

### Exercise 1

Let  $f : A \rightarrow B$  be a function. Recall that for any sets  $X \subseteq A$  and  $Y \subseteq B$ , we defined the *image*  $f(X) = \{f(x) \mid x \in X\}$  and preimage  $f^{-1}(Y) = \{x \in A \mid f(x) \in Y\}$ .

- Prove:  $f(X \cup X') = f(X) \cup f(X')$
- If you attempt to prove  $f(X \cap X') = f(X) \cap f(X')$  (which does not hold in general) with the same approach you took in a), where does your proof attempt fail?
- Show that  $f^{-1}(f(X)) \subseteq X$  does not hold in general by giving a counterexample.
- Prove:  $f(X \cap X') = f(X) \cap f(X')$  holds for all  $X, X' \subseteq A$  if and only if  $f$  is injective.

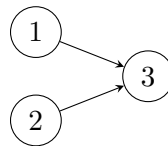
*Solution:*

- Direction  $\subseteq$ :** Let  $y$  be arbitrary with  $y \in f(X \cup X')$ . Then by definition of the image there exists  $x \in X \cup X'$  with  $f(x) = y$ . If  $x \in X$  we then have  $y \in f(X)$  and thus  $y \in f(X) \cup f(X')$ . If on the other hand  $x \in X'$  we have  $y \in f(X')$  and thus also  $y \in f(X) \cup f(X')$ .

**Direction  $\supseteq$ :** Let  $y$  be arbitrary with  $y \in f(X) \cup f(X')$ . If  $y \in f(X)$  then there exists  $x \in X$  with  $y = f(x)$  and hence also  $x \in X \cup X'$  and thus  $y \in f(X \cup X')$ . If on the other hand  $y \in f(X')$  there exists  $x \in X'$  with  $y = f(x)$  and hence also  $x \in X \cup X'$  and thus  $y \in f(X \cup X')$ .

- The direction  $\subseteq$  works just the same as before, but the direction  $\supseteq$  does not: the problem is that the assumption  $y \in f(X) \cap f(X')$  gives us  $y \in f(X)$  and  $y \in f(X')$  and from that  $x \in X$  and  $x' \in X'$  with  $f(x) = f(x') = y$ . But  $x$  and  $x'$  are not necessarily the same, so we cannot conclude  $x \in X \cap X'$ , and thus cannot move forward.
- What we need to achieve is the existence of an  $x' \in A \setminus X$  such that  $x' \in f^{-1}(f(X))$ , i.e.  $f(x') \in f(X)$ . In other words: if we draw the graph of  $f$  as a relation, then we need to find an  $x' \in A \setminus X$  that can be reached from some  $x \in X$  by taking an arrow to the right and then taking an arrow in reverse direction. That means there has to be some  $y \in B$  such that  $f(x) = y = f(x')$ .

We can easily construct such a situation by choosing e.g.  $A = \{1, 2\}$ ,  $B = \{3\}$ ,  $X = \{1\}$ , and  $f : \{1, 2\} \rightarrow \{3\}$ ,  $f(x) = 3$ . The graph representation of  $f$  as a relation then looks like this:



Now  $f(X) = \{f(1)\} = \{3\}$  and  $f^{-1}(f(X)) = \{1, 2\} \neq \{1\} = X$ .

- Direction  $\rightarrow$ :** Assume that  $f(X \cap X') = f(X) \cap f(X')$  holds for all  $X, X' \subseteq A$ . We have to prove that  $f$  is then injective. We proceed by contradiction.

Suppose  $f$  were not injective. Then there exist  $x, x' \in A$  such that  $x \neq x'$  and  $f(x) = f(x')$ . By our assumption, it holds that  $f(\{x\} \cap \{x'\}) = f(\{x\}) \cap f(\{x'\})$ . But because  $x \neq x'$  we have  $f(\{x\} \cap \{x'\}) = f(\emptyset) = \emptyset$ . Moreover, because  $f(x) = f(x')$  we have  $f(\{x\}) \cap f(\{x'\}) = \{f(x)\} \cap \{f(x')\} = \{f(x)\} \neq \emptyset$ . This is a contradiction.

**Direction  $\leftarrow$ :** Assume that  $f$  is injective. Let  $X, X' \subseteq A$  be arbitrary. We must prove that  $f(X \cap X') = f(X) \cap f(X')$ . We proceed by proving that either side is a subset of the other.

**Direction  $\subseteq$ :** Let  $y \in f(X \cap X')$ . Then there exists an  $x \in X \cap X'$  with  $f(x) = y$ . Then also  $y \in f(X)$  and  $y \in f(X')$  and thus  $y \in f(X) \cap f(X')$  follows.

