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```
rope<char> *root[10]; // nsqrt(n)
root[0] = new rope<char>();
root[1] = new rope<char>(*root[0]);
// root[1]->insert(pos, 'a');
// root[1]->at(pos); 0-base
// root[1]->erase(pos, size);
}
// __int128_t,__float128_t
// for (int i = bs._Find_first(); i < bs.size(); i = bs
._Find_next(i));</pre>
```

### 2 Graph

#### 2.1 BCC Vertex\*

```
vector<int> G[N]; // 1-base
vector<int> nG[N * 2], bcc[N];
int low[N], dfn[N], Time;
int bcc_id[N], bcc_cnt; // 1-base
bool is_cut[N]; // whether is av
bool cir[N * 2];
int st[N], top;
void dfs(int u, int pa = -1) {
  int child = 0;
  low[u] = dfn[u] = ++Time;
  st[top++] = u;
  for (int v : G[u])
    if (!dfn[v]) {
      dfs(v, u), ++child;
      low[u] = min(low[u], low[v]);
      if (dfn[u] <= low[v]) {</pre>
        is_cut[u] = 1;
        bcc[++bcc_cnt].clear();
        int t;
        do {
          bcc_id[t = st[--top]] = bcc_cnt;
          bcc[bcc_cnt].push_back(t);
        } while (t != v);
        bcc_id[u] = bcc_cnt;
        bcc[bcc_cnt].pb(u);
    } else if (dfn[v] < dfn[u] && v != pa)</pre>
      low[u] = min(low[u], dfn[v]);
  if (pa == -1 && child < 2) is_cut[u] = 0;</pre>
void bcc_init(int n) { // TODO: init {nG, cir}[1..2n]
 Time = bcc_cnt = top = 0;
  for (int i = 1; i <= n; ++i)</pre>
    G[i].clear(), dfn[i] = bcc_id[i] = is_cut[i] = 0;
void bcc_solve(int n) {
 for (int i = 1; i <= n; ++i)
    if (!dfn[i]) dfs(i);
  // block-cut tree
  for (int i = 1; i <= n; ++i)</pre>
    if (is_cut[i])
      bcc_id[i] = ++bcc_cnt, cir[bcc_cnt] = 1;
  for (int i = 1; i <= bcc_cnt && !cir[i]; ++i)</pre>
    for (int j : bcc[i])
      if (is_cut[j])
        nG[i].pb(bcc_id[j]), nG[bcc_id[j]].pb(i);
```

### 2.2 Bridge\*

```
int low[N], dfn[N], Time; // 1-base
vector<pii> G[N], edge;
vector<bool> is_bridge;

void init(int n) {
   Time = 0;
   for (int i = 1; i <= n; ++i)
      G[i].clear(), low[i] = dfn[i] = 0;
}

void add_edge(int a, int b) {</pre>
```

```
G[a].pb(pii(b, SZ(edge))), G[b].pb(pii(a, SZ(edge)));
edge.pb(pii(a, b));
}

void dfs(int u, int f) {
    dfn[u] = low[u] = ++Time;
    for (auto i : G[u])
        if (!dfn[i.X])
            dfs(i.X, i.Y), low[u] = min(low[u], low[i.X]);
        else if (i.Y != f) low[u] = min(low[u], dfn[i.X]);
    if (low[u] == dfn[u] && f != -1) is_bridge[f] = 1;
}

void solve(int n) {
    is_bridge.resize(SZ(edge));
    for (int i = 1; i <= n; ++i)
        if (!dfn[i]) dfs(i, -1);
}</pre>
```

### 2.3 2SAT (SCC)\*

```
struct SAT { // 0-base
   int low[N], dfn[N], bln[N], n, Time, nScc;
   bool instack[N], istrue[N];
   stack<int> st;
   vector<int> G[N], SCC[N];
   void init(int _n) {
     n = _n; // assert(n * 2 <= N);
     for (int i = 0; i < n + n; ++i) G[i].clear();</pre>
   void add_edge(int a, int b) { G[a].pb(b); }
   int rv(int a) {
     if (a >= n) return a - n;
     return a + n;
   void add_clause(int a, int b) {
     add_edge(rv(a), b), add_edge(rv(b), a);
   void dfs(int u) {
     dfn[u] = low[u] = ++Time;
     instack[u] = 1, st.push(u);
     for (int i : G[u])
       if (!dfn[i])
         dfs(i), low[u] = min(low[i], low[u]);
       else if (instack[i] && dfn[i] < dfn[u])</pre>
         low[u] = min(low[u], dfn[i]);
     if (low[u] == dfn[u]) {
       int tmp;
       do {
         tmp = st.top(), st.pop();
instack[tmp] = 0, bln[tmp] = nScc;
       } while (tmp != u);
       ++nScc;
     }
   bool solve() {
     Time = nScc = 0;
     for (int i = 0; i < n + n; ++i)</pre>
       SCC[i].clear(), low[i] = dfn[i] = bln[i] = 0;
     for (int i = 0; i < n + n; ++i)</pre>
       if (!dfn[i]) dfs(i);
     for (int i = 0; i < n + n; ++i) SCC[bln[i]].pb(i);</pre>
     for (int i = 0; i < n; ++i) {</pre>
       if (bln[i] == bln[i + n]) return false;
       istrue[i] = bln[i] < bln[i + n];</pre>
       istrue[i + n] = !istrue[i];
     return true;
  }
};
```

#### 2.4 MinimumMeanCycle\*

```
11 road[N][N]; // input here
struct MinimumMeanCycle {
   11 dp[N + 5][N], n;
   pl1 solve() {
     11 a = -1, b = -1, L = n + 1;
     for (int i = 2; i <= L; ++i)
        for (int k = 0; k < n; ++k)
        for (int j = 0; j < n; ++j)</pre>
```

```
dp[i][i] =
          min(dp[i - 1][k] + road[k][j], dp[i][j]);
  for (int i = 0; i < n; ++i) {</pre>
   if (dp[L][i] >= INF) continue;
   11 ta = 0, tb = 1;
   for (int j = 1; j < n; ++j)
      if (dp[j][i] < INF &&</pre>
        ta * (L - j) < (dp[L][i] - dp[j][i]) * tb)
        ta = dp[L][i] - dp[j][i], tb = L - j;
    if (ta == 0) continue;
   if (a == -1 || a * tb > ta * b) a = ta, b = tb;
 if (a != -1) {
   ll g = \_gcd(a, b);
   return pll(a / g, b / g);
 return pll(-1LL, -1LL);
void init(int _n) {
 n = _n;
 for (int i = 0; i < n; ++i)</pre>
    for (int j = 0; j < n; ++j) dp[i + 2][j] = INF;
```

#### 2.5 Virtual Tree\*

```
vector<int> vG[N];
int top, st[N];
void insert(int u) {
 if (top == -1) return st[++top] = u, void();
  int p = LCA(st[top], u);
 if (p == st[top]) return st[++top] = u, void();
 while (top >= 1 && dep[st[top - 1]] >= dep[p])
    vG[st[top - 1]].pb(st[top]), --top;
 if (st[top] != p)
   vG[p].pb(st[top]), --top, st[++top] = p;
 st[++top] = u;
}
void reset(int u) {
 for (int i : vG[u]) reset(i);
  vG[u].clear();
void solve(vector<int> &v) {
 top = -1;
 sort(ALL(v),
   [&](int a, int b) { return dfn[a] < dfn[b]; });</pre>
  for (int i : v) insert(i);
 while (top > 0) vG[st[top - 1]].pb(st[top]), --top;
 // do somethina
  reset(v[0]);
```

#### 2.6 Maximum Clique Dyn\*

```
struct MaxClique { // fast when N <= 100</pre>
 bitset<N> G[N], cs[N];
  int ans, sol[N], q, cur[N], d[N], n;
  void init(int _n) {
    n = _n;
    for (int i = 0; i < n; ++i) G[i].reset();</pre>
  void add_edge(int u, int v) {
    G[u][v] = G[v][u] = 1;
  void pre_dfs(vector<int> &r, int 1, bitset<N> mask) {
    if (1 < 4) {</pre>
      for (int i : r) d[i] = (G[i] & mask).count();
      sort(ALL(r), [\&](int x, int y) \{ return d[x] > d[
    vector<int> c(SZ(r));
    int lft = max(ans - q + 1, 1), rgt = 1, tp = 0;
    cs[1].reset(), cs[2].reset();
    for (int p : r) {
      int k = 1;
      while ((cs[k] & G[p]).any()) ++k;
      if (k > rgt) cs[++rgt + 1].reset();
```

```
cs[k][p] = 1;
      if (k < lft) r[tp++] = p;
    for (int k = 1ft; k <= rgt; ++k)</pre>
      for (int p = cs[k]._Find_first(); p < N; p = cs[k</pre>
           ]._Find_next(p))
        r[tp] = p, c[tp] = k, ++tp;
    dfs(r, c, l + 1, mask);
  void dfs(vector<int> &r, vector<int> &c, int 1,
      bitset<N> mask) {
    while (!r.empty()) {
      int p = r.back();
      r.pop_back(), mask[p] = 0;
      if (q + c.back() <= ans) return;</pre>
      cur[q++] = p;
      vector<int> nr;
      for (int i : r) if (G[p][i]) nr.pb(i);
      if (!nr.empty()) pre_dfs(nr, 1, mask & G[p]);
      else if (q > ans) ans = q, copy_n(cur, q, sol);
      c.pop_back(), --q;
  int solve() {
    vector<int> r(n);
    ans = q = 0, iota(ALL(r), 0);
    pre_dfs(r, 0, bitset<N>(string(n, '1')));
    return ans;
  }
};
```

### 2.7 Minimum Steiner Tree\*

```
struct SteinerTree { // 0-base
  int n, dst[N][N], dp[1 << T][N], tdst[N];</pre>
  int vcst[N]; // the cost of vertexs
  void init(int _n) {
    n = _n;
for (int i = 0; i < n; ++i) {</pre>
      fill_n(dst[i], n, INF);
      dst[i][i] = vcst[i] = 0;
  void chmin(int &x, int val) {
    x = min(x, val);
  void add_edge(int ui, int vi, int wi) {
    chmin(dst[ui][vi], wi);
  void shortest_path() {
    for (int k = 0; k < n; ++k)
      for (int i = 0; i < n; ++i)</pre>
        for (int j = 0; j < n; ++j)</pre>
           chmin(dst[i][j], dst[i][k] + dst[k][j]);
  int solve(const vector<int>& ter) {
    shortest_path();
    int t = SZ(ter), full = (1 << t) - 1;</pre>
    for (int i = 0; i <= full; ++i)</pre>
      fill_n(dp[i], n, INF);
    copy_n(vcst, n, dp[0]);
    for (int msk = 1; msk <= full; ++msk) {</pre>
      if (!(msk & (msk - 1))) {
        int who = __lg(msk);
for (int i = 0; i < n; ++i)</pre>
           dp[msk][i] = vcst[ter[who]] + dst[ter[who]][i
               ];
      for (int i = 0; i < n; ++i)</pre>
         for (int sub = (msk - 1) & msk; sub; sub = (sub
               - 1) & msk)
           chmin(dp[msk][i], dp[sub][i] + dp[msk ^ sub][
              i] - vcst[i]);
      for (int i = 0; i < n; ++i) {
         tdst[i] = INF;
         for (int j = 0; j < n; ++j)</pre>
           chmin(tdst[i], dp[msk][j] + dst[j][i]);
      copy_n(tdst, n, dp[msk]);
    return *min_element(dp[full], dp[full] + n);
```

```
| \ \}; // O(V 3^T + V^2 2^T)
```

### 2.8 Dominator Tree\*

```
struct dominator_tree { // 1-base
  vector<int> G[N], rG[N];
  int n, pa[N], dfn[N], id[N], Time;
  int semi[N], idom[N], best[N];
  vector<int> tree[N]; // dominator_tree
  void init(int _n) {
    n = _n;
    for (int i = 1; i <= n; ++i)</pre>
      G[i].clear(), rG[i].clear();
  void add edge(int u, int v) {
    G[u].pb(v), rG[v].pb(u);
  void dfs(int u) {
    id[dfn[u] = ++Time] = u;
    for (auto v : G[u])
      if (!dfn[v]) dfs(v), pa[dfn[v]] = dfn[u];
  int find(int y, int x) {
    if (y <= x) return y;</pre>
    int tmp = find(pa[y], x);
    if (semi[best[y]] > semi[best[pa[y]]])
      best[y] = best[pa[y]];
    return pa[y] = tmp;
  void tarjan(int root) {
    Time = ∅;
    for (int i = 1; i <= n; ++i) {</pre>
      dfn[i] = idom[i] = 0;
      tree[i].clear();
      best[i] = semi[i] = i;
    dfs(root);
    for (int i = Time; i > 1; --i) {
      int u = id[i];
      for (auto v : rG[u])
        if (v = dfn[v]) {
          find(v, i);
          semi[i] = min(semi[i], semi[best[v]]);
      tree[semi[i]].pb(i);
      for (auto v : tree[pa[i]]) {
        find(v, pa[i]);
        idom[v] =
          semi[best[v]] == pa[i] ? pa[i] : best[v];
      tree[pa[i]].clear();
    for (int i = 2; i <= Time; ++i) {
      if (idom[i] != semi[i]) idom[i] = idom[idom[i]];
      tree[id[idom[i]]].pb(id[i]);
 }
}:
```

### 2.9 Minimum Arborescence\*

```
/* TODO
DSU: disjoint set
- DSU(n), .boss(x), .Union(x, y)
min_heap<T, Info>: min heap for type {T, Info} with
    lazy tag
- .push({w, i}), .top(), .join(heap), .pop(), .empty(),
     .add_lazy(v)
struct E { int s, t; ll w; }; // O-base
vector<int> dmst(const vector<E> &e, int n, int root) {
 vector<min_heap<11, int>> h(n * 2);
  for (int i = 0; i < SZ(e); ++i)</pre>
    h[e[i].t].push({e[i].w, i});
  DSU dsu(n * 2);
  vector<int> v(n * 2, -1), pa(n * 2, -1), r(n * 2);
  v[root] = n + 1;
 int pc = n;
  for (int i = 0; i < n; ++i) if (v[i] == -1) {</pre>
    for (int p = i; v[p] == -1 \mid \mid v[p] == i; p = dsu.
        boss(e[r[p]].s)) {
```

```
if (v[p] == i) {
        int q = p; p = pc++;
          h[q].add_lazy(-h[q].top().X);
          pa[q] = p, dsu.Union(p, q), h[p].join(h[q]);
        } while ((q = dsu.boss(e[r[q]].s)) != p);
      v[p] = i;
      while (!h[p].empty() \&\& dsu.boss(e[h[p].top().Y].
          s) == p
        h[p].pop();
      if (h[p].empty()) return {}; // no solution
      r[p] = h[p].top().Y;
   }
  }
  vector<int> ans;
  for (int i = pc - 1; i >= 0; i--) if (i != root && v[
      i] != n) {
    for (int f = e[r[i]].t; ~f && v[f] != n; f = pa[f])
     v[f] = n;
    ans.pb(r[i]);
  }
  return ans; // default minimize, returns edgeid array
} // O(Ef(E)), f(E) from min_heap
```

#### 2.10 Vizing's theorem\*

```
namespace vizing { // returns edge coloring in adjacent
     matrix G. 1 - based
const int N = 105;
int C[N][N], G[N][N], X[N], vst[N], n;
void init(int _n) { n = _n;
  for (int i = 0; i <= n; ++i)
    for (int j = 0; j <= n; ++j)</pre>
      C[i][j] = G[i][j] = 0;
void solve(vector<pii> &E) {
  auto update = [&](int u)
  { for (X[u] = 1; C[u][X[u]]; ++X[u]); };
  auto color = [&](int u, int v, int c) {
    int p = G[u][v];
    G[u][v] = G[v][u] = c;
    C[u][c] = v, C[v][c] = u;
    C[u][p] = C[v][p] = 0;
    if (p) X[u] = X[v] = p;
    else update(u), update(v);
    return p;
  }:
  auto flip = [&](int u, int c1, int c2) {
    int p = C[u][c1];
    swap(C[u][c1], C[u][c2]);
    if (p) G[u][p] = G[p][u] = c2;
    if (!C[u][c1]) X[u] = c1;
    if (!C[u][c2]) X[u] = c2;
    return p;
  fill_n(X + 1, n, 1);
  for (int t = 0; t < SZ(E); ++t) {</pre>
    int u = E[t].X, v0 = E[t].Y, v = v0, c0 = X[u], c =
         c0, d;
    vector<pii> L;
    fill_n(vst + 1, n, 0);
    while (!G[u][v0]) {
      L.emplace_back(v, d = X[v]);
      if (!C[v][c]) for (int a = SZ(L) - 1; a >= 0; --a
          ) c = color(u, L[a].X, c);
      else if (!C[u][d]) for (int a = SZ(L) - 1; a >=
      0; --a) color(u, L[a].X, L[a].Y);
else if (vst[d]) break;
      else vst[d] = 1, v = C[u][d];
    if (!G[u][v0]) {
      for (; v; v = flip(v, c, d), swap(c, d));
      if (int a; C[u][c0]) {
        for (a = SZ(L) - 2; a >= 0 && L[a].Y != c; --a)
        for (; a >= 0; --a) color(u, L[a].X, L[a].Y);
      else --t;
    }
  }
```

