int a, int b)
{ return a / b - (a % b && (a < 0) ^ (b < 0)); }
int ceil(int a, int b)
{ return a / b + (a % b && (a < 0) ^ (b > 0)); }
```

6.3 Floor Enumeration

```
// enumerating x = floor(n / i), [l, r]
for (int l = 1, r; l <= n; l = r + 1) {
   int x = n / l;
   r = n / x;
}</pre>
```

6.4 Mod Min

```
// min{k | l <= ((ak) mod m) <= r}, no solution -> -1
int mod_min(int a, int m, int l, int r) {
  if (a == 0) return l ? -1 : 0;
  if (int k = (l + a - 1) / a; k * a <= r)
    return k;
  int b = m / a, c = m % a;
  if (int y = mod_min(c, a, a - r % a, a - l % a))
    return (l + y * c + a - 1) / a + y * b;
  return -1;
}</pre>
```

6.5 Gaussian integer qcd

6.6 floor sum*

```
int floor_sum(int n, int m, int a, int b) {
   int ans = 0;
   if (a >= m)
        ans += (n - 1) * n * (a / m) / 2, a %= m;
   if (b >= m)
        ans += n * (b / m), b %= m;
   int y_max
        = (a * n + b) / m, x_max = (y_max * m - b);
   if (y_max == 0) return ans;
   ans += (n - (x_max + a - 1) / a) * y_max;
   ans += floor_sum(y_max, a, m, (a - x_max % a) % a);
   return ans;
}// sum^{{
      n-1}_0 floor((a * i + b) / m) in log(n + m + a + b)
```

6.7 Miller Rabin*

6.8 Simultaneous Equations

```
struct matrix { //m variables, n equations
   fraction M[maxn][maxn + 1], sol[maxn];
   int solve() { //-1: inconsistent, >= 0: rank
     for (int i = 0; i < n; ++i) {</pre>
       int piv = 0;
       while (piv < m && !M[i][piv].n) ++piv;</pre>
       if (piv == m) continue;
for (int j = 0; j < n; ++j) {</pre>
         if (i == j) continue;
          fraction tmp = -M[j][piv] / M[i][piv];
         for (int k = 0; k <=</pre>
               m; ++k) M[j][k] = tmp * M[i][k] + M[j][k];
      }
     int rank = 0;
     for (int i = 0; i < n; ++i) {</pre>
       int piv = 0;
       while (piv < m && !M[i][piv].n) ++piv;</pre>
       if (piv == m && M[i][m].n) return -1;
       else if (piv
             < m) ++rank, sol[piv] = M[i][m] / M[i][piv];</pre>
     return rank;
};
```

6.9 Pollard Rho*

6.10 Simplex Algorithm

```
const int maxn = 11000, maxm = 405;
const double eps = 1E-10;
double a[maxn][maxm], b[maxn], c[maxm];
double d[maxn][maxm], x[maxm];
int ix[maxn + maxm]; // !!! array all indexed from 0
// max{cx} subject to {Ax<=b,x>=0}
// n: constraints, m: vars !!!
// x[] is the optimal solution vector
// usage :
// value = simplex(a, b, c, N, M);
double simplex(int n, int m){
    ++m;
    fill_n(d[n], m + 1, 0);
```

```
fill_n(d[n + 1], m + 1, 0);
iota(ix, ix + n + m, 0);
int r = n, s = m - 1;
for (int i = 0; i < n; ++i) {</pre>
   for (int j = 0; j < m - 1; ++j) d[i][j] = -a[i][j];</pre>
   d[i][m - 1] = 1;
   d[i][m] = b[i];
   if (d[r][m] > d[i][m]) r = i;
copy_n(c, m - 1, d[n]);
d[n + 1][m - 1] = -1;
for (double dd;; ) {
   if (r < n) {
      swap(ix[s], ix[r + m]);
     d[r][s] = 1.0 / d[r][s];
for (int j = 0; j <= m; ++j)
  if (j != s) d[r][j] *= -d[r][s];</pre>
     for (int i = 0; i <= n + 1; ++i) if (i != r) {
  for (int j = 0; j <= m; ++j) if (j != s)
    d[i][j] += d[r][j] * d[i][s];</pre>
         d[i][s] *= d[r][s];
     }
  }
  for (int j = 0; j < m; ++j)
  if (s < 0 || ix[s] > ix[j]) {
        if (d[n + 1][j] > eps ||
              (d[n + 1][j] > -eps && d[n][j] > eps))
  if (s < 0) break;
for (int i = 0; i < n; ++i) if (d[i][s] < -eps) {</pre>
     if (r < 0 ||
           (dd = d[r][m]
           / d[r][s] - d[i][m] / d[i][s]) < -eps ||
(dd < eps && ix[r + m] > ix[i + m]))
         r = i:
   if (r < 0) return -1; // not bounded
if (d[n + 1][m] < -eps) return -1; // not executable</pre>
double ans = 0:
fill_n(x, m, 0);
for (int i = m; i <</pre>
  n + m; ++i) { // the missing enumerated x[i] = 0
if (ix[i] < m - 1){
ans += d[i - m][m] * c[ix[i]];
     x[ix[i]] = d[i-m][m];
}
return ans;
```

6.10.1 Construction

Standard form: maximize $\mathbf{c}^T\mathbf{x}$ subject to $A\mathbf{x} \leq \mathbf{b}$ and $\mathbf{x} \geq 0$. Dual LP: minimize $\mathbf{b}^T\mathbf{y}$ subject to $A^T\mathbf{y} \geq \mathbf{c}$ and $\mathbf{y} \geq 0$. $\bar{\mathbf{x}}$ and $\bar{\mathbf{y}}$ are optimal if and only if for all $i \in [1,n]$, either $\bar{x}_i = 0$ or $\sum_{j=1}^m A_{ji}\bar{y}_j = c_i$ holds and for all $i \in [1,m]$ either $\bar{y}_i = 0$ or $\sum_{j=1}^n A_{ij}\bar{x}_j = b_j$ holds.

```
1. In case of minimization, let c_i' = -c_i
```

```
2. \sum_{1 \leq i \leq n} A_{ji} x_i \geq b_j \rightarrow \sum_{1 \leq i \leq n} -A_{ji} x_i \leq -b_j
```

- $3. \sum_{1 \le i \le n}^{-} A_{ji} x_i = b_j$
 - $\sum_{1 \leq i \leq n} A_{ji} x_i \leq b_j$
 - $\sum_{1 \le i \le n}^{-} A_{ji} x_i \ge b_j$
- 4. If x_i has no lower bound, replace x_i with $x_i x_i'$

6.11 chineseRemainder

```
int solve(int x1, int m1, int x2, int m2) {
  int g = gcd(m1, m2);
  if ((x2 - x1) % g) return -1; // no sol
  m1 /= g; m2 /= g;
  pii p = exgcd(m1, m2);
  int lcm = m1 * m2 * g;
  int res = p.first * (x2 - x1) * m1 + x1;
  // be careful with overflow
  return (res % lcm + lcm) % lcm;
}
```

6.12 Factorial without prime factor*

```
// O(p^k + log^2 n), pk = p^k
int prod[maxp];
int fac_no_p(int n, int p, int pk) {
  prod[0] = 1;
  for (int i = 1; i <= pk; ++i)
    if (i % p) prod[i] = prod[i - 1] * i % pk;
    else prod[i] = prod[i - 1];
int rt = 1;
for (; n; n /= p) {
    rt = rt * mpow(prod[pk], n / pk, pk) % pk;
    rt = rt * prod[n % pk] % pk;
}
return rt;
} // (n! without factor p) % p^k</pre>
```

6.13 QuadraticResidue*

```
int Jacobi(int a, int m) {
  int s = 1;
  for (; m > 1; ) {
    a %= m;
    if (a == 0) return 0;
    const int r = __builtin_ctz(a);
    if ((r \& 1) \&\& ((m + 2) \& 4)) s = -s;
    if (a & m & 2) s = -s;
    swap(a, m);
int QuadraticResidue(int a, int p) {
  if (p == 2) return a & 1;
  const int jc = Jacobi(a, p);
  if (jc == 0) return 0;
  if (jc == -1) return -1;
  int b, d;
  for (; ; ) {
   b = rand() % p;
d = (1LL * b * b + p - a) % p;
    if (Jacobi(d, p) == -1) break;
  int f0 = b, f1 = 1, g0 = 1, g1 = 0, tmp;
  for (int e = (1LL + p) >> 1; e; e >>= 1) {
    if (e & 1) {
      tmp = (1LL *
      g0 * f0 + 1LL * d * (1LL * g1 * f1 % p)) % p;
g1 = (1LL * g0 * f1 + 1LL * g1 * f0) % p;
    tmp = (1LL)
        * f0 * f0 + 1LL * d * (1LL * f1 * f1 % p)) % p;
    f1 = (2LL * f0 * f1) % p;
    f0 = tmp;
  return g0;
```

6.14 PiCount*

```
int PrimeCount(int n) { // n ~ 10^13 => < 2s</pre>
 if (n <= 1) return 0;
  int v = sqrt(n), s = (v + 1) / 2, pc = 0;
  vector < int > smalls(v + 1), skip(v + 1), roughs(s);
  vector<int> larges(s);
  for (int i = 2; i <= v; ++i) smalls[i] = (i + 1) / 2;</pre>
  for (int i = 0; i < s; ++i) {
  roughs[i] = 2 * i + 1;</pre>
    larges[i] = (n / (2 * i + 1) + 1) / 2;
  for (int p = 3; p <= v; ++p) {</pre>
    if (smalls[p] > smalls[p - 1]) {
      int q = p * p;
      ++pc;
      if (1LL * q * q > n) break;
      skip[p] = 1;
      for (int i = q; i <= v; i += 2 * p) skip[i] = 1;</pre>
      int ns = 0;
      for (int k = 0; k < s; ++k) {</pre>
        int i = roughs[k];
         if (skip[i]) continue;
         int d = 1LL * i * p;
        larges[ns] = larges[k] - (d <= v ? larges</pre>
             [smalls[d] - pc] : smalls[n / d]) + pc;
```

