# Algorithms in Invariant Theory

With 0 Figures

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# **Preface**

A core interest of mine (at the time of writing this) is algorithms in the context of commutative algebra and algebraic geometry. As such, Bernd Sturmfel's text *Algorithms in Invariant Theory* [Str08] is a good resource for first learning this stuff. Perhaps in the future I'll include notes and some source code

### Chapter 1

### Introduction

### 1.1 Symmetric Polynomials

**Solution:** [Str08] 1.1.5: From [Page 23, 2.11'][Mac98], we have the following identities

$$\begin{aligned} 1 \cdot \sigma_1 &= p_1, \\ 2 \cdot \sigma_2 &= p_1 \sigma_1 - p_2, \\ 3 \cdot \sigma_3 &= p_1 \sigma_2 - p_2 \sigma_1 + p_3, \\ &\vdots \\ k \cdot \sigma_k &= \sum_{r=1}^k (-1)^{r-1} p_r \sigma_{k-r}. \end{aligned}$$

Treating the  $\sigma_i$  as indeterminants, we can re-express the above system of equations:

$$\begin{split} p_1 &= 1 \cdot \sigma_1, \\ p_2 &= p_1 \sigma_1 - 2 \cdot \sigma_2, \\ p_3 &= p_2 \sigma_1 - p_1 \sigma_2 + 3 \cdot \sigma_3, \\ &\vdots \\ p_k &= \left(\sum_{r=1}^{k-1} (-1)^{r-1} p_{k-r} \sigma_r\right) + (-1)^{k-1} \cdot k \sigma_k. \end{split}$$

Consider  $\sigma_1, -\sigma_2, \sigma_3, \dots, (-1)^n \sigma_{n-1}, \sigma_n$  as indeterminants. Thus, we obtain the following matrix equation:

$$\begin{pmatrix} 1 & 0 & 0 & \cdots & 0 \\ p_1 & 2 & 0 & \cdots & 0 \\ p_2 & p_1 & 3 & \cdots & 0 \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ p_{k-1} & p_{k-2} & \cdots & \cdots & (-1)^k \cdot k \end{pmatrix} \begin{pmatrix} \sigma_1 \\ -\sigma_2 \\ \vdots \\ (-1)^k \sigma_{k-1} \\ \sigma_k \end{pmatrix} = \begin{pmatrix} p_1 \\ p_2 \\ \vdots \\ p_{k-1} \\ p_k \end{pmatrix}.$$

Then, Cramer's rule yields that

$$\sigma_{k} = \det \begin{pmatrix} 1 & 0 & 0 & \cdots & p_{1} \\ p_{1} & 2 & 0 & \cdots & p_{2} \\ p_{2} & p_{1} & 3 & \cdots & p_{3} \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ p_{k-1} & p_{k-2} & \cdots & \cdots & p_{k} \end{pmatrix} \det \begin{pmatrix} 1 & 0 & 0 & \cdots & 0 \\ p_{1} & 2 & 0 & \cdots & 0 \\ p_{2} & p_{1} & 3 & \cdots & 0 \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ p_{k-1} & p_{k-2} & \cdots & \cdots & (-1)^{k} \cdot k \end{pmatrix}^{-1}$$

$$= \frac{(-1)^{k}}{k!} \det \begin{pmatrix} 1 & 0 & 0 & \cdots & p_{1} \\ p_{1} & 2 & 0 & \cdots & p_{2} \\ p_{2} & p_{1} & 3 & \cdots & p_{3} \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ p_{k-1} & p_{k-2} & \cdots & \cdots & p_{k} \end{pmatrix}$$

$$= \frac{(-1)^{k}}{k!} \cdot (-1)^{k} \det \begin{pmatrix} p_{1} & 1 & 0 & \cdots & 0 \\ p_{2} & p_{1} & 2 & \cdots & 0 \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ p_{k-1} & p_{k-2} & \cdots & p_{1} & k-1 \\ p_{k} & p_{k-1} & \cdots & \cdots & p_{1} \end{pmatrix} = \frac{1}{k!} \det \begin{pmatrix} p_{1} & 1 & 0 & \cdots & 0 \\ p_{2} & p_{1} & 2 & \cdots & 0 \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ p_{k-1} & p_{k-2} & \cdots & p_{1} & k-1 \\ p_{k} & p_{k-1} & \cdots & \cdots & p_{1} \end{pmatrix}$$

