



Tutorial on Widening (and Narrowing)

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Goal of Abstract Interpretation

```
int P() {  
    int x = 2;  
    while (x >= 0) {  
        // Question: which value can “x” have here?  
        x = x + 1;  
    }  
    return(x);  
}
```

- An abstract interpretation answers this question by “**executing**” a program “**abstractly**”.
- The answer is an upper approximation.



Usual Abstract Interpretation

$x = 2; \text{ while } (x \geq 0) \{ /* \text{ i2I } */ \ x = x+1; \}$

- First, define a **finite** lattice $\langle L, ?, t, u, \geq \rangle$.
 - Each element of L denotes a set of integers that “ x ” can have in the loop.
- Second, define a monotone function $F: L \rightarrow L$.
 - F is the abstraction of the loop body.
- Finally, compute the least fixpoint $\text{fix}(F)$.
 - $\text{fix}(F)$ is the limit of $?, F(?), F^2(?), \dots$
 - Since L is finite and $F^n(?)$ is increasing, the sequence “terminates”.
- Example: $L = \{?, +, 0, -, >\}$ and $F(X) = (+ \ t \ (X \ \odot \ +))$



Infinite Abstract Domain

- Question: Can we use an infinite lattice L ?
 - Using a larger L , we can obtain a better estimate of a program invariant.
 - But, $\text{fix}(F)$ might not be computable.
- Answer: Yes, we can, if we have widening.
 - Intuitively, widening works by picking a different finite subset of L for each program F .



Goal of this Talk

- My goal is to demystify widening and narrowing:
 1. What are widening and narrowing?
 2. How do they allow us to use an infinite lattice?
 3. How to design widening and narrowing?
 4. Are they really necessary? Can we just live with finite lattices, or lattices with no strictly increasing chain?
- I'll mostly focus on widening.



Problem Mathematically

- Question: Given an infinite lattice $\langle L, ?, t, u, \rangle$ and monotone $F: L \rightarrow L$, compute $\text{fix}(F)$.
 - $?, F(?), F^2(?), \dots$ converges to $\text{fix}(F)$.
 - But, the sequence might strictly increase.
- Widening approach asks a different question:
Find an upper approximation “a” of $\text{fix}(F)$,
 $\text{fix}(F) \vee a$.
 1. First, approximate the sequence $\{F^n(?)\}$ by a “terminating” sequence $\{a_n\}$: $F^n(?) \vee a_n$.
 2. Then, compute the limit “a” of $\{a_n\}$.



Widening “ $r: L \sqsubseteq L \rightarrow L$ ”