```
roughs[ns++] = i;
        }
         s = ns;
         for (int j = v / p; j >= p; --j) {
           int c =
           smalls[j] - pc, e = min(j * p + p, v + 1); \\ \mbox{for (int } i = j * p; i < e; ++i) smalls[i] -= c; \\ \mbox{}
        }
      }
   for (int k = 1; k < s; ++k) {
   const int m = n / roughs[k];</pre>
      int t = larges[k] - (pc + k - 1);
      for (int l = 1; l < k; ++l) {</pre>
         int p = roughs[l];
         if (1LL * p * p > m) break;
         t = smalls[m / p] - (pc + l - 1);
      larges[0] -= t;
   return larges[0];
}
```

6.15 Discrete Log*

```
int DiscreteLog(int s, int x, int y, int m) {
   constexpr int kStep = 32000;
   unordered_map < int , int > p;
   int b = 1:
   for (int i = 0; i < kStep; ++i) {</pre>
    p[y] = i;
     y = 1LL * y * x % m;
     b = 1LL * b * x % m;
  for (int i = 0; i < m + 10; i += kStep) {
    s = 1LL * s * b % m;</pre>
     if (p.find(s) != p.end()) return i + kStep - p[s];
   }
   return -1:
int DiscreteLog(int x, int y, int m) {
   if (m == 1) return 0;
   int s = 1;
   for (int i = 0; i < 100; ++i) {</pre>
     if (s == y) return i;
s = 1LL * s * x % m;
   if (s == y) return 100;
   int p = 100 + DiscreteLog(s, x, y, m);
   if (fpow(x, p, m) != y) return -1;
   return p;
```

6.16 Berlekamp Massey

```
template <typename T>
vector<T> BerlekampMassey(const vector<T> &output) {
  vector<T> d(SZ(output) + 1), me, he;
  for (int f = 0, i = 1; i <= SZ(output); ++i) {
    for (int j = 0; j < SZ(me); ++j)
        d[i] += output[i - j - 2] * me[j];
    if ((d[i] -= output[i - 1]) == 0) continue;
    if (me.empty()) {
        me.resize(f = i);
        continue;
    }
    vector<T> o(i - f - 1);
    T k = -d[i] / d[f]; o.eb(-k);
    for (T x : he) o.eb(x * k);
        o.resize(max(SZ(o), SZ(me)));
    for (int j = 0; j < SZ(me); ++j) o[j] += me[j];
    if (i - f + SZ(he) >= SZ(me)) he = me, f = i;
    me = o;
    }
    return me;
}
```

6.17 Primes

6.18 Theorem

· Cramer's rule

$$ax+by=e \Rightarrow x = \frac{ed-bf}{ad-bc}$$

$$cx+dy=f \Rightarrow y = \frac{af-ec}{ad-bc}$$

Vandermonde's Identity

$$C(n+m,k) = \sum_{i=0}^{k} C(n,i)C(m,k-i)$$

· Kirchhoff's Theorem

Denote L be a $n \times n$ matrix as the Laplacian matrix of graph G, where $L_{ii} = d(i)$, $L_{ij} = -c$ where c is the number of edge (i,j) in G.

- The number of undirected spanning in G is $|\det(\tilde{L}_{11})|$.
- The number of directed spanning tree rooted at r in G is $|\det(\tilde{L}_{rr})|$.
- · Tutte's Matrix

Let D be a n imes n matrix, where $d_{ij} = x_{ij}$ (x_{ij} is chosen uniformly at random) if i < j and $(i,j) \in E$, otherwise $d_{ij} = -d_{ji}$. $\frac{rank(D)}{2}$ is the maximum matching on G.

- Cayley's Formula
 - Given a degree sequence $d_1, d_2, ..., d_n$ for each labeled vertices, there are $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$ spanning trees.
 - Let $T_{n,k}$ be the number of labeled forests on n vertices with k components, such that vertex $1,2,\ldots,k$ belong to different components. Then $T_{n,k}\!=\!kn^{n-k-1}$.
- Erdős–Gallai theorem

A sequence of nonnegative integers $d_1 \ge \cdots \ge d_n$ can be represented as the degree sequence of a finite simple graph on n vertices if and only if

$$d_1+\dots+d_n \text{ is even and } \sum_{i=1}^k d_i \leq k(k-1)+\sum_{i=k+1}^n \min(d_i,k) \text{ holds for every } 1\leq k\leq n.$$

Gale-Ryser theorem

A pair of sequences of nonnegative integers $a_1 \ge \cdots \ge a_n$ and b_1, \dots, b_n is bigraphic if and only if $\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$ and $\sum_{i=1}^k a_i \leq \sum_{i=1}^n \min(b_i,k)$ holds for

Fulkerson-Chen-Anstee theorem

A sequence $(a_1,\ b_1),\ ...\ ,\ (a_n,\ b_n)$ of nonnegative integer pairs with $a_1 \geq \cdots \geq a_n$ is digraphic if and only if $\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$ and

$$\sum_{i=1}^k a_i \leq \sum_{i=1}^k \min(b_i,k-1) + \sum_{i=k+1}^n \min(b_i,k) \text{ holds for every } 1 \leq k \leq n.$$
 • Möbius inversion formula

- - $f(n) = \sum_{d|n} g(d) \Leftrightarrow g(n) = \sum_{d|n} \mu(d) f(\frac{n}{d})$
 - $f(n) = \sum_{n|d} g(d) \Leftrightarrow g(n) = \sum_{n|d} \mu(\frac{d}{n}) f(d)$
- Spherical cap
 - A portion of a sphere cut off by a plane.
 - r: sphere radius, a: radius of the base of the cap, h: height of the cap,
 - Volume = $\pi h^2 (3r h)/3 = \pi h (3a^2 + h^2)/6 = \pi r^3 (2 + \cos \theta)(1 \theta)$ $\cos\theta)^2/3$.
 - $\label{eq:Area} \text{Area} \stackrel{\cdot}{=} 2\pi r h = \pi (a^2 + h^2) = 2\pi r^2 (1 \cos \theta).$
- · Lagrange multiplier
 - Optimize $f(x_1,...,x_n)$ when k constraints $g_i(x_1,...,x_n) = 0$.
 - Lagrangian function $\mathcal{L}(x_1,\ldots,x_n,\lambda_1,\ldots,\lambda_k)=f(x_1,\ldots,x_n)=$ $\sum_{i=1}^k \lambda_i g_i(x_1,...,x_n).$
 - The solution corresponding to the original constrained optimization is always a saddle point of the Lagrangian function.

6.19 Estimation

- Estimation
 - The number of divisors of n is at most around 100 for n < 5e4, 500 for n < 1e7, 2000 for n < 1e10, 200000 for n < 1e19.
 - The number of ways of writing n as a sum of positive integers, disregarding the order of the summands. 1,1,2,3,5,7,11,15,22,30 for $n = 0 \sim 9$, 627 for n = 20, $\sim 2e5$ for n = 50, $\sim 2e8$ for n = 100.
 - Total number of partitions of n distinct elements: B(n)1,1,2,5,15,52,203,877,4140,21147,115975,678570,4213597,27644437,190899322,....

6.20 Euclidean Algorithms

- $m = \lfloor \frac{an+b}{a} \rfloor$
- Time complexity: $O(\log n)$

$$\begin{split} f(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)}{2} + \lfloor \frac{b}{c} \rfloor \cdot (n+1) \\ + f(a \operatorname{mod} c, b \operatorname{mod} c, c, n), & a \geq c \vee b \geq c \\ 0, & n < 0 \vee a = 0 \\ nm - f(c, c - b - 1, a, m - 1), & \text{otherwise} \end{cases} \end{split}$$

$$\begin{split} g(a,b,c,n) &= \sum_{i=0}^{n} i \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor \cdot \frac{n(n+1)}{2} \\ +g(a \bmod c, b \bmod c, c, n), & a \geq c \vee b \geq c \\ 0, & n < 0 \vee a = 0 \\ \frac{1}{2} \cdot (n(n+1)m - f(c, c - b - 1, a, m - 1) \\ -h(c, c - b - 1, a, m - 1)), & \text{otherwise} \end{cases} \end{split}$$

$$\begin{split} h(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor^2 \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor^2 \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor^2 \cdot (n+1) \\ + \lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1) \\ + h(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{a}{c} \rfloor \cdot g(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{b}{c} \rfloor \cdot f(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ nm(m+1) - 2g(c, c - b - 1, a, m - 1) \\ - 2f(c, c - b - 1, a, m - 1) - f(a, b, c, n), & \text{otherwise} \end{cases} \end{split}$$

General Purpose Numbers

$$\begin{split} & \text{Bernoulli numbers} \\ & B_0 - 1, B_1^{\pm} = \pm \frac{1}{2}, B_2 = \frac{1}{6}, B_3 = 0 \\ & \sum_{j=0}^m \binom{m+1}{j} B_j = 0, \text{EGF is } B(x) = \frac{x}{e^x - 1} = \sum_{n=0}^\infty B_n \frac{x^n}{n!}. \\ & S_m(n) = \sum_{k=1}^n k^m = \frac{1}{m+1} \sum_{k=0}^m \binom{m+1}{k} B_k^+ n^{m+1-k}. \end{split}$$

ullet Stirling numbers of the second kind Partitions of n distinct elements into exactly k groups.

$$S(n,k) = S(n-1,k-1) + kS(n-1,k), S(n,1) = S(n,n) = 1$$

$$S(n,k) = \frac{1}{k!} \sum_{i=0}^{k} (-1)^{k-i} {k \choose i} i^n$$

$$x^n = \sum_{i=0}^{n} S(n,i)(x)_i$$
 • Pentagonal number theorem

$$\prod_{n=1}^{\infty} (1-x^n) = 1 + \sum_{k=1}^{\infty} (-1)^k \left(x^{k(3k+1)/2} + x^{k(3k-1)/2} \right)$$
 • Catalan numbers
$$C_n^{(k)} = \frac{1}{(k-1)n+1} \binom{kn}{n}$$

$$C_n^{(k)} = \frac{1}{(k-1)n+1} {kn \choose n}$$
$$C^{(k)}(x) = 1 + x[C^{(k)}(x)]^k$$

Number of permutations $\pi \in S_n$ in which exactly k elements are greater than the previous element. k j:s s.t. $\pi(j) > \pi(j+1)$, k+1 j:s s.t. $\pi(j) \ge j$, k j:s s.t. $\pi(j) > j$.

$$E(n,k) = (n-k)E(n-1,k-1) + (k+1)E(n-1,k)$$

$$E(n,0) = E(n,n-1) = 1$$

$$E(n,k) = \sum_{j=0}^{k} (-1)^{j} {n+1 \choose j} (k+1-j)^{n}$$

Tips for Generating Functions

- Ordinary Generating Function $A(x) = \sum_{i>0} a_i x^i$
 - $A(rx) \Rightarrow r^n a_n$
 - $A(x) + B(x) \Rightarrow a_n + b_n$
 - $A(x)B(x) \Rightarrow \sum_{i=0}^{n} a_i b_{n-i}$
 - $A(x)^k \Rightarrow \sum_{i_1+i_2+\dots+i_k=n} a_{i_1} a_{i_2} \dots a_{i_k}$
 - $xA(x)' \Rightarrow na_n$
 - $-\frac{A(x)}{1-x} \Rightarrow \sum_{i=0}^{n} a_i$
- Exponential Generating Function $A(x) = \sum_{i>0} \frac{a_i}{i!} x_i$

 - $-A(x)+B(x)\Rightarrow a_n+b_n$ $-A^{(k)}(x)\Rightarrow a_{n+k}$ $-A(x)B(x)\Rightarrow \sum_{i=0}^{n}\binom{n}{i}a_ib_{n-i}$ $-A(x)^k\Rightarrow \sum_{i_1+i_2+\dots+i_k=n}^{n}\binom{n}{i_1,i_2,\dots,i_k}a_{i_1}a_{i_2}\dots a_{i_k}$
- Special Generating Function $(1+x)^n = \sum_{i \ge 0} \binom{n}{i} x^i$

 - $-\frac{1}{(1-x)^n} = \sum_{i>0} {i \choose n-1} x^i$

7 Polynomial

7.1 Fast Fourier Transform

