

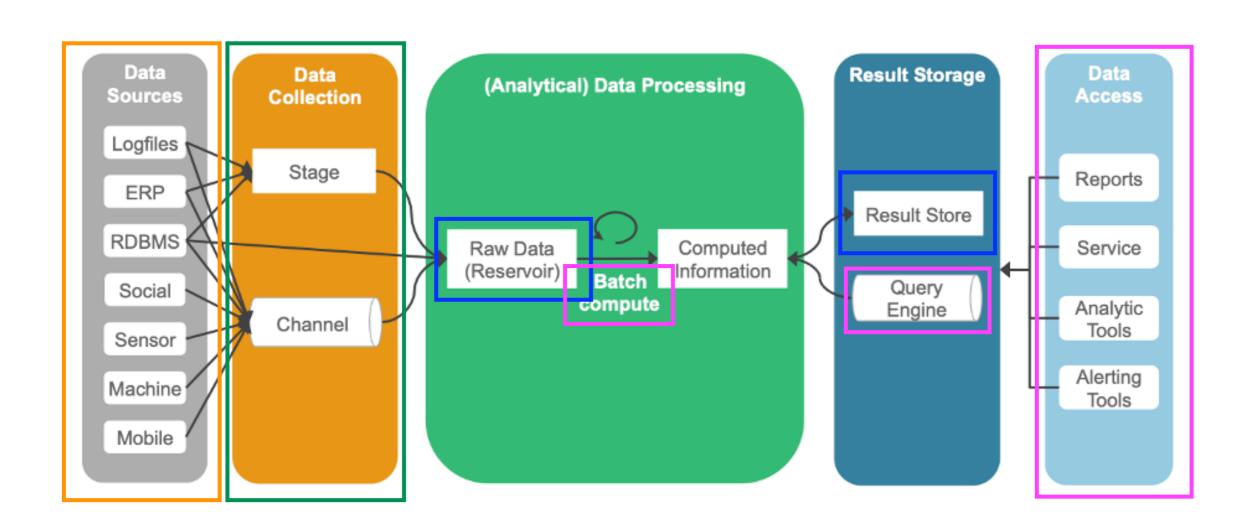
2110446 - Data Science and Data Engineering

Data Storages

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Simple Big Data Analytic Architecture