```
|} // namespace vizing
```

### 2.11 Minimum Clique Cover\*

```
struct Clique_Cover { // 0-base, 0(n2^n)
  int co[1 << N], n, E[N];</pre>
  int dp[1 << N];</pre>
  void init(int _n) {
    n = _n, fill_n(dp, 1 << n, 0);
    fill_n(E, n, 0), fill_n(co, 1 << n, 0);
  void add_edge(int u, int v) {
    E[u] = 1 << v, E[v] = 1 << u;
  int solve() {
     for (int i = 0; i < n; ++i)</pre>
       co[1 << i] = E[i] | (1 << i);
     co[0] = (1 << n) - 1;

dp[0] = (n & 1) * 2 - 1;
     for (int i = 1; i < (1 << n); ++i) {</pre>
       int t = i & -i;
dp[i] = -dp[i ^ t];
       co[i] = co[i ^ t] & co[t];
     for (int i = 0; i < (1 << n); ++i)</pre>
       co[i] = (co[i] & i) == i;
     fwt(co, 1 << n, 1);
    for (int ans = 1; ans < n; ++ans) {
  int sum = 0; // probabilistic</pre>
       for (int i = 0; i < (1 << n); ++i)</pre>
         sum += (dp[i] *= co[i]);
       if (sum) return ans;
     return n;
  }
};
```

### 2.12 NumberofMaximalClique\*

```
struct BronKerbosch { // 1-base
   int n, a[N], g[N][N];
   int S, all[N][N], some[N][N], none[N][N];
   void init(int _n) {
    n = _n;
for (int i = 1; i <= n; ++i)</pre>
       for (int j = 1; j <= n; ++j) g[i][j] = 0;</pre>
   void add_edge(int u, int v) {
     g[u][v] = g[v][u] = 1;
   void dfs(int d, int an, int sn, int nn) {
  if (S > 1000) return; // pruning
     if (sn == 0 && nn == 0) ++S;
     int u = some[d][0];
     for (int i = 0; i < sn; ++i) {</pre>
       int v = some[d][i];
       if (g[u][v]) continue;
       int tsn = 0, tnn = 0;
       copy_n(all[d], an, all[d + 1]);
       all[d + 1][an] = v;
       for (int j = 0; j < sn; ++j)
         if (g[v][some[d][j]])
            some[d + 1][tsn++] = some[d][j];
       for (int j = 0; j < nn; ++j)
  if (g[v][none[d][j]])</pre>
            none[d + 1][tnn++] = none[d][j];
       dfs(d + 1, an + 1, tsn, tnn);
       some[d][i] = 0, none[d][nn++] = v;
     }
   int solve() {
     iota(some[0], some[0] + n, 1);
     S = 0, dfs(0, 0, n, 0);
     return S;
  }
};
```

#### Data Structure 3

#### 3.1 Discrete Trick

```
vector<int> val;
// build
sort(ALL(val)), val.resize(unique(ALL(val)) - val.begin
    ());
// index of x
upper_bound(ALL(val), x) - val.begin();
// max idx <= x
upper_bound(ALL(val), x) - val.begin();
// max idx < x
lower_bound(ALL(val), x) - val.begin();
3.2 Leftist Tree
struct node {
 11 v, data, sz, sum;
node *1, *r;
  node(ll k)
    : v(0), data(k), sz(1), l(0), r(0), sum(k) {}
11 sz(node *p) { return p ? p->sz : 0; }
11 V(node *p) { return p ? p->v : -1; }
11 sum(node *p) { return p ? p->sum : 0; }
node *merge(node *a, node *b) {
 if (!a || !b) return a ? a : b;
```

### 3.3 Heavy light Decomposition

if (a->data < b->data) swap(a, b);

if  $(V(a\rightarrow r) \rightarrow V(a\rightarrow l))$  swap $(a\rightarrow r, a\rightarrow l)$ ;

 $a\rightarrow sum = sum(a\rightarrow 1) + sum(a\rightarrow r) + a\rightarrow data;$ 

a -> v = V(a -> r) + 1, a -> sz = sz(a -> 1) + sz(a -> r) + 1;

a->r = merge(a->r, b);

o = merge(o->1, o->r);

return a;

delete tmp;

void pop(node \*&o) { node \*tmp = o;

}

}

```
struct Heavy_light_Decomposition { // 1-base
   int n, ulink[N], deep[N], mxson[N], w[N], pa[N];
   int t, pl[N], data[N], dt[N], bln[N], edge[N], et;
   vector<pii> G[N];
   void init(int _n) {
    n = _n, t = 0, et = 1;
for (int i = 1; i <= n; ++i)
       G[i].clear(), mxson[i] = 0;
   void add_edge(int a, int b, int w) {
     G[a].pb(pii(b, et));
     G[b].pb(pii(a, et));
     edge[et++] = w;
  void dfs(int u, int f, int d) {
  w[u] = 1, pa[u] = f, deep[u] = d++;
     for (auto &i : G[u])
       if (i.X != f) {
         dfs(i.X, u, d), w[u] += w[i.X];
         if (w[mxson[u]] < w[i.X]) mxson[u] = i.X;</pre>
       } else bln[i.Y] = u, dt[u] = edge[i.Y];
   void cut(int u, int link) {
     data[pl[u] = t++] = dt[u], ulink[u] = link;
     if (!mxson[u]) return;
     cut(mxson[u], link);
     for (auto i : G[u])
       if (i.X != pa[u] && i.X != mxson[u])
         cut(i.X, i.X);
   void build() { dfs(1, 1, 1), cut(1, 1), /*build*/; }
   int query(int a, int b) {
     int ta = ulink[a], tb = ulink[b], re = 0;
     while (ta != tb)
       if (deep[ta] < deep[tb])</pre>
           *query*/, tb = ulink[b = pa[tb]];
       else /*query*/, ta = ulink[a = pa[ta]];
     if (a == b) return re;
     if (pl[a] > pl[b]) swap(a, b);
     /*query*/
     return re;
};
```

### 3.4 Centroid Decomposition\*

```
struct Cent_Dec { // 1-base
  vector<pll> G[N];
  pll info[N]; // store info. of itself
  pll upinfo[N]; // store info. of climbing up
  int n, pa[N], layer[N], sz[N], done[N];
  11 dis[__lg(N) + 1][N];
  void init(int _n) {
    n = _n, layer[0] = -1;
    fill_n(pa + 1, n, 0), fill_n(done + 1, n, 0);
    for (int i = 1; i <= n; ++i) G[i].clear();</pre>
  void add_edge(int a, int b, int w) {
    G[a].pb(pll(b, w)), G[b].pb(pll(a, w));
  void get_cent(
    int u, int f, int &mx, int &c, int num) {
    int mxsz = 0;
    sz[u] = 1;
    for (pll e : G[u])
      if (!done[e.X] && e.X != f) {
        get_cent(e.X, u, mx, c, num);
        sz[u] += sz[e.X], mxsz = max(mxsz, sz[e.X]);
    if (mx > max(mxsz, num - sz[u]))
      mx = max(mxsz, num - sz[u]), c = u;
  void dfs(int u, int f, ll d, int org) {
    // if required, add self info or climbing info
    dis[layer[org]][u] = d;
    for (pll e : G[u])
      if (!done[e.X] && e.X != f)
        dfs(e.X, u, d + e.Y, org);
  int cut(int u, int f, int num) {
    int mx = 1e9, c = 0, lc;
    get_cent(u, f, mx, c, num);
    done[c] = 1, pa[c] = f, layer[c] = layer[f] + 1;
    for (pll e : G[c])
      if (!done[e.X]) {
        if (sz[e.X] > sz[c])
          lc = cut(e.X, c, num - sz[c]);
        else lc = cut(e.X, c, sz[e.X]);
        upinfo[lc] = pll(), dfs(e.X, c, e.Y, c);
    return done[c] = 0, c;
  void build() { cut(1, 0, n); }
  void modify(int u) {
    for (int a = u, ly = layer[a]; a;
         a = pa[a],
                    --ly) {
      info[a].X += dis[ly][u], ++info[a].Y;
      if (pa[a])
        upinfo[a].X += dis[ly - 1][u], ++upinfo[a].Y;
    }
  11 query(int u) {
    11 \text{ rt} = 0;
    for (int a = u, ly = layer[a]; a;
         a = pa[a], --ly) {
      rt += info[a].X + info[a].Y * dis[ly][u];
      if (pa[a])
          upinfo[a].X + upinfo[a].Y * dis[ly - 1][u];
    return rt;
  }
};
```

#### 3.5 Link cut tree\*

```
struct Splay { // xor-sum
    static Splay nil;
    Splay *ch[2], *f;
    int val, sum, rev, size;
    Splay(int _val = 0)
        : val(_val), sum(_val), rev(0), size(1) {
        f = ch[0] = ch[1] = &nil;
    }
    bool isr() {
        return f->ch[0] != this && f->ch[1] != this;
}
```

```
int dir() { return f->ch[0] == this ? 0 : 1; }
  void setCh(Splay *c, int d) {
    ch[d] = c;
    if (c != &nil) c->f = this;
    pull();
  void give_tag(int r)
  { if (r) swap(ch[0], ch[1]), rev ^= 1; }
  void push() {
    if (ch[0] != &nil) ch[0]->give_tag(rev);
    if (ch[1] != &nil) ch[1]->give_tag(rev);
  void pull() {
    // take care of the nil!
    size = ch[0] -> size + ch[1] -> size + 1;
    sum = ch[0] \rightarrow sum ^ ch[1] \rightarrow sum ^ val;
    if (ch[0] != &nil) ch[0]->f = this;
    if (ch[1] != &nil) ch[1]->f = this;
} Splay::nil;
Splay *nil = &Splay::nil;
void rotate(Splay *x) {
  Splay *p = x->f;
  int d = x->dir();
  if (!p->isr()) p->f->setCh(x, p->dir());
  else x->f = p->f;
  p->setCh(x->ch[!d], d);
  x->setCh(p, !d);
  p->pull(), x->pull();
void splav(Splav *x) {
  vector<Splay *> splayVec;
  for (Splay *q = x;; q = q \rightarrow f) {
    splayVec.pb(q);
    if (q->isr()) break;
  reverse(ALL(splayVec));
  for (auto it : splayVec) it->push();
  while (!x->isr()) {
    if (x->f->isr()) rotate(x);
    else if (x->dir() == x->f->dir())
      rotate(x->f), rotate(x);
    else rotate(x), rotate(x);
  }
Splay *access(Splay *x) {
  Splay *q = nil;
  for (; x != nil; x = x->f)
    splay(x), x \rightarrow setCh(q, 1), q = x;
  return q;
void root_path(Splay *x) { access(x), splay(x); }
void chroot(Splay *x) {
  root_path(x), x->rev ^= 1;
  x->push(), x->pull();
void split(Splay *x, Splay *y) {
  chroot(x), root_path(y);
void link(Splay *x, Splay *y) {
  root_path(x), chroot(y);
  x->setCh(y, 1);
void cut(Splay *x, Splay *y) {
  split(x, y);
  if (y->size != 5) return;
  y->push();
  y - ch[0] = y - ch[0] - f = nil;
Splay *get_root(Splay *x) {
  for (root_path(x); x\rightarrow ch[0] != nil; x = x\rightarrow ch[0])
   x->push();
  splay(x);
  return x;
bool conn(Splay *x, Splay *y) {
  return get_root(x) == get_root(y);
Splay *lca(Splay *x, Splay *y) {
```

access(x), root\_path(y);

```
if (y->f == nil) return y;
  return y->f;
}
void change(Splay *x, int val) {
  splay(x), x->val = val, x->pull();
}
int query(Splay *x, Splay *y) {
  split(x, y);
  return y->sum;
}
```

#### 3.6 KDTree

```
namespace kdt {
int root, lc[maxn], rc[maxn], xl[maxn], xr[maxn],
 yl[maxn], yr[maxn];
point p[maxn];
int build(int 1, int r, int dep = 0) {
  if (1 == r) return -1;
  function<bool(const point &, const point &)> f =
    [dep](const point &a, const point &b) {
      if (dep & 1) return a.x < b.x;</pre>
      else return a.y < b.y;</pre>
    };
  int m = (1 + r) >> 1;
  nth_element(p + 1, p + m, p + r, f);
  x1[m] = xr[m] = p[m].x;
  yl[m] = yr[m] = p[m].y;
  lc[m] = build(1, m, dep + 1);
  if (~lc[m]) {
    xl[m] = min(xl[m], xl[lc[m]]);
    xr[m] = max(xr[m], xr[lc[m]]);
    yl[m] = min(yl[m], yl[lc[m]]);
    yr[m] = max(yr[m], yr[lc[m]]);
  rc[m] = build(m + 1, r, dep + 1);
  if (~rc[m]) {
   xl[m] = min(xl[m], xl[rc[m]]);
    xr[m] = max(xr[m], xr[rc[m]]);
   yl[m] = min(yl[m], yl[rc[m]]);
yr[m] = max(yr[m], yr[rc[m]]);
  return m;
bool bound(const point &q, int o, long long d) {
  double ds = sqrt(d + 1.0);
  if (q.x < x1[o] - ds || q.x > xr[o] + ds ||
    q.y < yl[o] - ds || q.y > yr[o] + ds)
    return false;
  return true;
long long dist(const point &a, const point &b) {
  return (a.x - b.x) * 111 * (a.x - b.x) +
    (a.y - b.y) * 111 * (a.y - b.y);
void dfs(
  const point &q, long long &d, int o, int dep = 0) {
  if (!bound(q, o, d)) return;
  long long cd = dist(p[o], q);
  if (cd != 0) d = min(d, cd);
  if ((dep & 1) && q.x < p[o].x ||</pre>
    !(dep & 1) && q.y < p[o].y) {
    if (~lc[o]) dfs(q, d, lc[o], dep + 1);
    if (~rc[o]) dfs(q, d, rc[o], dep + 1);
  } else {
    if (~rc[o]) dfs(q, d, rc[o], dep + 1);
    if (~lc[o]) dfs(q, d, lc[o], dep + 1);
  }
void init(const vector<point> &v) {
  for (int i = 0; i < v.size(); ++i) p[i] = v[i];</pre>
  root = build(0, v.size());
long long nearest(const point &q) {
  long long res = 1e18;
  dfs(q, res, root);
  return res;
} // namespace kdt
```

# 4 Flow/Matching

#### 4.1 Kuhn Munkres\*

```
struct KM { // 0-base
  11 w[N][N], h1[N], hr[N], s1k[N];
  int f1[N], fr[N], pre[N], qu[N], q1, qr, n;
  bool v1[N], vr[N];
  void init(int _n) {
    n = _n;
    for (int i = 0; i < n; ++i)</pre>
      fill_n(w[i], n, -INF);
  void add_edge(int a, int b, ll wei) {
    w[a][b] = wei;
  bool Check(int x) {
    if (vl[x] = 1, \sim fl[x])
      return vr[qu[qr++] = fl[x]] = 1;
    while (\sim x) swap(x, fr[fl[x] = pre[x]]);
    return 0;
  void bfs(int s) {
    fill_n(slk, n, INF), fill_n(vl, n, 0), fill_n(vr, n
         , 0);
    ql = qr = 0, qu[qr++] = s, vr[s] = 1;
    for (11 d;;) {
      while (ql < qr)
for (int x = 0, y = qu[ql++]; x < n; ++x)</pre>
           if (!v1[x] \&\& s1k[x] >= (d = h1[x] + hr[y] -
               w[x][y])) {
             if (pre[x] = y, d) slk[x] = d;
             else if (!Check(x)) return;
      d = INF;
      for (int x = 0; x < n; ++x)
        if (!v1[x] && d > s1k[x]) d = s1k[x];
       for (int x = 0; x < n; ++x) {
        if (v1[x]) h1[x] += d;
         else slk[x] -= d;
        if (vr[x]) hr[x] -= d;
      for (int x = 0; x < n; ++x)
         if (!v1[x] && !slk[x] && !Check(x)) return;
    }
  11 solve() {
    fill_n(fl, n, -1), fill_n(fr, n, -1), fill_n(hr, n,
          0);
    for (int i = 0; i < n; ++i)</pre>
      hl[i] = *max_element(w[i], w[i] + n);
     for (int i = 0; i < n; ++i) bfs(i);</pre>
    11 res = 0;
    for (int i = 0; i < n; ++i) res += w[i][fl[i]];</pre>
    return res;
  }
};
```

#### 4.2 MincostMaxflow\*

```
struct MinCostMaxFlow { // 0-base
  struct Edge {
    11 from, to, cap, flow, cost, rev;
  } *past[N];
  vector<Edge> G[N];
  int inq[N], n, s, t;
  11 dis[N], up[N], pot[N];
  bool BellmanFord() {
    fill_n(dis, n, INF), fill_n(inq, n, 0);
    queue<int> q;
    auto relax = [&](int u, ll d, ll cap, Edge *e) {
      if (cap > 0 && dis[u] > d) {
        dis[u] = d, up[u] = cap, past[u] = e;
        if (!inq[u]) inq[u] = 1, q.push(u);
     }
    relax(s, 0, INF, 0);
    while (!q.empty()) {
      int u = q.front();
      q.pop(), inq[u] = 0;
      for (auto &e : G[u]) {
```

```
11 d2 = dis[u] + e.cost + pot[u] - pot[e.to];
relax(e.to, d2, min(up[u], e.cap - e.flow), &e)
      }
    return dis[t] != INF;
  void solve(int _s, int _t, ll &flow, ll &cost, bool
      neg = true) {
    s = _s, t = _t, flow = 0, cost = 0;
    if (neg) BellmanFord(), copy_n(dis, n, pot);
    for (; BellmanFord(); copy_n(dis, n, pot)) {
      for (int i = 0; i < n; ++i) dis[i] += pot[i] -</pre>
          pot[s];
      flow += up[t], cost += up[t] * dis[t];
      for (int i = t; past[i]; i = past[i]->from) {
        auto &e = *past[i];
        e.flow += up[t], G[e.to][e.rev].flow -= up[t];
      }
    }
  void init(int _n) {
    n = _n, fill_n(pot, n, 0);
    for (int i = 0; i < n; ++i) G[i].clear();</pre>
  void add_edge(ll a, ll b, ll cap, ll cost) {
    G[a].pb(Edge{a, b, cap, 0, cost, SZ(G[b])});
    G[b].pb(Edge{b, a, 0, 0, -cost, SZ(G[a]) - 1});
  }
};
```

### 4.3 Maximum Simple Graph Matching\*