```
template < int maxn >
struct FFT {
    using val_t = complex < double >;
    const double PI = acos(-1);
    val_t w[maxn];
    FFT() {
        for (int i = 0; i < maxn; ++i) {
            double arg = 2 * PI * i / maxn;
            w[i] = val_t(cos(arg), sin(arg));
        }
    void bitrev(val_t *a, int n); // see NTT
    void trans
        (val_t *a, int n, bool inv = false); // see NTT;
    // remember to replace LL with val_t
};</pre>
```

7.2 Number Theory Transform*

```
//(2^16)+1, 65537, 3
//7*17*(2^23)+1, 998244353, 3
//1255*(2^20)+1, 1315962881,
//51*(2^25)+1, 1711276033, 29
template < int maxn, int P, int RT > //maxn must be 2^k
struct NTT {
  int w[maxn];
  int mpow(int a, int n);
  int minv(int a) { return mpow(a, P - 2); }
  NTT() {
    int dw = mpow(RT, (P - 1) / maxn);
    w[0] = 1;
    for (int
         i = 1; i < maxn; ++i) w[i] = w[i - 1] * dw % P;
  void bitrev(int *a, int n) {
    int i = 0;
    for (int j = 1; j < n - 1; ++j) {
  for (int k = n >> 1; (i ^= k) < k; k >>= 1);
      if (j < i) swap(a[i], a[j]);</pre>
    }
  }
  void operator()(int
        *a, int n, bool inv = false) { //0 <= a[i] < P
    bitrev(a, n);
    for (int L = 2; L <= n; L <<= 1) {</pre>
      int dx = maxn / L, dl = L >> 1;
      for (int i = 0; i < n; i += L) {</pre>
         for (int
              j = i, x = 0; j < i + dl; ++j, x += dx) {
           int tmp = a[j + dl] * w[x] % P;
           if ((a[j
                 + dl] = a[j] - tmp) < 0) a[j + dl] += P;
           if ((a[j] += tmp) >= P) a[j] -= P;
        }
      }
    if (inv) {
      reverse(a + 1, a + n);
      int invn = minv(n);
      for (int
            i = 0; i < n; ++i) a[i] = a[i] * invn % P;
  }
};
```

7.3 Fast Walsh Transform*

```
/* x: a[j], y: a[j + (L >> 1)]
or: (y += x * op), and: (x += y * op)
xor: (x, y = (x + y) * op, (x - y) * op)
invop: or, and, xor = -1, -1, 1/2 */
void fwt(int *a, int n, int op) { //or
for (int L = 2; L <= n; L <<= 1)
    for (int i = 0; i < n; i += L)
    for (int j = i; j < i + (L >> 1); ++j)
        a[j + (L >> 1)] += a[j] * op;
}
const int N = 21;
int f[
    N][1 << N], g[N][1 << N], h[N][1 << N], ct[1 << N];
void
    subset_convolution(int *a, int *b, int *c, int L) {</pre>
```

```
// c_k = \sum_{{i | j = k, i & j = 0}} a_i * b_j
int n = 1 << L;
for (int i = 1; i < n; ++i)
    ct[i] = ct[i & (i - 1)] + 1;
for (int i = 0; i < n; ++i)
    f[ct[i]][i] = a[i], g[ct[i]][i] = b[i];
for (int i = 0; i <= L; ++i)
    fwt(f[i], n, 1), fwt(g[i], n, 1);
for (int i = 0; i <= L; ++i)
    for (int j = 0; j <= i; ++j)
    for (int x = 0; x < n; ++x)
        h[i][x] += f[j][x] * g[i - j][x];
for (int i = 0; i <= L; ++i)
    fwt(h[i], n, -1);
for (int i = 0; i < n; ++i)
    c[i] = h[ct[i]][i];
}</pre>
```

7.4 Polynomial Operation

```
fi(s, n) for (int i = (int)(s); i < (int)(n); ++i) template<int maxn, int P, int RT> // maxn = 2^k
struct Poly : vector<int> { // coefficients in [0, P)
  using vector<int>::vector;
  static NTT<maxn, P, RT> ntt;
  int n() const { return (int)size(); } // n() >= 1
  Poly(const Poly &p, int m) : vector<int>(m) {
    copy_n(p.data(), min(p.n(), m), data());
  Poly& irev()
       { return reverse(data(), data() + n()), *this; }
  Poly& isz(int m) { return resize(m), *this; }
  Poly& iadd(const Poly &rhs) { // n() == rhs.n()
    fi(0, n()) if
         (((*this)[i] += rhs[i]) >= P) (*this)[i] -= P;
    return *this;
  Poly& imul(int k) {
    fi(0, n()) (*this)[i] = (*this)[i] * k % P;
    return *this;
  Poly Mul(const Poly &rhs) const {
    int m = 1;
    while (m < n() + rhs.n() - 1) m <<= 1;</pre>
    Poly X(*this, m), Y(rhs, m);
    ntt(X.data(), m), ntt(Y.data(), m);
    fi(0, m) X[i] = X[i] * Y[i] % P;
    ntt(X.data(), m, true);
    return X.isz(n() + rhs.n() - 1):
  Poly Inv() const { // (*this)[0] != 0, 1e5/95ms
    if (n() == 1) return {ntt.minv((*this)[0])};
    int m = 1;
    while (m < n() * 2) m <<= 1;</pre>
    Poly Xi = Poly(*this, (n() + 1) / 2).Inv().isz(m);
    Poly Y(*this, m);
    ntt(Xi.data(), m), ntt(Y.data(), m);
    fi(0, m) {
    Xi[i] *= (2 - Xi[i] * Y[i]) % P;
      if ((Xi[i] %= P) < 0) Xi[i] += P;</pre>
    ntt(Xi.data(), m, true);
    return Xi.isz(n());
  Poly Sqrt()
       const { // Jacobi((*this)[0], P) = 1, 1e5/235ms
    if (n()
        == 1) return {QuadraticResidue((*this)[0], P)};
    Poly
        X = Poly(*this, (n() + 1) / 2).Sqrt().isz(n());
         X.iadd(Mul(X.Inv()).isz(n())).imul(P / 2 + 1);
  pair < Poly , Poly > DivMod
      (const Poly &rhs) const { // (rhs.)back() != 0
    if (n() < rhs.n()) return \{\{\theta\}, *this\};
    const int m = n() - rhs.n() + 1;
    Poly X(rhs); X.irev().isz(m);
    Poly Y(*this); Y.irev().isz(m);
    Poly Q = Y.Mul(X.Inv()).isz(m).irev();
    X = rhs.Mul(Q), Y = *this
    fi(0, n()) if ((Y[i] -= X[i]) < 0) Y[i] += P;
    return {Q, Y.isz(max(1, rhs.n() - 1))};
 Poly Dx() const {
```

Poly ret(n() - 1);