Data Generation

Data Ingestion

Data Storage

Data Analytics

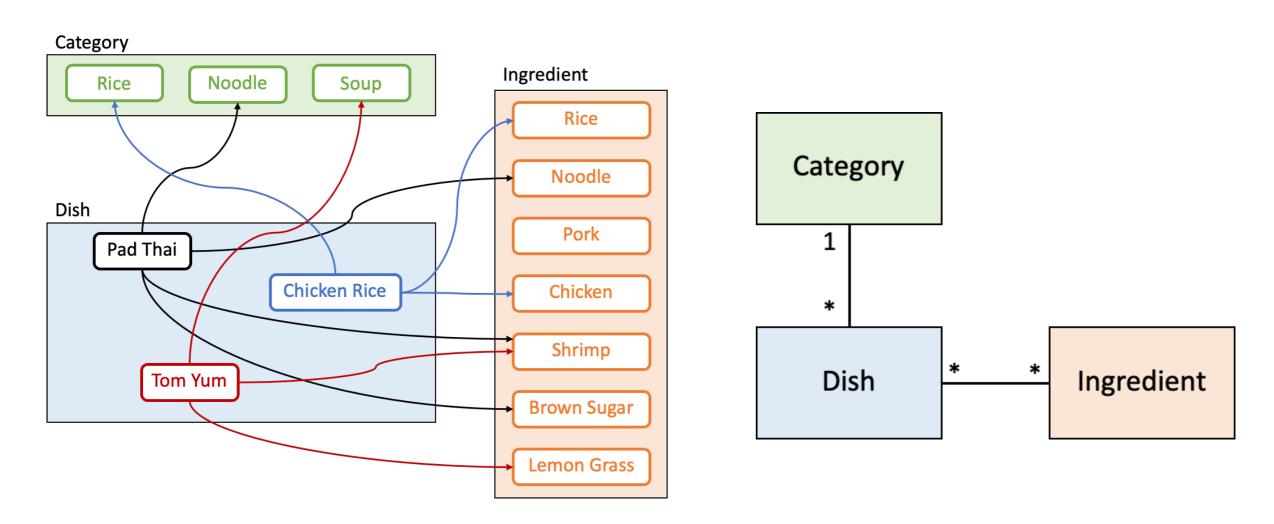
Traditional Database

- Based on "relational model"
 - Data is split and stored into tables
 - Tables can be processed together using set-like operations
 - Data model is usually normalized to remove duplication
- Very suitable for OLTP or transaction systems
 - Provide lots of complicated SQL operations
 - Lots of inserts and updates

Problems of Data Science Storage

- There are several needs for data analytics purposes e.g. traditional data store, caching, feature store
- Data is historical data and its volume can be huge
- Scalability is extremely important and "Relational + Consistency" can limit scalability
- SQL command can be very complex and time-consuming
 - It requires the synchronization of data accessing between multiple tables
 - It will be poor when using on more than a few servers in the same cluster

Relational Database: Normalized Data Model

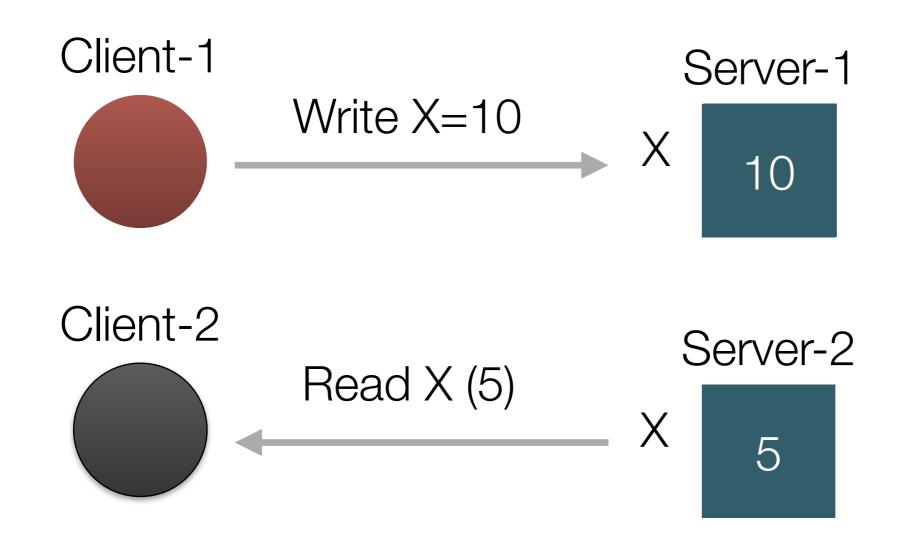


How can we split these tables to lots of machines?

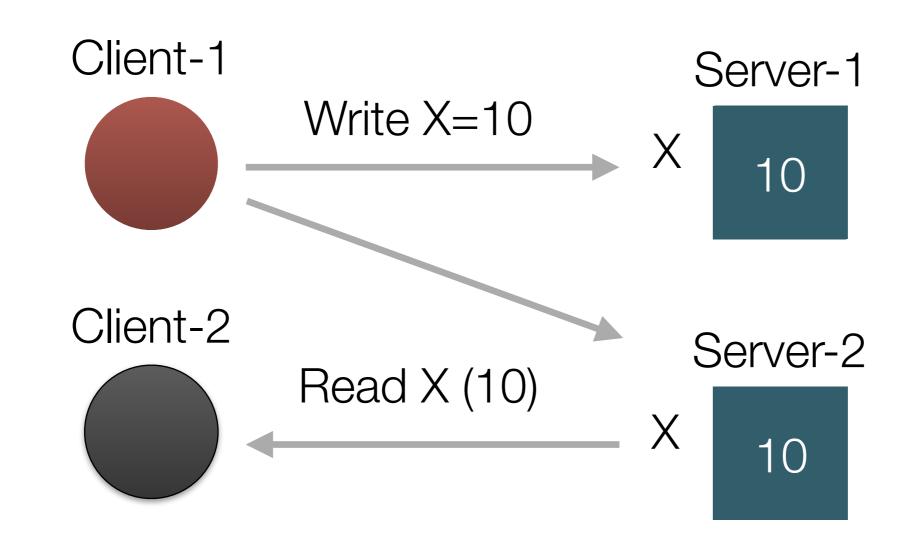
- Dishes of the same category in the same machine
- How about dishes and ingredients?