**Direction  $\supseteq$ :** Let  $y \in f(X) \cap f(X')$ . Then there exist  $x \in X$  and  $x' \in X'$  with  $f(x) = f(x') = y$ . By injectivity of  $f$  we then also have  $x = x'$  and thus  $x \in X \cap X'$  and from this  $y \in f(X \cap X')$ .

**Note:** This is essentially the proof you tried to do in b), but now we *can* conclude that  $x = x'$  because of the injectivity assumption.

## Exercise 2

Let  $\Sigma$  be a finite non-empty set and  $\Sigma^*$  the set of words over  $\Sigma$ . We will look at the relation  $\sqsubseteq$  on  $\Sigma^*$  where  $w_1 \sqsubseteq w_2 \iff$  (there exist  $u, v \in \Sigma^*$  such that  $w_2 = uw_1v$ ).

That is:  $w_1 \sqsubseteq w_2$  holds if  $w_1$  is a *subword* of  $w_2$ . E.g.  $\varepsilon, abc, bc, e$  are all  $\sqsubseteq abcde$ .

- Show that  $(\Sigma^*, \sqsubseteq)$  is a poset. Is it also totally ordered?
- Draw Hasse diagrams of the following posets:
  - $([5], \leq_{\mathbb{N}})$
  - $(\{a, b\}^*, \sqsubseteq)$  with  $\sqsubseteq$  defined as above (only draw it for words up to length 3)
  - $([3] \times [3], \leq_{\text{comp}})$ , where  $\leq_{\text{comp}}$  is the product of the natural order  $\leq_{\mathbb{N}}$  on  $[3]$  with itself.
  - $(A, \sqsubseteq)$  where  $A = \{X \subseteq [4] \mid 1 \in X \text{ or } 2 \in X\}$ .
- How can you read off the minimal, maximal, least, greatest elements from a Hasse diagram? Mark them in the diagrams you drew!
- Which of these orders are total? Can you tell from the Hasse diagrams? For non-total orders, indicate two incomparable elements.

*Solution:*

- Reflexivity:** For any  $w \in \Sigma^*$ , we have  $w = \varepsilon w \varepsilon$  and thus  $w \sqsubseteq w$  (by choosing  $u = v = \varepsilon$ ).

**Transitivity:** Suppose  $w_1 \sqsubseteq w_2 \sqsubseteq w_3$ . Then there exist  $u_1, u_2, v_1, v_2$  such that  $w_2 = u_1 w_1 v_1$  and  $w_3 = u_2 w_2 v_2$ . Let  $u = u_2 u_1$  and  $v = v_1 v_2$ . Then  $w_3 = u_2 u_1 w_1 v_1 v_2 = u w_1 v$  and thereby  $w_1 \sqsubseteq w_3$  as desired.

**Antisymmetry:** Suppose  $w_1 \sqsubseteq w_2$  and  $w_2 \sqsubseteq w_1$ . We must now show that  $w_1 = w_2$ .

There are two observations that can make this proof a bit easier:

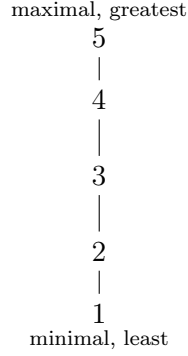
- First, if  $x \sqsubseteq y$  for two words, then  $|x| \leq |y|$ . This is obvious since if  $y = uxv$  then  $|y| = |u| + |x| + |v| \geq |x|$ .
- Second, if  $x \sqsubseteq y$  and  $|x| = |y|$  then  $x = y$ . This is also not too hard to see: if  $y = uxv$  and  $u \neq \varepsilon$  or  $v \neq \varepsilon$ , we would have  $|y| > |x|$ , which contradicts  $|x| = |y|$ . Thus  $u = v = \varepsilon$  and then  $x = y$ .

With the first of these observations,  $w_1 \sqsubseteq w_2$  and  $w_2 \sqsubseteq w_1$  directly imply  $|w_1| = |w_2|$  and then with the second observation  $w_1 = w_2$ .