```
fi(0,
      ret.n()) ret[i] = (i + 1) * (*this)[i + 1] % P;
  return ret.isz(max(1, ret.n()));
Poly Sx() const {
  Poly ret(n() + 1);
  fi(0, n())
       ret[i + 1] = ntt.minv(i + 1) * (*this)[i] % P;
  return ret;
     _tmul(int nn, const Poly &rhs) const {
Polv
  Poly Y = Mul(rhs).isz(n() + nn - 1);
  return Poly(Y.data() + n() - 1, Y.data() + Y.n());
vector<int> _eval(const
     vector<int> &x, const vector<Poly> &up) const {
  const int m = (int)x.size();
  if (!m) return {};
  vector<Poly> down(m * 2);
  // down[1] = DivMod(up[1]).second;
  // fi(2, m *
       2) down[i] = down[i / 2].DivMod(up[i]).second;
  down[1] = Poly(up[1])
      .irev().isz(n()).Inv().irev()._tmul(m, *this);
  fi(2, m * 2) down[i]
       = up[i ^ 1]._tmul(up[i].n() - 1, down[i / 2]);
  vector<int> y(m);
  fi(0, m) y[i] = down[m + i][0];
  return v:
}
static vector<Poly>
                     _tree1(const vector<int> &x) {
  const int m = (int)x.size();
  vector<Poly> up(m * 2);
  fi(0, m) up[m + i] = \{(x[i] ? P - x[i] : 0), 1\};
  for (int i = m - 1; i
      > 0; --i) up[i] = up[i * 2].Mul(up[i * 2 + 1]);
vector<int
    > Eval(const vector<int> &x) const { // 1e5, 1s
  auto up = _tree1(x); return _eval(x, up);
static Poly Interpolate(const vector
    <int> &x, const vector<int> &y) { // 1e5, 1.4s
  const int m = (int)x.size();
  vector<Poly> up = _tree1(x), down(m * 2);
  vector < int > z = up[1].Dx()._eval(x, up);
  fi(0, m) z[i] = y[i] * ntt.minv(z[i]) % P;
  fi(0, m) down[m + i] = {z[i]};
  for (int i = m -
       1; i > 0; --i) down[i] = down[i * 2].Mul(up[i
      * 2 + 1]).iadd(down[i * 2 + 1].Mul(up[i * 2]));
  return down[1];
Poly Ln() const { // (*this)[0] == 1, 1e5/170ms
  return Dx().Mul(Inv()).Sx().isz(n());
Poly Exp() const { // (*this)[0] == 0, 1e5/360ms
 if (n() == 1) return {1};
  Poly X = Poly(*this, (n() + 1) / 2).Exp().isz(n());
  Poly Y = X.Ln(); Y[0] = P - 1;
  fi(0, n())
       if ((Y[i] = (*this)[i] - Y[i]) < 0) Y[i] += P;</pre>
  return X.Mul(Y).isz(n());
}
^{-}// M := P(P - 1). If k >= M, k := k % M + M.
Poly Pow(int k) const {
  int nz = 0;
  while (nz < n() && !(*this)[nz]) ++nz;</pre>
  if (nz * min(k, (int)n()) >= n()) return Poly(n());
  if (!k) return Poly(Poly {1}, n());
  Poly X(data() + nz, data() + nz + n() - nz * k);
  const int c = ntt.mpow(X[0], k % (P - 1));
  return X.Ln().imul
      (k % P).Exp().imul(c).irev().isz(n()).irev();
static int LinearRecursion
    (const vector <int> &a, const vector <int> &coef, int n) { // a_n = |sum c_j| a_(n-j)
  const int k = (int)a.size();
  assert((int)coef.size() == k + 1);
  Poly C(k + 1), W(Poly \{1\}, k), M = \{0, 1\};
  fi(1, k + 1) C[k - i] = coef[i] ? P - coef[i] : 0;
  C[k] = 1:
  while (n) {
```

```
if (n % 2) W = W.Mul(M).DivMod(C).second;
    n /= 2, M = M.Mul(M).DivMod(C).second;
}
int ret = 0;
fi(0, k) ret = (ret + W[i] * a[i]) % P;
return ret;
}
};
#undef fi
using Poly_t = Poly<131072 * 2, 998244353, 3>;
template<> decltype(Poly_t::ntt) Poly_t::ntt = {};
```

7.5 Value Polynomial

```
struct Poly {
   mint base; // f(x) = poly[x - base]
   vector<mint> poly;
   Poly(mint b = 0, mint x = 0): base(b), poly(1, x) {}
   mint get_val(const mint &x) {
     if (x >= base && x < base + SZ(poly))</pre>
       return poly[x - base];
     mint rt = 0;
     vector<mint> lmul(SZ(poly), 1), rmul(SZ(poly), 1);
     for (int i = 1; i < SZ(poly); ++i)
  lmul[i] = lmul[i - 1] * (x - (base + i - 1));</pre>
     for (int i = SZ(poly) - 2; i >= 0; --i)
       rmul[i] = rmul[i + 1] * (x - (base + i + 1));
     for (int i = 0; i < SZ(poly); ++i)</pre>
       rt += poly[i] * ifac[i] * inegfac
           [SZ(poly) - 1 - i] * lmul[i] * rmul[i];
     return rt;
   void raise() { // g(x) = sigma\{base:x\} f(x)
     if (SZ(poly) == 1 && poly[0] == 0)
       return:
     mint nw = get_val(base + SZ(poly));
     poly.eb(nw);
     for (int i = 1; i < SZ(poly); ++i)</pre>
       poly[i] += poly[i - 1];
  }
};
```

7.6 Newton's Method

Given F(x) where

$$F(x) = \sum_{i=0}^{\infty} \alpha_i (x - \beta)^i$$

for β being some constant. Polynomial P such that F(P)=0 can be found iteratively. Denote by Q_k the polynomial such that $F(Q_k)=0$ (mod x^{2^k}), then

$$Q_{k+1}\!=\!Q_k\!-\!\frac{F(Q_k)}{F'(Q_k)}\pmod{x^{2^{k+1}}}$$

8 Geometry 8.1 Default Code

typedef pair < double, double > pdd; typedef pair < pdd, pdd > Line; struct Cir{ pdd 0; double R; }; const double eps = 1e-8; pdd operator + (pdd a, pdd b) { return pdd(a.X + b.X, a.Y + b.Y); pdd operator - (pdd a, pdd b)

```
{ return pdd(a.X + b.X, a.Y + b.Y); }
{ return pdd(a.X - b.X, a.Y - b.Y); }
pdd operator*(pdd a, double b)
{ return pdd(a.X * b, a.Y * b); }
pdd operator/(pdd a, double b)
{ return pdd(a.X / b, a.Y / b); }
double dot(pdd a, pdd b)
{ return a.X * b.X + a.Y * b.Y; }
double cross(pdd a, pdd b)
{ return a.X * b.Y - a.Y * b.X; }
double abs2(pdd a)
{ return dot(a, a); }
double abs(pdd a)
{ return sqrt(dot(a, a)); }
int sign(double a)
{ return fabs(a) < eps ? 0 : a > 0 ? 1 : -1; }
int ori(pdd a, pdd b, pdd c)
{ return sign(cross(b - a, c
                             - a)): }
bool collinearity(pdd p1, pdd p2, pdd p3)
{ return sign(cross(p1 - p3, p2 - p3)) == 0; }
bool btw(pdd p1, pdd p2, pdd p3) {
```

```
if (!collinearity(p1, p2, p3)) return 0;
  return sign(dot(p1 - p3, p2 - p3)) <= 0;</pre>
bool seg_intersect(pdd p1, pdd p2, pdd p3, pdd p4) {
  int a123 = ori(p1, p2, p3);
  int a124 = ori(p1, p2, p4);
  int a341 = ori(p3, p4, p1);
  int a342 = ori(p3, p4, p2);
  if (a123 == 0 && a124 == 0)
    return btw(p1, p2, p3) || btw(p1, p2, p4) ||
  btw(p3, p4, p1) || btw(p3, p4, p2);
return a123 * a124 <= 0 && a341 * a342 <= 0;
pdd intersect(pdd p1, pdd p2, pdd p3, pdd p4) {
  double a123 = cross(p2 - p1, p3 - p1);
  double a124 = cross(p2 - p1, p4 - p1);
  return (p4
       * a123 - p3 * a124) / (a123 - a124); // C^3 / C^2
pdd perp(pdd p1)
{ return pdd(-p1.Y, p1.X); }
pdd projection(pdd p1, pdd p2, pdd p3)
{ return p1 + (
    p2 - p1) * dot(p3 - p1, p2 - p1) / abs2(p2 - p1); }
pdd reflection(pdd p1, pdd p2, pdd p3)
{ return p3 + perp(p2 - p1
        ) * cross(p3 - p1, p2 - p1) / abs2(p2 - p1) * 2; }
pdd linearTransformation
    (pdd p0, pdd p1, pdd q0, pdd q1, pdd r) {
  pdd dp = p1 - p0
       , dq = q1 - q0, num(cross(dp, dq), dot(dp, dq));
  return q0 + pdd(
       cross(r - p0, num), dot(r - p0, num)) / abs2(dp);
\} // from line p0--p1 to q0--q1, apply to r
```

8.2 PointSegDist*

8.3 Heart

```
pdd circenter
    (pdd p0, pdd p1, pdd p2) { // radius = abs(center)
  p1 = p1 - p0, p2 = p2 - p0;
  double x1 = p1.X, y1 = p1.Y, x2 = p2.X, y2 = p2.Y;
  double m = 2. * (x1 * y2 - y1 * x2);
 center.X = (x1 * x1)
       * y2 - x2 * x2 * y1 + y1 * y2 * (y1 - y2)) / m;
  center.Y = (x1 * x2)
       * (x2 - x1) - y1 * y1 * x2 + x1 * y2 * y2) / m;
  return center + p0;
pdd incenter
    (pdd p1, pdd p2, pdd p3) { // radius = area / s * 2
  double a =
      abs(p2 - p3), b = abs(p1 - p3), c = abs(p1 - p2);
  double s = a + b + c;
  return (a * p1 + b * p2 + c * p3) / s;
pdd masscenter(pdd p1, pdd p2, pdd p3)
{ return (p1 + p2 + p3) / 3; }
pdd orthcenter(pdd p1, pdd p2, pdd p3)
{ return masscenter
    (p1, p2, p3) * 3 - circenter(p1, p2, p3) * 2; }
```

8.4 point in circle

```
__int128 det = (__int128)-u13 * u22 * u31
	+ (__int128)u12 * u23 * u31 + (__int128)u13 *
		u21 * u32 - (__int128)u11 * u23 * u32 - (__int128
	)u12 * u21 * u33 + (__int128)u11 * u22 * u33;
	return det > eps;
}
```

8.5 Convex hull*

8.6 PointInConvex*

```
bool PointInConvex
    (const vector < pii> &C, pii p, bool strict = true) {
    int a = 1, b = SZ(C) - 1, r = !strict;
    if (SZ(C) == 0) return false;
    if (SZ(C) < 3) return r && btw(C[0], C.back(), p);
    if (ori(C[0], C[a], C[b]) > 0) swap(a, b);
    if (ori
        (C[0], C[a], p) >= r || ori(C[0], C[b], p) <= -r)
        return false;
    while (abs(a - b) > 1) {
        int c = (a + b) / 2;
        (ori(C[0], C[c], p) > 0 ? b : a) = c;
    }
    return ori(C[a], C[b], p) < r;
}</pre>
```

8.7 TangentPointToHull*