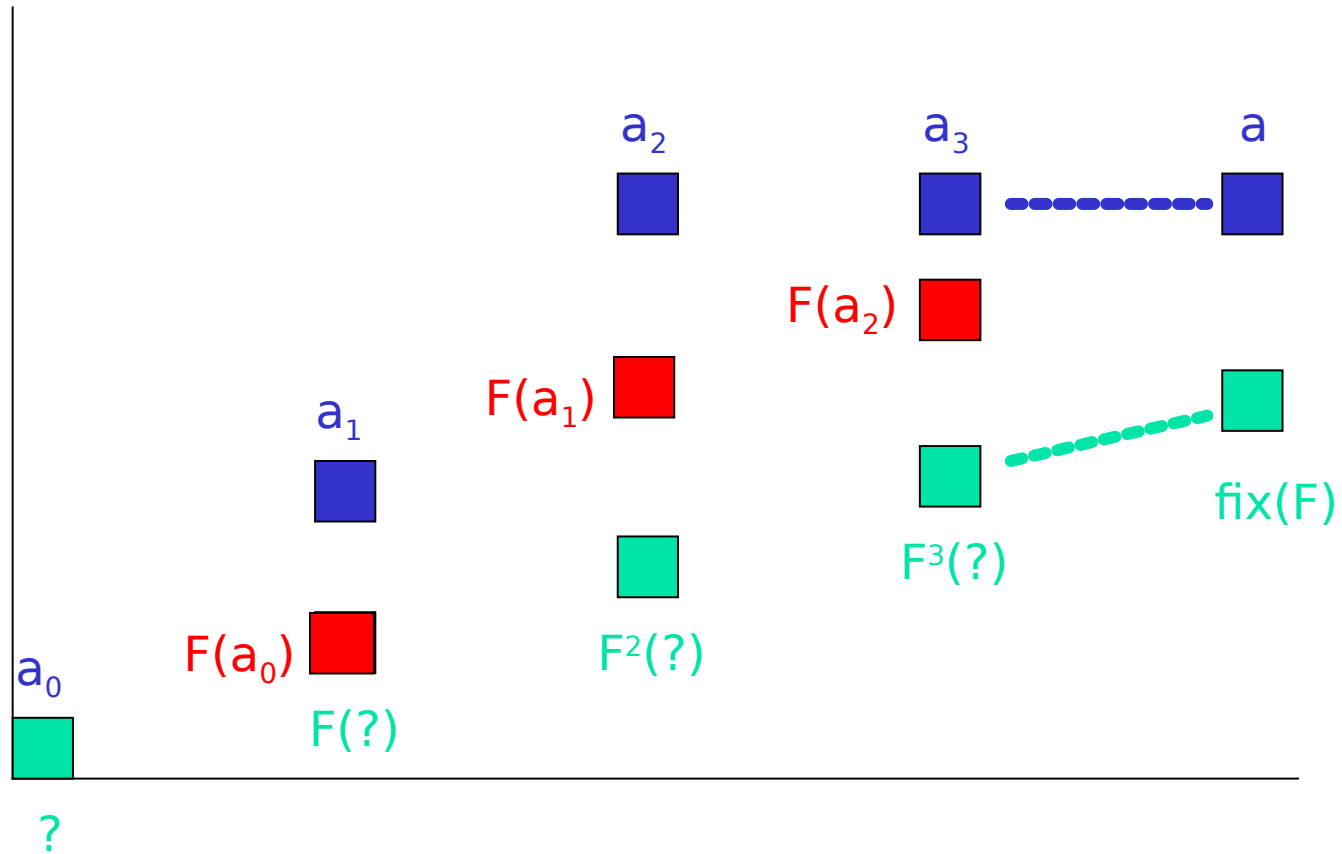
- General Dfn: Widening is what gives us $\{a_n\}$ for every F .
 - For every monotone function F , the below sequence $\{a_n\}$ approximates $\{F^n(?)\}$ and terminates:

$$a_0 = ? \qquad a_{n+1} = a_n \sqcup F(a_n)$$

- Specific Dfn: Widening $r: L \sqsubseteq L \rightarrow L$ is a function s.t.
 1. $x \sqcup x \sqcup r y$ and $y \sqcup y \sqcup r x$; and
 2. for all increasing sequences $\{x_n\}_n$, the “widened” sequence y_n terminates:

$$y_0 = x_0 \qquad y_{n+1} = y_n \sqcup x_{n+1}$$

Widening





Example

- Find $\text{fix}(F)$ in the interval domain L :

$$L = \{?\} \sqcup \{[l, u] \mid l, u \in (\mathbb{Z} \cup \{-\infty, \infty\}) \wedge l \leq u\}$$

$$F(X) = [0, 0] \sqcup (X + [1, 1])$$
- $F^n(?) = [0, n-1]$. Thus, $\{F^n(?)\}$ is strictly increasing.
- Use the following widening r :

$$[l, u] \sqcup ? = ? \sqcup [l, u] = [l, u]$$

$$[l, u] \sqcup [l', u'] = [\text{if } l' < l \text{ then } -\infty \text{ else } l, \text{if } u' > u \text{ then } \infty \text{ else } u]$$
- What is the limit of $\{a_n\}$?



Technique for Designing Widening

- For each non-bottom x ($\neq \perp$),
 - decide a finite lattice L_x (μL), and
 - define $(x\text{-})$ as an “abstraction” fn from L to L_x .
$$(x\text{-}) : L \rightarrow L_x : \text{id}$$
- Then, define $\text{?rx} = x$.
- Example: for r in the previous slide,
$$L_{[l,u]} = \{\perp, [l,u], [l,\text{inf}], [-\text{inf},u], [-\text{inf},\text{inf}]\}$$
- This kind of widening lets us pick a different finite subset L_F (μ) for each F .
- Exercise: $F(X) = ([1,1] \sqcap (X + [-1,1])) \sqcup [-\text{inf},4]$. Design a widening whose “a” for F is $[-\text{inf},4]$.



