## **MIT Art Design and Technology University**

## MIT School of Computing, Pune

## **Department of Computer Science and Engineering**

Second Year B. Tech

Academic Year 2022-2023. (SEM-II)

**Subject: Advance Data Structures Laboratory** 

## **Assignment 11**

**Assignment Title:** A Dictionary stores keywords & its meanings. Provide facility for adding new keywords, deleting keywords, updating values of any entry. Provide facility to display whole data sorted in ascending/ Descending order. Also find how many maximum comparisons may require for finding any keyword. Use Height balance tree and find the complexity for finding a keyword.

Aim: Implement Dictionary using AVL tree.

## **Prerequisite:**

- 1. Basic concepts of Dictionary.
- 2. Basic concepts of AVL tree and its rotations.
- 3. Procedure to create height balanced tree.

### **Objectives:**

Implement a Program in python for the following operations:

- 1. Create a Dictionary using any data structure.
- 2. Use AVL tree for performing different operations on dictionary.

#### **Outcomes:**

## Upon Completion of the assignment the students will be able to

- 1. Create dictionary using any data structure.
- 2. Understand and analyse creation of dictionary using AVL tree.

## Theory:

AVL tree is a self-balancing BST, where the difference between heights of left and right subtrees cannot be more than one for all nodes. Named after their inventors, Adelson-Velskii and Landis.

Searching of desired node is faster due to balancing of tree height.

## **Definition-**

- ➤ An AVL tree is a balanced binary search tree.
- ➤ In an AVL tree, balance factor of every node is either -1, 0 or +1.
- > Every subtree is an AVL tree.

## Balance Factor = Height of Left Subtree - Height of Right Subtree

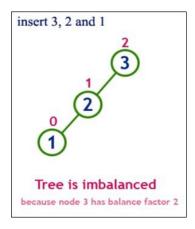
### **AVL Tree Rotations:**

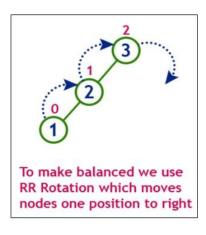
Rotation is the process of moving nodes either to left or to right to make the tree balanced.

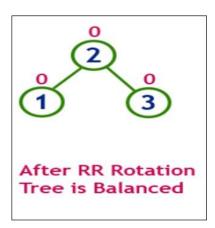
- ➤ Left Rotation (LL Rotation)
- ➤ Right Rotation (RR Rotation)
- ➤ Left Right Rotation (LR Rotation)
- ➤ Right Left Rotation (RL Rotation)

## LL Rotations - Single Right Rotation

- ➤ When tree is unbalanced because of adding left child in the left subtree of an unbalanced node, we need to use LL rotations to make balanced tree.
- In LL Rotation, every node moves one position to right from the current position.

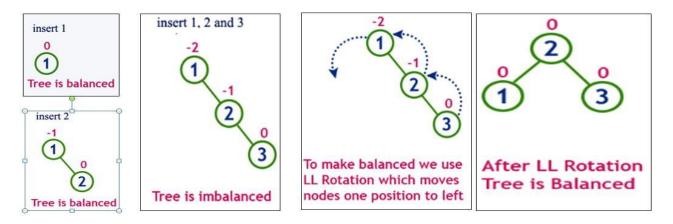






## **RR Rotations – Single Left Rotation**

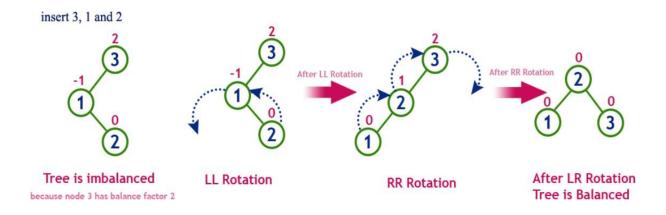
- ➤ When tree is unbalanced because of adding right child in the right subtree of an unbalanced node, we need to use RR rotations to make balanced tree.
- ➤ In RR Rotation, every node moves one position to left from the current position.