#### 1.2 Gröbner Bases

### Lemma 1.2.1:

Let  $R = \mathbb{C}[x_1, \dots, x_n]$ . Then with the usual grading, let  $H(R, z) \coloneqq \sum_{d=0}^{\infty} \dim_{\mathbb{C}}(R_d) z^d$ . We have that

$$H(R,z) := \sum_{d=0}^{\infty} \dim_{\mathbb{C}}(R_d) z^d = \sum_{d=0}^{\infty} \binom{d+n-1}{n-1} = \frac{1}{(1-z)^n}.$$

**Proof:** To see that  $H(R,z) = \sum_{d=0}^{\infty} {d+n-1 \choose n-1}$ , just count the number of monomials of degree d in n variables. The value  ${d+n-1 \choose n-1}$  is the number of non-negative integer solutions to  $a_1+\cdots+a_n=d$ . Each solution corresponds to a monomial  $x_1^{a_1}\cdots x_n^{a_n}$ . Then to see that  $H(R,z)=\frac{1}{(1-z)^n}$ , consider the product of infinite sums  $(1+z+z^2+\cdots)\cdots(1+z+z^2+\cdots)$  a total of n-times. Then the coefficient of  $z^d$  again corresponds to the number of such non-negative integer solutions. Since  $\frac{1}{1-z}=1+z+z^2+\cdots$ , we obtain the desired equality.

### Lemma 1.2.2:

For  $1 \le k \le n$ , we have that

$$h_k(x_k,...,x_n) + \sum_{i=1}^k (-1)^i h_{k-i}(x_k,...,x_n) \sigma_i(x_1,...,x_n) = 0.$$

**Proof**: Using the generating functions for the  $h_i$  and  $\sigma_i$ , we have that the above expression is the coefficient of  $t^k$  in the product

$$\prod_{i=k}^{n} (1 - x_i t)^{-1} \cdot \prod_{i=1}^{n} (1 - x_i t) = \prod_{i=1}^{k-1} (1 - x_i t).$$

However, the right-hand side of this has degree k-1 in t. Thus, the coefficient of  $t^k$  is indeed 0.

**Solution:** [Str08] 1.2.1: Let M be a set of monomial generators for  $\operatorname{init}(I)$  and let m be minimally nonstandard. Since m is a monomial and in  $\operatorname{init}(I)$ , we have that  $m' \mid m$  for some monomial  $m' \in M$ . However, note that  $m' \in \operatorname{init}(I)$  and furthermore by the fact that m is minimally nonstandard, we cannot have that m' strictly divides m. Thus, m' = m and  $m \in M$ . Then by Dickson's Lemma, we have that M is a finite set and thus there are finitely many minimally nonstandard monomials.

**Solution:** [Str08] 1.2.2: This is [CLO15, Chapter 2,  $\S$ 7, Theorem 5].

**Solution:** [Str08] 1.2.3: This is [CLO15, Chapter 3,  $\S$ 1, Theorem 2].

**Solution:** [Str08] 1.2.4: This is following [Rob85] and [GP07]. Let  $\geq$  be a monomial ordering on  $\mathbb{C}[x_1,\ldots,x_n]$ . This is equivalent to a total semigroup ordering  $\geq$  on  $\mathbb{Z}^n$ . Such a semigroup ordering gives a unique total ordering on  $\mathbb{Q}^n$ . To see this, for  $\overline{q}=(q_1,\ldots,q_n)\in\mathbb{Q}^n$ , let  $m\in\mathbb{Z}$  such that  $m\cdot q_i\in\mathbb{Z}$  for all i. Then say that  $\overline{q}\geq 0$  if and only if  $m\cdot \overline{q}\geq 0$  where the latter ordering is in  $\mathbb{Z}^n$ .

