Lab Assignment 2 & 3

Team Number: **50**

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Q.1

```
Command to Compile: g++ q1.cpp -o q1.exe -lmysql
```

Q.2

```
Command to Compile: g++ q2.cpp -o q2.exe -lmysql
```

Q.3

```
Command to Compile:
```

```
g++ q3_create_table.cpp -o q3_create_table.exe -lmysql
g++ q3_a.cpp -o q3_a.exe -lmysql
g++ q3_b.cpp -o q3_b.exe -lmysql
g++ q3_c.cpp -o q3_c.exe -lmysql
```

In this question, you first need to run the file *q3_create_table.exe* as it creates the tables which are required in the other subparts.

The table so created has the following columns:

```
← T→ salesman_id name address_city coverage_city commission
```

A)
After entering some dummy data into the table and running all queries given in Q3 part
A, the table looks like this:

←T			~	salesman_id	name	address_city	coverage_city	commission	date_of_employment	date_of_release
	<i> </i>	≩ Copy	Delete	1	raj	mumbai	delhi	750	NULL	NULL
	Ø Edit	≩ Copy	Delete	2	tom	bhopal	pune	450	NULL	NULL
	<i> </i>	≩ Сору	Delete	3	manya	delhi	kolkata	700	NULL	NULL
	<i>⊘</i> Edit	≩ Copy	Delete	4	abc	mumbai	chennai	600	NULL	NULL
	<i> </i>	≩ Copy	Delete	5	bhim	surat	bhopal	500	NULL	NULL
	<i> </i>	≩ Copy	Delete	6	sara	delhi	mumbai	600	NULL	NULL
	<i> </i>	≩ Copy	Delete	7	raj	chennai	bhopal	450	NULL	NULL
	<i> </i>	≩ Copy	Delete	8	manu	chennai	kolkata	850	NULL	NULL
	<i> </i>	≩ € Copy	Delete	9	garry	mumbai	mumbai	450	NULL	NULL
	Edit	≩ Copy	Delete	10	binny	delhi	delhi	500	NULL	NULL
	<i> </i>	≩ € Copy	Delete	11	geeta	bhopal	pune	450	NULL	NULL
	Edit	≩ Copy	Delete	12	blue	kolkata	kolkata	500	NULL	NULL
	<i> </i>	≩ € Copy	Delete	13	abc	pune	chennai	500	NULL	NULL
	<i> </i>	≩ Copy	Delete	14	xyz	pune	bhopal	500	NULL	NULL
	Edit	≩ € Copy	Delete	15	feeli	delhi	mumbai	700	NULL	NULL
	<i> </i>	≩ Сору	Delete	16	raj	mumbai	mumbai	450	NULL	NULL
	<i> </i>	≩ Copy	Delete	17	raj	mumbai	kolkata	900	NULL	NULL
	<i> </i>	≩ € Copy	Delete	18	raj	pune	pune	450	NULL	NULL
	<i>⊘</i> Edit	≩ Copy	Delete	19	raj	delhi	chennai	850	NULL	NULL

Here date_of_employment and release columns have been added after data entry, so they currently contain null values. So the employee entries made before adding the columns have null values, but now whenever any entry for an employee is made, date_of_emplomeny and release can be specified.

```
B)
The functional dependencies formed in the table after Q3-A-(k) are:-
salesman_id → {'name', 'address_city', 'coverage_city', 'commission'}
('address_city', 'coverage_city') → {'commission'}
```

As we can see that there is a transitive dependency present in the table, thus it cannot be in 3NF.

Also, as there are no partial dependencies present in the table as the primary key is non-composite, thus it can be in 2NF.

Therefore, the normal form of the table after subpart A(k) is 2NF.

We can improve upon it by changing it to 3NF by removing the transitive dependency.

To do this, we have split our original table into 2 tables. First, we have a table called **commissions** consisting of 'salesman_id', 'address_city', 'coverage_city', and 'commission' in which the primary key is ('salesman_id', 'address_city', 'comission_city'). Secondly, the original table **salesmen** now only contain 'salesman_id', 'name', 'date_of_employment' and 'date_of_release', with the primary key 'salesman_id'.