## **LR Rotations – Double Rotations**

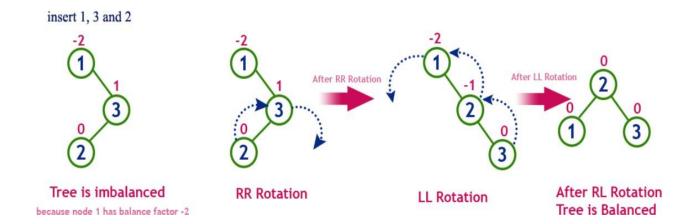
The LR Rotation is sequence of single left rotation followed by single right rotation.

- ➤ When tree is unbalanced because of adding right child in the left subtree of an unbalanced node, we need to use LR rotations to make balanced tree.
- Tree is unbalanced because of adding right child in the left subtree of an unbalanced node. The LR Rotation is sequence of single left rotation followed by single rightrotation.



#### **RL** Rotations – Double Rotations

- ➤ When tree is unbalanced because of adding left child in the right subtree of an unbalanced node, we need to use RL rotations to make balanced tree.
- ➤ The RL Rotation is sequence of single right rotation followed by single left rotation.



#### **Creation of AVL tree:**

- Insert the element in the AVL tree in the same way the insertion is performed in BST.
- After insertion, check the balance factor of each node of the resulting tree and consider following cases for same.
- > Case 1: Balance factor of each node IS {-1,0,1}
  - > The tree is balanced.
  - > Conclude the operation.
  - > Insert the next element if any.
- > Case2: Balance factor of each node is NOT {-1,0,1}
  - > The tree is imbalanced.
  - Perform the suitable **ROTATION(s)** to balance the tree.
  - After the tree is balanced, insert the next element if any.

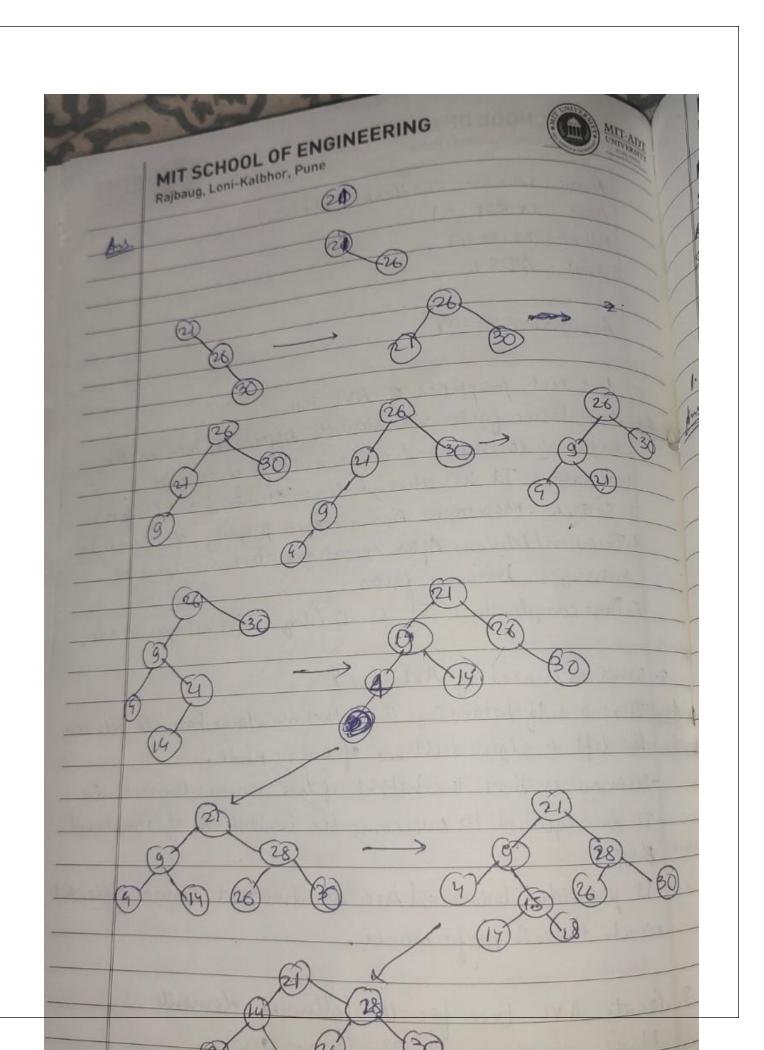
## **Conclusion:**

We have created dictionary using AVL tree.