```
struct GenMatch { // 1-base
 int V, pr[N];
  bool el[N][N], inq[N], inp[N], inb[N];
  int st, ed, nb, bk[N], djs[N], ans;
 void init(int _V) {
   V = V;
    for (int i = 0; i <= V; ++i) {</pre>
      for (int j = 0; j <= V; ++j) el[i][j] = 0;</pre>
      pr[i] = bk[i] = djs[i] = 0;
      inq[i] = inp[i] = inb[i] = 0;
   }
  void add_edge(int u, int v) {
   el[u][v] = el[v][u] = 1;
 int lca(int u, int v) {
   fill_n(inp, V + 1, 0);
   while (1)
      if (u = djs[u], inp[u] = true, u == st) break;
      else u = bk[pr[u]];
    while (1)
      if (v = djs[v], inp[v]) return v;
      else v = bk[pr[v]];
    return v;
  void upd(int u) {
   for (int v; djs[u] != nb;) {
      v = pr[u], inb[djs[u]] = inb[djs[v]] = true;
      u = bk[v];
      if (djs[u] != nb) bk[u] = v;
   }
 }
  void blo(int u, int v, queue<int> &qe) {
   nb = lca(u, v), fill_n(inb, V + 1, 0);
    upd(u), upd(v);
    if (djs[u] != nb) bk[u] = v;
    if (djs[v] != nb) bk[v] = u;
    for (int tu = 1; tu <= V; ++tu)</pre>
      if (inb[djs[tu]])
        if (djs[tu] = nb, !inq[tu])
          qe.push(tu), inq[tu] = 1;
  void flow() {
    fill_n(inq + 1, V, 0), fill_n(bk + 1, V, 0);
    iota(djs + 1, djs + V + 1, 1);
    queue<int> qe;
    qe.push(st), inq[st] = 1, ed = 0;
    while (!qe.empty()) {
      int u = qe.front();
```

```
qe.pop();
      for (int v = 1; v <= V; ++v)</pre>
        if (el[u][v] && djs[u] != djs[v] &&
          pr[u] != v) {
           if ((v == st) ||
             (pr[v] > 0 \&\& bk[pr[v]] > 0)) {
            blo(u, v, qe);
           } else if (!bk[v]) {
             if (bk[v] = u, pr[v] > 0) {
               if (!inq[pr[v]]) qe.push(pr[v]);
             } else {
               return ed = v, void();
          }
        }
    }
  void aug() {
    for (int u = ed, v, w; u > 0;)
      v = bk[u], w = pr[v], pr[v] = u, pr[u] = v,
  int solve() {
    fill_n(pr, V + 1, 0), ans = 0;
    for (int u = 1; u <= V; ++u)</pre>
      if (!pr[u])
        if (st = u, flow(), ed > 0) aug(), ++ans;
    return ans;
  }
};
```

### 4.4 Minimum Weight Matching (Clique version)\*

```
struct Graph { // 0-base (Perfect Match), n is even
   int n, match[N], onstk[N], stk[N], tp;
   11 edge[N][N], dis[N];
   void init(int _n) {
     n = _n, tp = 0;
     for (int i = 0; i < n; ++i) fill_n(edge[i], n, 0);</pre>
   void add_edge(int u, int v, ll w) {
     edge[u][v] = edge[v][u] = w;
   bool SPFA(int u) {
     stk[tp++] = u, onstk[u] = 1;
     for (int v = 0; v < n; ++v)
  if (!onstk[v] && match[u] != v) {</pre>
         int m = match[v];
         if (dis[m] >
           dis[u] - edge[v][m] + edge[u][v]) {
           dis[m] = dis[u] - edge[v][m] + edge[u][v];
           onstk[v] = 1, stk[tp++] = v;
           if (onstk[m] || SPFA(m)) return 1;
            --tp, onstk[v] = 0;
       }
     onstk[u] = 0, --tp;
     return 0;
   11 solve() { // find a match
     for (int i = 0; i < n; ++i) match[i] = i ^ 1;</pre>
     while (1) {
  int found = 0;
       fill_n(dis, n, 0);
       fill_n(onstk, n, 0);
       for (int i = 0; i < n; ++i)</pre>
         if (tp = 0, !onstk[i] && SPFA(i))
           for (found = 1; tp >= 2;) {
             int u = stk[--tp];
             int v = stk[--tp];
             match[u] = v, match[v] = u;
       if (!found) break;
     11 ret = 0;
     for (int i = 0; i < n; ++i)</pre>
       ret += edge[i][match[i]];
     return ret >> 1:
| };
```

#### 4.5 SW-mincut

```
struct SW{ // global min cut, O(V^3)
#define REP for (int i = 0; i < n; ++i)
  static const int MXN = 514, INF = 2147483647;
  int vst[MXN], edge[MXN][MXN], wei[MXN];
  void init(int n) {
   REP fill_n(edge[i], n, 0);
  void addEdge(int u, int v, int w){
    edge[u][v] += w; edge[v][u] += w;
  int search(int &s, int &t, int n){
   fill_n(vst, n, 0), fill_n(wei, n, 0);
    s = t = -1;
    int mx, cur;
    for (int j = 0; j < n; ++j) {</pre>
      mx = -1, cur = 0;
      REP if (wei[i] > mx) cur = i, mx = wei[i];
      vst[cur] = 1, wei[cur] = -1;
      s = t; t = cur;
      REP if (!vst[i]) wei[i] += edge[cur][i];
    return mx;
  int solve(int n) {
    int res = INF;
    for (int x, y; n > 1; n--){
      res = min(res, search(x, y, n));
      REP edge[i][x] = (edge[x][i] += edge[y][i]);
      REP {
        edge[y][i] = edge[n - 1][i];
        edge[i][y] = edge[i][n - 1];
      return res;
 }
} sw;
```

### 4.6 BoundedFlow\*(Dinic\*)

```
struct BoundedFlow { // 0-base
 struct edge {
   int to, cap, flow, rev;
  vector<edge> G[N];
 int n, s, t, dis[N], cur[N], cnt[N];
 void init(int _n) {
   n = _n;
for (int i = 0; i < n + 2; ++i)</pre>
      G[i].clear(), cnt[i] = 0;
  void add_edge(int u, int v, int lcap, int rcap) {
    cnt[u] -= lcap, cnt[v] += lcap;
    G[u].pb(edge{v, rcap, lcap, SZ(G[v])});
    G[v].pb(edge{u, 0, 0, SZ(G[u]) - 1});
  void add_edge(int u, int v, int cap) {
   G[u].pb(edge{v, cap, 0, SZ(G[v])});
   G[v].pb(edge{u, 0, 0, SZ(G[u]) - 1});
  int dfs(int u, int cap) {
    if (u == t || !cap) return cap;
    for (int &i = cur[u]; i < SZ(G[u]); ++i) {</pre>
      edge &e = G[u][i];
      if (dis[e.to] == dis[u] + 1 && e.cap != e.flow) {
        int df = dfs(e.to, min(e.cap - e.flow, cap));
        if (df) {
          e.flow += df, G[e.to][e.rev].flow -= df;
          return df;
        }
     }
    dis[u] = -1;
   return 0;
  bool bfs() {
   fill_n(dis, n + 3, -1);
    queue<int> q;
    q.push(s), dis[s] = 0;
    while (!q.empty()) {
      int u = q.front();
```

```
q.pop();
       for (edge &e : G[u])
         if (!~dis[e.to] && e.flow != e.cap)
           q.push(e.to), dis[e.to] = dis[u] + 1;
    return dis[t] != -1;
  int maxflow(int _s, int _t) {
    s = _s, t = _t;
     int flow = 0, df;
     while (bfs()) {
      fill_n(cur, n + 3, 0);
       while ((df = dfs(s, INF))) flow += df;
    return flow;
  bool solve() {
    int sum = 0;
     for (int i = 0; i < n; ++i)</pre>
      if (cnt[i] > 0)
         add_edge(n + 1, i, cnt[i]), sum += cnt[i];
       else if (cnt[i] < 0) add_edge(i, n + 2, -cnt[i]);</pre>
     if (sum != maxflow(n + 1, n + 2)) sum = -1;
     for (int i = 0; i < n; ++i)</pre>
       if (cnt[i] > 0)
         G[n + 1].pop_back(), G[i].pop_back();
       else if (cnt[i] < 0)</pre>
        G[i].pop_back(), G[n + 2].pop_back();
    return sum != -1;
  int solve(int _s, int
                          _t) {
     add_edge(_t, _s, INF);
     if (!solve()) return -1; // invalid flow
    int x = G[_t].back().flow;
     return G[_t].pop_back(), G[_s].pop_back(), x;
};
```

### 4.7 Gomory Hu tree\*

```
MaxFlow Dinic;
int g[MAXN];
void GomoryHu(int n) { // 0-base
  fill_n(g, n, 0);
  for (int i = 1; i < n; ++i) {
    Dinic.reset();
    add_edge(i, g[i], Dinic.maxflow(i, g[i]));
    for (int j = i + 1; j <= n; ++j)
        if (g[j] == g[i] && ~Dinic.dis[j])
        g[j] = i;
  }
}</pre>
```

#### 4.8 Minimum Cost Circulation\*

```
struct MinCostCirculation { // 0-base
  struct Edge {
    ll from, to, cap, fcap, flow, cost, rev;
  } *past[N];
  vector<Edge> G[N];
  11 dis[N], inq[N], n;
  void BellmanFord(int s) {
    fill_n(dis, n, INF), fill_n(inq, n, 0);
    queue<int> q;
    auto relax = [&](int u, ll d, Edge *e) {
      if (dis[u] > d) {
        dis[u] = d, past[u] = e;
        if (!inq[u]) inq[u] = 1, q.push(u);
     }
    };
    relax(s, 0, 0);
    while (!q.empty()) {
      int u = q.front();
      q.pop(), inq[u] = 0;
      for (auto &e : G[u])
        if (e.cap > e.flow)
          relax(e.to, dis[u] + e.cost, &e);
    }
  void try_edge(Edge &cur) {
    if (cur.cap > cur.flow) return ++cur.cap, void();
```

```
BellmanFord(cur.to);
    if (dis[cur.from] + cur.cost < 0) {</pre>
      ++cur.flow, --G[cur.to][cur.rev].flow;
      for (int i = cur.from; past[i]; i = past[i]->from
        auto &e = *past[i];
        ++e.flow, --G[e.to][e.rev].flow;
    }
    ++cur.cap;
  void solve(int mxlg) {
    for (int b = mxlg; b >= 0; --b) {
      for (int i = 0; i < n; ++i)</pre>
        for (auto &e : G[i])
         e.cap *= 2, e.flow *= 2;
      for (int i = 0; i < n; ++i)
        for (auto &e : G[i])
          if (e.fcap >> b & 1)
            try_edge(e);
   }
  }
  void init(int _n) { n = _n;
    for (int i = 0; i < n; ++i) G[i].clear();</pre>
  void add_edge(ll a, ll b, ll cap, ll cost) {
    G[a].pb(Edge{a, b, 0, cap, 0, cost, SZ(G[b]) + (a
        == b)});
    G[b].pb(Edge{b, a, 0, 0, 0, -cost, SZ(G[a]) - 1});
} mcmf; // O(VE * ElogC)
```

#### 4.9 Flow Models

- Maximum/Minimum flow with lower bound / Circulation problem
  - 1. Construct super source  $\boldsymbol{S}$  and sink T.
  - 2. For each edge (x,y,l,u), connect  $x \to y$  with capacity u-l. 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
  - 4. If in(v)>0, connect  $S\to v$  with capacity in(v), otherwise, connect v o T with capacity -in(v).
    - To maximize, connect t o s with capacity  $\infty$  (skip this in circulation problem), and let f be the maximum flow from S to T. If  $f \neq \sum_{v \in V, in(v) > 0} in(v)$ , there's no solution. Otherwise, the maximum flow from s to t is
    - To minimize, let f be the maximum flow from S to T. Connect  $t \to s$  with capacity  $\infty$  and let the flow from S to T be f'. If  $f+f' \neq \sum_{v \in V, in(v)>0} in(v)$ , there's no solution. Otherwise, f' is the answer.
  - 5. The solution of each edge e is  $l_e+f_e$  , where  $f_e$  corresponds to the flow of edge  $\boldsymbol{e}$  on the graph.
- $\bullet$  Construct minimum vertex cover from maximum matching M on bipartite graph (X,Y)
  - 1. Redirect every edge:  $y \to x$  if  $(x,y) \in M$ ,  $x \to y$  otherwise.
  - 2. DFS from unmatched vertices in X.
  - 3.  $x \in X$  is chosen iff x is unvisited. 4.  $y \in Y$  is chosen iff y is visited.
- Minimum cost cyclic flow
  - 1. Consruct super source  ${\cal S}$  and sink  ${\cal T}$
  - 2. For each edge (x,y,c), connect  $x \to y$  with (cost,cap) = (c,1)
  - if c>0, otherwise connect  $y\to x$  with (cost, cap)=(-c,1)3. For each edge with c<0, sum these cost as K, then increase d(y) by 1, decrease d(x) by 1
    4. For each vertex v with d(v)>0, connect  $S\to v$  with
  - (cost, cap) = (0, d(v))
  - For each vertex v with d(v) < 0, connect v o T with (cost, cap) = (0, -d(v)) 6. Flow from S to T , the answer is the cost of the flow C+K
- Maximum density induced subgraph
  - 1. Binary search on answer, suppose we're checking answer  ${\cal T}$

  - 2. Construct a max flow model, let K be the sum of all weights 3. Connect source  $s \to v$ ,  $v \in G$  with capacity K 4. For each edge (u,v,w) in G, connect  $u \to v$  and  $v \to u$  with capacity w
  - For  $v \in G$ , connect it with sink  $v \to t$  with capacity  $K+2T-(\sum_{e \in E(v)} w(e))-2w(v)$
  - 6. T is a valid answer if the maximum flow f < K |V|
- Minimum weight edge cover
  - 1. For each  $v \in V$  create a copy v', and connect  $u' \to v'$  with
  - weight w(u,v). 2. Connect  $v \to v'$  with weight  $2\mu(v)$ , where  $\mu(v)$  is the cost of the cheapest edge incident to  $\boldsymbol{v}$
  - 3. Find the minimum weight perfect matching on  $G^{\prime}$ .
- Project selection problem

- 1. If  $p_v>0$ , create edge (s,v) with capacity  $p_v$ ; otherwise, create edge (v,t) with capacity  $-p_v$
- 2. Create edge (u,v) with capacity w with w being the cost of choosing u without choosing v.
- 3. The mincut is equivalent to the maximum profit of a subset of projects.
- Dual of minimum cost maximum flow
  - 1. Capacity  $c_{uv}$  , Flow  $f_{uv}$  , Cost  $w_{uv}$  , Required Flow difference for vertex  $b_u$  .
  - 2. If all  $w_{uv}$  are integers, then optimal solution can happen when all  $p_u$  are integers.

$$\begin{split} \min \sum_{uv} w_{uv} f_{uv} \\ -f_{uv} \geq -c_{uv} &\Leftrightarrow \min \sum_{u} b_{u} p_{u} + \sum_{uv} c_{uv} \max(0, p_{v} - p_{u} - w_{uv}) \\ \sum_{v} f_{vu} - \sum_{v} f_{uv} = -b_{u} \end{split}$$

# String

#### 5.1 **KMP**

```
int F[MAXN];
vector<int> match(string A, string B) {
  vector<int> ans;
  F[0] = -1, F[1] = 0;
  for (int i = 1, j = 0; i < SZ(B); F[++i] = ++j) {
    if (B[i] == B[j]) F[i] = F[j]; // optimize
    while (j != -1 \&\& B[i] != B[j]) j = F[j];
  for (int i = 0, j = 0; i < SZ(A); ++i) {
    while (j != -1 && A[i] != B[j]) j = F[j];
if (++j == SZ(B)) ans.pb(i + 1 - j), j = F[j];
  return ans;
```

### 5.2 Z-value\*

```
int z[MAXn];
void make_z(const string &s) {
  int 1 = 0, r = 0;
  for (int i = 1; i < SZ(s); ++i) {</pre>
    for (z[i] = max(0, min(r - i + 1, z[i - 1]));
         i + z[i] < SZ(s) && s[i + z[i]] == s[z[i]];
         ++z[i])
    if (i + z[i] - 1 > r) l = i, r = i + z[i] - 1;
```

#### 5.3 Manacher\*

```
int z[MAXN]; // 0-base
/* center i: radius z[i * 2 + 1] / 2
   center i, i + 1: radius z[i * 2 + 2] / 2
   both aba, abba have radius 2 */
void Manacher(string tmp) {
  string s = "%";
  int l = 0, r = 0;
  for (char c : tmp) s.pb(c), s.pb('%');
  for (int i = 0; i < SZ(s); ++i) {
  z[i] = r > i ? min(z[2 * 1 - i], r - i) : 1;
    while (i - z[i] >= 0 \&\& i + z[i] < SZ(s)
            && s[i + z[i]] == s[i - z[i]]) ++z[i];
    if (z[i] + i > r) r = z[i] + i, l = i;
  }
```

#### 5.4 SAIS\*

```
namespace sfx {
bool _t[N * 2];
int SA[N * 2], H[N], RA[N];
int _s[N * 2], _c[N * 2], x[N], _p[N], _q[N * 2];
// zero based, string content MUST > 0
// SA[i]: SA[i]-th suffix is the i-th lexigraphically
      smallest suffix.
// H[i]: longest common prefix of suffix SA[i] and
     suffix SA[i - 1].
```