New F.Ds are:-

```
<u>Salesmen</u>: salesman_id → {'name', 'date_of_employment', 'date_of_release'}

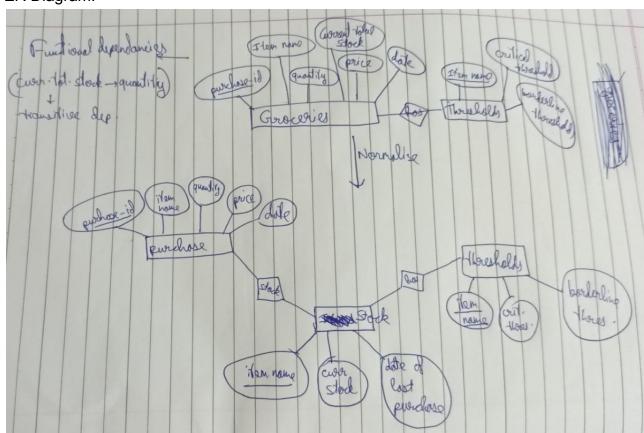
<u>Commissions</u>: ('salesman_id', 'address_city', 'coverage_city') → {'commission'}
```

Q.4

```
Command to Compile:

g++ q4_a.cpp -o q4_a.exe -lmysql
g++ q4_b.cpp -o q4_b.exe -lmysql
g++ q4_c.cpp -o q4_c.exe -lmysql
```

ER Diagram:-



B)

Currently, the table has the following dependencies:-

```
'purchase_id' → {'item_name', 'price', 'date_of_purchase', 'quantity'} 'quantity' → 'current_total_stock'
```

As we can see that there is a transitive dependency present in the table, thus it cannot be in 3NF.

Also, as there are no partial dependencies present in the table as the primary key is non-composite, thus it can be in 2NF.

Thus, the table is not in BCNF.

Now, to normalize the table into BCNF, we need to decompose it into 2 tables namely, 'groceries' and 'stock'. **groceries** consist of 'purchase_id', 'item_name', 'quantity', 'date_of_purchase' and 'price', with the primary key as 'purchase_id'. **stock** consists of 'item_name', 'last_purchase_date' and 'current_total_stock', with the primary key as 'item_name'.

New F.Ds are:-

```
<u>groceries</u>: purchase_id \rightarrow {'item_name', 'date_of_purchase', 'quantity', 'price'} 

<u>stock</u>: item_name \rightarrow {'last_purchase_date', 'current_total_stock'}
```

Since now all the functional dependencies have a primary key on the LHS, the table is in BCNF.

C)

Un-Normalized Tables -> Query of Q4-C-(v)

```
Showing rows 0 - 24 (59 total, 0 in query, Query took 0.0004 seconds.)

SELECT * FROM groceries INNER JOIN thresholds ON groceries.item_name = thresholds.item_name WHERE groceries.cur_stock > thresholds.borderline_threshold;
```

Normalized Tables -> Query of Q4-C-(v)

```
Showing rows 0 - 0 (1 total, Query took 0.0003 seconds.)

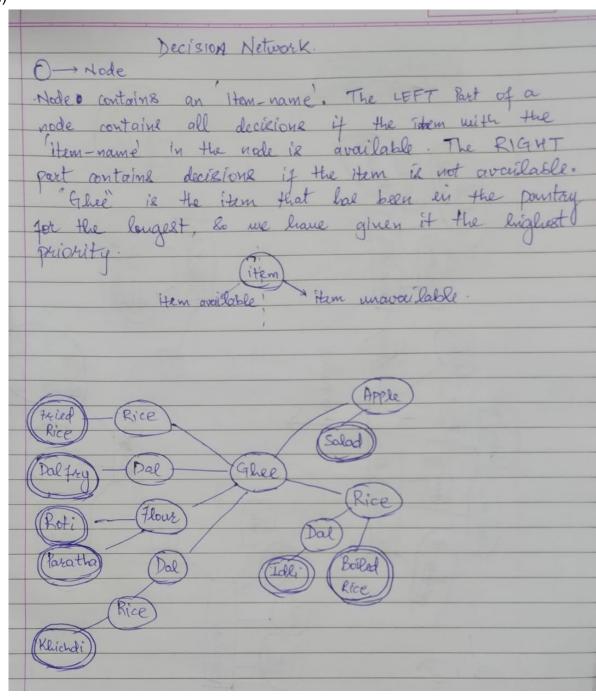
SELECT * FROM stock INNER JOIN thresholds ON stock.item_name = thresholds.item_name WHERE stock.cur_stock > thresholds.borderline_threshold;
```