```
/* The point should be strictly out of hull
  return arbitrary point on the tangent line */
pii get_tangent(vector<pii> &C, pii p) {
  auto gao = [&](int s) {
    return cyc_tsearch(SZ(C), [&](int x, int y)
      { return ori(p, C[x], C[y]) == s; });
  };
  return pii(gao(1), gao(-1));
} // return (a, b), ori(p, C[a], C[b]) >= 0
```

8.8 Intersection of line and convex

```
int TangentDir(vector<pii> &C, pii dir) {
  return cyc_tsearch(SZ(C), [&](int a, int b) {
    return cross(dir, C[a]) > cross(dir, C[b]);
#define cmpL(i) sign(cross(C[i] - a, b - a))
pii lineHull(pii a, pii b, vector<pii> &C) {
  int A = TangentDir(C, a - b);
   int B = TangentDir(C, b - a);
   int n = SZ(C);
  if (cmpL(A) < \theta \mid \mid cmpL(B) > \theta)
    return pii(-1, -1); // no collision
   auto gao = [&](int l, int r) {
    for (int t = l; (l + 1) % n != r; ) {
       int m = ((l + r + (l < r? 0 : n)) / 2) % n;
       (cmpL(m) = cmpL(t) ? l : r) = m;
    }
     return (l + !cmpL(r)) % n;
  pii res = pii(gao(B, A), gao(A, B)); // (i, j)
  if (res.X == res.Y) // touching the corner i
     return pii(res.X, -1);
   if (!
       cmpL(res.X) && !cmpL(res.Y)) // along side i, i+1
     switch ((res.X - res.Y + n + 1) % n) {
      case 0: return pii(res.X, res.X);
       case 2: return pii(res.Y, res.Y);
  /* crossing sides (i, i+1) and (j, j+1)
  crossing corner i is treated as side (i, i+1)
  returned
        in the same order as the line hits the convex */
  return res;
```

8.9 minMaxEnclosingRectangle*

```
const double INF = 1e18, qi = acos(-1) / 2 * 3;
pdd solve(vector<pii> &dots) {
#define diff(u, v) (dots[u] - dots[v])
#define vec(v) (dots[v] - dots[i])
 hull(dots);
  double Max = 0, Min = INF, deg;
  int n = SZ(dots);
  dots.eb(dots[0]);
  for (int i = 0, u = 1, r = 1, l = 1; i < n; ++i) {</pre>
   pii nw = vec(i + 1);
    while (cross(nw, vec(u + 1)) > cross(nw, vec(u)))
      u = (u + 1) \% n;
    while (dot(nw, vec(r + 1)) > dot(nw, vec(r)))
   r = (r + 1) \% n;
if (!i) l = (r + 1) \% n;
    while (dot(nw, vec(l + 1)) < dot(nw, vec(l)))
     l = (l + 1) \% n;
   Min = min(Min, (double)(dot(nw, vec(r)) - dot
        (nw, vec(l))) * cross(nw, vec(u)) / abs2(nw));
    deg = acos(dot(diff(r
   , l), vec(u)) / abs(diff(r, l)) / abs(vec(u)));
deg = (qi - deg) / 2;
   }
  return pdd(Min, Max);
```

8.10 VectorInPoly*

8.11 PolyUnion*

```
double rat(pii a, pii b) {
  return sign
      (b.X) ? (double)a.X / b.X : (double)a.Y / b.Y;
     all poly. should be ccw
double polyUnion(vector<vector<pii>>> &poly) {
  double res = 0;
  for (auto &p : poly)
    for (int a = 0; a < SZ(p); ++a) {</pre>
      pii A = p[a], B = p[(a + 1) \% SZ(p)];
      vector
      <pair < double, int >> segs = {{0, 0}, {1, 0}};
for (auto &q : poly) {
        if (&p == &q) continue;
        for (int b = 0; b < SZ(q); ++b) {
           pii C = q[b], D = q[(b + 1) \% SZ(q)];
          int sc = ori(A, B, C), sd = ori(A, B, D);
if (sc != sd && min(sc, sd) < 0) {</pre>
             double sa = cross(D
                   - C, A - C), sb = cross(D - C, B - C);
             segs.emplace_back
                 (sa / (sa - sb), sign(sc - sd));
           if (!sc && !sd &&
               &q < &p && sign(dot(B - A, D - C)) > \theta) {
             segs.emplace_back(rat(C - A, B - A), 1);
             segs.emplace_back(rat(D - A, B - A), -1);
          }
        }
      sort(ALL(segs));
      for (auto &s : segs) s.X = clamp(s.X, 0.0, 1.0);
      double sum = 0;
      int cnt = segs[0].second;
      for (int j = 1; j < SZ(segs); ++j) {</pre>
        if (!cnt) sum += segs[j].X - segs[j - 1].X;
        cnt += segs[j].Y;
```

```
res += cross(A, B) * sum;
}
return res / 2;
}
```

8.12 PolyCut

8.13 Polar Angle Sort*

```
int cmp(pii a, pii b, bool same = true) {
#define is_neg(k) (
    sign(k.Y) < 0 || (sign(k.Y) == 0 && sign(k.X) < 0))
int A = is_neg(a), B = is_neg(b);
if (A != B)
    return A < B;
if (sign(cross(a, b)) == 0)
    return same ? abs2(a) < abs2(b) : -1;
return sign(cross(a, b)) > 0;
```

8.14 Half plane intersection*

```
pii area_pair(Line a, Line b)
{ return pii(cross(a.Y
       - a.X, b.X - a.X), cross(a.Y - a.X, b.Y - a.X)); }
bool isin(Line l0, Line l1, Line l2) {
  // Check inter(l1, l2) strictly in l0
  auto [a02X, a02Y] = area_pair(l0, l2);
auto [a12X, a12Y] = area_pair(l1, l2);
  if (a12X - a12Y < \theta) a12X *= -1, a12Y *= -1;
  return (__int128)
a02Y * a12X - (__int128) a02X * a12Y > 0; // C^4
/* Having solution, check size > 2 */
/* --^-- Line.X --^-- Line.Y --^-- */
vector<Line> halfPlaneInter(vector<Line> arr) {
  sort(ALL(arr), [&](Line a, Line b) -> int {
  if (cmp(a.Y - a.X, b.Y - b.X, 0) != -1)
    return cmp(a.Y - a.X, b.Y - b.X, 0);
     return ori(a.X, a.Y, b.Y) < 0;</pre>
  });
  deque<Line> dq(1, arr[0]);
  for (auto p : arr) {
     if (cmp(
          dq.back().Y - dq.back().X, p.Y - p.X, 0) == -1)
       continue:
     while (SZ(dq
         ) >= 2 && !isin(p, dq[SZ(dq) - 2], dq.back()))
       dq.pb();
     while (SZ(dq) \ge 2 \&\& !isin(p, dq[0], dq[1]))
       dq.pop_front();
     dq.eb(p);
  while (SZ(dq)
        >= 3 && !isin(dq[0], dq[SZ(dq) - 2], dq.back()))
     dq.pb();
  while (SZ(dq) >= 3 \&\& !isin(dq.back(), dq[0], dq[1]))
     da.pop front():
  return vector < Line > (ALL(dq));
```

8.15 RotatingSweepLine

```
void rotatingSweepLine(vector<pii> &ps) {
   int n = SZ(ps), m = 0;
   vector<int> id(n), pos(n);
   vector<pii> line(n * (n - 1));
   for (int i = 0; i < n; ++i)
        for (int j = 0; j < n; ++j)
        if (i != j) line[m++] = pii(i, j);
   sort(ALL(line), [&](pii a, pii b) {
        return cmp(ps[a.Y] - ps[a.X], ps[b.Y] - ps[b.X]);
   }
}</pre>
```

8.16 Minimum Enclosing Circle*

```
pdd Minimum_Enclosing_Circle
    (vector<pdd> dots, double &r) {
  pdd cent:
  random_shuffle(ALL(dots));
  cent = dots[0], r = 0;
for (int i = 1; i < SZ(dots); ++i)</pre>
    if (abs(dots[i] - cent) > r) {
       cent = dots[i], r = 0;
for (int j = 0; j < i; ++j)
  if (abs(dots[j] - cent) > r) {
            cent = (dots[i] + dots[j]) / 2;
            r = abs(dots[i] - cent);
            for(int k = 0; k < j; ++k)
              if(abs(dots[k] - cent) > r)
                 cent = excenter
                      (dots[i], dots[j], dots[k], r);
         }
    }
  return cent;
```

8.17 Intersection of two circles*

8.18 Intersection of polygon and circle*

```
// Divides into multiple triangle, and sum up
const double PI=acos(-1);
double _area(pdd pa, pdd eb, double r){
  if(abs(pa)<abs(eb)) swap(pa, eb);</pre>
  if(abs(eb)<eps) return 0;
double S, h, theta;</pre>
  double a=abs(eb),b=abs(pa),c=abs(eb-pa);
  double cosB = dot(eb,eb-pa) / a / c, B = acos(cosB);
  double cosC = dot(pa,eb) / a / b, C = acos(cosC);
  if(a > r){
    S = (C/2)*r*r;
    h = a*b*sin(C)/c;
    if (h < r && B
        < PI/2) S -= (acos(h/r)*r*r - h*sqrt(r*r-h*h));
  else if(b > r){
    theta = PI - B - asin(sin(B)/r*a);
    S = .5*a*r*sin(theta) + (C-theta)/2*r*r;
  else S = .5*sin(C)*a*b;
  return S:
double area_poly_circle(const
     vector<pdd> poly,const pdd &0,const double r){
  double S=0;
  for(int i=0;i<SZ(poly);++i)</pre>
    S += _area(poly[i]-0,poly[(i+1)%SZ(poly
        )]-0,r)*ori(0,poly[i],poly[(i+1)%SZ(poly)]);
  return fabs(S):
```

8.19 Intersection of line and circle*

8.20 Tangent line of two circles