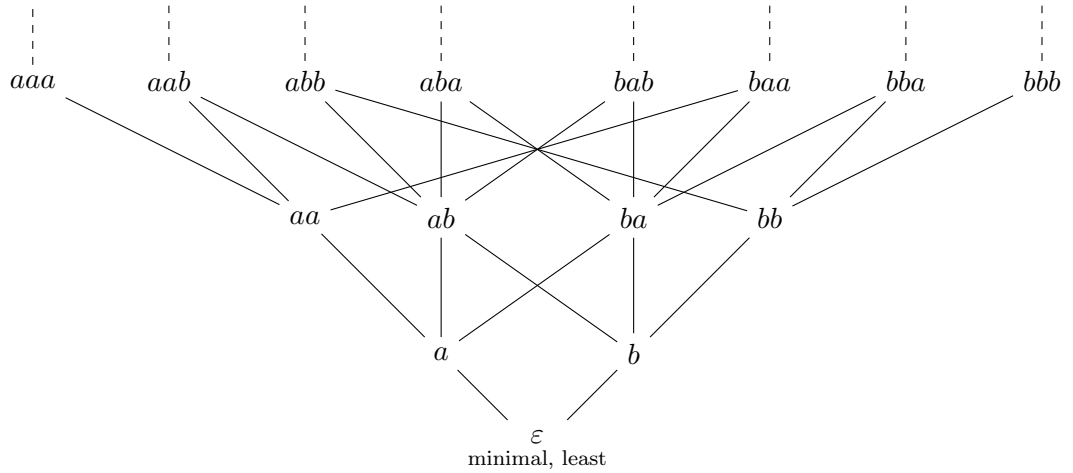
**Totality:** If  $\Sigma$  is a singleton set (say  $\Sigma = \{a\}$ ) it is a total order, since we can then write any word in the form  $a^n$  for some  $n \in \mathbb{N}$  and  $w_1 \sqsubseteq w_2$  holds if and only if  $|w_1| \leq |w_2|$ , which can easily be seen to be total.

If on the other hand  $\Sigma$  is not a singleton set (and, as we assumed, nonempty), then there exist  $a, b \in \Sigma$  with  $a \neq b$  and e.g. the words  $a$  and  $b$  are incomparable. That these two are incomparable can be seen from Observation 2 from our antisymmetry proof: since  $|a| = |b| = 1$ , if  $a \sqsubseteq b$  or  $b \sqsubseteq a$  then  $a = b$ , which is a contradiction.

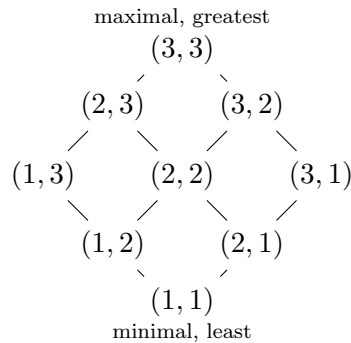
- b) i) There are no incomparable elements since this is a total order.



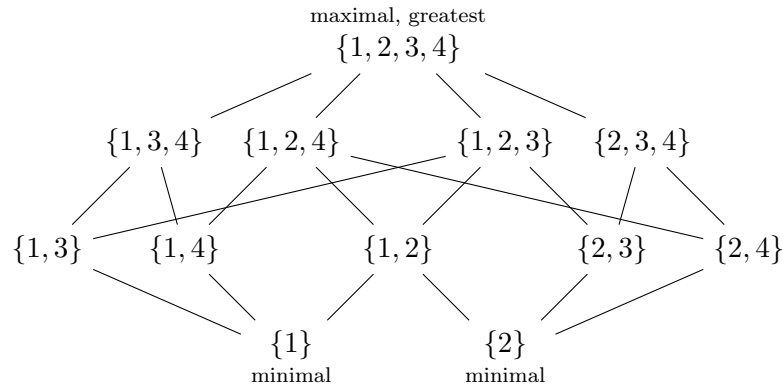
- ii) E.g.  $a$  and  $b$  or  $aa$  and  $bab$  are incomparable.



- iii) E.g.  $(1, 2)$  and  $(2, 1)$  or  $(1, 2)$  and  $(3, 1)$  are incomparable.



iv) E.g.  $\{1, 3\}$  and  $\{1, 2\}$  or  $\{1, 3\}$  and  $\{2, 3, 4\}$  are incomparable.



- c) The minimal elements are those that have no edges going down. A least element  $x$  additionally has to be connected to every other element, i.e. for every other element  $y$  there must be a path from  $x$  to  $y$  that only ever goes up.

For maximal/greatest element, the same holds switching all occurrences of ‘up’ and ‘down’.

- d) Only i) is total. There, the elements are arranged in one long line so that for every pair of elements  $x$  and  $y$ , we can reach one from the other by only going up or only going down. This is also why total orders are sometimes called *linear orders*.

In the others we have pairs of elements that cannot be reached from one another just by going only up or only down. Such pairs of elements are incomparable.