Finite Domain on the Fly (My Observation)

Suppose that for each non-bottom x ($\neq \perp$), there is a subset L_x of L such that

1. L_x with the order of L is a lattice with finite height;
2. “ $\alpha_x : L \rightarrow L_x : \text{id}$ ” is a Galois embedding; and
3. $\alpha_x(\alpha_x(x)) = x = \perp_{L_x}$.

Proposition1: If for all x and all y in L_x , $L_y \sqsubseteq L_x$, then “ $\text{urw} = \alpha_u$ (w)” and “ $\text{?rw} = w$ ” define a widening operator.

Proposition2: Moreover, if for all x and y in L_x , $\alpha_y(z) = \alpha_x(z) \sqcup_x y$, then for every non-strict monotone function F , this widening picks $L_{F(\perp)}$ in the computation: the $n+1$ -th term a_{n+1} of the widened sequence is the n -th iterate of $(\alpha_{F(\perp)} \circ F)$: $L_{F(\perp)} \sqsubseteq L_{F(\perp)}$.



Widening Sequence $\{r_i\}$

- Is there a widening that uses $F^k(?)$ to decide a finite subset f or F ?
 - Need to use different “widening” r_i for a_i .
- For every monotone function F , the below sequence $\{a_n\}$ approximates $\{F^n(?)\}$ and terminates:

$$a_0 = ? \qquad a_{n+1} = a_n r_n F(a_n)$$

- A **widening sequence** $\{r_n\}$ is a sequence of fns s.t.
 1. $x \vee x r_n y$ and $y \vee x r_n y$; and
 2. for all increasing sequences $\{x_n\}_n$, the “widened” sequence y_n terminates:

$$y_0 = x_0 \qquad y_{n+1} = y_n r_n x_{n+1}$$

- Design Technique: Usually, $r_0 = \dots = r_{k-1} = t$, and $r_k = r_{k+1} = \dots$
- Exercise: $F(X) = ([1,1] \ t \ (X + [-10,10])) \cup [-\text{inf}, 30]$.

