# Hash Index vs B-tree Index

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## What is Indexing?

To understand what indexing is, first let us consider a sample Database having a table called Employee which has Name, Age and Sex as the columns. Assume that this table has thousands of rows. Suppose we want to find all employees who have name as "Ram".

Without an index, the Databse software would have to sequentially scan every single row in the table to see if the employee name matches 'Ram'. This is called a full table scan.

The whole point of indexing is to speed up search queries by cutting down the number of records/rows that needs to be examined.

PostgreSQL provides the index methods B-tree, hash, GiST, and GIN. By default, the CREATE INDEX command creates B-tree indexes, which works well with all the comparison operators, or with operators such as BETWEEN, IN, IS NULL, LIKE etc. Hash indexes can only handle equality comparisons.

#### More about Hash Indexes:

The reason hash indexes are used is because hash tables are extremely efficient when it comes to just looking up values. So, queries that compare for equality to a string work well. For instance, the query where we fetch employees whose name is 'RAM' could benefit from a hash index created on the Name column.

The way a hash index would work is that the column value will be the key into the hash table and the actual value mapped to that key would just be a pointer to the row data in the table. Since a hash table is basically an associative array, a typical entry would look something like "RAM => 0×28939", where 0×28939 is a reference to the table row where RAM is stored in memory. Looking up a value like "RAM" in a hash table index and getting back a reference to the row in memory is obviously a lot faster than scanning the table to find all the rows with a value of "RAM" in the Name column.

## Reason for choosing this topic:

While going through the source code and the mailing lists, i came across the following lines

"Note: Testing has shown PostgreSQL's hash indexes to perform no better than B-tree indexes, and the index size and build time for hash indexes is much worse. For these reasons, hash index use is presently discouraged."

While it's evident that Hash index operations are not currently WAL logged and also that it works only for statements that check for equality, hence the usage being discouraged. However, I decided to check for myself how it compares against B-Tree indexes in terms of speed.

Here, I try to present the results of my comparison between the two indexing techniques.

## Creation of Tables, Comparison:

First of all, initialize a schema test (As was shown in the postgres\_installation).

Enter the schema

./bin/psql test

Now create a test\_table with random\_string as a column.

Create table test\_table (random\_string text);

To test the index creation time, we would need atleast a million rows. Let's generate that using a perl script. (insertdata.pl)

```
The script is as follows:
```

```
#!/usr/bin/perl -w
use strict;
print generate_random_string() . "\n" for 1 .. 10_000_000;
exit;
sub generate_random_string {
    my $word_count = 2 + int(rand() * 4);
    return join(' ', map { generate_random_word() } 1..$word_count);
}
```

```
sub generate_random_word {
  my $len = 3 + int(rand() * 5);
  my @chars = ( "a".."z", "A".."Z", "0".."9" );
  my @word_chars = map { $chars[rand @chars] } 1.. $len;
  return join ", @word_chars;
}
```

The above script generates a million random strings each of which has a few words. The output of this file has to be redirected to a .lst file. That is dont by using

Perl insertdata.pl > data.lst

This will redirect the output to the data.lst file.

Next step is to import the data from the data.lst file to the test\_table relation.

That is done by using: test=#\copy test table from 'data.lst';

Now the table test\_table has 10M rows.

We can check this by using Select count(\*) from test\_table;

Also, \timing can be used to turn on the timing.

The results until this point are as seen in the following snapshot:

```
sourabh@sourabh-HP-ENVY-m6-Notebook-PC: ~/tarballs/postgresql-9.3.5/pgsql
                                                                                                        📆 📗 🔼 🗔 (3:37, 84%) 🜒 7:43 PM 🛱 Sourabh
sourabh@sourabh-HP-ENVY-m6-Notebook-PC:~/tarballs/postgresql-9.3.5/pgsql$ ./bin/psql test
psql (9.3.5)
Type "help" for help.
test=# \list
                         List of databases
          | Owner | Encoding | Collate | Ctype | Access privileges
  Name
postgres | sourabh | UTF8
                              en_IN en_IN
template0 | sourabh | UTF8
                                en_IN
                                         en_IN | =c/sourabh
                                                | sourabh=CTc/sourabh
template1 | sourabh | UTF8
                                en_IN
                                        | en_IN | =c/sourabh
                                                 sourabh=CTc/sourabh
          | sourabh | UTF8
                              en_IN en_IN
(4 rows)
test=# \d
          List of relations
         Name | Type | Owner
public | test_table | table | sourabh
test=# \d+ test_table
                      Table "public.test_table"
              | Type | Modifiers | Storage | Stats target | Description
random_string | text |
                              | extended |
Has OIDs: no
test=# \timing
Timing is on.
test=# select count(*) from test_table ;
10000001
Time: 1259.103 ms
test=#
```

Now the next step is to check the speed for index creation.

