

Sorting

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Outline

- Sorting-concept
- Sorting Terms
- Bubble sort
- Insertion sort
- Counting sort
- Sorting applications





Sorting

- Sorting is any process of arranging items systematically in a particular order
 - Sorting in ascending order :arrange n keys in such a way that key_i < key_i for any i & j such that i<j
 - Sorting in descending order: arrange n keys in such a way that key_i > key_j for any i & j any i & j such that i<j





Sorting Terms

- Stable sort
- Inplace sort
- Number of Passes





Bubble Sort

- Compares adjacent array elements
 - Exchanges their values if they are <u>out of order</u>

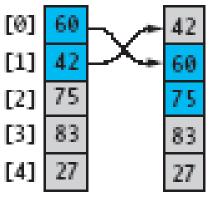
- Smaller values <u>bubble up</u> to the top of the array
 - Larger values sink to the bottom





FIGURE 10.1

One Pass of Bubble Sort



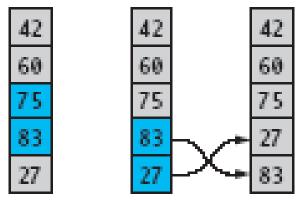


FIGURE 10.2

Array After Completion of Each Pass

[6]	42
[1]	60
[2]	75
[3]	27
[41]	83

42	
60	
27	
75	
83	

42	
27	
60	
75	
83	

27
42
60
75
83





Bubble Sort Algorithm

- 1. do
- 2. for each pair of adjacent array elements
- 3. if values are out of order
- 4. Exchange the values
- 5. while the array is not sorted





Bubble Sort Algorithm, Refined

- 1. do
- 2. Initialize exchanges to false
- 3. for each pair of adjacent array elements
- 4. if values are out of order
- 5. Exchange the values
- 6. Set exchanges to true
- 7. while exchanges





- Excellent performance <u>in some cases</u>
 - But <u>very poor</u> performance in others!
- Works best when array is nearly sorted to begin with
- Worst case number of comparisons: O(n²)
- Worst case number of exchanges: O(n²)
- <u>Best case</u> occurs when the array is already sorted:
 - O(n) comparisons
 - O(1) exchanges (none actually)



```
bubbleSort(int arr[], int n)
  int i, j;
  for (i = 0; i < n-1; i++)
  // Last i elements are already in place
  for (j = 0; j < n-i-1; j++)
     if (arr[j] > arr[j+1])
       swap(&arr[j], &arr[j+1]);
```





Insertion Sort

- Based on technique of card players to arrange a hand
 - Player keeps cards picked up so far in <u>sorted order</u>
 - When the player picks up a new card











K J Somaiya College of English Sertion sort algorithm

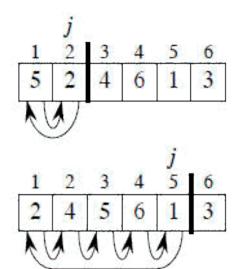
INSERTION_SORT takes as parameters an array A[1...n] and the length n of the array. The array A is sorted in place: the numbers are rearranged within the array, with at most a constant number outside the array at any time.

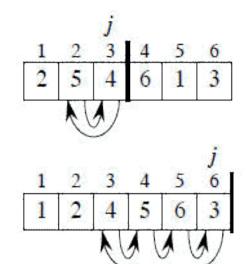
```
INSERTION SORT (A)
      FOR j \leftarrow 2 TO length[A]
2.
            DO key \leftarrow A[i]
3.
                 {Put A[j] into the sorted sequence A[1..j-1]}
4.
                 i \leftarrow i - 1
5.
                 WHILE i > 0 and A[i] > \text{key}
                           DO A[i+1] \leftarrow A[i]
6.
                                 i \leftarrow i - 1
7.
8.
                  A[i+1] \leftarrow \text{key}
```

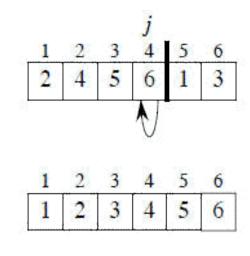




K J Somaiya College of Englishern sertion sort algorithm







Stability:

Since multiple keys with the same value are placed in the sorted array in the same order that they appear in the input array, Insertion sort is stable.





Counting sort

- sorting is based on keys between a specific range.
- It works by counting the number of objects having distinct key values
- Followed by computation of position of each object in the output sequence.