D)

As we can see that the execution time for the case on unnormalised data is longer than that in the normalised table. This is the case when the table has entries that are less than 100. Now when we deal with bigger databases, when the number of entries will be in millions, this difference will amplify, and will definitely be considerable in real-time

usage. Besides, Normalised data makes the schema of the database lucid and structured.

E)

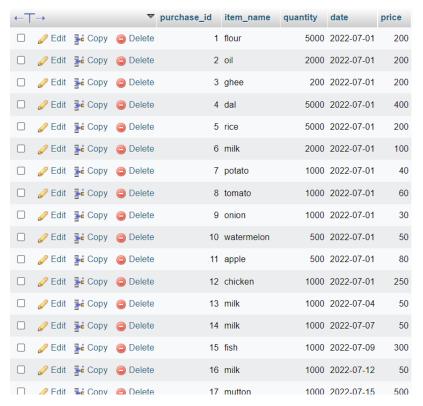


Groceries table (un-normalized table)

←Τ	_→		∇	purchase_id	item_name	quantity_bought	cur_stock	date	price
	<i></i> €dit	≩ Сору	Delete	1	flour	5000	5000	2022-07-01	200
	<i></i> Edit	≩ € Сору	Delete	2	oil	2000	2000	2022-07-01	200
	<i></i> €dit	≩ € Сору	Delete	3	ghee	200	200	2022-07-01	200
		З Сору	Delete	4	dal	5000	5000	2022-07-01	400
	<i> </i>	≩ Сору	Delete	5	rice	5000	5000	2022-07-01	200
		≩ € Сору	Delete	6	milk	2000	2000	2022-07-01	100
	<i> </i>	≩ Сору	Delete	7	potato	1000	1000	2022-07-01	40
		≩ € Сору	Delete	8	tomato	1000	1000	2022-07-01	60
	<i> </i>	≩ € Сору	Delete	9	onion	1000	1000	2022-07-01	30
		≩ € Сору	Delete	10	watermelon	500	500	2022-07-01	50
	<i> </i>	≩ Сору	Delete	11	apple	500	500	2022-07-01	80
		≩ € Сору	Delete	12	chicken	1000	1000	2022-07-01	250
	<i> </i>	≩ Сору	Delete	13	milk	1000	1000	2022-07-04	50
		З Сору	Delete	14	milk	1000	1000	2022-07-07	50
	<i> </i>	≩ Сору	Delete	15	fish	1000	1000	2022-07-09	300
		≩ € Сору	Delete	16	milk	1000	1000	2022-07-12	50
	<i>A</i> Edit	₹ Conv	ΠαΙατα	17	mutton	1000	1000	2022 07 15	500

Tables formed after normalization to BCNF:-

Purchase table



Stock table



Thresholds table (For Q4 c (iii, iv and v))

Since different items will have different thresholds for when we need to/may/need not buy them.



Contributions

Soumik Roy:

ER Diagram, Code for Problem 1, Problem 2 and Parts of Problem 3. Kulkarni Tanmay Shreevallabh:

Decision Network, Schema for problem 3, Code for problem 3, Problem 4 Yash Bhargava:

Code for Problem 4, Report writing, Decision Network.