```
void pre(int *sa, int *c, int n, int z)
{ fill_n(sa, n, 0), copy_n(c, z, x); }
void induce(int *sa, int *c, int *s, bool *t, int n,
     int z) {
   copy_n(c, z - 1, x + 1);
  for (int i = 0; i < n; ++i)
  if (sa[i] && !t[sa[i] - 1])</pre>
       sa[x[s[sa[i] - 1]]++] = sa[i] - 1;
  copy_n(c, z, x);
  for (int i = n - 1; i >= 0; --i)
     if (sa[i] && t[sa[i] - 1])
       sa[--x[s[sa[i] - 1]]] = sa[i] - 1;
void sais(int *s, int *sa, int *p, int *q, bool *t, int
      *c, int n, int z) {
  bool uniq = t[n - 1] = true;
  int nn = 0, nmxz = -1, *nsa = sa + n, *ns = s + n,
    last = -1;
  fill_n(c, z, 0);
  for (int i = 0; i < n; ++i) uniq &= ++c[s[i]] < 2;</pre>
   partial_sum(c, c + z, c);
   if (uniq) {
     for (int i = 0; i < n; ++i) sa[--c[s[i]]] = i;</pre>
     return;
   for (int i = n - 2; i >= 0; --i)
     t[i] = (s[i] == s[i + 1] ? t[i + 1] : s[i] < s[i +
         1]);
  pre(sa, c, n, z);
   for (int i = 1; i <= n - 1; ++i)</pre>
     if (t[i] && !t[i - 1])
       sa[--x[s[i]]] = p[q[i] = nn++] = i;
  induce(sa, c, s, t, n, z);
for (int i = 0; i < n; ++i)
     if (sa[i] && t[sa[i]] && !t[sa[i] - 1]) {
       bool neq = last < 0 \mid | !equal(s + sa[i], s + p[q[
            sa[i]] + 1], s + last);
       ns[q[last = sa[i]]] = nmxz += neq;
   sais(ns, nsa, p + nn, q + n, t + n, c + z, nn, nmxz +
        1);
   pre(sa, c, n, z);
   for (int i = nn - 1; i >= 0; --i)
    sa[--x[s[p[nsa[i]]]]] = p[nsa[i]];
  induce(sa, c, s, t, n, z);
void mkhei(int n) {
  for (int i = 0, j = 0; i < n; ++i) {
     if (RA[i])
     for (; _s[i + j] == _s[SA[RA[i] - 1] + j]; ++j);
H[RA[i]] = j, j = max(0, j - 1);
  }
void build(int *s, int n) {
  copy_n(s, n, _s), _s[n] = 0;
sais(_s, SA, _p, _q, _t, _c, n + 1, 256);
copy_n(SA + 1, n, SA);
  for (int i = 0; i < n; ++i) RA[SA[i]] = i;</pre>
  mkhei(n);
| } }
```

#### 5.5 Aho-Corasick Automatan

```
const int len = 400000, sigma = 26;
struct AC_Automatan {
   int nx[len][sigma], fl[len], cnt[len], pri[len], top;
   int newnode() {
      fill(nx[top], nx[top] + sigma, -1);
      return top++;
   }
   void init() { top = 1, newnode(); }
   int input(
      string &s) { // return the end_node of string
      int X = 1;
      for (char c : s) {
        if (!~nx[X][c - 'a']) nx[X][c - 'a'] = newnode();
        X = nx[X][c - 'a'];
   }
   return X;
}
   void make_fl() {
      queue<int> q;
```

```
q.push(1), fl[1] = 0;
    for (int t = 0; !q.empty();) {
      int R = q.front();
      q.pop(), pri[t++] = R;
      for (int i = 0; i < sigma; ++i)</pre>
        if (~nx[R][i]) {
           int X = nx[R][i], Z = f1[R];
           for (; Z && !~nx[Z][i];) Z = f1[Z];
          fl[X] = Z ? nx[Z][i] : 1, q.push(X);
    }
  }
  void get_v(string &s) {
    int X = 1;
    fill(cnt, cnt + top, 0);
    for (char c : s) {
      while (X && !\sim nx[X][c - 'a']) X = fl[X];
      X = X ? nx[X][c - 'a'] : 1, ++cnt[X];
    for (int i = top - 2; i > 0; --i)
      cnt[fl[pri[i]]] += cnt[pri[i]];
};
```

#### 5.6 Smallest Rotation

```
string mcp(string s) {
  int n = SZ(s), i = 0, j = 1;
  s += s;
  while (i < n && j < n) {
    int k = 0;
    while (k < n && s[i + k] == s[j + k]) ++k;
    if (s[i + k] <= s[j + k]) j += k + 1;
    else i += k + 1;
    if (i == j) ++j;
  }
  int ans = i < n ? i : j;
  return s.substr(ans, n);
}</pre>
```

#### 5.7 De Bruijn sequence\*

```
constexpr int MAXC = 10, MAXN = 1e5 + 10;
struct DBSeq {
  int C, N, K, L, buf[MAXC * MAXN]; // K <= C^N</pre>
  void dfs(int *out, int t, int p, int &ptr) {
     if (ptr >= L) return;
     if (t > N) {
       if (N % p) return;
       for (int i = 1; i <= p && ptr < L; ++i)</pre>
         out[ptr++] = buf[i];
     } else ·
       buf[t] = buf[t - p], dfs(out, t + 1, p, ptr);
       for (int j = buf[t - p] + 1; j < C; ++j)</pre>
         buf[t] = j, dfs(out, t + 1, t, ptr);
    }
  }
  void solve(int _c, int _n, int _k, int *out) {
     int p = 0;
     C = _c, N = _n, K = _k, L = N + K - 1;
dfs(out, 1, 1, p);
     if (p < L) fill(out + p, out + L, 0);</pre>
} dbs;
```

#### 5.8 Extended SAM\*

```
struct exSAM {
  int len[N * 2], link[N * 2]; // maxlength, suflink
  int next[N * 2][CNUM], tot; // [0, tot), root = 0
  int lenSorted[N * 2]; // topo. order
  int cnt[N * 2]; // occurence
  int newnode() {
    fill_n(next[tot], CNUM, 0);
    len[tot] = cnt[tot] = link[tot] = 0;
    return tot++;
}

void init() { tot = 0, newnode(), link[0] = -1; }
  int insertSAM(int last, int c) {
    int cur = next[last][c];
    len[cur] = len[last] + 1;
```

```
int p = link[last];
    while (p != -1 && !next[p][c])
      next[p][c] = cur, p = link[p];
    if (p == -1) return link[cur] = 0, cur;
    int q = next[p][c];
    if (len[p] + 1 == len[q]) return link[cur] = q, cur
    int clone = newnode();
    for (int i = 0; i < CNUM; ++i)</pre>
      next[clone][i] = len[next[q][i]] ? next[q][i] :
    len[clone] = len[p] + 1;
    while (p != -1 && next[p][c] == q)
      next[p][c] = clone, p = link[p];
    link[link[cur] = clone] = link[q];
    link[q] = clone;
    return cur:
  void insert(const string &s) {
    int cur = 0;
    for (auto ch : s) {
      int &nxt = next[cur][int(ch - 'a')];
      if (!nxt) nxt = newnode();
      cnt[cur = nxt] += 1;
    }
  void build() {
    queue<int> q;
    q.push(0);
    while (!q.empty()) {
      int cur = q.front();
      q.pop();
      for (int i = 0; i < CNUM; ++i)</pre>
        if (next[cur][i])
          q.push(insertSAM(cur, i));
    vector<int> lc(tot);
    for (int i = 1; i < tot; ++i) ++lc[len[i]];</pre>
    partial_sum(ALL(lc), lc.begin());
    for (int i = 1; i < tot; ++i) lenSorted[--lc[len[i</pre>
        ]]] = i;
  void solve() {
    for (int i = tot - 2; i >= 0; --i)
      cnt[link[lenSorted[i]]] += cnt[lenSorted[i]];
};
```

### 5.9 PalTree\*

```
struct palindromic_tree {
  struct node {
    int next[26], fail, len;
    int cnt, num; // cnt: appear times, num: number of
                   // pal. suf.
    node(int 1 = 0) : fail(0), len(1), cnt(0), num(0) {
      for (int i = 0; i < 26; ++i) next[i] = 0;</pre>
   }
  };
  vector<node> St;
  vector<char> s;
  int last, n;
  palindromic_tree() : St(2), last(1), n(0) {
    St[0].fail = 1, St[1].len = -1, s.pb(-1);
  inline void clear() {
    St.clear(), s.clear(), last = 1, n = 0;
    St.pb(0), St.pb(-1);
    St[0].fail = 1, s.pb(-1);
  inline int get_fail(int x) {
    while (s[n - St[x].len - 1] != s[n])
     x = St[x].fail;
    return x;
 inline void add(int c) {
  s.push_back(c -= 'a'), ++n;
    int cur = get_fail(last);
    if (!St[cur].next[c]) {
      int now = SZ(St);
      St.pb(St[cur].len + 2);
      St[now].fail =
```

```
St[get_fail(St[cur].fail)].next[c];
St[cur].next[c] = now;
St[now].num = St[St[now].fail].num + 1;
}
last = St[cur].next[c], ++St[last].cnt;
}
inline void count() { // counting cnt
auto i = St.rbegin();
for (; i != St.rend(); ++i) {
St[i->fail].cnt += i->cnt;
}
}
inline int size() { // The number of diff. pal.
return SZ(St) - 2;
}
};
```

### 6 Math

### 6.1 ax+by=gcd(only exgcd \*)

```
pll exgcd(ll a, ll b) {
   if (b == 0) return pll(1, 0);
   ll p = a / b;
   pll q = exgcd(b, a % b);
   return pll(q.Y, q.X - q.Y * p);
}
/* ax+by=res, let x be minimum non-negative
g, p = gcd(a, b), exgcd(a, b) * res / g
if p.X < 0: t = (abs(p.X) + b / g - 1) / (b / g)
else: t = -(p.X / (b / g))
p += (b / g, -a / g) * t */</pre>
```

#### 6.2 Floor and Ceil

```
int floor(int a, int b)
{ return a / b - (a % b && (a < 0) ^ (b < 0)); }
int ceil(int a, int b)
{ return a / b + (a % b && (a < 0) ^ (b > 0)); }
```

#### 6.3 Floor Enumeration

```
// enumerating x = floor(n / i), [l, r]
for (int l = 1, r; l <= n; l = r + 1) {
  int x = n / l;
  r = n / x;
}</pre>
```

#### 6.4 Mod Min

```
// min{k | L <= ((ak) mod m) <= r}, no solution -> -1
11 mod_min(11 a, 11 m, 11 1, 11 r) {
   if (a == 0) return 1 ? -1 : 0;
   if (11 k = (1 + a - 1) / a; k * a <= r)
      return k;
   11 b = m / a, c = m % a;
   if (11 y = mod_min(c, a, a - r % a, a - 1 % a))
      return (1 + y * c + a - 1) / a + y * b;
   return -1;
}</pre>
```

#### 6.5 Gaussian integer gcd

#### 6.6 Miller Rabin\*

#### 6.7 Simultaneous Equations

```
struct matrix { //m variables, n equations
  int n, m;
  fraction M[MAXN][MAXN + 1], sol[MAXN];
  int solve() { //-1: inconsistent, >= 0: rank
    for (int i = 0; i < n; ++i) {</pre>
      int piv = 0;
       while (piv < m && !M[i][piv].n) ++piv;</pre>
       if (piv == m) continue;
      for (int j = 0; j < n; ++j) {
         if (i == j) continue;
         fraction tmp = -M[j][piv] / M[i][piv];
         for (int k = 0; k \leftarrow m; ++k) M[j][k] = tmp * M[
             i][k] + M[j][k];
      }
    int rank = 0;
    for (int i = 0; i < n; ++i) {</pre>
      int piv = 0;
      while (piv < m && !M[i][piv].n) ++piv;</pre>
      if (piv == m && M[i][m].n) return -1;
       else if (piv < m) ++rank, sol[piv] = M[i][m] / M[</pre>
           i][piv];
    return rank;
  }
};
```

#### 6.8 Pollard Rho\*

```
map<ll, int> cnt;
void PollardRho(ll n) {
  if (n == 1) return;
  if (prime(n)) return ++cnt[n], void();
  if (n % 2 == 0) return PollardRho(n / 2), ++cnt[2],
      void();
  11 x = 2, y = 2, d = 1, p = 1;
  #define f(x, n, p) ((mul(x, x, n) + p) % n) while (true) {
    if (d != n && d != 1) {
      PollardRho(n / d);
      PollardRho(d);
      return;
    if (d == n) ++p;
    x = f(x, n, p), y = f(f(y, n, p), n, p);
    d = gcd(abs(x - y), n);
}
```

### 6.9 Simplex Algorithm

```
const int MAXN = 11000, MAXM = 405;
const double eps = 1E-10;
double a[MAXN][MAXM], b[MAXN], c[MAXM];
double d[MAXN][MAXM], x[MAXM];
int ix[MAXN + MAXM]; // !!! array all indexed from 0
// max{cx} subject to {Ax<=b,x>=0}
// n: constraints, m: vars !!!
// x[] is the optimal solution vector
```

```
// usage :
// value = simplex(a, b, c, N, M);
double simplex(int n, int m){
  ++m:
  fill_n(d[n], m + 1, 0);
  fill_n(d[n + 1], m + 1, 0);
  iota(ix, ix + n + m, \theta);
  int r = n, s = m - 1;
  for (int i = 0; i < n; ++i) {</pre>
    for (int j = 0; j < m - 1; ++j) d[i][j] = -a[i][j];</pre>
    d[i][m - 1] = 1;
    d[i][m] = b[i];
    if (d[r][m] > d[i][m]) r = i;
  copy_n(c, m - 1, d[n]);
  d[n + 1][m - 1] = -1;
  for (double dd;; ) {
    if (r < n) {
       swap(ix[s], ix[r + m]);
      d[r][s] = 1.0 / d[r][s];
      for (int j = 0; j <= m; ++j)
  if (j != s) d[r][j] *= -d[r][s];</pre>
      for (int i = 0; i <= n + 1; ++i) if (i != r) {</pre>
         for (int j = 0; j <= m; ++j) if (j != s)</pre>
           d[i][j] += d[r][j] * d[i][s];
         d[i][s] *= d[r][s];
    }
    r = s = -1;
    for (int j = 0; j < m; ++j)
  if (s < 0 || ix[s] > ix[j]) {
         if (d[n + 1][j] > eps ||
             (d[n + 1][j] > -eps && d[n][j] > eps))
    if (s < 0) break;</pre>
    for (int i = 0; i < n; ++i) if (d[i][s] < -eps) {</pre>
      if (r < 0 ||
           (dd = d[r][m] / d[r][s] - d[i][m] / d[i][s])
                < -eps ||
           (dd < eps && ix[r + m] > ix[i + m]))
         r = i;
    if (r < 0) return -1; // not bounded
  if (d[n + 1][m] < -eps) return -1; // not executable</pre>
  double ans = 0:
  fill_n(x, m, 0);
  for (int i = m; i < n + m; ++i) { // the missing</pre>
       enumerated x[i] = 0
    if (ix[i] < m - 1){</pre>
      ans += d[i - m][m] * c[ix[i]];
       x[ix[i]] = d[i-m][m];
  }
  return ans;
```

### 6.9.1 Construction

Standard form: maximize  $\mathbf{c}^T\mathbf{x}$  subject to  $A\mathbf{x} \leq \mathbf{b}$  and  $\mathbf{x} \geq 0$ . Dual LP: minimize  $\mathbf{b}^T\mathbf{y}$  subject to  $A^T\mathbf{y} \geq \mathbf{c}$  and  $\mathbf{y} \geq 0$ .  $\bar{\mathbf{x}}$  and  $\bar{\mathbf{y}}$  are optimal if and only if for all  $i \in [1,n]$ , either  $\bar{x}_i = 0$  or  $\sum_{j=1}^m A_{ji}\bar{y}_j = c_i$  holds and for all  $i \in [1,m]$  either  $\bar{y}_i = 0$  or  $\sum_{j=1}^n A_{ij}\bar{x}_j = b_j$  holds.