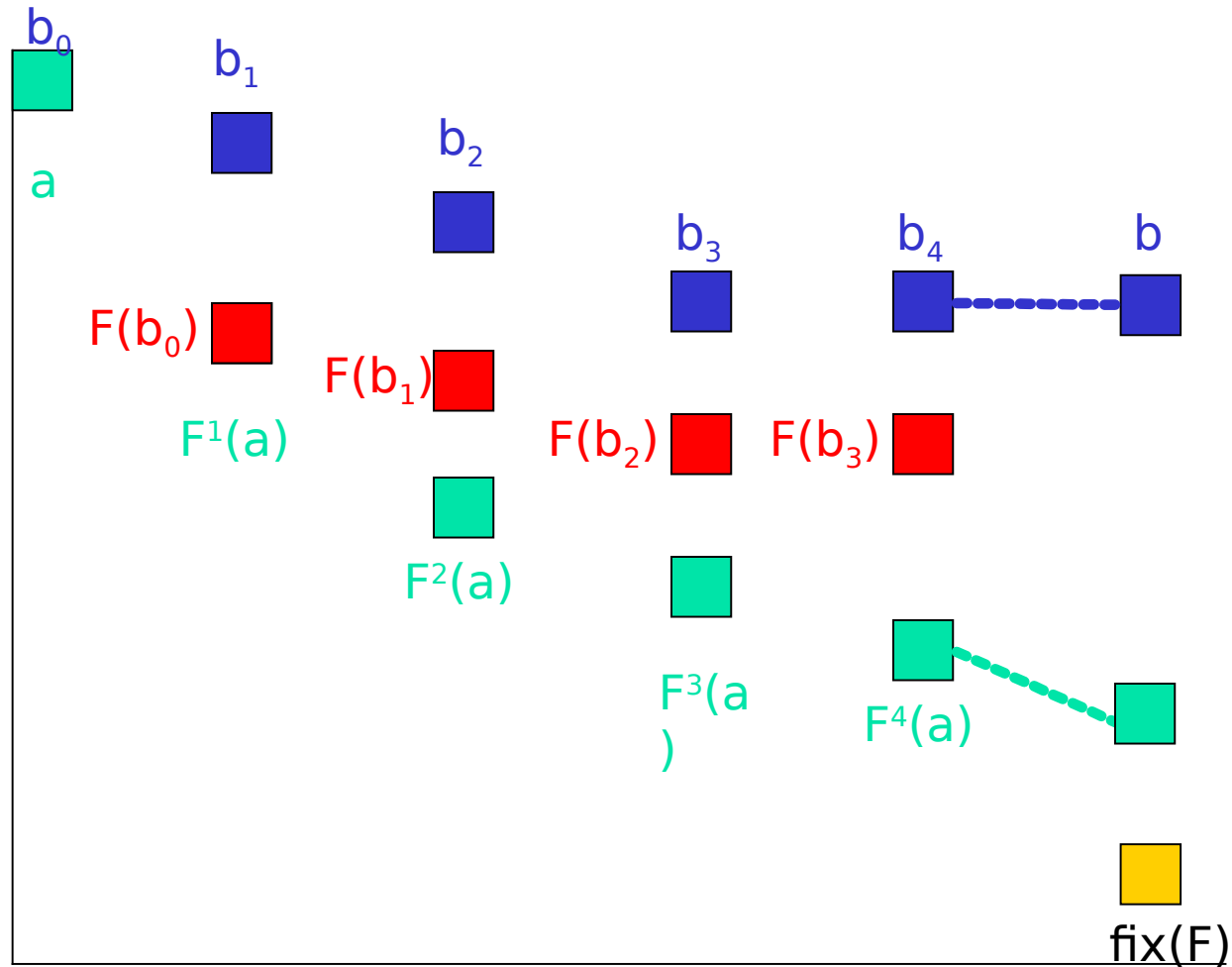


Increase the Precision

- Problem:
 - Suppose we obtained an approximation “a” of $\text{fix}(F)$ via widening.
 - From “a”, find a more precise approximation.
- Not-computable “solution”:
 - $a, F(a), F^2(a), F^3(a), \dots$
 - For each n , $\text{fix}(F) \sqsubseteq F^n(a)$ and $F^{n+1} \sqsubseteq F^n(a)$. (* a w $F(a)$)
- Narrowing approach:
 1. Approximate the sequence $\{F^n(a)\}$ by a “terminating” decreasing sequence $\{b_n\}$: $F^n(a) \sqsupseteq b_n$.
 2. Then, b is the limit of $\{b_n\}$.



Narrowing





Narrowing

- General Dfn: Narrowing is what gives us $\{b_n\}$ for every F .
 - For every monotone function F , the below sequence $\{b_n\}$ approximates $\{F^n(a)\}$, is decreasing and terminates:

$$b_0 = a \qquad b_{n+1} = b_n \sqcap F(b_n)$$

- Specific Dfn: Narrowing $\Delta: L \times L \rightarrow L$ is a function such that
 1. $x \sqsupseteq x \Delta y \sqsupseteq y$, and
 2. for all decreasing sequences $\{x_n\}_n$, the “narrowed” sequence $\{y_n\}$ terminates:

$$y_0 = x_0 \qquad y_{n+1} = y_n \Delta x_{n+1}$$



Widening/Narrowing Example

- $F(X) = ([1,1] \sqcup (X + [1,1])) \sqcup [-\infty, 100]$
- Widening:
 - $?rx = xr? = ?$
 - $[l,u] \sqcup [l',u'] = [\text{if } l > l' \text{ then } -\infty \text{ else } l, \text{if } u < u' \text{ then } \infty \text{ else } u]$
- Narrowing:
 - $? \Delta x = x \Delta ? = ?$
 - $[l,u] \Delta [l',u'] = [\text{if } l = -\infty \text{ then } l' \text{ else } l, \text{if } u = \infty \text{ then } u' \text{ else } u]$
- Exercise: Compute an approximation of $\text{fix}(F)$.



Finite Lattices Are Not as Powerful as Widening.

```
int Pnm() {  
    int i = n;  
    while (i <= m) { // i ∈ [n,m]  
        i := i + 1; }  
    return(i);  
}
```

- With a single “finite” abstract domain, we cannot find the invariant of “P_{nm}()” for all n and m.
 - $\{[n,m] \mid n \cdot m\}$ has a strictly increasing sequence.
- We can find the invariant of all “P_{nm}()”s, using the interval domain and the widening and narrowing in the previous slide.



Conclusion

- Widening lets us use an infinite lattice in designing an abstract interpreter.
 - It picks a finite subset depending on an input program.
- A common technique for designing widening is to compute $F^k(?)$ for some fixed k , and to use the result to build a finite subset.
- Why don't you design an abstract interpreter based on an infinite lattice and widening?