```
vector < Line
   > go( const Cir& c1 , const Cir& c2 , int sign1 ){
// sign1 = 1 for outer tang, -1 for inter tang
   vector<Line> ret;
   double d_sq = abs2(c1.0 - c2.0);
   if (sign(d_sq) == 0) return ret;
   double d = sqrt(d_sq);
   pdd v = (c2.0 - c1.0) / d;
   double c = (c1.R - sign1 * c2.R) / d;
   if (c * c > 1) return ret;
   double h = sqrt(max(0.0, 1.0 - c * c));
   for (int sign2 = 1; sign2 >= -1; sign2 -= 2) {
  pdd n = pdd(v.X * c - sign2 * h * v.Y,
       v.Y * c + sign2 * h * v.X);
     pdd p1 = c1.0 + n * c1.R;
     pdd p2 = c2.0 + n * (c2.R * sign1);
     if (sign(p1.X - p2.X) == 0 and
sign(p1.Y - p2.Y) == 0)
       p2 = p1 + perp(c2.0 - c1.0);
     ret.eb(Line(p1, p2));
   return ret;
}
```

8.21 CircleCover*

```
const int N = 1021;
struct CircleCover {
  int C;
  Cir c[N];
  bool g[N][N], overlap[N][N];
  // Area[i] : area covered by at least i circles
  double Area[ N ];
  void init(int _C){ C = _C;}
  struct Teve {
     pdd p; double ang; int add;
     Teve() {}
     Teve(pdd
    , double _b, int _c):p(_a), ang(_b), add(_c){}
bool operator<(const Teve &a)const</pre>
     {return ang < a.ang;}
  }eve[N * 2];
  // strict: x = 0, otherwise x = -1
  bool disjuct(Cir &a, Cir &b, int x)
{return sign(abs(a.0 - b.0) - a.R - b.R) > x;}
bool contain(Cir &a, Cir &b, int x)
  {return sign(a.R - b.R - abs(a.0 - b.0)) > x;}
  bool contain(int i, int j) {
    /* c[j] is non-strictly in c[i]. */
     return (sign
         (c[i].R - c[j].R) > 0 \mid | (sign(c[i].R - c[j].
         R) == 0 \&\& i < j)) \&\& contain(c[i], c[j], -1);
  void solve(){
    fill_n(Area, C + 2, 0);
for(int i = 0; i < C; ++i)
       for(int j = 0; j < C; ++j)</pre>
         overlap[i][j] = contain(i, j);
     for(int i = 0; i < C; ++i)</pre>
       for(int j = 0; j < C; ++j)
         g[i][j] = !(overlap[i][j] || overlap[j][i] ||
              disjuct(c[i], c[j], -1));
     for(int i = 0; i < C; ++i){</pre>
       int E = 0, cnt = 1;
       for(int j = 0; j < C; ++j)
         if(j != i && overlap[j][i])
            ++cnt;
       for(int j = 0; j < C; ++j)</pre>
         if(i != j && g[i][j]) {
            pdd aa, bb;
            CCinter(c[i], c[j], aa, bb);
```

```
double A =
                atan2(aa.Y - c[i].0.Y, aa.X - c[i].0.X);
           double B =
                atan2(bb.Y - c[i].0.Y, bb.X - c[i].0.X);
           eve[E++] = Teve
               (bb, B, 1), eve[E++] = Teve(aa, A, -1);
           if(B > A) ++cnt;
      if(E == 0) Area[cnt] += pi * c[i].R * c[i].R;
      else{
        sort(eve, eve + E);
         eve[E] = eve[0];
         for(int j = 0; j < E; ++j){</pre>
           cnt += eve[j].add;
           Area[cnt
               ] += cross(eve[j].p, eve[j + 1].p) * .5;
           double theta = eve[j + 1].ang - eve[j].ang;
           if (theta < 0) theta += \frac{1}{2}. * pi;
           Area[cnt] += (theta
                - sin(theta)) * c[i].R * c[i].R * .5;
      }
    }
  }
};
```

8.22 3Dpoint*

```
struct Point {
  double x, y, z;
  Point(double _x = 0, double _y = 0, double _z = 0): x(_x), y(_y), z(_z){}
  Point(pdd p) { x = p.X, y = p.Y, z = abs2(p); }
Point operator - (Point p1, Point p2)
     Point(p1.x - p2.x, p1.y - p2.y, p1.z - p2.z); }
Point operator+(Point p1, Point p2)
{ return
     Point(p1.x + p2.x, p1.y + p2.y, p1.z + p2.z); }
Point operator*(Point p1, double v)
{ return Point(p1.x * v, p1.y * v, p1.z * v); }
Point operator/(Point p1, double v)
{ return Point(p1.x / v, p1.y / v, p1.z / v); }
Point cross(Point p1, Point p2) { return Point(p1.y * p2.z - p1.z * p2.y, p1.z
    * p2.x - p1.x * p2.z, p1.x * p2.y - p1.y * p2.x); }
double dot(Point p1, Point p2)
{ return p1.x * p2.x + p1.y * p2.y + p1.z * p2.z; }
double abs(Point a)
{ return sqrt(dot(a, a)); }
Point cross3(Point a, Point b, Point c)
{ return cross(b - a, c - a); }
double area(Point a, Point b, Point c)
{ return abs(cross3(a, b, c)); }
double volume(Point a, Point b, Point c, Point d)
{ return dot(cross3(a, b, c), d - a); }
//Azimuthal
     angle (longitude) to x-axis in interval [-pi, pi]
double phi(Point p) { return atan2(p.y, p.x); }
     angle (latitude) to the z-axis in interval [0, pi]
double theta(Point p)
    { return atan2(sqrt(p.x * p.x + p.y * p.y), p.z); }
Point masscenter(Point a, Point b, Point c, Point d)
{ return (a + b + c + d) / 4; }
pdd proj(Point a, Point b, Point c, Point u) {
// proj. u to the plane of a, b, and c
  Point e1 = b - a;
  Point e2 = c - a;
  e1 = e1 / abs(e1);
  e2 = e2 - e1 * dot(e2, e1);
  e2 = e2 / abs(e2);
  Point p = u - a;
  return pdd(dot(p, e1), dot(p, e2));
Point
     rotate_around(Point p, double angle, Point axis) {
  double s = sin(angle), c = cos(angle);
  Point u = axis / abs(axis);
        * dot(u, p) * (1 - c) + p * c + cross(u, p) * s;
```

```
struct convex hull 3D {
struct Face {
  int a, b, c;
  Face(int ta, int tb, int tc): a(ta), b(tb), c(tc) {}
}; // return the faces with pt indexes
vector < Face > res;
vector<Point> P:
convex_hull_3D(const vector<Point> &_P): res(), P(_P) {
// all points coplanar case will WA, O(n^2)
  int n = SZ(P);
  if (n <= 2) return; // be careful about edge case</pre>
  // ensure first 4 points are not coplanar
  swap(P[1], *find_if(ALL(P), [&](
      auto p) { return sign(abs2(P[0] - p)) != 0; }));
  swap(P[2], *find_if(ALL(P), [&](auto p) { return
       sign(abs2(cross3(p, P[0], P[1]))) != 0; }));
  swap(P[3], *find_if(ALL(P), [&](auto p) { return
       sign(volume(P[0], P[1], P[2], p)) != 0; }));
  vector<vector<int>> flag(n, vector<int>(n));
  res.emplace_back(0, 1, 2); res.emplace_back(2, 1, 0);
  for (int i = 3; i < n; ++i) {</pre>
    vector < Face > next;
    for (auto f : res) {
      int d
           = sign(volume(P[f.a], P[f.b], P[f.c], P[i]));
      if (d <= 0) next.eb(f);</pre>
      int ff = (d > 0) - (d < 0);
      flag[f.a][
          f.b] = flag[f.b][f.c] = flag[f.c][f.a] = ff;
    for (auto f : res) {
      auto F = [&](int x, int y) {
        if (flag[x][y] > 0 && flag[y][x] <= 0)
          next.emplace_back(x, y, i);
      F(f.a, f.b); F(f.b, f.c); F(f.c, f.a);
    }
    res = next;
  }
bool same(Face s, Face t) {
  if (sign(volume
      (P[s.a], P[s.b], P[s.c], P[t.a])) != 0) return 0;
  if (sign(volume
      (P[s.a], P[s.b], P[s.c], P[t.b])) != 0) return 0;
  if (sign(volume
      (P[s.a], P[s.b], P[s.c], P[t.c])) != 0) return 0;
int polygon_face_num() {
  int ans = 0;
  for (int i = 0; i < SZ(res); ++i)</pre>
    ans += none_of(res.begin(), res.begin()
         + i, [&](Face g) { return same(res[i], g); });
  return ans;
double get_volume() {
  double ans = 0;
  for (auto f : res)
    ans +=
        volume(Point(0, 0, 0), P[f.a], P[f.b], P[f.c]);
  return fabs(ans / 6);
double get_dis(Point p, Face f) {
  Point p1 = P[f.a], p2 = P[f.b], p3 = P[f.c];
  double a = (p2.y - p1.y)
       * (p3.z - p1.z) - (p2.z - p1.z) * (p3.y - p1.y);
  double b = (p2.z - p1.z)
       * (p3.x - p1.x) - (p2.x - p1.x) * (p3.z - p1.z);
  double c = (p2.x - p1.x)
       * (p3.y - p1.y) - (p2.y - p1.y) * (p3.x - p1.x);
  double d = 0 - (a * p1.x + b * p1.y + c * p1.z);
  return fabs(a * p.x + b *
      p.y + c * p.z + d) / sqrt(a * a + b * b + c * c);
// n^2 delaunay: facets with negative z normal of
// convexhull of (x, y, x^2 + y^2), use a pseudo-point
// (0, 0, inf) to avoid degenerate case
```

8.24 DelaunayTriangulation*

```
/* Delaunay Triangulation:
Given a sets of points on 2D plane, find a
triangulation such that no points will strictly
inside circumcircle of any triangle.
```