### Exercise 3

The following are two different posets, but they are also ‘the same’ in some sense because one can obtain one from the other by simply renaming elements ( $1 \mapsto 3, 2 \mapsto 1, 3 \mapsto 4, 4 \mapsto 2$ ):



While the specific elements are different, the orders themselves have ‘the same shape’. One says that these two posets are *isomorphic*. If we only care about different ‘isomorphism types’, we can draw the above partial order e.g. like this:



Draw diagrams like these for all posets of size  $n$  for  $n = 0, 1, 2, 3, 4$ . Be careful not to draw the same isomorphism type twice by accident!

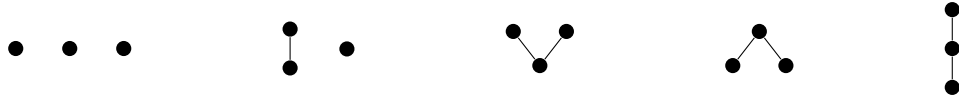
**Hint:** Just so that you can be sure not to have missed one: the number of posets up to isomorphism for  $n = 0, 1, 2, 3, 4$  is 1, 1, 2, 5, 16. (Sequence A000112 on OEIS)

*Solution:* The only partial order on  $\emptyset$  is the empty relation (whose diagram is empty). The only partial order on a singleton set  $\{x\}$  is  $\text{Id}_{\{x\}}$ , whose diagram is just a single  $\bullet$ . As for the rest:

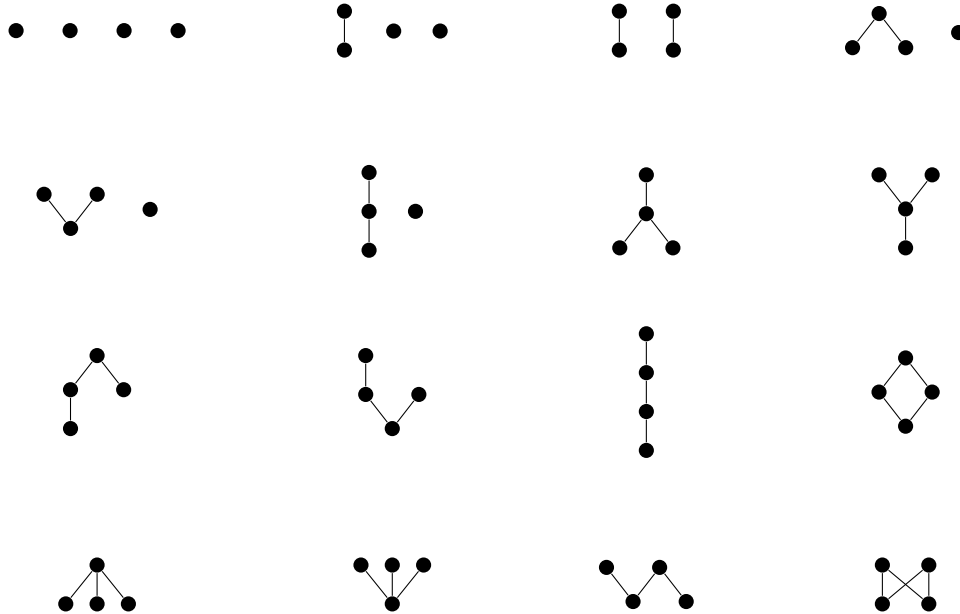
- $n = 2$ :



- $n = 3$ :



- $n = 4$ :



One good strategy to come up with all of these for some size  $n$  is to first list all the ones that are simply disjoint unions of smaller partial orders (see the Bonus Exercise for a precise definition of what a disjoint union of two orders is).

E.g. for  $n = 4$ , we can decompose  $4 = 1 + 1 + 1 + 1$  or  $4 = 2 + 1 + 1$  or  $4 = 2 + 2$  or  $4 = 3 + 1$ . This accounts for the first 6 diagrams.

Next, we look at all diagrams of size 3 that are not decomposable into disjoint unions of smaller orders (there are 3 of those) and go through all the different places where we can add a new element. However, one has to be very careful about eliminating duplicates (it may be difficult to tell at first glance if two diagrams are actually the same).