For this, i will need to create the following indexes:

- Hash Index
- B-Tree Index
- Unique B-Tree Index
- Unique Hash Index

The command to create an index for a column is:

CREATE [UNIQUE] INDEX [index-name] ON table [USING method](column-name(s));

where,

index name can be anything unique hash is not implemented method can either be hash or btree column names are the column(s) on which the indexes are to be created.

The names given are as follows: btree\_index for B-Tree u\_btree\_index for Unique B-Tree hash\_index for Hash

sourabh@sourabh-HP-ENVY-m6-Notebook-PC: ~/tarballs/postgresql-9.3.5/pgsql

The snapshot of the Database at that time is as follows:

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```
test=# \d+ test_table
                       Table "public.test_table"
              | Type | Modifiers | Storage | Stats target | Description
random_string | text | extended |
Has OIDs: no
test=# \timing
Timing is on.
test=# select count(*) from test_table ;
count
10000001
(1 row)
Time: 1259.103 ms
test=# create index hash_index on test_table using hash(random_string);
CREATE INDEX
Time: 22687.794 ms
test=# create index btree_index on test_table using btree(random_string);
CREATE INDEX
Time: 106681.112 ms
test=# create unique index u_hash_index on test_table using hash(random_string);
ERROR: access method "hash" does not support unique indexes
STATEMENT: create unique index u_hash_index on test_table using hash(random_string);
ERROR: access method "hash" does not support unique indexes
Time: 0.484 ms
test=# create unique index u_btree_index on test_table using btree(random_string);
CREATE INDEX
Time: 105360.166 ms
test=# \di
                 List of relations
             Name | Type | Owner | Table
Schema |
public | btree_index | index | sourabh | test_table
public | hash_index
                       | index | sourabh | test_table
public | u_btree_index | index | sourabh | test_table
(3 rows)
test=#
```

As we can observe from the above snapshot, the index creation time for hash index is much lesser than the time taken to create a B-Tree index. This was observed to be true even when repeated again after dropping the index.

What about the size of the indexes created? As per what we checked, the size of the Hash index was much smaller than the size of the B-tree index. Here's a snapshot comparing the size in MB. sourabh@sourabh-HP-ENVY-m6-Notebook-PC: ~/tarballs/postgresql-9.3.5/pgsql

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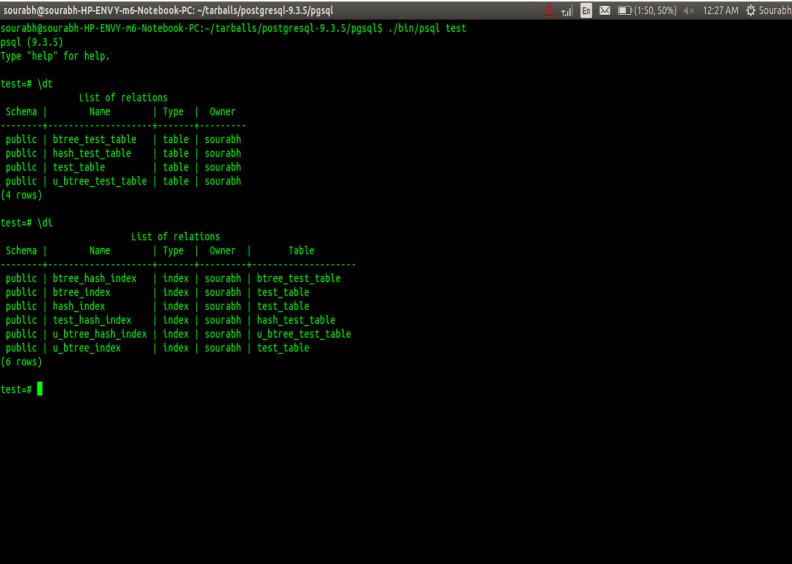
```
pg_shdepend_depender_index
pg_shdepend_reference_index
pg_shdescription_o_c_index
Time: 0.996 ms
test=# SELECT relname, relpages * 8 / 1024 AS "MB" FROM pg_class ORDER BY relpages DESC;
                                             MB
test_table
u_btree_index
btree_index
hash_index
pg_proc
pg_depend
pg_attribute
pg_toast_2618
pg_description
pg_depend_depender_index
pg_depend_reference_index
pg_proc_proname_args_nsp_index
pg_description_o_c_o_index
pg_attribute_relid_attnam_index
pg_statistic
pg_operator
pg_rewrite
pg_attribute_relid_attnum_index
pg_proc_oid_index
pg_type
pg_class
sql_features
pg_operator_oprname_l_r_n_index
pg_operator_oid_index
pg_amproc_fam_proc_index
pg_amop_opr_fam_index
pg_amop_fam_strat_index
pg_amop
pg_ts_config_map_index
pg_type_typname_nsp_index
pg_class_relname_nsp_index
pg_statistic_relid_att_inh_index
pg_amop_oid_index
```

pg\_conversion

How do they fare against each other when it comes to statements that test for equality?

We try this out for 2 types of queries. One having values present in the table and one which doesn't.

To enable simultaneous execution of queries, I copied the same table thrice and renamed them and created respective indexes as shown below.