```
A Triangle is in triangulation iff. its has_chd is 0.
Region of triangle u: iterate each u.edge[i].tri,
each points are u.p[(i+1)\%3], u.p[(i+2)\%3]
Voronoi diagram: for each triangle in triangulation,
the bisector of all its edges will split the region.
nearest point will belong to the triangle containing it
    int inf = maxc * maxc * 100; // lower_bound unknown
struct Tri:
struct Edge {
  Tri* tri; int side;
  Edge(): tri(0), side(0){}
  Edge(Tri* _tri, int _side): tri(_tri), side(_side){}
struct Tri {
  pii p[3];
  Edge edge[3];
  Tri() {}
  Tri(const pii& p0, const pii& p1, const pii& p2) {
    p[0] = p0; p[1] = p1; p[2] = p2;
chd[0] = chd[1] = chd[2] = 0;
  bool has_chd() const { return chd[0] != 0; }
  int num_chd() const {
    return !!chd[0] + !!chd[1] + !!chd[2];
  bool contains(pii const& q) const {
  for (int i = 0; i < 3; ++i)</pre>
      if (ori(p[i], p[(i + 1) % 3], q) < 0)</pre>
        return 0;
    return 1;
  }
} pool[N * 10], *tris;
void edge(Edge a, Edge b) {
  if(a.tri) a.tri->edge[a.side] = b;
  if(b.tri) b.tri->edge[b.side] = a;
struct Trig { // Triangulation
  Trig() {
    the_root
         = // Tri should at least contain all points
      new(tris++) Tri(pii(-inf, -inf),
            pii(inf + inf, -inf), pii(-inf, inf + inf));
  Tri* find(pii p) { return find(the_root, p); }
  void add_point(const
       pii &p) { add_point(find(the_root, p), p); }
  Tri* the_root;
  static Tri* find(Tri* root, const pii &p) {
    while (1) {
      if (!root->has_chd())
        return root;
      for (int i = 0; i < 3 && root->chd[i]; ++i)
        if (root->chd[i]->contains(p)) {
          root = root->chd[i];
          break;
        }
    assert(0); // "point not found"
  void add_point(Tri* root, pii const& p) {
    Tri* t[3];
     /* split it into three triangles */
    for (int i = 0; i < 3; ++i)</pre>
      t[i] = new(tris
           ++) Tri(root->p[i], root->p[(i + 1) % 3], p);
    for (int i = 0; i < 3; ++i)</pre>
      edge(Edge(t[i], 0), Edge(t[(i + 1) % 3], 1));
    for (int i = 0; i < 3; ++i)</pre>
    edge(Edge(t[i], 2), root->edge[(i + 2) % 3]);
for (int i = 0; i < 3; ++i)</pre>
      root->chd[i] = t[i];
    for (int i = 0; i < 3; ++i)</pre>
      flip(t[i], 2);
  void flip(Tri* tri, int pi) {
    Tri* trj = tri->edge[pi].tri;
int pj = tri->edge[pi].side;
    if (!trj) return;
    if (!in_cc(tri->p
        [0], tri->p[1], tri->p[2], trj->p[pj])) return;
    /* flip edge between tri,trj */
```

```
Tri* trk = new(tris++) Tri
         (tri->p[(pi + 1) % 3], trj->p[pj], tri->p[pi]);
    Tri* trl = new(tris++) Tri
         (trj->p[(pj + 1) % 3], tri->p[pi], trj->p[pj]);
     edge(Edge(trk, 0), Edge(trl, 0));
    edge(Edge(trk, 1), tri->edge[(pi + 2) % 3]);
    edge(Edge(trk, 2), trj->edge[(pj + 1) % 3]);
     edge(Edge(trl, 1), trj->edge[(pj + 2) % 3]);
     edge(Edge(trl, 2), tri->edge[(pi + 1) % 3]);
     tri->chd
        [0] = trk; tri->chd[1] = trl; tri->chd[2] = 0;
     trj->chd
        [0] = trk; trj->chd[1] = trl; trj->chd[2] = 0;
     flip(trk, 1); flip(trk, 2);
     flip(trl, 1); flip(trl, 2);
  }
};
vector<Tri*> triang; // vector of all triangle
set<Tri*> vst;
void go(Tri* now) { // store all tri into triang
  if (vst.find(now) != vst.end())
    return;
  vst.insert(now);
  if (!now->has_chd())
    return triang.eb(now);
  for (int i = 0; i < now->num_chd(); ++i)
    go(now->chd[i]);
void build(int n, pii* ps) { // build triangulation
  tris = pool; triang.clear(); vst.clear();
  random_shuffle(ps, ps + n);
  Trig tri; // the triangulation structure
  for (int i = 0; i < n; ++i)</pre>
    tri.add_point(ps[i]);
  qo(tri.the root);
}
```

8.25 Triangulation Vonoroi*

```
vector<Line> ls[N];
pii arr[N];
Line make_line(pdd p, Line l) {
  pdd d = l.Y - l.X; d = perp(d);
  pdd m = (l.X + l.Y) / 2;
  l = Line(m, m + d);
  if (ori(l.X, l.Y, p) < 0)</pre>
    l = Line(m + d, m);
  return l;
double calc_area(int id) {
  // use to calculate
  the area of point "strictly in the convex hull" vector<Line> hpi = halfPlaneInter(ls[id]);
  vector<pdd> ps;
  for (int i = 0; i < SZ(hpi); ++i)</pre>
    ps.eb(intersect(hpi[i].X, hpi[i].Y, hpi[(i
          + 1) % SZ(hpi)].X, hpi[(i + 1) % SZ(hpi)].Y));
  double rt = 0;
  for (int i = 0; i < SZ(ps); ++i)</pre>
    rt += cross(ps[i], ps[(i + 1) % SZ(ps)]);
  return fabs(rt) / 2;
void solve(int n, pii *oarr) {
  map<pii, int> mp;
  for (int i = 0; i < n; ++i)</pre>
    arr[i] = pii(oarr[i].X, oarr[i].Y), mp[arr[i]] = i;
  build(n, arr); // Triangulation
for (auto *t : triang) {
    vector<int> p;
     for (int i = 0; i < 3; ++i)</pre>
       if (mp.find(t->p[i]) != mp.end())
         p.eb(mp[t->p[i]]);
    for (int i = 0; i < SZ(p); ++i)
  for (int j = i + 1; j < SZ(p); ++j) {</pre>
         Line l(oarr[p[i]], oarr[p[j]]);
         ls[p[i]].eb(make_line(oarr[p[i]], l));
         ls[p[j]].eb(make_line(oarr[p[j]], l));
       }
  }
}
```

8.26 Minkowski Sum*

```
vector<pii> Minkowski(vector<pii> A, vector<pii> B) {
  hull(A), hull(B);
  vector<pii> C(1, A[0] + B[0]), s1, s2;
```

```
for (int i = 0; i < SZ(A); ++i)
  s1.eb(A[(i + 1) % SZ(A)] - A[i]);</pre>
for (int i = 0; i < SZ(B); i++)</pre>
  s2.eb(B[(i + 1) % SZ(B)] - B[i]);
for (int i = 0, j = 0; i < SZ(A) || j < SZ(B);)</pre>
  if (j >= SZ
       (B) || (i < SZ(A) && cross(s1[i], s2[j]) >= 0))
    C.eb(B[j % SZ(B)] + A[i++]);
    C.eb(A[i % SZ(A)] + B[j++]);
return hull(C), C;
```

9 Else

9.1 Cyclic Ternary Search*

```
* bool pred(int a, int b);
f(0) \sim f(n-1) is a cyclic-shift U-function return idx s.t. pred(x, idx) is false forall x*/
int cyc_tsearch(int n, auto pred) {
  if (n == 1) return 0;
  int l = 0, r = n; bool rv = pred(1, 0);
while (r - l > 1) {
  int m = (l + r) / 2;
     if (pred(0, m) ? rv: pred(m, (m + 1) % n)) r = m;
     else l = m:
  return pred(l, r % n) ? l : r % n;
```

9.2 Mo's Alogrithm(With modification)

```
Mo's Algorithm With modification
Block: N^{2/3}, Complexity: N^{5/3}
struct Query {
   int L, R, LBid, RBid, T;
  Query(int l, int r, int t):
   L(l), R(r), LBid(l / blk), RBid(r / blk), T(t) {}
   bool operator < (const Query &q) const {</pre>
     if (LBid != q.LBid) return LBid < q.LBid;</pre>
     if (RBid != q.RBid) return RBid < q.RBid;</pre>
     return T < b.T;</pre>
  }
void solve(vector<Query> query) {
  sort(ALL(query));
  int L=0, R=0, T=-1;
for (auto q : query) {
     while (T < q.T) addTime(L, R, ++T); // TODO
while (T > q.T) subTime(L, R, T--); // TODO
     while (R < q.R) add(arr[++R]); // TODO
while (L > q.L) add(arr[--L]); // TODO
     while (R > q.R) sub(arr[R--]); // TODO
     while (L < q.L) sub(arr[L++]); // TODO</pre>
      // answer query
}
```

Mo's Alogrithm On Tree

```
Mo's Algorithm On Tree
Preprocess:
1) LCA
2) dfs with in[u] = dft++, out[u] = dft++
3) ord[in[u]] = ord[out[u]] = u
4) bitset<maxn> inset
struct Query {
  int L, R, LBid, lca;
  Query(int u, int v) {
    int c = LCA(u, v);
    if (c == u || c == v)
      q.lca = -1, q.L = out[c ^ u ^ v], q.R = out[c];
    else if (out[u] < in[v])</pre>
      q.lca = c, q.L = out[u], q.R = in[v];
    else
      q.lca = c, q.L = out[v], q.R = in[u];
    q.Lid = q.L / blk;
  bool operator<(const Query &q) const {</pre>
    if (LBid != q.LBid) return LBid < q.LBid;</pre>
    return R < q.R;</pre>
```

