

0.1 BINARY NUMERATION SYSTEM

Jean-Pierre Deschamps

University Rovira i Virgili, Tarragona, Spain

0.1 UAF

de Barcelona

1. Computer information representation



A computer receives, stores, processes, transmits data.

Data types: numbers, characters, sounds (audio), pictures (video), etc.

Data encoding: strings of zeroes and ones.

Computer technology is based on electronic circuits able to process vectors of 0's and 1's (the so-called digital electronic circuits). For that reason all data are encoded by strings of 0's and 1's.

This type of information encoding is called **binary encoding system**.



Most used systems:

- Decimal system
- Binary system
- Hexadecimal system

Conversion methods will be presented..





2.1 Decimal system

• Uses ten digits:

• **Positional** system: a weight is associated to every digit position so that position is relevant.

Example: 653

(weights) 10^2 10^1 10^0

$$653 = 6 \cdot 10^2 + 5 \cdot 10^1 + 3 \cdot 10^0$$

6 hundreds, 5 tens, 3 units



2.2 Binary system

- Uses two digits (binary digits, bits): 0, 1
- Positional system.

(weights)
$$2^3$$
 2^2 2^1 2^0

• To compute the decimal representation, add up the weights corresponding to the 1's of the binary representation:

$$(1101)_2 = 1 \cdot 2^3 + 1 \cdot 2^2 + 0 \cdot 2^1 + 1 \cdot 2^0 = 8 + 4 + 0 + 1 = 13$$

Exercise

Compute the decimal representation of the following binary number: (101001)₂

6

de Barcelona

Exercise (solution)



Compute the decimal representation of the following binary number: (101001)₂

weights
$$2^5$$
 2^4 2^3 2^2 2^1 2^0

$$(101001)_2 = 1 \cdot 2^5 + 0 \cdot 2^4 + 1 \cdot 2^3 + 0 \cdot 2^2 + 0 \cdot 2^1 + 1 \cdot 2^0 = 32 + 8 + 1 = 41$$

UAB Universitat Autònoma de Barcelona

2.2 Binary system: representation range

• Pure binary system: **non-negative number** representation.

• With *n* bits: **2**^{*n*} distinct values.

• Representation range: **0 to 2**ⁿ - **1**.

EXAMPLE:	Binary	Decimal	Binary	Decimal
n = 4 bits	0000	0	1000	8
16 different combinations	0001	1	1001	9
from 0 to 15 = 2^4 -1	0010	2	1010	10
110111 0 to 13 - 2 -1	0011	3	1011	11
	0100	4	1100	12
	0101	5	1101	13
	0110	6	1110	14
	0111	7	1111	15

8



2.2 Binary system: representation range

- ✓ If $n = 3 : 2^3 = 8$ representable numbers, from 0 to 7.
- ✓ If $n = 4 : 2^4 = 16$ representable, from 0 to 15.
- ✓ If $n = 5 : 2^5 = 32$ representable, from 0 to 31.
- ✓ If $n = 6 : 2^6 = 64$ representable, from 0 to 63 ...

Example: how many bits do we need to represent 48?

$$31 \le 48 \le 63$$

=> to represent decimal number 48_{10} we need 6 bits: 110000

2.3 Hexadecimal system

• Uses sixteen dígits:

Positional system

• To compute the decimal representation, add up the digits multiplied by the corresponding weights:

$$(3A9F)_{16} = 3.16^3 + 10.16^2 + 9.16^1 + 15.16^0 = 15007_{10}$$

(weights)



