Optimal Broadcast Auctions with Costly Actions

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1. INTRODUCTION

2. COST MODEL AND SETTINGS

DEFINITION 1. In our settings, there's only one seller selling one item to n i.i.d. bidders whose valuation distribution (CDF) is continuous over [0,1]. The seller can broadcast a message to all bidders with a cost b (the broadcast cost). A bidder may remain silent after that broadcast with no cost. If not, it costs c_0 if the bidder's bidding(reply) doesn't imply positive valuation. Otherwise, if the bidding implies a positive valuation, a greater bidding cost $c > c_0$ is charged. Such bidding cost c main contain two parts, c_1, c_2 ($c_1 + c_2 = c$). The first part c_1 is charged to the seller while the second is charged to the corresponding bidder.

3. OPTIMAL MECHANISMS WITH EFFI-CIENCY AND ONLY SELLER'S COST

3.1 Optimality of Multi-round Vickrey Auctions (MVA)

Theorem 1. Suppose there's one seller selling one item to n i.i.d. bidders whose valuation distribution (CDF) is continuous over [0,1]. The seller can broadcast a message to all bidders with a cost b. A bidder may remain silent after that broadcast with no cost. If not, it costs c_0 if the bidder's bidding(reply) doesn't imply positive valuation. Otherwise, if the bidding implies a positive valuation, a greater bidding cost $c_1 > c_0$ is charged. Among all mechanisms that can include multiple rounds of broadcasts and are required to be efficient (allocate the item to the bidder with highest valuation), Multi-round Vickrey Auctions (MVAs) are of minimum cost.

COROLLARY 1. If all broadcast costs and bidding costs are charged to sellers, MVA gives the maximum utility to sellers if efficiency is required. Such optimal MVA is the one that minimizes the overall cost.

3.2 Cost Minimized α-MVA

Appears in: Proceedings of the 11th International Conference on Autonomous Agents and Multiagent Systems (AA-MAS 2012), Conitzer, Winikoff, Padgham, and van der Hoek (eds.), June, 4–8, 2012, Valencia, Spain.

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3.3 Approximation of α and Experiments

4. OPTIMAL MECHANISMS WITH BOTH SELLER'S AND BIDDER'S COST

4.1 Payment Equivalence Theorem and Revenue Optimization Strategy

THEOREM 2. If two mechanisms satisfy:

- 1. They only allocate the item to bidders whose valuation are at least l (the lowest type is l)
- 2. If they will allocate the item, they will always allocate the item to the bidder with highest valuation.

They will have the same overall payment from all bidders (including bidders' bidding costs and payments to sellers).

COROLLARY 2. For mechanisms that have a fixed lowest type l and will always allocate the item to the bidder with highest valuation if it's at least l, the maximum utility for sellers is achieved when the mechanism minize the cost.

Theorem 3. Under the same circumstances of theorem 1, except that the efficiency constraint is replaced by lowest type l and the bidding cost is now also charged to corresponding bidder, MVAs are of the minimum cost among all those mechanisms that have a fixed lowest type l and will always allocate the item to the bidder with highest valuation if it's at least l.

COROLLARY 3. MVAs are optimal in terms of maximizing seller's utility.

COROLLARY 4. Shall we increase lowest type l a little above Myerson's optimal lowest type to trade payment with cost? Or is it optimal to use Myerson's l and then minimize cost according to that?

- **4.2** Experiments to Discover Optimal MVA with a Given Lowest Type
- 4.3 Analysis of Optimal MVA with Lowest Type
- 4.4 Using Piecewise Linear MVA to Approximate Optimal MVA
- 4.5 Choosing Lowest Type?

5. CONCLUSION