



#### GANAK: A Scalable Probabilistic Exact Model Counter

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### Propositional Model Counting

- Given:
  - Propositional formula F (CNF) over a set of variables X
- Propositional Model Counting (#SAT):
  - Compute the number of satisfying assignments of F
- #SAT is a #P complete problem

### Propositional Model Counting

- Probabilistic Exact Model Counting
  - Given a propositional formula F (CNF) and confidence  $\delta \in (0,1]$ , counter returns count such that:

$$\mathsf{Pr}ig[|\mathsf{Solutions}\;\mathsf{of}\;\mathsf{F}|=\mathsf{count}ig]\geq 1-\delta$$

<sup>&</sup>lt;sup>1</sup>Chakraborty et al., 2019

# Propositional Model Counting

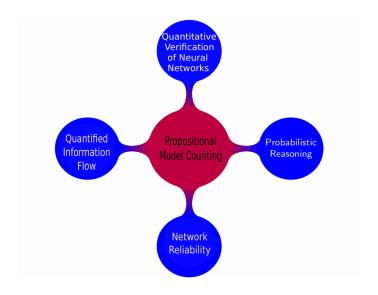
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 Probabilistic Exact Model Counting is almost as hard as Exact Model Counting<sup>1</sup>

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#### Applications of Propositional Model Counting



• Decision Process:

$$\begin{array}{ll} - & (F \wedge I) \vee (F \wedge \neg I) \\ - & \#(F) = \#(F \wedge I) + \#(F \wedge \neg I) \end{array}$$
 mutually inconsistent

Decision Process:

$$-(F \wedge I) \vee (F \wedge \neg I)$$
 mutually inconsistent 
$$-\#(F) = \#(F \wedge I) + \#(F \wedge \neg I)$$

• Component Decomposition:

$$-F = \Delta_1 \wedge \Delta_2 \cdots \Delta_n \quad \Delta_1 \cdots \Delta_n \text{ does not share any variables}$$

$$-\#(F) = \#(\Delta_1) \times \#(\Delta_2) \cdots \times \#(\Delta_n) \quad \text{mutually disjoint}$$

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• Component Caching:

Key Value	
$\Delta_1$	$\#(\Delta_1)$
$\Delta_2$	$\#(\Delta_2)$

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Conflict Driven Clause Learning

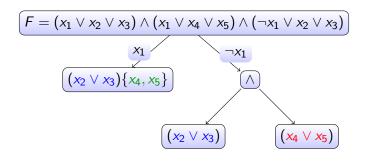
$$F = (x_1 \lor x_2 \lor x_3) \land (x_1 \lor x_4 \lor x_5) \land (\neg x_1 \lor x_2 \lor x_3)$$

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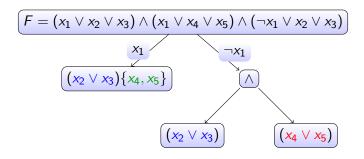
$$x_1$$

$$(x_2 \lor x_3)\{x_4, x_5\}$$

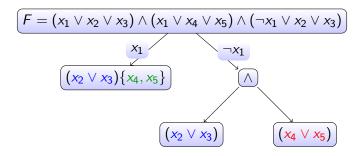
Key	Value	
$(x_2 \lor x_3)$	3	
$(x_2 \lor x_3)\{x_4, x_5\}$	12	



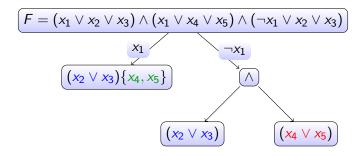
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$(x_2 \lor x_3)\{x_4, x_5\}$	12
$(x_4 \vee x_5)$	3
$(x_2 \lor x_3) \land (x_4 \lor x_5)$	9



Key	
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$(x_2 \lor x_3)\{x_4, x_5\}$	12
$(x_4 \lor x_5)$	3
$(x_2 \lor x_3) \land (x_4 \lor x_5)$	9
$F = (x_1 \lor x_2 \lor x_3) \land (x_1 \lor x_4 \lor x_5) \land (\neg x_1 \lor x_2 \lor x_3)$	21

#### Our Contribution

- Probabilistic Component Caching (PCC)
- Variable Branching Heuristic (CSVSADS)
- Phase Selection Heuristic (PC)
- Independent Support (IS)
- Restarts (LSO)
- Exponentially Decaying Randomness (EDR)