BINARY HEXA 0 0 0 0 0 0 0 0 1 1 0 0 1 0 2 0 0 1 1 3 0 1 0 0 4 0 1 0 0 4 0 1 0 1 5 0 1 1 0 6 0 1 1 1 7 1 0 0 0 8 1 0 0 0 8 1 0 1 0 A 1 0 1 0 A 1 0 1 1 B 1 1 0 1 D 1 1 1 0 E 1 1 1 1 1			I	Iniversitat	Autònoma
0 0 0 1 1 0 0 1 0 2 0 0 1 1 3 0 1 0 0 4 0 1 0 1 5 0 1 1 0 6 0 1 1 1 7 1 0 0 0 8 1 0 0 0 8 1 0 0 1 9 1 0 1 0 A 1 0 1 1 B 1 0 1 0 C 1 1 0 1 D 1 1 1 0 E		BIN			
0 0 1 0 2 0 0 1 1 3 0 1 0 0 4 0 1 0 1 5 0 1 1 0 6 0 1 1 1 7 1 0 0 0 8 1 0 0 0 8 1 0 0 1 9 1 0 1 0 A 1 0 1 1 B 1 0 0 C C 1 1 0 1 D 1 1 1 0 E	0	0	0	0	0
0 0 1 1 3 0 1 0 0 4 0 1 0 1 5 0 1 1 0 6 0 1 1 1 7 1 0 0 0 8 1 0 0 0 8 1 0 0 1 9 1 0 1 0 A 1 0 1 1 B 1 1 0 0 C 1 1 0 1 D 1 1 1 0 E	0	0	0	1	1
0 1 0 0 4 0 1 0 1 5 0 1 1 0 6 0 1 1 1 7 1 0 0 0 8 1 0 0 1 9 1 0 1 0 A 1 0 1 1 B 1 1 0 0 C 1 1 0 1 D 1 1 1 0 E	0	0	1	0	2
0 1 0 1 5 0 1 1 0 6 0 1 1 1 7 1 0 0 0 8 1 0 0 1 9 1 0 1 0 A 1 0 1 1 B 1 1 0 0 C 1 1 0 1 D 1 1 1 0 E	0	0	1	1	3
0 1 1 0 6 0 1 1 1 7 1 0 0 8 1 0 0 1 9 1 0 1 0 A 1 0 1 B 1 1 0 0 C 1 1 0 1 D 1 1 0 E	0	1	0	0	4
0 1 1 7 1 0 0 0 8 1 0 0 1 9 1 0 1 0 A 1 0 1 1 B 1 1 0 0 C 1 1 0 1 D 1 1 1 0 E	0	1	0	1	5
1 0 0 0 8 1 0 0 1 9 1 0 1 0 A 1 0 1 1 B 1 1 0 0 C 1 1 0 1 D 1 1 1 0 E	0	1	1	0	6
1 0 0 1 9 1 0 1 0 A 1 0 1 1 B 1 1 0 0 C 1 1 0 1 D 1 1 1 0 E	0	1	1	1	7
1 0 1 0 A 1 0 1 B 1 1 0 0 C 1 1 0 1 D 1 1 0 E	1	0	0	0	8
1 0 1 1 B 1 1 0 0 C 1 1 D 1 1 0 E	1	0	0	1	9
1 1 0 0 C 1 1 0 1 D 1 1 1 0 E	1	0	1	0	Α
1 1 0 1 D 1 1 1 0 E	1	0	1	1	В
1 1 1 0 E	1	1	0	0	C
	1	1	0	1	D
1 1 1 F	1	1	1	0	E
	1	1	1	1	F



3. Base conversion (HEXADECIMAL to BINARY) (BINARY to HEXADECIMAL)

• HEXADECIMAL to BINARY: 1 hexadecimal digit \rightarrow 4 bits.

• BINARY to HEXADECIMAL: 4 bits \rightarrow 1 hexadecimal digit (starting from the four rightmost bits)

11





	BIN	ARY		HEXA
0	0	0	0	0
0	0	0	1	1
0	0	1	0	2
0	0	1	1	3
0	1	0	0	4
0	1	0	1	5
0	1	1	0	6
0	1	1	1	7
1	0	0	0	8
1	0	0	1	9
1	0	1	0	Α
1	0	1	1	В
1	1	0	0	С
1	1	0	1	D
1	1	1	0	E
1	1	1	1	F

Compute the hexadecimal representation:

$$110110101001100_2 = ????_{16}$$

Compute the binary representation:

$$5F2C_{16} = ????_{2}$$





	BIN	ARY		HEXA
0	0	0	0	0
0	0	0	1	1
0	0	1	0	2
0	0	1	1	3
0	1	0	0	4
0	1	0	1	5
0	1	1	0	6
0	1	1	1	7
1	0	0	0	8
1	0	0	1	9
1	0	1	0	Α
1	0	1	1	В
1	1	0	0	С
1	1	0	1	D
1	1	1	0	E
1	1	1	1	F