```
1. In case of minimization, let c_i'=-c_i   
2. \sum_{1\leq i\leq n}A_{ji}x_i\geq b_j\to \sum_{1\leq i\leq n}-A_{ji}x_i\leq -b_j   
3. \sum_{1\leq i\leq n}A_{ji}x_i=b_j
```

- $\sum_{1 \le i \le n} A_{ji} x_i \le b_j$   $\sum_{1 \le i \le n} A_{ji} x_i \ge b_j$
- 4. If  $x_i$  has no lower bound, replace  $x_i$  with  $x_i x_i^\prime$

#### 6.10 chineseRemainder

```
11 solve(11 x1, 11 m1, 11 x2, 11 m2) {
    11 g = gcd(m1, m2);
    if ((x2 - x1) % g) return -1; // no sol
    m1 /= g; m2 /= g;
    p11 p = exgcd(m1, m2);
```

```
11 lcm = m1 * m2 * g;
11 res = p.first * (x2 - x1) * m1 + x1;
// be careful with overflow
return (res % lcm + lcm) % lcm;
}
```

#### 6.11 Factorial without prime factor\*

```
// O(p^k + Log^2 n), pk = p^k
ll prod[MAXP];
ll fac_no_p(ll n, ll p, ll pk) {
  prod[0] = 1;
  for (int i = 1; i <= pk; ++i)
    if (i % p) prod[i] = prod[i - 1] * i % pk;
    else prod[i] = prod[i - 1];
ll rt = 1;
for (; n; n /= p) {
    rt = rt * mpow(prod[pk], n / pk, pk) % pk;
    rt = rt * prod[n % pk] % pk;
}
return rt;
} // (n! without factor p) % p^k</pre>
```

### 6.12 QuadraticResidue\*

```
int Jacobi(int a, int m) {
 int s = 1;
  for (; m > 1; ) {
    a %= m;
   if (a == 0) return 0;
    const int r = __builtin_ctz(a);
    if ((r \& 1) \&\& ((m + 2) \& 4)) s = -s;
    a >>= r;
    if (a \& m \& 2) s = -s;
   swap(a, m);
 return s:
int QuadraticResidue(int a, int p) {
 if (p == 2) return a & 1;
  const int jc = Jacobi(a, p);
  if (jc == 0) return 0;
 if (jc == -1) return -1;
 int b, d;
 for (; ; ) {
   b = rand() % p;
   d = (1LL * b * b + p - a) \% p;
    if (Jacobi(d, p) == -1) break;
 int f0 = b, f1 = 1, g0 = 1, g1 = 0, tmp;
  for (int e = (1LL + p) >> 1; e; e >>= 1) {
    if (e & 1) {
      tmp = (1LL * g0 * f0 + 1LL * d * (1LL * g1 * f1 %
           p)) % p;
      g1 = (1LL * g0 * f1 + 1LL * g1 * f0) % p;
     g0 = tmp;
   tmp = (1LL * f0 * f0 + 1LL * d * (1LL * f1 * f1 % p)
       )) % p;
   f1 = (2LL * f0 * f1) % p;
    f0 = tmp;
  return g0;
```

### 6.13 PiCount\*

```
ll PrimeCount(ll n) { // n ~ 10^13 => < 2s
    if (n <= 1) return 0;
    int v = sqrt(n), s = (v + 1) / 2, pc = 0;
    vector<int> smalls(v + 1), skip(v + 1), roughs(s);
    vector<ll> larges(s);
    for (int i = 2; i <= v; ++i) smalls[i] = (i + 1) / 2;
    for (int i = 0; i < s; ++i) {
        roughs[i] = 2 * i + 1;
        larges[i] = (n / (2 * i + 1) + 1) / 2;
    }
    for (int p = 3; p <= v; ++p) {
        if (smalls[p] > smalls[p - 1]) {
            int q = p * p;
        }
}
```

```
++pc;
      if (1LL * q * q > n) break;
      skip[p] = 1;
      for (int i = q; i <= v; i += 2 * p) skip[i] = 1;</pre>
      int ns = 0;
      for (int k = 0; k < s; ++k) {
        int i = roughs[k];
        if (skip[i]) continue;
        ll d = 1LL * i * p;
        larges[ns] = larges[k] - (d \leftarrow v ? larges[
             smalls[d] - pc] : smalls[n / d]) + pc;
        roughs[ns++] = i;
      s = ns;
      for (int j = v / p; j >= p; --j) {
        int c = smalls[j] - pc, e = min(j * p + p, v + p
             1);
        for (int i = j * p; i < e; ++i) smalls[i] -= c;</pre>
      }
    }
  for (int k = 1; k < s; ++k) {
    const 11 m = n / roughs[k];
    ll t = larges[k] - (pc + k - 1);
    for (int 1 = 1; 1 < k; ++1) {
      int p = roughs[1];
      if (1LL * p * p > m) break;
      t = smalls[m / p] - (pc + 1 - 1);
    larges[0] -= t;
  return larges[0];
}
```

### 6.14 Discrete Log\*

```
int DiscreteLog(int s, int x, int y, int m) {
   constexpr int kStep = 32000;
   unordered_map<int, int> p;
   int b = 1;
   for (int i = 0; i < kStep; ++i) {</pre>
    p[y] = i;
y = 1LL * y * x % m;
     b = 1LL * b * x % m;
   for (int i = 0; i < m + 10; i += kStep) {</pre>
     s = 1LL * s * b % m;
     if (p.find(s) != p.end()) return i + kStep - p[s];
   return -1:
 int DiscreteLog(int x, int y, int m) {
  if (m == 1) return 0;
   int s = 1;
   for (int i = 0; i < 100; ++i) {</pre>
     if (s == y) return i;
     s = 1LL * s * x % m;
   if (s == y) return 100;
   int p = 100 + DiscreteLog(s, x, y, m);
   if (fpow(x, p, m) != y) return -1;
   return p;
```

#### 6.15 Berlekamp Massey

```
template <typename T>
vector<T> BerlekampMassey(const vector<T> &output) {
  vector<T> d(SZ(output) + 1), me, he;
  for (int f = 0, i = 1; i <= SZ(output); ++i) {
    for (int j = 0; j < SZ(me); ++j)
        d[i] += output[i - j - 2] * me[j];
    if ((d[i] -= output[i - 1]) == 0) continue;
    if (me.empty()) {
        me.resize(f = i);
        continue;
    }
    vector<T> o(i - f - 1);
    T k = -d[i] / d[f]; o.pb(-k);
    for (T x : he) o.pb(x * k);
        o.resize(max(SZ(o), SZ(me)));
```

```
for (int j = 0; j < SZ(me); ++j) o[j] += me[j];
if (i - f + SZ(he) >= SZ(me)) he = me, f = i;
}
return me;
```

#### 6.16 Primes

```
/* 12721 13331 14341 75577 123457 222557 556679 999983
    1097774749 1076767633 100102021 999997771
    1001010013 1000512343 987654361 999991231 999888733
     98789101 987777733 999991921 1010101333 1010102101
     100000000039 100000000000037 2305843009213693951
     4611686018427387847 9223372036854775783
    18446744073709551557 */
```

#### 6.17 Theorem

• Cramer's rule

$$\begin{array}{l} ax+by=e & x=\frac{ed-bf}{ad-bc} \\ cx+dy=f \Rightarrow & y=\frac{af-ec}{ad-bc} \end{array}$$

• Vandermonde's Identity

$$C(n+m,k) = \sum_{i=0}^{k} C(n,i)C(m,k-i)$$

· Kirchhoff's Theorem

Denote L be a  $n\times n$  matrix as the Laplacian matrix of graph G , where  $L_{ii}=d(i)$  ,  $L_{ij}=-c$  where c is the number of edge (i,j) in

- The number of undirected spanning in G is  $|\det(\tilde{L}_{11})|$ . The number of directed spanning tree rooted at r in G is  $|\det(\tilde{L}_{rr})|$ .
- Tutte's Matrix

Let D be a n imes n matrix, where  $d_{ij} = x_{ij}$  ( $x_{ij}$  is chosen uniformly at random) if i < j and  $(i,j) \in E$ , otherwise  $d_{ij} = -d_{ji}$ .  $rac{rank(D)}{2}$ is the maximum matching on G.

- Cayley's Formula
  - Given a degree sequence  $d_1,d_2,\ldots,d_n$  for each labeled vertices, there are  $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_p-1)!}$  spanning trees. Let  $T_{n,k}$  be the number of labeled forests on n vertices with
  - k components, such that vertex  $1,2,\ldots,k$  belong to different components. Then  $T_{n,k}=kn^{n-k-1}$  .
- Erdős-Gallai theorem

A sequence of nonnegative integers  $d_1 \geq \cdots \geq d_n$  can be represented as the degree sequence of a finite simple graph on  $\boldsymbol{n}$  vertices if and only if  $d_1+\cdots+d_n$  is even and  $\sum_{i=1}^k d_i \leq k(k-1)+\sum_{i=k+1}^n \min(d_i,k)$ holds for every  $1 \le k \le n$ .

• Gale-Ryser theorem

A pair of sequences of nonnegative integers  $a_1 \geq \cdots \geq a_n$  and  $b_1,\dots,b_n$  is bigraphic if and only if  $\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$  and  $\sum_{i=1}^k a_i \le$  $\sum_{i=1}^n \min(b_i,k)$  holds for every  $1 \leq k \leq n$  .

• Fulkerson-Chen-Anstee theorem

A sequence  $(a_1,b_1),\ldots,(a_n,b_n)$  of nonnegative integer pairs with A sequence  $(a_1, a_1), \dots, (a_n, a_n)$  .  $a_1 \geq \dots \geq a_n \text{ is digraphic if and only if } \sum_{i=1}^n a_i = \sum_{i=1}^n b_i \text{ and } a_i = \sum_{i=1}^n b_i$  $\sum_{i=1}^{^{\kappa}}a_i\leq \sum_{i=1}^{^{\kappa}}\min(b_i,k-1)+\sum_{i=k+1}^{n}\min(b_i,k) \text{ holds for every } 1\leq k\leq n\text{.}$ 

• Möbius inversion formula

- 
$$f(n) = \sum_{d|n} g(d) \Leftrightarrow g(n) = \sum_{d|n} \mu(d) f(\frac{n}{d})$$
  
-  $f(n) = \sum_{n|d} g(d) \Leftrightarrow g(n) = \sum_{n|d} \mu(\frac{d}{n}) f(d)$ 

- Spherical cap
  - A portion of a sphere cut off by a plane. r: sphere radius, a: radius of the base of the cap, h: height of the cap,  $\theta$ :  $\arcsin(a/r)$ .

- Volume  $=\pi h^2(3r-h)/3=\pi h(3a^2+h^2)/6=\pi r^3(2+\cos\theta)(1-\cos\theta)$
- $\cos \theta)^2/3$ . Area  $= 2\pi r h = \pi(a^2 + h^2) = 2\pi r^2(1 \cos \theta)$ .
- Lagrange multiplier

  - Optimize  $f(x_1,\ldots,x_n)$  when k constraints  $g_i(x_1,\ldots,x_n)=0$ . Lagrangian function  $\mathcal{L}(x_1,\ldots,x_n,\lambda_1,\ldots,\lambda_k)=f(x_1,\ldots,x_n)=\sum_{i=1}^k\lambda_ig_i(x_1,\ldots,x_n)$ . The solution corresponding to the original constrained optimization is always a saddle point of the Lagrangian function.

#### 6.18 Estimation

- Estimation
  - The number of divisors of n is at most around  $100\ \mathrm{for}$  $n\,<\,5e4$  ,  $\,500$  for  $n\,<\,1e7$  ,  $\,2000$  for  $n\,<\,1e10$  ,  $\,200000$  for n < 1e19.
  - n < 1e19.

    The number of ways of writing n as a sum of positive integers, disregarding the order of the summands. 1,1,2,3,5,7,11,15,22,30 for  $n=0\sim 9$ , 627 for n=20,  $\sim 2e5$  for n=50,  $\sim 2e8$  for n=100.

    Total number of partitions of n distinct elements: B(n)=1
  - 1, 1, 2, 5, 15, 52, 203, 877, 4140, 21147, 115975, 678570, 4213597,27644437, 190899322, . . . .

### 6.19 Euclidean Algorithms

- $m = \lfloor \frac{an+b}{a} \rfloor$
- Time complexity:  $O(\log n)$

$$\begin{split} f(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)}{2} + \lfloor \frac{b}{c} \rfloor \cdot (n+1) \\ +f(a \text{ mod } c,b \text{ mod } c,c,n), & a \geq c \vee b \geq c \\ 0, & n < 0 \vee a = 0 \\ nm - f(c,c-b-1,a,m-1), & \text{otherwise} \end{cases} \end{split}$$

$$\begin{split} g(a,b,c,n) &= \sum_{i=0}^n i \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor \cdot \frac{n(n+1)}{2} \\ +g(a \bmod c, b \bmod c, c, n), & a \geq c \vee b \geq c \\ 0, & n < 0 \vee a = 0 \end{cases} \\ &= \frac{\frac{1}{2} \cdot (n(n+1)m - f(c, c-b-1, a, m-1)}{-h(c, c-b-1, a, m-1)),} & \text{otherwise} \end{split}$$

$$\begin{split} h(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor^2 \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor^2 \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor^2 \cdot (n+1) \\ + \lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1) \\ + h(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{a}{c} \rfloor \cdot g(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{b}{c} \rfloor \cdot f(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ nm(m+1) - 2g(c, c-b-1, a, m-1) \\ - 2f(c, c-b-1, a, m-1) - f(a, b, c, n), & \text{otherwise} \end{cases} \end{split}$$

#### 6.20 General Purpose Numbers

• Bernoulli numbers

$$\begin{split} B_0 - 1, B_1^{\pm} &= \pm \frac{1}{2}, B_2 = \frac{1}{6}, B_3 = 0 \\ \sum_{j=0}^m {m+1 \choose j} B_j &= 0 \text{, EGF is } B(x) = \frac{x}{e^x - 1} = \sum_{n=0}^\infty B_n \frac{x^n}{n!} \,. \\ S_m(n) &= \sum_{k=1}^n k^m = \frac{1}{m+1} \sum_{k=0}^m {m+1 \choose k} B_k^+ n^{m+1-k} \end{split}$$

ullet Stirling numbers of the second kind Partitions of n distinct elements into exactly k groups.

$$\begin{split} S(n,k) &= S(n-1,k-1) + kS(n-1,k), S(n,1) = S(n,n) = 1 \\ S(n,k) &= \frac{1}{k!} \sum_{i=0}^k (-1)^{k-i} {k \choose i} i^n \\ x^n &= \sum_{i=0}^n S(n,i)(x)_i \end{split}$$

• Pentagonal number theorem

$$\prod_{n=1}^{\infty} (1 - x^n) = 1 + \sum_{k=1}^{\infty} (-1)^k \left( x^{k(3k+1)/2} + x^{k(3k-1)/2} \right)$$

• Catalan numbers

$$C_n^{(k)} = \frac{1}{(k-1)n+1} {kn \choose n}$$
$$C^{(k)}(x) = 1 + x[C^{(k)}(x)]^k$$

```
• Eulerian numbers  \begin{array}{l} \text{Number of permutations } \pi \in S_n \text{ in which exactly } k \text{ elements are greater than the previous element. } k j \text{:s s.t. } \pi(j) > \pi(j+1), \\ k+1 j \text{:s s.t. } \pi(j) \geq j, \ k j \text{:s s.t. } \pi(j) > j. \\ E(n,k) = (n-k)E(n-1,k-1) + (k+1)E(n-1,k) \\ E(n,0) = E(n,n-1) = 1 \\ E(n,k) = \sum_{j=0}^k (-1)^j \binom{n+1}{j} (k+1-j)^n \end{array}
```

### 6.21 Tips for Generating Functions

```
• Ordinary Generating Function A(x) = \sum_{i \geq 0} a_i x^i
```

```
\begin{array}{l} - \ A(rx) \Rightarrow r^n a_n \\ - \ A(x) + B(x) \Rightarrow a_n + b_n \\ - \ A(x) B(x) \Rightarrow \sum_{i=0}^n a_i b_{n-i} \\ - \ A(x)^k \Rightarrow \sum_{i_1 + i_2 + \dots + i_k = n} a_{i_1} a_{i_2} \dots a_{i_k} \\ - \ x A(x)' \Rightarrow n a_n \\ - \ \frac{A(x)}{1 - x} \Rightarrow \sum_{i=0}^n a_i \end{array}
```

• Exponential Generating Function  $A(x) = \sum_{i > 0} \frac{a_i}{i!} x_i$ 

```
- A(x) + B(x) \Rightarrow a_n + b_n

- A^{(k)}(x) \Rightarrow a_{n+k_n}

- A(x)B(x) \Rightarrow \sum_{i=0}^{n} \binom{n}{i} a_i b_{n-i}

- A(x)^k \Rightarrow \sum_{i_1+i_2+\dots+i_k=n}^{n} \binom{n}{i_1,i_2,\dots,i_k} a_{i_1} a_{i_2} \dots a_{i_k}

- xA(x) \Rightarrow na_n
```

• Special Generating Function

```
- (1+x)^n = \sum_{i\geq 0} \binom{n}{i} x^i
- \frac{1}{(1-x)^n} = \sum_{i\geq 0} \binom{n}{n-1} x^i
```

# 7 Polynomial

//(2^16)+1, 65537, 3

#### 7.1 Fast Fourier Transform

```
template<int MAXN>
struct FFT {
   using val_t = complex<double>;
   const double PI = acos(-1);
   val_t w[MAXN];
   FFT() {
      for (int i = 0; i < MAXN; ++i) {
           double arg = 2 * PI * i / MAXN;
           w[i] = val_t(cos(arg), sin(arg));
      }
   }
   void bitrev(val_t *a, int n); // see NTT
   void trans(val_t *a, int n, bool inv = false); // see NTT;
   // remember to replace LL with val_t
};</pre>
```

#### 7.2 Number Theory Transform\*

```
//7*17*(2^23)+1, 998244353, 3
//1255*(2^20)+1, 1315962881, 3
//51*(2^25)+1, 1711276033, 29
template<int MAXN, 11 P, 11 RT> //MAXN must be 2^k
struct NTT {
  11 w[MAXN];
 11 mpow(ll a, ll n);
 11 minv(ll a) { return mpow(a, P - 2); }
 NTT() {
   11 dw = mpow(RT, (P - 1) / MAXN);
    w[0] = 1;
    for (int i = 1; i < MAXN; ++i) w[i] = w[i - 1] * dw
         % P;
 void bitrev(ll *a, int n) {
    int i = 0;
    for (int j = 1; j < n - 1; ++j) {
      for (int k = n >> 1; (i ^= k) \langle k; k >>= 1);
      if (j < i) swap(a[i], a[j]);</pre>
   }
  void operator()(ll *a, int n, bool inv = false) { //0
       <= a[i] < P
    bitrev(a, n);
    for (int L = 2; L <= n; L <<= 1) {</pre>
      int dx = MAXN / L, dl = L >> 1;
```

#### 7.3 Fast Walsh Transform\*