```
Now type the following perl script and run it
(compare.pl).
#!/usr/bin/perl -w
use strict:
use Benchmark qw( cmpthese );
use DBI;
my @existing strings = (
  'PDVfHsH 2zm 3Lls ncdh Xf006',
  'VfR5L HIVWFt P2H5 yWxC6',
  'hyh ru6z19',
  'nJkC pU4V7 l4f 1ewgmxd 6BN',
  '6Mmn D1NeP IMOT2j',
  '53xWE6g fop0OM moBi zOP',
  '2EgFHwr AwpD9tt 5GY zewDok OLIE',
  '4YAQAF B87GBC7 GtIPEBC IgML awGO',
  'uxJKdN dQela AKC OlygiTL X1govd',
  'uzEFodF oFdH b9zHQ HWYQp1p',
);
my @nonexisting_strings = (
  '3zwy PJe 314P',
  'pqo6d n6j 5t1V',
  'XEyoX ciKRh01d iJrSS',
  'pa3r29P V16Gv1gy ayEv',
  'TB0f9 5TE wbP1GI',
  'T0btpG aK6R1hF x1mXlo',
```

'2T2 xdMf6k1b WUuZi'.

```
'dPV xipP4 az1F',
  't3IC Qibtx5V H186',
  'XLyiZa6 Wj1 XK1Kz',
);
my $dbh = DBI->connect(
  'dbi:Pg:dbname=test;host=127.0.0.1;port=5432', "sourabh", "",
  {
     'AutoCommit' => 1,
     'RaiseError' => 1,
     'PrintError' => 1,
  }
);
print "Testing existing keys:\n";
cmpthese(
  50000,
  {
     'existing hash' => sub
{ fetch from( 'hash test table',
                                \@existing strings); },
     'existing btree'
                          => sub
{ fetch from('btree test table', \@existing strings ); },
     'existing_unique_btree' => sub
{ fetch_from( 'u_btree_test_table', \@existing_strings ); },
  }
);
print "\nTesting nonexisting keys:\n";
```

```
cmpthese(
  50000,
  {
    'nonexisting hash' => sub { fetch from( 'hash test table',
\@nonexisting strings); },
    'nonexisting btree' => sub
{ fetch from('btree test table', \@nonexisting strings ); },
     'nonexisting unique btree' => sub
{ fetch from( 'u btree test table', \@nonexisting strings ); },
  }
);
exit;
sub fetch from {
  my ($table, $values) = @;
  my $sql = 'SELECT * FROM ' . $table . ' WHERE random_string =
ANY(?);
  my $rows = $dbh->selectall arrayref( $sql, undef, $values );
  return;
}
```

This creates a benchmark comparison between the three indexes. The results are as shown in the snapshot that follows.

sourabh@sourabh-HP-El	NVY-m6-Notebook-PC:	~/tarballs/postgresql-9.3	.5/pgsql			En 🖂	<b>■</b> (1:17, 43%) 《×	7:35 AM 🔱 Sourabh
sourabh@sourabh-HP-E Testing existing key	ENVY-m6-Notebook-PC			l\$ perl/co	mpare.pl			
Rate existing_unique_btree existing_hash existing_btree								
existing_unique_btre			-1%	- 9%				
existing_hash	3934/s	1%		-8%				
existing_btree	4277/s	10%	9%					
Testing nonexisting								
		.sting_unique_btree						
nonexisting_unique_t				-1%	- 4%			
nonexisting_btree	4888/s	1%			- 3%			
nonexisting_hash	5015/s	4%		3%				
sourabh@sourabh-HP-E	ENVY-MO-NO LEDOOK-PC	.:~/tarbatts/postgre	sqr-9.3.5/pgsq	LŞ				

#### Conclusion:

In my Comparison, i found out that Hash index creation need not always be slower as compared to B-Tree index creation. Also, in the test where we check for non existing strings, Hash index is 4%-6% faster than B-Tree index. While this may be true, it is still a marginal difference and also does not justify the usage of hash indexes over B-Tree indexes which can handle much more types of Queries.