Compute the hexadecimal representation:

$$110110101001100_2 = 6D4C_{16}$$

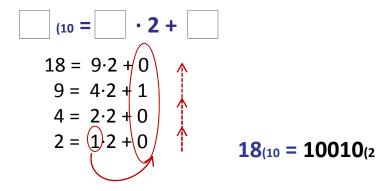
Compute the binary representation:



4. Base conversion (DECIMAL to BINARY)

- Divide the decimal number by 2. Divide the obtained quotient by 2. Keep dividing the obtained quotients by 2 until the obtained quotient is equal to 1.
- The base 2 number consists of the last quotient 1 and the set of previously obtained remainders.

Example: $18_{(10} = ?$



Exercise

43₍₁₀ = binary number?



Exercise (solution)

UAB
Universitat Autònoma
de Barcelona

$$43 = 21 \cdot 2 + 1$$

$$21 = 10 \cdot 2 + 1$$

$$10 = 5 \cdot 2 + 0$$

$$5 = 2 \cdot 2 + 1$$

$$2 = 1 \cdot 2 + 0$$



6. Sum and difference of binary numbers

Sum of 2 bits:

$$0 + 0 = 0$$

$$0 + 1 = 1$$

$$1 + 0 = 1$$

Sum example:

$$\begin{array}{c} 1 \ 1 \ 1 \\ 10 \ 10 \ 0 \ 10 \ 1 = A \\ + \ 10 \ 10 \ 11 \ 1 \\ 11 \ 11 \ 11 \ 0 \ 0 \end{array}$$

Difference of 2 bits:

$$0 - 0 = 0$$

$$1 - 0 = 1$$

$$1 - 1 = 0$$

$$0 - 1 = 11$$
 (current step bit: 1 borrow to the next step: 1)

Difference example:

$$10100101 = A$$

$$-1010111 = B$$

de Barcelona



Exercise

de Barcelona



Exercise (solution)





- Computer information representation.
- Numeration systems (decimal, binary, hexadecimal).
- Pure binary system and representation range.
- Base conversion.
- Sum and difference of binary numbers.



0.2 ALGORITHM REPRESENTATION IN PSEUDOCODE

Jean-Pierre Deschamps

University Rovira i Virgili, Tarragona, Spain





Algorithm

- Sequence of operations whose objective is the solution of some problem such as: complex computation, control of some process, etc.
- The result of the algorithm execution must be independent of the chosen type of representation.
- Some common representation methods are: natural language, pseudocode, flow diagrams and programming languages.

Pseudocode

Similar to programming language but more informal. It uses a mix of

- natural language sentences,
- programming language instructions,
- key words that define basic structures.

2. Operations and control structures

- ASSIGNMENTS
- OPERATORS
 - Comparison
 - Logic operations
 - Arithmetic operations
- SELECTION STRUCTURES (DECISIONS)
 - Simple, Double, Multiple, Case
- ITERATION STRUCTURES (CYCLES)
 - While, For
- PROCEDURES



2. 1 Assignments and types of operators

UAB Universitat Autònoma de Barcelona

variable <= expression;</pre>

ASSIGNMENTS

Assignment instruction <= : allows to store a value within a variable. Examples:

$$x \le 15$$
, $x \le 2x + y + z$, etc.

where x, y and z are variables.