```
/* x: a[j], y: a[j + (L >> 1)]
or: (y += x * op), and: (x += y * op)
xor: (x, y = (x + y) * op, (x - y) * op)
invop: or, and, xor = -1, -1, 1/2 */
void fwt(int *a, int n, int op) { //or
  for (int L = 2; L <= n; L <<= 1)
    for (int i = 0; i < n; i += L)</pre>
      for (int j = i; j < i + (L >> 1); ++j)
        a[j + (L >> 1)] += a[j] * op;
const int N = 21;
int f[N][1 << N], g[N][1 << N], h[N][1 << N], ct[1 << N</pre>
    1;
void subset_convolution(int *a, int *b, int *c, int L)
  // c_k = \sum_{i | j = k, i & j = 0} a_i * b_j
  int n = 1 << L;
  for (int i = 1; i < n; ++i)</pre>
   ct[i] = ct[i & (i - 1)] + 1;
  for (int i = 0; i < n; ++i)</pre>
    f[ct[i]][i] = a[i], g[ct[i]][i] = b[i];
  for (int i = 0; i <= L; ++i)</pre>
    fwt(f[i], n, 1), fwt(g[i], n, 1);
  for (int i = 0; i <= L; ++i)</pre>
    for (int j = 0; j <= i; ++j)</pre>
      for (int x = 0; x < n; ++x)
        h[i][x] += f[j][x] * g[i - j][x];
  for (int i = 0; i <= L; ++i)</pre>
    fwt(h[i], n, -1);
  for (int i = 0; i < n; ++i)</pre>
    c[i] = h[ct[i]][i];
```

#### 7.4 Polynomial Operation

```
#define fi(s, n) for (int i = (int)(s); i < (int)(n);
    ++i)
template<int MAXN, 11 P, 11 RT> // MAXN = 2^k
struct Poly : vector<ll> { // coefficients in [0, P)
  using vector<11>::vector;
  static NTT<MAXN, P, RT> ntt;
  int n() const { return (int)size(); } // n() >= 1
  Poly(const Poly &p, int m) : vector<ll>(m) {
    copy_n(p.data(), min(p.n(), m), data());
  Poly& irev() { return reverse(data(), data() + n()),
      *this; }
  Poly& isz(int m) { return resize(m), *this; }
  Poly& iadd(const Poly &rhs) { // n() == rhs.n()
    fi(0, n()) if (((*this)[i] += rhs[i]) >= P) (*this)
    [i] -= P;
return *this;
  Poly& imul(ll k) {
    fi(0, n()) (*this)[i] = (*this)[i] * k % P;
    return *this;
  Poly Mul(const Poly &rhs) const {
    int m = 1;
    while (m < n() + rhs.n() - 1) m <<= 1;</pre>
```

```
Poly X(*this, m), Y(rhs, m);
  ntt(X.data(), m), ntt(Y.data(), m);
  fi(0, m) X[i] = X[i] * Y[i] % P;
  ntt(X.data(), m, true);
  return X.isz(n() + rhs.n() - 1);
Poly Inv() const { // (*this)[0] != 0, 1e5/95ms
  if (n() == 1) return {ntt.minv((*this)[0])};
  int m = 1:
  while (m < n() * 2) m <<= 1;</pre>
  Poly Xi = Poly(*this, (n() + 1) / 2).Inv().isz(m);
  Poly Y(*this, m);
  ntt(Xi.data(), m), ntt(Y.data(), m);
  fi(0, m) {
    Xi[i] *= (2 - Xi[i] * Y[i]) % P;
    if ((Xi[i] %= P) < 0) Xi[i] += P;</pre>
  ntt(Xi.data(), m, true);
  return Xi.isz(n());
Poly Sqrt() const { // Jacobi((*this)[0], P) = 1, 1e5
  if (n() == 1) return {QuadraticResidue((*this)[0],
      P)};
  Poly X = Poly(*this, (n() + 1) / 2).Sqrt().isz(n())
  return X.iadd(Mul(X.Inv()).isz(n())).imul(P / 2 +
      1);
}
pair<Poly, Poly> DivMod(const Poly &rhs) const { // (
   rhs.)back() != 0
  if (n() < rhs.n()) return {{0}, *this};</pre>
  const int m = n() - rhs.n() + 1;
  Poly X(rhs); X.irev().isz(m);
  Poly Y(*this); Y.irev().isz(m);
  Poly Q = Y.Mul(X.Inv()).isz(m).irev();
 X = rhs.Mul(Q), Y = *this;
  fi(0, n()) if ((Y[i] -= X[i]) < 0) Y[i] += P;
  return {Q, Y.isz(max(1, rhs.n() - 1))};
Poly Dx() const {
 Poly ret(n() - 1);
  fi(0, ret.n()) ret[i] = (i + 1) * (*this)[i + 1] %
  return ret.isz(max(1, ret.n()));
Poly Sx() const {
 Poly ret(n() + 1);
  fi(0, n()) ret[i + 1] = ntt.minv(i + 1) * (*this)[i
      ] % P;
  return ret;
Poly
     _tmul(int nn, const Poly &rhs) const {
  Poly Y = Mul(rhs).isz(n() + nn - 1);
  return Poly(Y.data() + n() - 1, Y.data() + Y.n());
vector<ll> _eval(const vector<ll> &x, const vector<</pre>
    Poly> &up) const {
  const int m = (int)x.size();
  if (!m) return {};
  vector<Poly> down(m * 2);
  // down[1] = DivMod(up[1]).second;
 // fi(2, m * 2) down[i] = down[i / 2].DivMod(up[i])
      .second;
  down[1] = Poly(up[1]).irev().isz(n()).Inv().irev().
      _tmul(m, *this);
  fi(2, m * 2) down[i] = up[i ^ 1]._tmul(up[i].n() -
      1, down[i / 2]);
  vector<ll> y(m);
  fi(0, m) y[i] = down[m + i][0];
  return y;
static vector<Poly> _tree1(const vector<ll> &x) {
 const int m = (int)x.size();
vector<Poly> up(m * 2);
  fi(0, m) up[m + i] = \{(x[i] ? P - x[i] : 0), 1\};
  for (int i = m - 1; i > 0; --i) up[i] = up[i * 2].
      Mul(up[i * 2 + 1]);
  return up;
}
vector<ll> Eval(const vector<ll> &x) const { // 1e5,
    15
```

```
auto up = _tree1(x); return _eval(x, up);
  static Poly Interpolate(const vector<11> &x, const
      vector<ll> &y) { // 1e5, 1.4s
    const int m = (int)x.size();
    vector<Poly> up = _tree1(x), down(m * 2);
    vector<11> z = up[1].Dx()._eval(x, up);
    fi(0, m) z[i] = y[i] * ntt.minv(z[i]) % P;
    fi(0, m) down[m + i] = {z[i]};
for (int i = m - 1; i > 0; --i) down[i] = down[i *
         2].Mul(up[i * 2 + 1]).iadd(down[i * 2 + 1].Mul(
        up[i * 2]));
    return down[1];
  Poly Ln() const { // (*this)[0] == 1, 1e5/170ms
    return Dx().Mul(Inv()).Sx().isz(n());
  Poly Exp() const { // (*this)[0] == 0, 1e5/360ms
    if (n() == 1) return {1};
    Poly X = Poly(*this, (n() + 1) / 2).Exp().isz(n());
    Poly Y = X.Ln(); Y[0] = P - 1;
    fi(0, n()) if ((Y[i] = (*this)[i] - Y[i]) < 0) Y[i]
         += P:
    return X.Mul(Y).isz(n());
  ^{\prime}// M := P(P - 1). If k >= M, k := k % M + M.
  Poly Pow(ll k) const {
    int nz = 0;
    while (nz < n() && !(*this)[nz]) ++nz;</pre>
    if (nz * min(k, (ll)n()) >= n()) return Poly(n());
if (!k) return Poly(Poly {1}, n());
    Poly X(data() + nz, data() + nz + n() - nz * k);
    const ll c = ntt.mpow(X[0], k % (P - 1));
    return X.Ln().imul(k % P).Exp().imul(c).irev().isz(
        n()).irev();
  static ll LinearRecursion(const vector<ll> &a, const
      vector<11> &coef, 11 n) { // a_n = \sum_{j=1}^{n} a_{j}(n-1)
      i)
    const int k = (int)a.size();
    assert((int)coef.size() == k + 1);
    Poly C(k + 1), W(Poly \{1\}, k), M = \{0, 1\};
    fi(1, k + 1) C[k - i] = coef[i] ? P - coef[i] : 0;
    C[k] = 1;
    while (n)
      if (n % 2) W = W.Mul(M).DivMod(C).second;
      n /= 2, M = M.Mul(M).DivMod(C).second;
    11 ret = 0;
    fi(0, k) ret = (ret + W[i] * a[i]) % P;
    return ret;
  }
};
#undef fi
using Poly_t = Poly<131072 * 2, 998244353, 3>;
template<> decltype(Poly_t::ntt) Poly_t::ntt = {};
7.5 Value Polynomial
```

```
struct Poly {
  mint base; // f(x) = poly[x - base]
  vector<mint> poly;
  Poly(mint b = 0, mint x = 0): base(b), poly(1, x) {}
  mint get_val(const mint &x) {
    if (x >= base \&\& x < base + SZ(poly))
      return poly[x - base];
    mint rt = 0;
    vector<mint> lmul(SZ(poly), 1), rmul(SZ(poly), 1);
    for (int i = 1; i < SZ(poly); ++i)</pre>
      lmul[i] = lmul[i - 1] * (x - (base + i - 1));
    for (int i = SZ(poly) - 2; i >= 0; --i)
      rmul[i] = rmul[i + 1] * (x - (base + i + 1));
    for (int i = 0; i < SZ(poly); ++i)
  rt += poly[i] * ifac[i] * inegfac[SZ(poly) - 1 -</pre>
          i] * lmul[i] * rmul[i];
    return rt;
  void raise() { // g(x) = sigma\{base:x\} f(x)
    if (SZ(poly) == 1 && poly[0] == 0)
    mint nw = get_val(base + SZ(poly));
    poly.pb(nw);
```

```
for (int i = 1; i < SZ(poly); ++i)
    poly[i] += poly[i - 1];
}
};</pre>
```

#### 7.6 Newton's Method

Given F(x) where

$$F(x) = \sum_{i=0}^{\infty} \alpha_i (x - \beta)^i$$

for  $\beta$  being some constant. Polynomial P such that F(P)=0 can be found iteratively. Denote by  $Q_k$  the polynomial such that  $F(Q_k)=0$  (mod  $x^{2^k}$ ), then

$$Q_{k+1} = Q_k - \frac{F(Q_k)}{F'(Q_k)} \pmod{x^{2^{k+1}}}$$

## 8 Geometry

#### 8.1 Default Code

```
typedef pair<double, double> pdd;
typedef pair<pdd, pdd> Line;
struct Cir{ pdd O; double R; };
const double eps = 1e-8;
pdd operator+(pdd a, pdd b)
{ return pdd(a.X + b.X, a.Y + b.Y); }
pdd operator-(pdd a, pdd b)
{ return pdd(a.X - b.X, a.Y - b.Y); }
pdd operator*(pdd a, double b)
{ return pdd(a.X * b, a.Y * b); }
pdd operator/(pdd a, double b)
{ return pdd(a.X / b, a.Y / b); }
double dot(pdd a, pdd b)
{ return a.X * b.X + a.Y * b.Y; }
double cross(pdd a, pdd b)
{ return a.X * b.Y - a.Y * b.X; }
double abs2(pdd a)
{ return dot(a, a); }
double abs(pdd a)
{ return sqrt(dot(a, a)); }
int sign(double a)
{ return fabs(a) < eps ? 0 : a > 0 ? 1 : -1; }
int ori(pdd a, pdd b, pdd c)
{ return sign(cross(b - a, c - a)); }
bool collinearity(pdd p1, pdd p2, pdd p3)
{ return sign(cross(p1 - p3, p2 - p3)) == 0; }
bool btw(pdd p1, pdd p2, pdd p3) {
  if (!collinearity(p1, p2, p3)) return 0;
  return sign(dot(p1 - p3, p2 - p3)) <= 0;</pre>
bool seg_intersect(pdd p1, pdd p2, pdd p3, pdd p4) {
  int a123 = ori(p1, p2, p3);
  int a124 = ori(p1, p2, p4);
  int a341 = ori(p3, p4, p1);
  int a342 = ori(p3, p4, p2);
  if (a123 == 0 && a124 == 0)
    return btw(p1, p2, p3) | btw(p1, p2, p4) ||
      btw(p3, p4, p1) || btw(p3, p4, p2);
  return a123 * a124 <= 0 && a341 * a342 <= 0;
pdd intersect(pdd p1, pdd p2, pdd p3, pdd p4) {
  double a123 = cross(p2 - p1, p3 - p1);
double a124 = cross(p2 - p1, p4 - p1);
  return (p4 * a123 - p3 * a124) / (a123 - a124); // C
^3 / C^2
pdd perp(pdd p1)
{ return pdd(-p1.Y, p1.X); }
pdd projection(pdd p1, pdd p2, pdd p3)
{ return p1 + (p2 - p1) * dot(p3 - p1, p2 - p1) / abs2(
    p2 - p1); }
pdd reflection(pdd p1, pdd p2, pdd p3)
{ return p3 + perp(p2 - p1) * cross(p3 - p1, p2 - p1) /
     abs2(p2 - p1) * 2; }
pdd linearTransformation(pdd p0, pdd p1, pdd q0, pdd q1
    , pdd r) {
```

### 8.2 PointSegDist\*

```
double PointSegDist(pdd q0, pdd q1, pdd p) {
  if (sign(abs(q0 - q1)) == 0) return abs(q0 - p);
  if (sign(dot(q1 - q0, p - q0)) >= 0 && sign(dot(q0 -
          q1, p - q1)) >= 0)
    return fabs(cross(q1 - q0, p - q0) / abs(q0 - q1));
  return min(abs(p - q0), abs(p - q1));
}
```

#### 8.3 Heart

```
pdd circenter(pdd p0, pdd p1, pdd p2) { // radius = abs
    (center)
  p1 = p1 - p0, p2 = p2 - p0;
  double x1 = p1.X, y1 = p1.Y, x2 = p2.X, y2 = p2.Y;
  double m = 2. * (x1 * y2 - y1 * x2);
  center.X = (x1 * x1 * y2 - x2 * x2 * y1 + y1 * y2 * (
  y1 - y2)) / m;
center.Y = (x1 * x2 * (x2 - x1) - y1 * y1 * x2 + x1 *
      y2 * y2) / m;
  return center + p0;
pdd incenter(pdd p1, pdd p2, pdd p3) { // radius = area
     / s * 2
  double a = abs(p2 - p3), b = abs(p1 - p3), c = abs(p1
      - p2);
  double s = a + b + c;
  return (a * p1 + b * p2 + c * p3) / s;
pdd masscenter(pdd p1, pdd p2, pdd p3)
{ return (p1 + p2 + p3) / 3; }
pdd orthcenter(pdd p1, pdd p2, pdd p3)
{ return masscenter(p1, p2, p3) * 3 - circenter(p1, p2,
     p3) * 2; }
```

#### 8.4 point in circle

```
// return p4 is strictly in circumcircle of tri(p1,p2,
11 sqr(l1 x) { return x * x; }
bool in_cc(const pll& p1, const pll& p2, const pll& p3,
      const pll& p4) {
  11 u11 = p1.X - p4.X; 11 u12 = p1.Y - p4.Y;
  11 u21 = p2.X - p4.X; 11 u22 = p2.Y - p4.Y;
  11 u31 = p3.X - p4.X; 11 u32 = p3.Y - p4.Y;
  11 u13 = sqr(p1.X) - sqr(p4.X) + sqr(p1.Y) - sqr(p4.Y)
  11 u23 = sqr(p2.X) - sqr(p4.X) + sqr(p2.Y) - sqr(p4.Y)
  11 u33 = sqr(p3.X) - sqr(p4.X) + sqr(p3.Y) - sqr(p4.Y)
      );
  __int128 det = (__int128)-u13 * u22 * u31 + (_
                                                     int128
       )u12 * u23 * u31 + (__int128)u13 * u21 * u32 - (__int128)u11 * u23 * u32 - (__int128)u12 * u21 *
       u33 + (__int128)u11 * u22 * u33;
  return det > eps;
}
```

#### 8.5 Convex hull\*

#### 8.6 PointInConvex\*

#### 8.7 TangentPointToHull\*

```
/* The point should be strictly out of hull
  return arbitrary point on the tangent line */
pii get_tangent(vector<pll> &C, pll p) {
  auto gao = [&](int s) {
    return cyc_tsearch(SZ(C),
    [&](int x, int y) { return ori(p, C[x], C[y]) == s;
    });
  };
  return pii(gao(1), gao(-1));
} // return (a, b), ori(p, C[a], C[b]) >= 0
```

### 8.8 minMaxEnclosingRectangle\*

```
const double INF = 1e18, qi = acos(-1) / 2 * 3;
pdd solve(vector<pll> &dots) {
#define diff(u, v) (dots[u] - dots[v])
#define vec(v) (dots[v] - dots[i])
  hull(dots);
  double Max = 0, Min = INF, deg;
  int n = SZ(dots);
  dots.pb(dots[0]);
  for (int i = 0, u = 1, r = 1, l = 1; i < n; ++i) {
    pll nw = vec(i + 1);
    while (cross(nw, vec(u + 1)) > cross(nw, vec(u)))
      u = (u + 1) \% n;
    while (dot(nw, vec(r + 1)) > dot(nw, vec(r)))
     r = (r + 1) \% n;
    if (!i) l = (r + 1) % n;
    while (dot(nw, vec(1 + 1)) < dot(nw, vec(1)))
      1 = (1 + 1) \% n;
   Min = min(Min, (double)(dot(nw, vec(r)) - dot(nw,
    vec(1))) * cross(nw, vec(u)) / abs2(nw));
    deg = acos(dot(diff(r, 1), vec(u)) / abs(diff(r, 1)
        ) / abs(vec(u)));
    deg = (qi - deg) / 2
    Max = max(Max, abs(diff(r, 1)) * abs(vec(u)) * sin(
        deg) * sin(deg));
  return pdd(Min, Max);
```

#### 8.9 VectorInPoly\*