```
}
};
void flip(int x) {
     if (inset[x]) sub(arr[x]); // TODO
     else add(arr[x]); // TODO
     inset[x] = ~inset[x];
void solve(vector<Query> query) {
   sort(ALL(query));
   int L = 0, R = 0;
   for (auto q : query) {
     while (R < q.R) flip(ord[++R]);</pre>
     while (L > q.L) flip(ord[--L]);
     while (R > q.R) flip(ord[R--]);
     while (L < q.L) flip(ord[L++]);</pre>
     if (~q.lca) add(arr[q.lca]);
     // answer query
     if (~q.lca) sub(arr[q.lca]);
}
```

9.4 Additional Mo's Algorithm Trick

- · Mo's Algorithm With Addition Only
 - Sort querys same as the normal Mo's algorithm.
 - For each query [l,r]:
 - If l/blk = r/blk, brute-force.
 - If $l/blk \neq curL/blk$, initialize $curL := (l/blk+1) \cdot blk$, curR := curL-1
 - If r > cur R, increase cur R
 - decrease curL to fit \emph{l} , and then undo after answering
- · Mo's Algorithm With Offline Second Time
 - Require: Changing answer \equiv adding f([l,r],r+1).
 - Require: f([l,r],r+1) = f([1,r],r+1) f([1,l),r+1). Part1: Answer all f([1,r],r+1) first.

 - Part2: Store $curR \rightarrow R$ for curL (reduce the space to O(N)), and then answer them by the second offline algorithm.
 - Note: You must do the above symmetrically for the left boundaries.

9.5 Hilbert Curve

```
int hilbert(int n, int x, int y) {
  int res = 0;
  for (int s = n / 2; s; s >>= 1) {
    int rx = (x \& s) > 0;
    int ry = (y & s) > 0;
res += s * 1ll * s * ((3 * rx) ^ ry);
    if (ry == 0) {
      if (rx == 1) x = s - 1 - x, y = s - 1 - y;
       swap(x, y);
    }
  }
  return res;
 // n = 2^k
```

9.6 DynamicConvexTrick*

```
// only works for integer coordinates!! maintain max
struct Line {
  mutable int a, b, p;
  bool operator
       <(const Line &rhs) const { return a < rhs.a; }
  bool operator<(int x) const { return p < x; }</pre>
struct DynamicHull : multiset<Line, less<>>> {
  static const int kInf = 1e18;
  int Div(int a,
       int b) { return a / b - ((a ^ b) < 0 && a % b); }</pre>
  bool isect(iterator x, iterator y) {
     if (y == end()) { x->p = kInf; return 0; }
     if (x
         ->a == y->a) x->p = x->b > y->b ? kInf : -kInf;
     else x -> p = Div(y -> b - x -> b, x -> a - y -> a);
     return x->p >= y->p;
  void addline(int a, int b) {
     auto z = insert({a, b, \theta}), y = z++, x = y;
     while (isect(y, z)) z = erase(z);
     if (x != begin
         () && isect(--x, y)) isect(x, y = erase(y));
     while ((y = x) != begin
         () && (--x)->p >= y->p) isect(x, erase(y));
   int query(int x) {
     auto l = *lower_bound(x);
     return l.a * x + l.b;
};
```

9.7 All LCS*

```
void all_lcs(string s, string t) { // 0-base
  vector < int > h(SZ(t));
  iota(ALL(h), 0);
  for (int a = 0; a < SZ(s); ++a) {
    int v = -1;
    for (int c = 0; c < SZ(t); ++c)
        if (s[a] == t[c] || h[c] < v)
            swap(h[c], v);
        // LCS(s[0, a], t[b, c]) =
        // c - b + 1 - sum([h[i] >= b] | i <= c)
        // h[i] might become -1 !!
    }
}</pre>
```

9.8 DLX*

```
#define TRAV(i, link, start)
     for (int i = link[start]; i != start; i = link[i])
template < bool A, bool B = !A> // A: Exact
struct DLX {
              rg[NN], up[NN], dn[NN
  int lt[NN],
      ], cl[NN], rw[NN], bt[NN], s[NN], head, sz, ans;
  int columns;
  bool vis[NN];
  void remove(int c) {
    if (A) lt[rg[c]] = lt[c], rg[lt[c]] = rg[c];
    TRAV(i, dn, c) {
      if (A) {
        TRAV(j, rg, i)
          up[dn[j]]
                = up[j], dn[up[j]] = dn[j], --s[cl[j]];
        lt[rg[i]] = lt[i], rg[lt[i]] = rg[i];
   }
  void restore(int c) {
    TRAV(i, up, c) {
      if (A) {
        TRAV(j, lt, i)
          ++s[cl[j]], up[dn[j]] = j, dn[up[j]] = j;
        lt[rg[i]] = rg[lt[i]] = i;
    if (A) lt[rg[c]] = c, rg[lt[c]] = c;
  void init(int c) {
    columns = c;
    for (int i = 0; i < c; ++i) {
  up[i] = dn[i] = bt[i] = i;</pre>
      lt[i] = i == 0 ? c : i - 1;
      rg[i] = i == c - 1 ? c : i + 1;
      s[i] = 0;
    }
    rg[c] = 0, lt[c] = c - 1;
    up[c] = dn[c] = -1;
    head = c, sz = c + 1;
  void insert(int r, const vector<int> &col) {
    if (col.empty()) return;
    int f = sz;
    for (int i = 0; i < (int)col.size(); ++i) {</pre>
      int c = col[i], v = sz++;
      dn[bt[c]] = v;
      up[v] = bt[c], bt[c] = v;
      rg[v] = (i +
                     == (int)col.size() ? f : v + 1);
      rw[v] = r, cl[v] = c;
      ++s[c];
      if (i > 0) lt[v] = v - 1;
    lt[f] = sz - 1;
  int h() {
    int ret = 0;
    memset(vis, 0, sizeof(bool) * sz);
TRAV(x, rg, head) {
      if (vis[x]) continue;
      vis[x] = true, ++ret;
      TRAV(i, dn, x) TRAV(j, rg, i) vis[cl[j]] = true;
    return ret;
  void dfs(int dep) {
```

```
if (dep + (A ? 0 : h()) >= ans) return;
    if (rg[head] == head) return ans = dep, void();
    if (dn[rg[head]] == rg[head]) return;
    int w = rg[head];
    TRAV(x, rg, head) if (s[x] < s[w]) w = x;
    if (A) remove(w);
    TRAV(i, dn, w) {
      if (B) remove(i);
      TRAV(j, rg, i) remove(A ? cl[j] : j);
      dfs(dep + 1);
      TRAV(j, lt, i) restore(A ? cl[j] : j);
      if (B) restore(i);
    if (A) restore(w);
  int solve() {
    for (int i = 0; i < columns; ++i)</pre>
      dn[bt[i]] = i, up[i] = bt[i];
    ans = 1e9, dfs(0);
    return ans:
  }
};
```

9.9 Matroid Intersection

Start from $S = \emptyset$. In each iteration, let

```
• Y_1 = \{x \notin S \mid S \cup \{x\} \in I_1\}
• Y_2 = \{x \notin S \mid S \cup \{x\} \in I_2\}
```

If there exists $x \in Y_1 \cap Y_2$, insert x into S. Otherwise for each $x \in S, y \not\in S$, create edges

```
• x \to y if S - \{x\} \cup \{y\} \in I_1.
• y \to x if S - \{x\} \cup \{y\} \in I_2.
```

Find a *shortest* path (with BFS) starting from a vertex in Y_1 and ending at a vertex in Y_2 which doesn't pass through any other vertices in Y_2 , and alternate the path. The size of S will be incremented by 1 in each iteration. For the weighted case, assign weight w(x) to vertex x if $x \in S$ and -w(x) if $x \notin S$. Find the path with the minimum number of edges among all minimum length paths and alternate it.

9.10 AdaptiveSimpson

```
using F_t = function < double (double) >;
pdd simpson(const F_t &f, double l, double r
  double fl, double fr, double fm = nan("")) {
  if (isnan(fm)) fm = f((l + r) / 2);
return {fm, (r - l) / 6 * (fl + 4 * fm + fr)};
double simpson_ada(const F_t &f, double l, double r,
  double fl, double fm, double fr, double eps) {
  double m = (l + r) / 2,
          s = simpson(f, l, r, fl, fr, fm).second;
  auto [flm, sl] = simpson(f, l, m, fl, fm);
auto [fmr, sr] = simpson(f, m, r, fm, fr);
  double delta = sl + sr - s;
  if (abs(delta) <= 15 * eps)</pre>
    return sl + sr + delta / 15;
  return simpson_ada(f, l, m, fl, flm, fm, eps / 2) +
    simpson_ada(f, m, r, fm, fmr, fr, eps / 2);
double simpson_ada(const F_t &f, double l, double r) {
  return simpson_ada(
    f, l, r, f(l), f((l + r) / 2), f(r), 1e-9 / 7122);
double simpson_ada2(const F_t &f, double l, double r) {
  double h = (r - l) / 7122, s = 0;
for (int i = 0; i < 7122; ++i, l += h)</pre>
    s += simpson_ada(f, l, l + h);
  return s;
```

9.11 Simulated Annealing

9.12 Tree Hash*

```
ull seed;
ull shift(ull x) {
    x ^= x << 13;
    x ^= x >> 7;
    x ^= x << 17;
    return x;
}
ull dfs(int u, int f) {
    ull sum = seed;
    for (int i : G[u])
        if (i != f)
        sum += shift(dfs(i, u));
    return sum;
}</pre>
```

9.13 Binary Search On Fraction

```
struct Q {
  int p, q;
Q go(Q b, int d) { return {p + b.p*d, q + b.q*d}; }
bool pred(Q);
// returns smallest p/q in [lo, hi] such that
// pred(p/q) is true, and 0 <= p,q <= N
Q frac_bs(int N) {
  Q lo{0, 1}, hi{1, 0};
  if (pred(lo)) return lo;
  assert(pred(hi));
  bool dir = 1, L = 1, H = 1;
for (; L || H; dir = !dir) {
    int len = 0, step = 1;
    for (int t = 0; t < 2 && (t ? step/=2 : step*=2);)</pre>
      if (Q mid = hi.go(lo, len + step);
          mid.p > N || mid.q > N || dir ^ pred(mid))
        t++:
      else len += step;
    swap(lo, hi = hi.go(lo, len));
    (dir ? L : H) = !!len;
  return dir ? hi : lo;
```

10 Python

10.1 Misc