COMPARISON OPERATORS

- (1) Sometimes we can use <=
- (2) Sometimes we can use >=
- (3) Sometimes we can use /=

Operator	Meaning
<	Smaller than
>	Greater than
=	Equal to
≤ ₍₁₎	Smaller than or equal to
≥ (2)	Greater than or equal to
≠ (3)	Different from

24

2. 1 Assignments and types of operators

LOGIC OPERATORS

Operator	Meaning
and	Logic product
or	Logic sum
not	Negation

ARITHMETIC OPERATORS

Operator	Meaning
+	Sum
-	Difference
*	Product
/	Division



de Barcelona

3. Control structures

SELECTION STRUCTURES

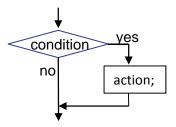
Simple

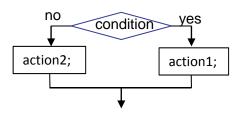
If condition **then** action(s);



Double

If condition then action1/s; else action2/s; end if;



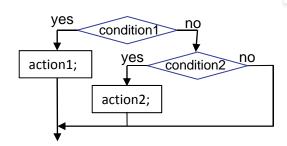




SELECTION STRUCTURES

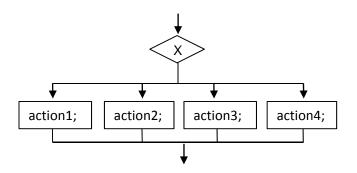
Multiple

```
If condition1 then action1/s;
        elsif condition2 then action2/s;
        end if;
end if;
```



Case

```
case x is
    when "value1" => action1;
    when "value2" => action2;
    when "value3" => action3;
    when "value4" => action4;
end case;
```



27

de Barcelona

UAB Universitat Autònoma de Barcelona

3. Control structures (example 1)

y = X/2, rounded down

Different cases:

- When X is even, the result of the division is exact (whether X is positive or negative)
- When X is odd, rounding must be calculated differently :
 - if X is negative: (X+1)/2
 - if X is positive (X-1)/2

```
If (X is even) then y<= X/2;
    elsif (X<0) then y <= (X+1)/2;
    else y <= (X-1)/2;
    end if;</pre>
```

3. Control structures (example 2)



X is a decimal digit: we calculate the binary representation

```
case x is
    when "0" => y <= 0000;
    when "1" => y <= 0001;
    ......
    when "9" => y <= 1001;
end case;</pre>
```

3. Control structures

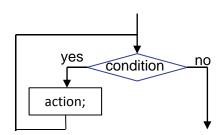
ITERATION STRUCTURES

While

While condition loop action/s; end loop;

For

for variable in min to max loop
 action/s;
end loop;



UAB
Universitat Autònoma
de Barcelona

UAB Universitat Autònoma de Barcelona

3. Control structures (example 3)

Given two vectors of 8 positions:

We want to do the following calculation:

```
y = a(0) \cdot x(0) + a(1) \cdot x(1) + a(2) \cdot x(2) + \dots + a(7) \cdot x(7)
acc <= 0;
for i in 0 to 7 loop
acc <= acc+ a(i) \cdot x(i);
end loop;
```

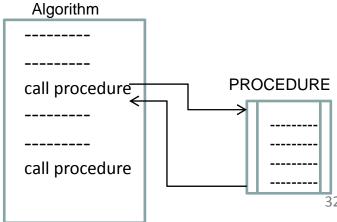
4. Procedures

UAB Universitat Autònoma de Barcelona

Procedure call name(parameters);

- **Procedure** (subroutine): sequence of instructions (operations and control structures) that execute some task (algorithm).
- A **name** and a set of **parameters** are associated to every procedure.

A procedure can be called one or several times within a program. When called, values are given to its parameters.



4. Procedures (example 1)



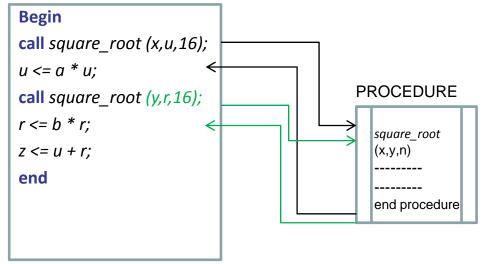
$$z = a \cdot \sqrt{x} + b \cdot \sqrt{y}$$

```
Procedure square_root(x, y, n) is
----
end procedure
```

Algorithm

```
call square_root (x, u, 16);
u <= a * u;
call square_root (y, r, 16);
r <= b * r;
z <= u + r;</pre>
```

ALGORITHM



33

SUMMARY

UAB
Universitat Autònoma
de Barcelona

- Algorithm representation
- Pseudocode
- Operations and control structures
- Procedures