Or we can replicate data — lead to data consistency problems

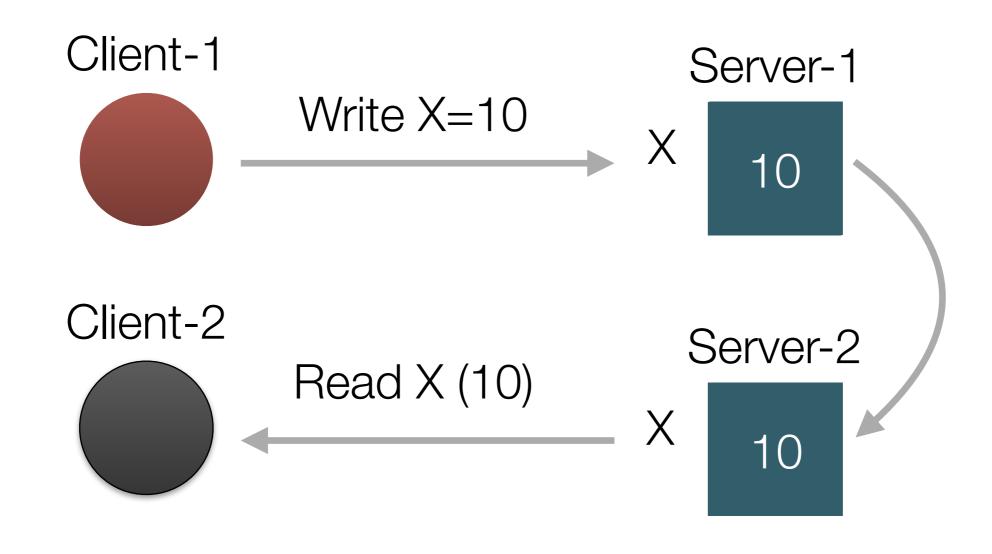
Replication and Data Consistency Problem



Solution to Data Consistency - Active Replication



Solution to Data Consistency - Passive Replication



Data Consistency Models

- · Strong consistency (ท่องตรงกันตองกอลา)
 - After update completes, any access will return the updated value
 - High costs for large scale system
- · Weak consistency (เมื่อถือเยื่อนใบหมือฐ มะตับตรยกัน)
 - Certain conditions must be met before the consistency is guaranteed
 - Time between update completion and guaranteed consistency is called inconsistency window

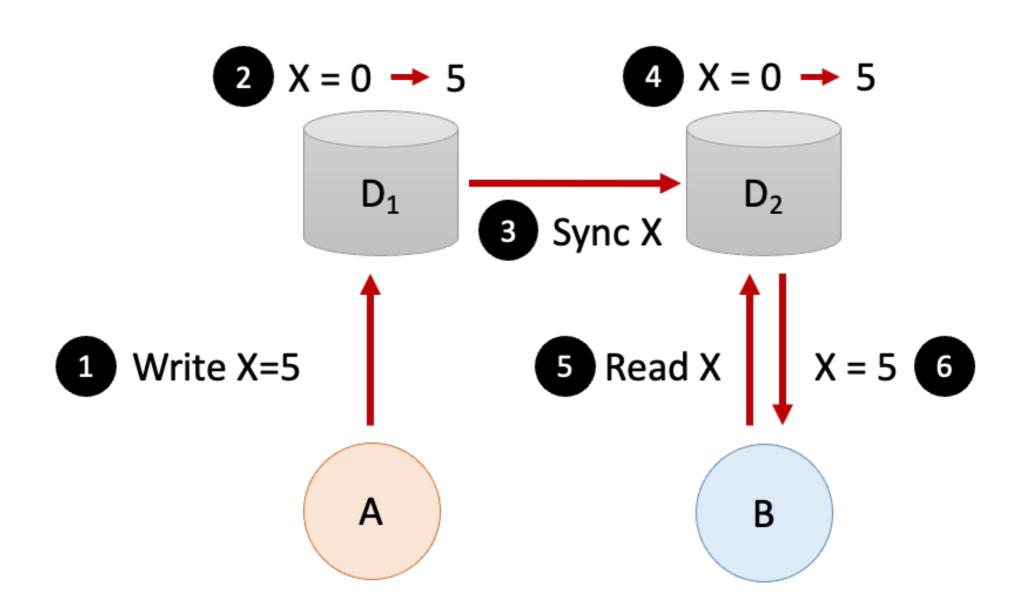
CAP Theorem (Brewer's Theorem)