```
// ori(a, b, c) >= 0, valid: "strict" angle from a-b to
    a-c
bool btwangle(pll a, pll b, pll c, pll p, int strict) {
    return ori(a, b, p) >= strict && ori(a, p, c) >=
        strict;
}
// whether vector{cur, p} in counter-clockwise order
    prv, cur, nxt
bool inside(pll prv, pll cur, pll nxt, pll p, int
    strict) {
    if (ori(cur, nxt, prv) >= 0)
        return btwangle(cur, nxt, prv, p, strict);
    return !btwangle(cur, prv, nxt, p, !strict);
}
```

### 8.10 PolyUnion\*

```
double rat(pll a, pll b) {
  return sign(b.X) ? (double)a.X / b.X : (double)a.Y /
      b.Y;
} // all poly. should be ccw
double polyUnion(vector<vector<pll>>> &poly) {
  double res = 0;
  for (auto &p : poly)
    for (int a = 0; a < SZ(p); ++a) {</pre>
      pll A = p[a], B = p[(a + 1) \% SZ(p)];
      vector<pair<double, int>> segs = {{0, 0}}, {1,
           ∅}};
      for (auto &q : poly) {
         if (&p == &q) continue;
         for (int b = 0; b < SZ(q); ++b) {</pre>
           pll C = q[b], D = q[(b + 1) \% SZ(q)];
           int sc = ori(A, B, C), sd = ori(A, B, D);
           if (sc != sd && min(sc, sd) < 0) {
  double sa = cross(D - C, A - C), sb = cross</pre>
                 (D - C, B - C);
             segs.emplace_back(sa / (sa - sb), sign(sc -
                  sd));
           if (!sc && !sd && &q < &p && sign(dot(B - A,
               D - C)) > 0) \{
             segs.emplace_back(rat(C - A, B - A), 1);
             segs.emplace_back(rat(D - A, B - A), -1);
        }
       sort(ALL(segs));
      for (auto &s : segs) s.X = clamp(s.X, 0.0, 1.0);
       double sum = 0;
      int cnt = segs[0].second;
       for (int j = 1; j < SZ(segs); ++j) {</pre>
         if (!cnt) sum += segs[j].X - segs[j - 1].X;
         cnt += segs[j].Y;
      res += cross(A, B) * sum;
    }
  return res / 2;
}
```

#### 8.11 Polar Angle Sort\*

#### 8.12 Half plane intersection\*

```
pll area_pair(Line a, Line b)
{ return pll(cross(a.Y - a.X, b.X - a.X), cross(a.Y - a
    .X, b.Y - a.X)); }
bool isin(Line 10, Line 11, Line 12) {
  // Check inter(l1, l2) strictly in l0
  auto [a02X, a02Y] = area_pair(10, 12);
  auto [a12X, a12Y] = area_pair(l1, l2);
  if (a12X - a12Y < 0) a12X *= -1, a12Y *= -1;</pre>
  return (__int128) a02Y * a12X - (__int128) a02X *
      a12Y > 0; // C^4
/* Having solution, check size > 2 */
/* --^-- Line.X --^-- Line.Y --^-- */
vector<Line> halfPlaneInter(vector<Line> arr) {
  sort(ALL(arr), [&](Line a, Line b) -> int {
  if (cmp(a.Y - a.X, b.Y - b.X, 0) != -1)
      return cmp(a.Y - a.X, b.Y - b.X, 0);
    return ori(a.X, a.Y, b.Y) < 0;</pre>
  });
  deque<Line> dq(1, arr[0]);
  for (auto p : arr) {
    if (cmp(dq.back().Y - dq.back().X, p.Y - p.X, 0) ==
          -1)
```

### 8.13 RotatingSweepLine

```
void rotatingSweepLine(vector<pii> &ps) {
  int n = SZ(ps), m = 0;
  vector<int> id(n), pos(n);
  vector<pii> line(n * (n - 1));
  for (int i = 0; i < n; ++i)</pre>
    for (int j = 0; j < n; ++j)</pre>
      if (i != j) line[m++] = pii(i, j);
  sort(ALL(line), [&](pii a, pii b) {
    return cmp(ps[a.Y] - ps[a.X], ps[b.Y] - ps[b.X]);
  }); // cmp(): polar angle compare
  iota(ALL(id), ∅);
  sort(ALL(id), [&](int a, int b) {
    if (ps[a].Y != ps[b].Y) return ps[a].Y < ps[b].Y;</pre>
    return ps[a] < ps[b];</pre>
  }); // initial order, since (1, 0) is the smallest
  for (int i = 0; i < n; ++i) pos[id[i]] = i;</pre>
  for (int i = 0; i < m; ++i) {</pre>
    auto l = line[i];
    // do something
    }
}
```

### 8.14 Minimum Enclosing Circle\*

```
pdd Minimum_Enclosing_Circle(vector<pdd> dots, double &
    r) {
  pdd cent;
 random_shuffle(ALL(dots));
  cent = dots[0], r = 0;
  for (int i = 1; i < SZ(dots); ++i)</pre>
   if (abs(dots[i] - cent) > r) {
      cent = dots[i], r = 0;
      for (int j = 0; j < i; ++j)</pre>
        if (abs(dots[j] - cent) > r) {
          cent = (dots[i] + dots[j]) / 2;
          r = abs(dots[i] - cent);
          for(int k = 0; k < j; ++k)
            if(abs(dots[k] - cent) > r)
              cent = excenter(dots[i], dots[j], dots[k
                   ], r);
        }
  return cent;
```

### 8.15 Intersection of two circles\*

### 8.16 Intersection of polygon and circle\*

```
// Divides into multiple triangle, and sum up
const double PI=acos(-1);
double _area(pdd pa, pdd pb, double r){
  if(abs(pa)<abs(pb)) swap(pa, pb);</pre>
  if(abs(pb)<eps) return 0;</pre>
  double S, h, theta;
  double a=abs(pb),b=abs(pa),c=abs(pb-pa);
  double cosB = dot(pb,pb-pa) / a / c, B = acos(cosB);
  double cosC = dot(pa,pb) / a / b, C = acos(cosC);
  if(a > r){
    S = (C/2)*r*r;
    h = a*b*sin(C)/c;
    if (h < r \&\& B < PI/2) S -= (acos(h/r)*r*r - h*sqrt
        (r*r-h*h));
  else if(b > r){
    theta = PI - B - asin(sin(B)/r*a);
    S = .5*a*r*sin(theta) + (C-theta)/2*r*r;
  else S = .5*sin(C)*a*b;
  return S:
double area_poly_circle(const vector<pdd> poly,const
    pdd &O, const double r){
  double S=0;
  for(int i=0;i<SZ(poly);++i)</pre>
    S+=\_area(poly[i]-0,poly[(i+1)\%SZ(poly)]-0,r)*ori(0,
        poly[i],poly[(i+1)%SZ(poly)]);
  return fabs(S);
```

#### 8.17 Intersection of line and circle\*

```
vector<pdd> circleLine(pdd c, double r, pdd a, pdd b) {
  pdd p = a + (b - a) * dot(c - a, b - a) / abs2(b - a)
  ;
  double s = cross(b - a, c - a), h2 = r * r - s * s /
      abs2(b - a);
  if (h2 < 0) return {};
  if (h2 == 0) return {p};
  pdd h = (b - a) / abs(b - a) * sqrt(h2);
  return {p - h, p + h};
}</pre>
```

### 8.18 Tangent line of two circles

```
vector<Line> go( const Cir& c1 , const Cir& c2 , int
     sign1 ){
  // sign1 = 1 for outer tang, -1 for inter tang
  vector<Line> ret;
  double d_sq = abs2(c1.0 - c2.0);
  if (sign(d_sq) == 0) return ret;
  double d = sqrt(d_sq);
  pdd v = (c2.0 - c1.0) / d;

double c = (c1.R - sign1 * c2.R) / d;
  if (c * c > 1) return ret;
  double h = sqrt(max(0.0, 1.0 - c * c));
  for (int sign2 = 1; sign2 >= -1; sign2 -= 2) {
  pdd n = pdd(v.X * c - sign2 * h * v.Y,
       v.Y * c + sign2 * h * v.X);
     pdd p1 = c1.0 + n * c1.R;
    pdd p2 = c2.0 + n * (c2.R * sign1);
    if (sign(p1.X - p2.X) == 0 and
    sign(p1.Y - p2.Y) == 0)
       p2 = p1 + perp(c2.0 - c1.0);
    ret.pb(Line(p1, p2));
  return ret;
```

#### 8.19 CircleCover\*

```
const int N = 1021;
struct CircleCover {
  int C;
  Cir c[N];
  bool g[N][N], overlap[N][N];
  // Area[i] : area covered by at least i circles
  double Area[ N ];
```

```
void init(int _C){ C = _C;}
  struct Teve {
    pdd p; double ang; int add;
    Teve() {}
    Teve(pdd _a, double _b, int _c):p(_a), ang(_b), add
         (_c){}
    bool operator<(const Teve &a)const</pre>
     {return ang < a.ang;}
  }eve[N * 2];
  // strict: x = 0, otherwise x = -1
  bool disjuct(Cir &a, Cir &b, int x)
  {return sign(abs(a.0 - b.0) - a.R - b.R) \rightarrow x;}
  bool contain(Cir &a, Cir &b, int x)
  {return sign(a.R - b.R - abs(a.0 - b.0)) > x;}
  bool contain(int i, int j) {
    /* c[j] is non-strictly in c[i]. */
    return (sign(c[i].R - c[j].R) > 0 || (sign(c[i].R -
          c[j].R) == 0 \&\& i < j)) \&\& contain(c[i], c[j],
          -1):
  void solve(){
    fill_n(Area, C + 2, 0);
    for(int i = 0; i < C; ++i)</pre>
       for(int j = 0; j < C; ++j)</pre>
         overlap[i][j] = contain(i, j);
    for(int i = 0; i < C; ++i)</pre>
       for(int j = 0; j < C; ++j)</pre>
         g[i][j] = !(overlap[i][j] || overlap[j][i] ||
             disjuct(c[i], c[j], -1));
    for(int i = 0; i < C; ++i){</pre>
      int E = 0, cnt = 1;
       for(int j = 0; j < C; ++j)</pre>
         if(j != i && overlap[j][i])
           ++cnt;
       for(int j = 0; j < C; ++j)</pre>
         if(i != j && g[i][j]) {
           pdd aa, bb;
           CCinter(c[i], c[j], aa, bb);
           double A = atan2(aa.Y - c[i].0.Y, aa.X - c[i
               ].0.X);
           double B = atan2(bb.Y - c[i].0.Y, bb.X - c[i
               ].O.X);
           eve[E++] = Teve(bb, B, 1), eve[E++] = Teve(aa)
               , A, -1);
           if(B > A) ++cnt;
      if(E == 0) Area[cnt] += pi * c[i].R * c[i].R;
         sort(eve, eve + E);
         eve[E] = eve[0];
         for(int j = 0; j < E; ++j){</pre>
           cnt += eve[j].add;
           Area[cnt] += cross(eve[j].p, eve[j + 1].p) *
           double theta = eve[j + 1].ang - eve[j].ang;
           if (theta < 0) theta += 2. * pi;</pre>
           Area[cnt] += (theta - sin(theta)) * c[i].R *
               c[i].R * .5;
      }
    }
  }
};
```

### 8.20 3Dpoint\*

```
{ return Point(p1.y * p2.z - p1.z * p2.y, p1.z * p2.x -
     p1.x * p2.z, p1.x * p2.y - p1.y * p2.x); }
double dot(const Point &p1, const Point &p2)
{ return p1.x * p2.x + p1.y * p2.y + p1.z * p2.z; }
double abs(const Point &a)
{ return sqrt(dot(a, a)); }
Point cross3(const Point &a, const Point &b, const
    Point &c)
{ return cross(b - a, c - a); }
double area(Point a, Point b, Point c)
{ return abs(cross3(a, b, c)); }
double volume(Point a, Point b, Point c, Point d)
{ return dot(cross3(a, b, c), d - a); }
Point masscenter(Point a, Point b, Point c, Point d)
{ return (a + b + c + d) / 4; }
pdd proj(Point a, Point b, Point c, Point u) {
// proj. u to the plane of a, b, and c
  Point e1 = b - a;
  Point e2 = c - a;
  e1 = e1 / abs(e1);
  e2 = e2 - e1 * dot(e2, e1);
  e2 = e2 / abs(e2);
  Point p = u - a;
  return pdd(dot(p, e1), dot(p, e2));
```

#### 8.21 Convexhull3D\*

```
struct convex hull 3D {
struct Face {
  int a, b, c;
  Face(int ta, int tb, int tc): a(ta), b(tb), c(tc) {}
}; // return the faces with pt indexes
vector<Face> res;
vector<Point> P;
convex_hull_3D(const vector<Point> &_P): res(), P(_P) {
// all points coplanar case will WA, O(n^2)
  int n = SZ(P);
  if (n <= 2) return; // be careful about edge case</pre>
  // ensure first 4 points are not coplanar
  swap(P[1], *find_if(ALL(P), [&](auto p) { return sign
       (abs2(P[0] - p)) != 0; }));
  swap(P[2], *find_if(ALL(P), [&](auto p) { return sign
      (abs2(cross3(p, P[0], P[1]))) != 0; }));
  swap(P[3], *find_if(ALL(P), [&](auto p) { return sign
      (volume(P[0], P[1], P[2], p)) != 0; }));
  vector<vector<int>> flag(n, vector<int>(n));
  res.emplace_back(0, 1, 2); res.emplace_back(2, 1, 0);
  for (int i = 3; i < n; ++i) {</pre>
    vector<Face> next;
    for (auto f : res) {
      int d = sign(volume(P[f.a], P[f.b], P[f.c], P[i])
      if (d <= 0) next.pb(f);</pre>
      int ff = (d > 0) - (d < 0);
      flag[f.a][f.b] = flag[f.b][f.c] = flag[f.c][f.a]
    for (auto f : res) {
      auto F = [\&](int x, int y) {
        if (flag[x][y] > 0 && flag[y][x] <= 0)</pre>
          next.emplace_back(x, y, i);
      F(f.a, f.b); F(f.b, f.c); F(f.c, f.a);
    }
    res = next;
  }
bool same(Face s, Face t) {
  if (sign(volume(P[s.a], P[s.b], P[s.c], P[t.a])) !=
      0) return 0;
  if (sign(volume(P[s.a], P[s.b], P[s.c], P[t.b])) !=
      0) return 0;
  if (sign(volume(P[s.a], P[s.b], P[s.c], P[t.c])) !=
      0) return 0;
  return 1;
int polygon_face_num() {
  int ans = 0;
  for (int i = 0; i < SZ(res); ++i)</pre>
    ans += none_of(res.begin(), res.begin() + i, [&](
        Face g) { return same(res[i], g); });
```

```
return ans:
double get_volume() {
   double ans = 0;
   for (auto f : res)
      ans += volume(Point(0, 0, 0), P[f.a], P[f.b], P[f.c
           ]);
   return fabs(ans / 6);
double get_dis(Point p, Face f) {
   Point p1 = P[f.a], p2 = P[f.b], p3 = P[f.c];

double a = (p2.y - p1.y) * (p3.z - p1.z) - (p2.z - p1
         .z) * (p3.y - p1.y);
   double b = (p2.z - p1.z) * (p3.x - p1.x) - (p2.x - p1
         .x) * (p3.z - p1.z);
   double c = (p2.x - p1.x) * (p3.y - p1.y) - (p2.y - p1
   .y) * (p3.x - p1.x);

double d = 0 - (a * p1.x + b * p1.y + c * p1.z);

return fabs(a * p.x + b * p.y + c * p.z + d) / sqrt(a * p.x + b * p.y + c * p.z + d)
          * a + b * b + c * c);
};
```

### 8.22 DelaunayTriangulation\*

```
/* Delaunay Triangulation:
Given a sets of points on 2D plane, find a
triangulation such that no points will strictly
inside circumcircle of any triangle.
find : return a triangle contain given point
add_point : add a point into triangulation
A Triangle is in triangulation iff. its has_chd is 0.
Region of triangle u: iterate each u.edge[i].tri,
each points are u.p[(i+1)\%3], u.p[(i+2)\%3]
Voronoi diagram: for each triangle in triangulation,
the bisector of all its edges will split the region.
nearest point will belong to the triangle containing it
const ll inf = MAXC * MAXC * 100; // Lower_bound
    unknown
struct Tri:
struct Edge {
  Tri* tri; int side;
 Edge(): tri(∅), side(∅){}
 Edge(Tri* _tri, int _side): tri(_tri), side(_side){}
struct Tri {
 pll p[3];
 Edge edge[3];
 Tri* chd[3];
 Tri() {}
 Tri(const pll& p0, const pll& p1, const pll& p2) {
   p[0] = p0; p[1] = p1; p[2] = p2;
    chd[0] = chd[1] = chd[2] = 0;
  bool has_chd() const { return chd[0] != 0; }
 int num_chd() const {
    return !!chd[0] + !!chd[1] + !!chd[2];
 bool contains(pll const& q) const {
    for (int i = 0; i < 3; ++i)
      if (ori(p[i], p[(i + 1) % 3], q) < 0)
       return 0;
    return 1:
} pool[N * 10], *tris;
void edge(Edge a, Edge b) {
 if(a.tri) a.tri->edge[a.side] = b;
  if(b.tri) b.tri->edge[b.side] = a;
struct Trig { // Triangulation
 Trig() {
   the_root = // Tri should at least contain all
        points
     new(tris++) Tri(pll(-inf, -inf), pll(inf + inf, -
          inf), pll(-inf, inf + inf));
  Tri* find(pll p) { return find(the_root, p); }
 void add_point(const pll &p) { add_point(find(
      the_root, p), p); }
  Tri* the root;
  static Tri* find(Tri* root, const pll &p) {
```