Let  $V \subseteq \mathbb{Q}^n$  be a  $\mathbb{Q}$ -vector space with  $\dim_{\mathbb{Q}}(V) = r$ . Then let

$$V_0 := \{ z \in \mathbb{R}^n \mid \forall \varepsilon > 0, \exists z_+(\varepsilon), z_-(\varepsilon) \in V \cap B_{\varepsilon}(z) \text{ such that } z_+(\varepsilon) > 0, z_-(\varepsilon) < 0 \}.$$

Then  $V_0$  is clearly a  $\mathbb{R}$ -subspace of  $\mathbb{R}^n$ . With the ordering  $\geq$  on  $\mathbb{Q}^n$ , we can define  $V_+$  and  $V_-$  depending on if  $\overline{q} \geq 0$  or  $\overline{q} < 0$ . We define a map  $\pi \colon V \to \{-1,1\}$ , where V has the Euclidean topology and  $\{-1,1\}$  has the discrete topology. Let  $\pi(q) = 1$  if there exists an open ball  $U_{\varepsilon}(q)$  such that  $U_{\varepsilon}(q) \cap V \subseteq V_+$  and  $\pi(q) = -1$  if there exists an open ball  $U_{\varepsilon}(q)$  such that  $U_{\varepsilon}(q) \cap V \subseteq V_-$ . Then  $\pi$  is continuous. Recall that topological vector spaces over  $\mathbb{R}$  are connected. Thus, we cannot have that  $\dim_{\mathbb{R}} V_0 < r-1$  as if it were, then  $V_{\mathbb{R}} \setminus V_0$  would be connected  $\langle \langle \text{why is this bad?} \rangle \rangle$ . Then suppose that  $\dim_{\mathbb{R}} V_0 = r$ . Then we have an ordered basis  $e_1, \ldots, e_r$  such that  $e_i > 0$  for all i. But then the linear combinations of the  $e_i$  with positive coefficients are a subspace of  $V_+$  which is a contradiction  $\langle \langle \text{why?} \rangle \rangle$ .

To construct the first row of the matrix, start with  $V = \mathbb{Q}^n$  and consider the obtained  $V_0$ . Then the dimension 1 subspace orthogonal to  $V_0$  in  $\mathbb{R}^n$  defines the first row of A. We can continue this construction inductively to obtain the full matrix A.

**Solution:** [Str08] 1.2.6: First, we recall some definitions. The S-polynomial of f and g is

$$S(f,g) := \frac{\operatorname{lcm}(\operatorname{LM}(f),\operatorname{LM}(g))}{\operatorname{LT}(f)} f - \frac{\operatorname{lcm}(\operatorname{LM}(f),\operatorname{LM}(g))}{\operatorname{LT}(g)} g.$$

For a set of polynomials  $\mathcal{F} = \{f_1, \dots, f_t\} \subseteq k[x_1, \dots, x_n]$ , we write  $f \to_{\mathcal{F}} 0$  if there exists  $a_1, \dots, a_t \in k[x_1, \dots, x_n]$  such that  $a_1f_1 + \dots + a_tf_t = 0$ . Then [CLO15, Chapter 2, §9, Theorem 3] says that a basis  $\mathcal{F} = \{f_1, \dots, f_t\}$  is a Gröbner basis for G if and only if  $S(f_i, f_j) \to_{\mathcal{F}} 0$  for all  $i \neq j$ . But [CLO15, Chapter 2, §9, Proposition 4] says that for  $f, g \in \mathcal{F}$  with relatively prime initial monomials, we have that  $S(f, g) \to_{\mathcal{F}} 0$ . This proves the claim.

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