- By Eric Brewer (University of California, Berkeley)
- It is impossible for a distributed computer system to simultaneously provide all three of the following guarantees: Consistency, Availability, Partition tolerance
- Scenario
 - Distributed system (clients and servers)
 - Multiple servers working together
 - Multiple clients may read or write on the same data at the same time

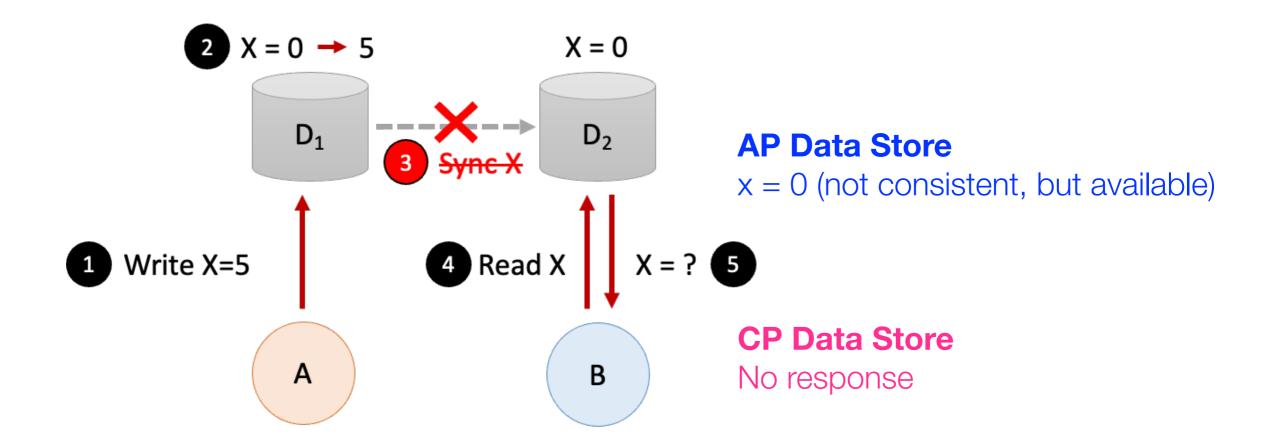
CAP Guarantees only 2 out of 3

- · Consistency (ข้อมูลเดียวกัน ห้องพรงกัน ในทุกฐที่)
 - Every read receives the most recent write or an error
- Availability (ถึงแม่สะบบ จะก่อม สะบบจะยงมีการ เทอบกัน เองใด หากมี server เนลืองย่า
 - Every request receives a (non-error) response without guarantee that it contains the most recent write
 - Response from any server is good
- - The system continues to operate despite arbitrary message loss or failure of part of the system
 - Lots of servers require long synchronization time causing servers to not be able to communicate among one another within reasonable time

CA Data Store



Given P, Choosing between C and A (at step 5)



The Landscape of NoSQL

- Alternatives to SQL database non-relational, distributed, and horizontally scalable
- Data is shared and distributed across multiple servers
- Typically use weak consistency model (but not always)
- Examples
 - Document: MongoDB, DynamoDB, CosmosDB, Couchbase, Firebase
 - Column: Cassandra, HBase, CosmosDB, Accumulo
 - Key-value: Redis, DynamoDB, CosmosDB, MemcacheDB
 - Graph: Neo4J, CosmosDB, ArrangoDB, OrientDB
 - · Search Engine: Elasticsearch, Splunk, Solr

How NoSQL can "Scale"