# Probabilistic Component Caching (PCC)

$$F = (\neg x_3 \lor \neg x_5 \lor x_6) \land (\neg x_1 \lor x_4 \lor \neg x_6) \land (x_2 \lor x_3 \lor x_6)$$

Schema	Key	Value
STD <sup>2</sup>	-3, -5, 6, 0, -1, 4, -6, 0, 2, 3, 6, 0	#(F)
HC <sup>3</sup>	1, 2, 3, 4, 5, 6, 1, 2, 3	#(F)
GANAK	Hash of HC/STD	#(F)

<sup>&</sup>lt;sup>2</sup>Sang et al., 2005

<sup>&</sup>lt;sup>3</sup>Thurley, 2006

# Variable Branching Heuristic (CSVSADS)

- $\bullet \ \, \mathsf{Score}(\mathsf{VSADS})^4 = \mathsf{p} \times \mathsf{Score}(\mathsf{VSIDS}) + \mathsf{q} \times \mathsf{Score}(\mathsf{DLCS})$ 
  - VSIDS: Prioritize variables present in recently generated conflict clauses
  - DLCS: Prioritize the highest occurring variable in the residual formula

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  - VSIDS: Prioritize variables present in recently generated conflict clauses
  - DLCS: Prioritize the highest occurring variable in the residual formula
- Score(CSVSADS) =  $\underline{\alpha \times CacheScore} + \beta \times Score(VSADS)$ 
  - Prioritize variables not present in the components which are recently added to the cache

<sup>&</sup>lt;sup>4</sup>Sang et al., 2005

# Phase Selection Heuristic (PC)

$$\mathsf{DLIS}^5 \left\{ \begin{array}{ll} I & |I| \geq |\neg I| \\ \neg I & \textit{otherwise} \end{array} \right.$$

<sup>&</sup>lt;sup>5</sup>Sang et al., 2005

# Phase Selection Heuristic (PC)

$$\mathsf{DLIS}^5 \left\{ \begin{array}{ll} I & |I| \geq |\neg I| \\ \neg I & \textit{otherwise} \end{array} \right.$$

• We reduce our trust on DLIS by adding randomness in DLIS if the difference in |I| and  $|\neg I|$  is not overwhelmingly high

<sup>&</sup>lt;sup>5</sup>Sang et al., 2005

#### Tool

- GANAK<sup>6</sup>: First Scalable Probabilistic Exact Model Counter
- Given a propositional formula F (CNF) and confidence  $\delta \in (0,1]$  GANAK $(F,\delta)$  returns count such that

$$\Pr[|Sol(F)| = \mathtt{count}] \ge 1 - \delta$$

Tool is available at: https://github.com/meelgroup/ganak

<sup>&</sup>lt;sup>6</sup>GANAK (गणक in Sanskrit) refers to a device that counts

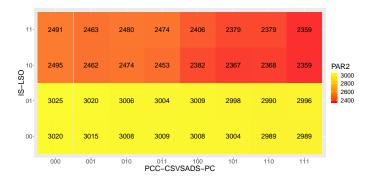
#### **Experimental Evaluation**

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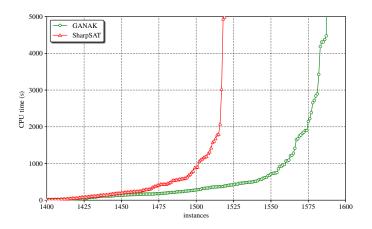
- Benchmarks arising from probabilistic reasoning, plan recognition, DQMR networks, ISCAS89 combinatorial circuits, quantified information flow, etc
- Objectives:
  - Study the impact of different configurations of heuristics
  - Study the performance of GANAK with respect to the state-of-the-art model counters

# Experimental Evaluation: Individual Analysis



• GANAK performed best when all the heuristics are turned on

### Experimental Evaluation: Comparison with other tools



- $\bullet$   $\delta = 0.05$ , Component Cache Size = 2 GB, Timeout=5000 secs
- In our experiments, the model count returned by GANAK was equal to the exact model count for all benchmarks

#### Thank You

Tool is available at: https://github.com/meelgroup/ganak