## **Frequently Asked Questions:**

- 1. List out properties of AVL tree.
- 2. What is a need of AVL tree?
- 3. Create AVL tree for following elements: 21, 26, 30, 9, 4, 14, 28, 18, 15, 10

# MIT SCHOOL OF ENGINEERING Rajbaug, Loni-Kalbhor, Pune Name: Souran Chailesh Toshniwal CLASS: SY CSE -1 Roll no : 2213047 Subject: ADS L Assignment - 17 1. List out properties of AVI tree. 1. Balance factor: The absolute difference between the height of its left and right subtrees is at most I 2 Height: It is always logarithmic i.e o(logn) 3. Search: Maintains Binary search property of O(logn) 4. Inscrtion/ Deletion: After operations, true maintains or rearrages bollance factor 5. Time complexity: It is O (log n) for all operations. 2 What is a need of AVL true? Ins It is a self-balancing BST that maintains balance between the left & right subtrees of enery node. - It ensures that the height of tree remains logarithmic - It is required to overcome the limitations of standard - It provides balanced tree structure that ensures efficient search times & performance 3. Greate AVL tree for the following elements: 21,26,30,9,9,19,28,18,15,70



#### Code:

```
#include <iostream>
#include <cstdio>
#include <sstream>
#include <algorithm>
#define pow2(n) (1 << (n))
using namespace std;

struct avl_node
{
    int data;
    struct avl_node *left;
    struct avl_node *right;
} *root;</pre>
```

```
class avlTree
public:
   int height(avl_node *);
   int diff(avl_node *);
   avl_node *rr_rotation(avl_node *);
   avl_node *ll_rotation(avl_node *);
   avl_node *lr_rotation(avl_node *);
   avl_node *rl_rotation(avl_node *);
   avl_node *balance(avl_node *);
   avl_node *insert(avl_node *, int);
   void display(avl_node *, int);
   void inorder(avl_node *);
   void preorder(avl_node *);
   void postorder(avl_node *);
   avlTree()
        root = \overline{NULL};
```

```
int main()
   int choice, item;
   avlTree avl;
        cout << "\n----" << endl;</pre>
        cout << "AVL Tree Implementation" << endl;</pre>
        cout << "\n----" << endl;</pre>
        cout << "1.Insert Element into the tree" << endl;</pre>
        cout << "2.Display Balanced AVL Tree" << endl;</pre>
        cout << "3.InOrder traversal" << endl;</pre>
        cout << "4.PreOrder traversal" << endl;</pre>
        cout << "5.PostOrder traversal" << endl;</pre>
        cout << "6.Exit" << endl;</pre>
        cin >> choice;
        switch (choice)
        case 1:
            cout << "Enter value to be inserted: ";</pre>
            cin >> item;
            root = avl.insert(root, item);
            break;
        case 2:
```

```
if (root == NULL)
                 cout << "Tree is Empty" << endl;</pre>
            cout << "Balanced AVL Tree:" << endl;</pre>
            avl.display(root, 1);
            break;
        case 3:
            cout << "Inorder Traversal:" << endl;</pre>
            avl.inorder(root);
            cout << endl;</pre>
            break;
        case 4:
            cout << "Preorder Traversal:" << endl;</pre>
            avl.preorder(root);
            cout << endl;</pre>
            break;
        case 5:
            cout << "Postorder Traversal:" << endl;</pre>
            avl.postorder(root);
            cout << endl;</pre>
            break;
        case 6:
            exit(1);
            break;
        default:
            cout << "Wrong Choice" << endl;</pre>
    return 0;