- Principle ideas
 - Split data into chunks or shards
 - Distribute data across multiple servers
 - Must require minimum synchronization
- Have to give up some traditional features
 - No complex relational model
 - Relax consistency
 - Duplicated information (not space optimized)
 - Fast to insert new record, but not so fast to update the existing one

CONSISTENCY

Clients will see the same data at the same time

Databases:

Relational/SQL

Achieved by:

Data in a write to one node is instantly forwarded or replicated across all other nodes

NOT POSSIBLE

Databases:

MongoDB Redis **HBase**

AVAILABILITY

Clients' data requests will always get a response even if nodes are down

Achieved by:

Replicating data across different servers

Databases:

Cassandra CouchDB

PARTITION TOLERANCE

Partition: a partial failure between two nodes

The system will continue to work despite a partition

Achieved by:

Data is sufficiently replicated across nodes and networks to maintain system during an outage

Source: https://dev.to/katkelly/cap-theorem-why-you-can-t-have-it-all-ga1

Redis Key-Value Store

Data Storage



Redis (เร็ว เหมาะกับ real time)

- Remote Dictionary Server
- In-memory data structure store with clustering, transactional, time-to-live limiting, and auto-failover capabilities
- Being used for database cache, message broker, streaming engine, feature store engine, etc.
- Support wide-range of data structure with lots of related operations for each structure
- Provide CLI and support many programming languages

Working with Redis

Redis CLI

- Standard client program to connect to any redis server
- Come with any redis installation (see: https://redis.io/docs/getting-started/)

```
redis-cli
redis-cli -h 34.143.227.66
```

You can type in redis command in the CLI input

Working with Redis

- Redis-Py
 - Standard python package for redis client

```
pip install redis
```