Counting sort

- Initialize count array of the size of input range
- Update the count array to store the count of each unique key.
- Further update the count array with cumulative additions of previous counts
- Shift the count array to right by one position; no circular shift
- Initialize sort array of the size of input sequence
- Update sort array by entering keys from input array at location from count array and increment the count by 1





K J Somaiya College of Engineering Counting sort example

• i/p:2312452154

 N= 10, rang 	e: 1:5					
Initialize count a	array of th	ne size of	input rar	nge		
count array	0	1	2	3	4	5
	0	0	0	0	0	0
Update the cou	nt array to	o store th	ne count d	of each u	nique key	'.
count array	0	1	2	3	4	5
	0	2	3	1	2	2
Further update the	e count arra	y with cum	nulative add	ditions of p	revious cou	ınts
count array	0	1	2	3	4	5
	0	2	5	6	8	10
Shift the count	array to ri	ght by or	ne positio	n; no circ	cular shift	
count array	0	1	2	3	4	5
		_	_	_		_



Initialize sort array of the size of input sequence

Sort Array 0 1 2 3 4 5 6 7 8 9

Update sort array by entering keys from input array at location from count array and increment the count by 1

i/p	2	3	1	2	4	5	2	1	5	4
count array	0	1	2	3	4	5				
	0	0	2	5	6	8				
Output	0	1	2	3	4	5	6	7	8	9
Sorted Sancina y	1	1	2	2	2	3	4	4	5	5



i/p	2	3	1	2	4	5	2	1	5	4
count array	0	1	2	3	4	5				
	0	0	2 3	5	6	8				
Sort Array	0	1	2	3	4	5	6	7	8	9
			2							
i/p	2	3	1	2	4	5	2	1	5	4
count array	0	1	2	3	4	5				
	0	0	2 3	5 6	6	8				
Sort Array	0	1	2	3	4	5	6	7	8	9
			2			3				





i/p	2	3	1	2	4	5	2	1	5	4
count array	0	1	2	3	4	5				
	0	0 1	2 3	5 6	6	8				
Sort Array	0	1	2	3	4	5	6	7	8	9
	1		2			3				
i/p	2	3	1	2	4	5	2	1	5	4
count array	0	1	2	3	4	5				
	0	0 1	2 3 4	5 6	6	8				
Sort Array	0	1	2	3	4	5	6	7	8	9
	1		2	2		3				





i/p	2	3	1	2	4	5	2	1	5	4
count array	0	1	2	3	4	5				
	0	0 1	2 3 4	5 6	6 7	8				
Sort Array	0	1	2	3	4	5	6	7	8	9
	1		2	2		3	4			
i/p	2	3	1	2	4	5	2	1	5	4
count array	0	1	2	3	4	5				
	0	0 1	2 3 4	5 6	6 7	8 9				
Sort Array	0	1	2	3	4	5	6	7	8	9
	1		2	2		3	4		5	





i/p	2	3	1	2	4	5	2	1	5	4
count array	0	1	2	3	4	5				
	0	0 1	2 3 4 5	5 6	6 7	8 9				
Sort Array	0	1	2	3	4	5	6	7	8	9
	1		2	2	2	3	4		5	
i/p	2	3	1	2	4	5	2	1	5	4
count array	0	1	2	3	4	5				
	0	0 1 2	2 3 4 5	5 6	6 7	8 9				
Sort Array	0	1	2	3	4	5	6	7	8	9
	1	1	2	2	2	3	4		5	





i/p	2	3	1	2	4	5	2	1	5	4
count array	0	1	2	3	4	5				
	0	0 1 2	2 3 4 5	5 6	6 7	8 9 10				
Sort Array	0	1	2	3	4	5	6	7	8	9
	1	1	2	2	2	3	4		5	5
i/p	2	3	1	2	4	5	2	1	5	4
count array	0	1	2	3	4	5				
	0	0 1 2	2 3 4 5	5 6	6 7 8	8 9 10				
Sort Array	0	1	2	3	4	5	6	7	8	9
ucha -	1	1	2	2	2	3	4	4	5	5



Analysis of sorting algorithms

Sr.	Algorithm	Stable?	Inplace?	#passes?
1	Bubble			
2	Insertion			
3	Counting			





Thank you