int avlTree::height(avl_node *temp)
    int h = 0;
        int l_height = height(temp->left);
        int r_height = height(temp->right);
        int max_height = max(l_height, r_height);
        h = max_height + 1;
    return h;
int avlTree::diff(avl_node *temp)
    int l_height = height(temp->left);
    int r_height = height(temp->right);
    int b_factor = l_height - r_height;
    return b_factor;
```

```
avl_node *avlTree::rr_rotation(avl_node *parent)
{
    avl_node *temp;
    temp = parent->right;
    parent->right = temp->left;
    temp->left = parent;
    return temp;
```

```
avl_node *avlTree::ll_rotation(avl_node *parent)
   avl_node *temp;
   temp = parent->left;
   parent->left = temp->right;
   temp->right = parent;
   return temp;
avl_node *avlTree::lr_rotation(avl_node *parent)
   avl_node *temp;
   temp = parent->left;
   parent->left = rr_rotation(temp);
   return ll_rotation(parent);
avl_node *avlTree::rl_rotation(avl_node *parent)
   avl_node *temp;
   temp = parent->right;
   parent->right = ll_rotation(temp);
   return rr_rotation(parent);
avl_node *avlTree::balance(avl_node *temp)
   int bal_factor = diff(temp);
   if (bal_factor > 1)
       if (diff(temp->left) > 0)
           temp = ll_rotation(temp);
       else
           temp = lr_rotation(temp);
   else if (bal_factor < -1)</pre>
       if (diff(temp->right) > 0)
           temp = rl_rotation(temp);
           temp = rr_rotation(temp);
avl_node *avlTree::insert(avl_node *root, int value)
   if (root == NULL)
       root = new avl_node;
       root->data = value;
       root->left = NULL;
       root->right = NULL;
   else if (value < root->data)
       root->left = insert(root->left, value);
       root = balance(root);
```

```
}
else if (value >= root->data)
{
    root->right = insert(root->right, value);
    root = balance(root);
}
return root;
}
```

```
void avlTree::display(avl_node *ptr, int level)
{
    int i;
    if (ptr != NULL)
    {
        display(ptr->right, level + 1);
        printf("\n");
        if (ptr == root)
            cout << "Root -> ";
        for (i = 0; i < level && ptr != root; i++)
            cout << " ";
        cout << ptr->data;
        display(ptr->left, level + 1);
    }
}
```

```
void avlTree::inorder(avl_node *tree)
{
   if (tree == NULL)
      return;
   inorder(tree->left);
   cout << tree->data << " ";
   inorder(tree->right);
}
```

```
void avlTree::preorder(avl_node *tree)
{
   if (tree == NULL)
      return;
   cout << tree->data << " ";
   preorder(tree->left);
   preorder(tree->right);
}
```

```
void avlTree::postorder(avl_node *tree)
{
    if (tree == NULL)
        return;
    postorder(tree->left);
    postorder(tree->right);
    cout << tree->data << " ";
}</pre>
```

Output:

```
PS C:\SOURAV\CODE\C++ language codes\ADS assignment> cd "c:\SOURAV\CODE\C++ language codes\ADS assignment\"; if ($?) { g++ A
AVL Tree Implementation
1.Insert Element into the tree
2.Display Balanced AVL Tree
3.InOrder traversal
4.PreOrder traversal
5.PostOrder traversal
6.Exit
Enter your Choice: 1
Enter value to be inserted: 2
AVL Tree Implementation
5.PostOrder traversal
6.Exit
Enter your Choice: 1
Enter value to be inserted: 5
AVL Tree Implementation
1.Insert Element into the tree
2.Display Balanced AVL Tree
3.InOrder traversal
4.PreOrder traversal
5.PostOrder traversal
6.Exit
Enter your Choice: 2
Balanced AVL Tree:
Root -> 2
```