Example

```
import redis

r = redis.Redis(host='hostname', port=port)
```

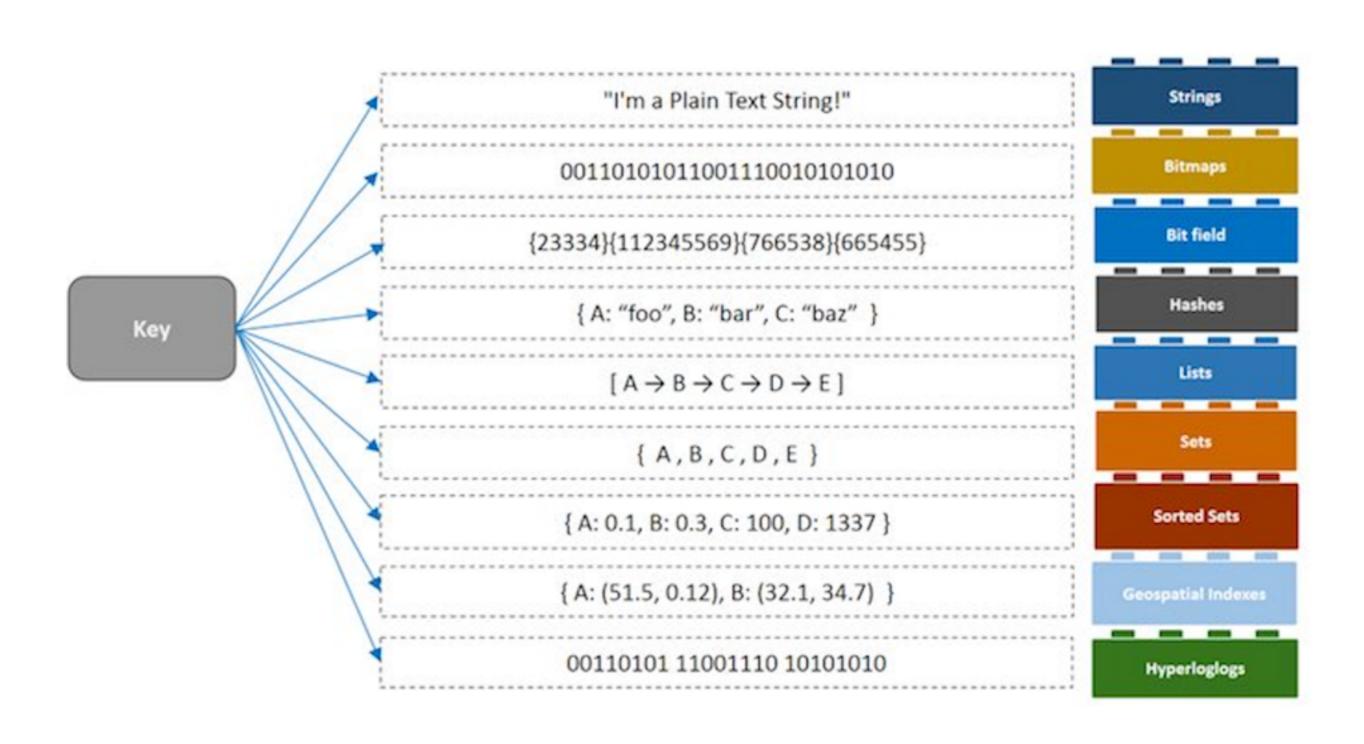
Running Redis Locally

The simplest way to run a redis instance is to use docker

```
docker pull redis
docker run -d --rm --name redis -p 6379:6379 redis
```

 This will start redis in your docker at port 6379 and map the port to your localhost

Redis Data Types



String

 Similar to Python or Java Strings, maximum length of 512MB

```
SET "user" "Natawut Nupairoj"
GET "user"
DEL "user"
```

- Use cases
 - Server-side object cache e.g. HTML fragments, shopping cart, user profile
 - Queues
 - Activity tracking

Other String Commands

APPEND INCR SET

DECR INCRBY SETEX

DECRBY INCRBYFLOAT SETNX

GET LCS SETRANGE

GETDEL MGET STRLEN

GETEX MSET SUBSTR

GETRANGE MSETNX

GETSET PSETEX

Useful Commands

Any item in Redis can be made to expire after or at a certain time

```
EXPIRE user 60. # in seconds
TTL user
```

- You can scan all index with scan command scan 0
- You can delete item or test its existence

```
DEL mykey
EXISTS mykey
```

List

- · List of strings, sorted by insertion order
- · Can be used as list, queue, stack

```
LPUSH mylist abc # mylist contains "abc"

LPUSH mylist xyz # mylist contains "xyz", "abc"

RPUSH mylist 123 # mylist contains "xyz", "abc", "123"
```

- Use cases
 - Queue
 - Timelines

List Commands

BLMOVE LMOVE LSET

BLMPOP LMPOP LTRIM

BLPOP LPOP RPOP

BRPOP LPOS RPOPLPUSH

BRPOPLPUSH LPUSH RPUSH

LINDEX LPUSHX RPUSHNX

LINSERT LRANGE

LLEN LREM

Set

- Powerful data types for unordered non-duplicated keys
- · Support many set operations e.g. intersection, union, etc.

```
SADD user_set natawut
SCARD user_set
SMEMBERS user set
```

- Use cases
 - Set of user profiles
 - Set of inappropriate words for inappropriate content filtering

Set Commands

SADD SISMEMBER SSCAN

SCARD SMEMBERS SUNION

SDIFF SMISMEMBER SUNIONSTORE

SDIFFSTORE SMOVE

SINTER SPOP

SINTERCARD SRANDMEMBER

SINTERSTORE SREM

Sorted Set

 Set of sorted items based on the score associated to each member

```
ZADD my_sortedset 5 data1
ZADD my_sortedset 1 data2 10 data3
ZRANGEBYSCORE my_sortedset 5. +inf WITHSCORES
```

- Use
 - Leader scoreboard
 - Priority queue

Sorted Set Commands

BZMPOP ZDIFFSTORE ZMSCORE

BZPOPMAX ZINCRBY ZPOPMAX

BZPOPMIN ZINTER ZPOPMIN

ZADD ZINTERCARD ZRANDMEMBER

ZCARD ZINTERSTORE ZRANGE

ZCOUNT ZRANGEBYLEX

ZDIFF ZMPOP ZRANGEBYSCORE

Hash

A container of unique fields and their values

```
HMSET profile:12345 user nnp id 12345 name "Natawut Nupairoj" balance 10
HGETALL profile:12345
HINCRBYprofile:12345 balance 5
```