### Bonus exercise

- Show that if  $(A, \preceq_1)$  and  $(B, \preceq_2)$  are posets that have a least element, then their product order  $(A \times B, \preceq_{\text{comp}})$  also has a least element.
- Let  $\preceq_1$  and  $\preceq_2$  be preorders on  $A$  and  $B$  with  $A \cap B = \emptyset$ . We define the preorder  $\preceq_{1 \uplus 2}$  as  $\preceq_{1 \uplus 2} = \preceq_1 \cup \preceq_2$ . That is:

$$x \preceq_{1 \uplus 2} y \iff (x, y \in A \text{ and } x \preceq_1 y) \text{ or } (x, y \in B \text{ and } x \preceq_2 y)$$

We call  $\preceq_{1 \uplus 2}$  the *disjoint union* of  $\preceq_1$  and  $\preceq_2$ .

Is  $\preceq_{1 \uplus 2}$  a partial order if  $\preceq_1$  and  $\preceq_2$  are partial orders? What about total? What are the minimal, maximal, least, greatest elements of  $\preceq_{1 \uplus 2}$ ? What does its Hasse diagram look like if  $\preceq_1$  and  $\preceq_2$  are partial orders?

- c) In the lecture we showed that if a partial order  $\preceq$  on a finite set has exactly one minimal element  $x$ , then  $x$  is also least. Show that this is not true if we do not assume finiteness.

*Solution:*

- a) Let  $x_1$  and  $x_2$  be least elements of  $(A, \preceq_1)$  and  $(B, \preceq_2)$ , respectively. Let  $(y_1, y_2) \in A \times B$  be arbitrary. Then we have  $x_1 \preceq_1 y_1$  and  $x_2 \preceq_2 y_2$  because  $x_1$  and  $x_2$  are least in their respective orders. Then, by definition of the product order,  $(x_1, x_2) \preceq_{\text{comp}} (y_1, y_2)$ .
- b) If  $A$  is empty, then  $\preceq_{1 \uplus 2}$  is just  $\preceq_2$  and thus inherits all its properties (including whether it is a partial/total order), and analogously if  $B$  is empty.

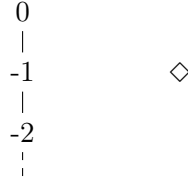
Otherwise, if  $A$  and  $B$  are both nonempty, the disjoint union order can never be total since any element from  $A$  is incomparable to any element from  $B$  by definition. For the same reason, there also cannot be any least or greatest elements (least/greatest elements must be comparable to every element). It is however a partial order (as you can check easily).

As for minimal/maximal elements, it is easy to see that any minimal element of  $\preceq_1$  is also minimal in  $\preceq_{1 \uplus 2}$ : Let  $x \in A$  be minimal in  $\preceq_1$ . To show that it is also minimal in  $\preceq_{1 \uplus 2}$ , let  $y \in A \cup B$  with  $y \prec_{1 \uplus 2} x$ . Then by definition of  $\preceq_{1 \uplus 2}$  (and the fact that  $x \in A$  and  $A \cap B = \emptyset$ ) we have  $y \in A$  and  $y \prec_1 x$ . But that is not possible because  $x$  is minimal in  $\prec_1$ .

By duality, it follows that maximal elements of  $\preceq_1$  are also preserved. Lastly, due to the commutativity of  $\cup$  and the definition of the disjoint union order, the minima and maxima of  $\preceq_2$  are then also preserved.

If  $\preceq_1$  and  $\preceq_2$  are preorders, the Hasse diagram of  $\preceq_{1 \uplus 2}$  can be obtained simply by writing those of  $\preceq_1$  and  $\preceq_2$  next to one another independently.

- c) Let  $\diamond$  be some object that is  $\notin \mathbb{Z}$ .  
Consider the poset  $(\mathbb{Z}_{\leq 0} \cup \{\diamond\}, \preceq)$  where  $\mathbb{Z}_{\leq 0} = \{0, -1, -2, \dots\}$  and  $\preceq$  as follows:



It is easy to see from this picture that  $\diamond$  is the only minimal element, but it is clearly not least (since it is incomparable to e.g. 0).

$\preceq$  can be formally defined as:

$$x \preceq y \iff x = y = \diamond \text{ or } (x, y \in \mathbb{Z} \text{ and } x \leq y)$$

That is: all the integers are ordered by their natural ordering in  $\mathbb{Z}$  and  $\diamond$  is incomparable to the integers.

An alternative view of this argument is that our  $\preceq$  is simply the disjoint union of the natural order  $(\mathbb{Z}_{\leq 0}, \leq)$  and the trivial order  $(\{\diamond\}, \text{id}_{\{\diamond\}})$ . And as we have seen in c), such an order inherits all the minimal/maximal elements from its constituent orders (in this case the minimal element  $\diamond$  and the maximal element 0) but has no least/greatest elements.