```
while (1) {
      if (!root->has_chd())
        return root;
      for (int i = 0; i < 3 && root->chd[i]; ++i)
        if (root->chd[i]->contains(p)) {
          root = root->chd[i];
          break;
    assert(0); // "point not found"
  void add_point(Tri* root, pll const& p) {
    Tri* t[3];
    /* split it into three triangles */
    for (int i = 0; i < 3; ++i)</pre>
      t[i] = new(tris++) Tri(root->p[i], root->p[(i +
          1) % 3], p);
    for (int i = 0; i < 3; ++i)
      edge(Edge(t[i], 0), Edge(t[(i + 1) % 3], 1));
    for (int i = 0; i < 3; ++i)
      edge(Edge(t[i], 2), root->edge[(i + 2) % 3]);
    for (int i = 0; i < 3; ++i)
      root->chd[i] = t[i];
    for (int i = 0; i < 3; ++i)
      flip(t[i], 2);
  void flip(Tri* tri, int pi) {
    Tri* trj = tri->edge[pi].tri;
    int pj = tri->edge[pi].side;
    if (!trj) return;
    if (!in_cc(tri->p[0], tri->p[1], tri->p[2], trj->p[
        pj])) return;
    /* flip edge between tri,trj *,
    Tri* trk = new(tris++) Tri(tri->p[(pi + 1) % 3],
        trj->p[pj], tri->p[pi]);
    Tri* trl = new(tris++) Tri(trj->p[(pj + 1) % 3],
         tri->p[pi], trj->p[pj]);
    edge(Edge(trk, 0), Edge(trl, 0));
    edge(Edge(trk, 1), tri->edge[(pi + 2) % 3]);
    edge(Edge(trk, 2), trj->edge[(pj + 1) % 3]);
    edge(Edge(trl, 1), trj->edge[(pj + 2) % 3]);
    edge(Edge(trl, 2), tri->edge[(pi + 1) % 3]);
    tri->chd[0] = trk; tri->chd[1] = trl; tri->chd[2] =
         0:
    trj->chd[0] = trk; trj->chd[1] = trl; trj->chd[2] =
    flip(trk, 1); flip(trk, 2);
    flip(trl, 1); flip(trl, 2);
  }
};
vector<Tri*> triang; // vector of all triangle
set<Tri*> vst;
void go(Tri* now) { // store all tri into triang
  if (vst.find(now) != vst.end())
    return:
  vst.insert(now);
  if (!now->has chd())
    return triang.pb(now);
  for (int i = 0; i < now->num_chd(); ++i)
    go(now->chd[i]);
void build(int n, pll* ps) { // build triangulation
  tris = pool; triang.clear(); vst.clear();
  random_shuffle(ps, ps + n);
  Trig tri; // the triangulation structure
  for (int i = 0; i < n; ++i)</pre>
    tri.add_point(ps[i]);
  go(tri.the_root);
}
8.23 Triangulation Vonoroi*
```

```
vector<Line> ls[N];
pll arr[N];
Line make_line(pdd p, Line 1) {
  pdd d = 1.Y - 1.X; d = perp(d);
  pdd m = (1.X + 1.Y) / 2;
  l = Line(m, m + d);
  if (ori(1.X, 1.Y, p) < 0)</pre>
    l = Line(m + d, m);
  return 1;
}
```

```
double calc_area(int id) {
  // use to calculate the area of point "strictly in
      the convex hull"
  vector<Line> hpi = halfPlaneInter(ls[id]);
  vector<pdd> ps;
  for (int i = 0; i < SZ(hpi); ++i)</pre>
    ps.pb(intersect(hpi[i].X, \, hpi[i].Y, \, hpi[(i\,+\,1)\,\,\%
        SZ(hpi)].X, hpi[(i + 1) % SZ(hpi)].Y));
  double rt = 0;
  for (int i = 0; i < SZ(ps); ++i)</pre>
    rt += cross(ps[i], ps[(i + 1) % SZ(ps)]);
  return fabs(rt) / 2;
void solve(int n, pii *oarr) {
  map<pll, int> mp;
  for (int i = 0; i < n; ++i)</pre>
    arr[i] = pll(oarr[i].X, oarr[i].Y), mp[arr[i]] = i;
  build(n, arr); // Triangulation
  for (auto *t : triang) {
    vector<int> p;
    for (int i = 0; i < 3; ++i)
      if (mp.find(t->p[i]) != mp.end())
        p.pb(mp[t->p[i]]);
    for (int i = 0; i < SZ(p); ++i)</pre>
      for (int j = i + 1; j < SZ(p); ++j) {</pre>
        Line l(oarr[p[i]], oarr[p[j]]);
        ls[p[i]].pb(make_line(oarr[p[i]], 1));
        ls[p[j]].pb(make_line(oarr[p[j]], 1));
  }
}
```

### 8.24 Minkowski Sum\*

#### 9 Else

### 9.1 Cyclic Ternary Search\*

```
/* bool pred(int a, int b);
f(0) ~ f(n - 1) is a cyclic-shift U-function
return idx s.t. pred(x, idx) is false forall x*/
int cyc_tsearch(int n, auto pred) {
   if (n == 1) return 0;
   int l = 0, r = n; bool rv = pred(1, 0);
   while (r - 1 > 1) {
      int m = (1 + r) / 2;
      if (pred(0, m) ? rv: pred(m, (m + 1) % n)) r = m;
      else l = m;
   }
   return pred(1, r % n) ? l : r % n;
}
```

#### 9.2 Mo's Alogrithm(With modification)

```
/*
Mo's Algorithm With modification
Block: N^{2/3}, Complexity: N^{5/3}
*/
struct Query {
  int L, R, LBid, RBid, T;
  Query(int 1, int r, int t):
   L(1), R(r), LBid(1 / blk), RBid(r / blk), T(t) {}
```

```
bool operator<(const Query &q) const {</pre>
     if (LBid != q.LBid) return LBid < q.LBid;</pre>
     if (RBid != q.RBid) return RBid < q.RBid;</pre>
     return T < b.T;</pre>
};
void solve(vector<Query> query) {
   sort(ALL(query));
   int L=0, R=0, T=-1;
   for (auto q : query) {
     while (T < q.T) addTime(L, R, ++T); // TODO</pre>
     while (T > q.T) subTime(L, R, T--); // TODO
     while (R < q.R) add(arr[++R]); // TODO</pre>
     while (L > q.L) add(arr[--L]); // TODO
while (R > q.R) sub(arr[R--]); // TODO
     while (L < q.L) sub(arr[L++]); // TODO</pre>
     // answer auerv
}
```

#### 9.3 Mo's Alogrithm On Tree

```
Mo's Algorithm On Tree
Preprocess:
1) LCA
2) dfs with in[u] = dft++, out[u] = dft++
3) ord[in[u]] = ord[out[u]] = u
4) bitset<MAXN> inset
struct Query {
  int L, R, LBid, lca;
  Query(int u, int v) {
     int c = LCA(u, v);
     if (c == u || c == v)
       q.lca = -1, q.L = out[c ^ u ^ v], q.R = out[c];
     else if (out[u] < in[v])</pre>
       q.lca = c, q.L = out[u], q.R = in[v];
     else
       q.lca = c, q.L = out[v], q.R = in[u];
     q.Lid = q.L / blk;
  bool operator<(const Query &q) const {</pre>
     if (LBid != q.LBid) return LBid < q.LBid;</pre>
     return R < q.R;
  }
void flip(int x) {
    if (inset[x]) sub(arr[x]); // TODO
     else add(arr[x]); // TODO
     inset[x] = ~inset[x];
void solve(vector<Query> query) {
  sort(ALL(query));
  int L = 0, R = 0;
  for (auto q : query) {
    while (R < q.R) flip(ord[++R]);</pre>
     while (L > q.L) flip(ord[--L]);
    while (R > q.R) flip(ord[R--]);
while (L < q.L) flip(ord[L++]);</pre>
     if (~q.lca) add(arr[q.lca]);
     // answer query
     if (~q.lca) sub(arr[q.lca]);
}
```

#### 9.4 Additional Mo's Algorithm Trick

• Mo's Algorithm With Addition Only

```
- Sort querys same as the normal Mo's algorithm. - For each query [l,r]: - If l/blk = r/blk, brute-force. - If l/blk \neq curL/blk, initialize curL := (l/blk + 1) \cdot blk, curR := curL - 1 - If r > curR, increase curR - decrease curL to fit l, and then undo after answering
```

• Mo's Algorithm With Offline Second Time

```
- Require: Changing answer \equiv adding f([l,r],r+1).
- Require: f([l,r],r+1) = f([l,r],r+1) - f([l,l),r+1).
- Part1: Answer all f([l,r],r+1) first.
- Part2: Store curR \to R for curL (reduce the space to O(N)), and then answer them by the second offline algorithm.
- Note: You must do the above symmetrically for the left boundaries.
```

#### 9.5 Hilbert Curve

```
ll hilbert(int n, int x, int y) {
    ll res = 0;
    for (int s = n / 2; s; s >>= 1) {
        int rx = (x & s) > 0;
        int ry = (y & s) > 0;
        res += s * 111 * s * ((3 * rx) ^ ry);
        if (ry == 0) {
            if (rx == 1) x = s - 1 - x, y = s - 1 - y;
            swap(x, y);
        }
    }
    return res;
} // n = 2^k
```

#### 9.6 DynamicConvexTrick\*

```
// only works for integer coordinates!! maintain max
struct Line {
 mutable 11 a, b, p;
  bool operator<(const Line &rhs) const { return a <</pre>
     rhs.a; }
 bool operator<(11 x) const { return p < x; }</pre>
struct DynamicHull : multiset<Line, less<>>> {
  static const ll kInf = 1e18;
 ll Div(ll a, ll b) { return a / b - ((a ^ b) < 0 && a
       % b); }
 bool isect(iterator x, iterator y) {
   if (y == end()) { x->p = kInf; return 0; }
    if (x->a == y->a) x->p = x->b > y->b ? kInf : -kInf
    else x -> p = Div(y -> b - x -> b, x -> a - y -> a);
    return x->p >= y->p;
 }
 void addline(ll a, ll b) {
    auto z = insert({a, b, 0}), y = z++, x = y;
    while (isect(y, z)) z = erase(z);
    if (x != begin() && isect(--x, y)) isect(x, y =
        erase(y));
    while ((y = x) != begin() && (--x)->p >= y->p)
        isect(x, erase(y));
 11 query(ll x) {
    auto 1 = *lower_bound(x);
    return 1.a * x + 1.b;
```

#### 9.7 All LCS\*

```
void all_lcs(string s, string t) { // 0-base
  vector<int> h(SZ(t));
  iota(ALL(h), 0);
  for (int a = 0; a < SZ(s); ++a) {
    int v = -1;
    for (int c = 0; c < SZ(t); ++c)
        if (s[a] == t[c] || h[c] < v)
            swap(h[c], v);
        // LCS(s[0, a], t[b, c]) =
        // c - b + 1 - sum([h[i] >= b] | i <= c)
        // h[i] might become -1 !!
    }
}</pre>
```

#### 9.8 DLX\*

```
up[dn[j]] = up[j], dn[up[j]] = dn[j], --s[cl[
               j]];
       } else {
        lt[rg[i]] = lt[i], rg[lt[i]] = rg[i];
    }
  void restore(int c) {
    TRAV(i, up, c) {
       if (A) {
        TRAV(j, lt, i)
          ++s[cl[j]], up[dn[j]] = j, dn[up[j]] = j;
       } else {
        lt[rg[i]] = rg[lt[i]] = i;
       }
     if (A) lt[rg[c]] = c, rg[lt[c]] = c;
  void init(int c) {
     columns = c;
     for (int i = 0; i < c; ++i) {</pre>
       up[i] = dn[i] = bt[i] = i;
       lt[i] = i == 0 ? c : i - 1;
       rg[i] = i == c - 1 ? c : i + 1;
       s[i] = 0;
     rg[c] = 0, lt[c] = c - 1;
    up[c] = dn[c] = -1;
    head = c, sz = c + 1;
  void insert(int r, const vector<int> &col) {
     if (col.empty()) return;
     int f = sz;
     for (int i = 0; i < (int)col.size(); ++i) {</pre>
       int c = col[i], v = sz++;
       dn[bt[c]] = v;
       up[v] = bt[c], bt[c] = v;
       rg[v] = (i + 1 == (int)col.size() ? f : v + 1);
       rw[v] = r, cl[v] = c;
       ++s[c];
       if (i > 0) lt[v] = v - 1;
     lt[f] = sz - 1;
  int h() {
    int ret = 0;
     memset(vis, 0, sizeof(bool) * sz);
     TRAV(x, rg, head) {
       if (vis[x]) continue;
       vis[x] = true, ++ret;
       TRAV(i, dn, x) TRAV(j, rg, i) vis[cl[j]] = true;
    }
    return ret;
  void dfs(int dep) {
    if (dep + (A ? 0 : h()) >= ans) return;
     if (rg[head] == head) return ans = dep, void();
     if (dn[rg[head]] == rg[head]) return;
     int w = rg[head];
    TRAV(x, rg, head) if (s[x] < s[w]) w = x;
     if (A) remove(w);
     TRAV(i, dn, w) {
       if (B) remove(i);
       TRAV(j, rg, i) remove(A ? cl[j] : j);
       dfs(dep + 1);
       TRAV(j, lt, i) restore(A ? cl[j] : j);
       if (B) restore(i);
     if (A) restore(w);
  int solve() {
    for (int i = 0; i < columns; ++i)</pre>
      dn[bt[i]] = i, up[i] = bt[i];
     ans = 1e9, dfs(0);
     return ans;
};
```

#### 9.9 Matroid Intersection

```
Start from S=\emptyset. In each iteration, let Y_1=\{x\not\in S\mid S\cup\{x\}\in I_1\}
```

```
• Y_2=\{x\not\in S\mid S\cup\{x\}\in I_2\} If there exists x\in Y_1\cap Y_2, insert x into S. Otherwise for each x\in S, y\not\in S, create edges • x\to y if S-\{x\}\cup\{y\}\in I_1. • y\to x if S-\{x\}\cup\{y\}\in I_2.
```

Find a shortest path (with BFS) starting from a vertex in  $Y_1$  and ending at a vertex in  $Y_2$  which doesn't pass through any other vertices in  $Y_2$ , and alternate the path. The size of S will be incremented by 1 in each iteration. For the weighted case, assign weight w(x) to vertex x if  $x \in S$  and -w(x) if  $x \not\in S$ . Find the path with the minimum number of edges among all minimum length paths and alternate it.

### 9.10 AdaptiveSimpson

```
using F_t = function<double(double)>;
pdd simpson(const F_t &f, double 1, double r,
  double f1, double fr, double fm = nan("")) {
 if (isnan(fm)) fm = f((1 + r) / 2);
return {fm, (r - 1) / 6 * (f1 + 4 * fm + fr)};
double simpson_ada(const F_t &f, double 1, double r,
  double f1, double fm, double fr, double eps) {
  double m = (1 + r) / 2,
  s = simpson(f, 1, r, fl, fr, fm).second;
auto [flm, sl] = simpson(f, 1, m, fl, fm);
  auto [fmr, sr] = simpson(f, m, r, fm, fr);
  double delta = sl + sr - s;
  if (abs(delta) <= 15 * eps)</pre>
    return sl + sr + delta / 15;
  return simpson_ada(f, 1, m, f1, f1m, fm, eps / 2) +
    simpson_ada(f, m, r, fm, fmr, fr, eps / 2);
double simpson_ada(const F_t &f, double 1, double r) {
  return simpson_ada(
    f, 1, r, f(1), f((1 + r) / 2), f(r), 1e-9 / 7122);
double simpson_ada2(const F_t &f, double 1, double r) {
    double h = (r - 1) / 7122, s = 0;
  for (int i = 0; i < 7122; ++i, l += h)</pre>
    s += simpson_ada(f, 1, 1 + h);
  return s;
```

### 9.11 Simulated Annealing

#### 9.12 Tree Hash\*

```
ull seed;
ull shift(ull x) {
    x ^= x << 13;
    x ^= x >> 7;
    x ^= x << 17;
    return x;
}
ull dfs(int u, int f) {
    ull sum = seed;
    for (int i : G[u])
        if (i != f)
            sum += shift(dfs(i, u));
    return sum;
}</pre>
```

#### 9.13 Binary Search On Fraction

```
struct Q {
    11 p, q;
    Q go(Q b, 11 d) { return {p + b.p*d, q + b.q*d}; }
};
bool pred(Q);
// returns smallest p/q in [lo, hi] such that
```

```
// pred(p/q) is true, and 0 <= p,q <= N
Q frac_bs(ll N) {
  Q lo{0, 1}, hi{1, 0};
  if (pred(lo)) return lo;
  assert(pred(hi));
  bool dir = 1, L = 1, H = 1;
  for (; L || H; dir = !dir) {
    ll len = 0, step = 1;
    for (int t = 0; t < 2 && (t ? step/=2 : step*=2);)</pre>
      if (Q mid = hi.go(lo, len + step);
          mid.p > N \mid\mid mid.q > N \mid\mid dir \land pred(mid))
        t++;
      else len += step;
    swap(lo, hi = hi.go(lo, len));
    (dir ? L : H) = !!len;
  return dir ? hi : lo;
```

## 10 Python

#### 10.1 Misc