- Use
 - User profile information
 - Post. Information

Hash Commands

HDEL HSTRLEN

HEXISTS HMGET HVALS

HGET HMSET

HGETALL HRANDFIELD

HINCRBY HSCAN

HINCRBYFLOAT HSET

HKEYS HSETNX

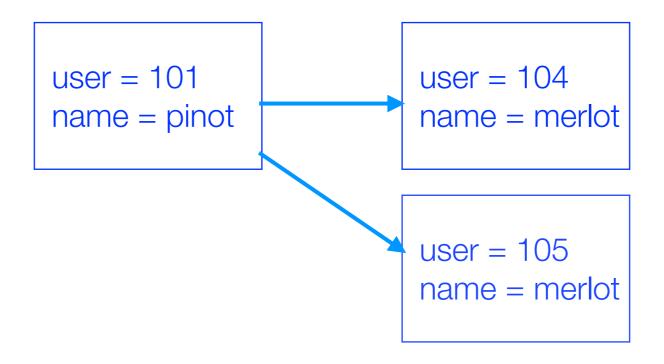
Example: Redis in Action

- Create data model for simple social network
- Data items
 - Users id, name, can follow others, can be followed
 - Posts id, content
- Relationship
 - 1 user can have many posts, each post can associate to only one user
 - User can follow one another

Users and Posts

```
post (id = 1, user=101, content="this is...")
 user = 101
 name = pinot
                            post (id = 2, user=105, content="...")
                            post (id = 8, user=101, content="...")
SET user:101:name pinot
                           # user id 101 with name "pinot"
                             # refer back to user-id
SET username:pinot 101
HMSET post: 1 user 101 content "this is the first post"
RPUSH user:101:post 1
RPUSH user:101:post 8
```

Users and Followers



SADD user:101:follows 104

SADD user:101:follows 105

SADD user:104:followed by 101

SADD user:105:followed by 101

Simple Redis-Py Example

This notebook contains simple redis python commands.

```
In [1]: import redis
         Connect to local server -- no hostname or ip is needed
In [2]: rd = redis.Redis(charset="utf-8", decode_responses=True)
In [3]: rd.set('user:101:name', 'pinot')
Out[3]: True
In [4]: rd.get('user:101:name')
Out[4]: 'pinot'
In [5]: rd.hset('post:1', 'user', 101)
    rd.hset('post:1', 'content', 'this is the first post')
Out[5]: 1
In [6]: rd.hgetall('post:1')
Out[6]: {'user': '101', 'content': 'this is the first post'}
In [7]: rd.rpush('user:101:post', 1)
         rd.rpush('user:101:post', 8)
Out[7]: 2
In [8]: rd.llen('user:101:post')
Out[8]: 2
```

```
In [9]: rd.lrange('user:101:post', 0, -1)
 Out[9]: ['1', '8']
In [10]: rd.sadd('user:101:follows', 104)
         rd.sadd('user:101:follows', 105)
Out[10]: 1
In [11]: rd.scard('user:101:follows')
Out[11]: 2
In [12]: rd.smembers('user:101:follows')
Out[12]: {'104', '105'}
In [13]: cursor = 0
         cursor, keys = rd.scan(cursor=cursor, match='user:*')
         while cursor > 0:
             for key in keys:
                 print('found: ', key)
             cursor, keys = rd.scan(cursor=cursor, match='username:*')
         for key in keys:
             print('found: ', key)
         found: user:101:follows
         found: user:101:name
         found: user:101:post
```

Additional Information - see redis.io

- Redis supports many level of persistence: no persistence, RDB(point-in-time snapshot), AOF (log every write), RDB+AOF
- Redis can be setup as a cluster of multiple Redis servers for horizontally scaling and weak consistency replication
- Many useful modules are available to extend the functionality of Redis core e.g. RedisJSON, RedisSearch, RedisTimeSeries, Redis OM, etc.

References

- ScaleGrid, "Top Redis Use Cases by. Core Data Structure Types", https://scalegrid.io/blog/top-redis-use-cases-by-core-data-structure-types/
- Jerry An, "The most important Redis data structures you must understand", https://medium.com/analytics-vidhya/the-most-important-redis-data-structures-you-must-understand-2e95b5cf2bce
- Brad Solomon, "How to use Redis with python", https://realpython.com/python-redis/