# Using Algebraic Geometry

With 0 Figures

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## **Preface**

At the time of writing this, I am starting my PhD at The Ohio State University. Currently a large part of my interests in algebra are about algorithms as they relate to polynomials and algebraic geometry. I've been doing a bunch of problems from *Ideals, Varieties, and Algorithms* [CLO15]. However, it seems that *Using Algebraic Geometry* [CLO05] Ex. moves through the material faster as it assumes you know more algebra. So I've moved onto working through this book as well as trying to comprehend Sturmfel's *Algorithms in Invariant Theory* [Str08].

### Chapter 1

### Introduction

### 1.1 Polynomials and Ideals

Solution: [CLO05] Ex. 1.1.1:

- (a) We have that  $x(x y^2) + y(xy) = x^2 xy^2 + xy^2 = x^2$ .
- (b) It suffices to check for generators. We have that  $x + (-1)(y^2) = x y^2$ , y(x) = xy, and  $y^2 = y^2$  showing that  $\langle x y^2, xy, y^2 \rangle \subseteq \langle x, y^2 \rangle$ . Then  $x y^2 + y^2 = x$  and  $y^2 = y^2$  shows the reverse containment and overall the ideals are equal.
- (c) We already know from 1. that  $x^2$  lives in  $\langle x-y^2, xy \rangle$ . Since xy=xy, we overall have that  $\langle x^2, xy \rangle \subseteq \langle x-y^2, xy \rangle$ . It remains to check if  $x-y^2 \in \langle x^2, xy \rangle$ . However, notice that every element of  $\langle x^2, xy \rangle$  is divisible by x while  $x-y^2$  is clearly not divisible by x. Thus  $x-y^2 \notin \langle x^2, xy \rangle$  and the two ideals are not equal.

Solution: [CLO05] Ex. 1.1.2:

Let  $f,g \in \langle f_1,\ldots,f_s \rangle$ . Then  $\exists p_1,\ldots,p_s,q_1,\ldots,q_s$  such that  $f=\sum_{i=1}^s p_i \cdot f_i$  and  $g=\sum_{i=1}^s q_i \cdot f_i$ . Thus  $f+g=\sum_{i=1}^s (p_i+q_i) \cdot f_i$  which shows that  $f+g \in \langle f_1,\ldots,f_s \rangle$ . Then let  $p \in k[x_1,\ldots,x_n]$ . We have that  $p \cdot f=p\sum_{i=1}^s p_i f_i=\sum_{i=1}^s (p \cdot p_i) \cdot f_i$  which shows that  $\langle f_1,\ldots,f_s \rangle$  is an ideal.

Solution: [CLO05] Ex. 1.1.3:

We already know that  $\langle f_1,\ldots,f_s\rangle$  is an ideal by [CLO05] Ex. 1.1.2. Now suppose that J is an ideal containing  $\{f_1,\ldots,f_s\}$ . Then, since ideals are closed under addition and scaling, we have that for all  $p_1,\ldots,p_s\in k[x_1,\ldots,x_n]$  that  $\sum_{i=1}^s p_i\cdot f_i\in J$ . Thus,  $\langle f_1,\ldots,f_s\rangle\subseteq J$ .

#### Solution: [CLO05] Ex. 1.1.4:

We claim that  $\langle f_1,\ldots,f_s\rangle=\langle g_1,\ldots,g_t\rangle$  if and only if  $\{g_1,\ldots,g_t\}\subseteq I$  and  $\{f_1,\ldots,f_s\}\subseteq J$ . The forward implication is immediate. Then by [CLO05] Ex. 1.1.3, if  $\{g_1,\ldots,g_t\}\subseteq I$  then  $J\subseteq I$ . Similarly,  $\{f_1,\ldots,f_s\}\subseteq J$   $I\subseteq J$  and overall I=J. This fact was used in [CLO05] Ex. 1.1.1 (b).

### Solution: [CLO05] Ex. 1.1.5:

It suffices to show that  $z-x^3 \in \langle y-x^2, z-xy \rangle$  and and  $z-xy \in \langle x-y^2, z-x^3 \rangle$ . Indeed we have that  $(z-xy)+x(y-x^2)=z-x^3$  which also yields that  $z-xy=z-x^3-x(y-x^2)$ .

### Solution: [CLO05] Ex. 1.1.6:

If  $I = \{0\}$  then  $I = \langle 0 \rangle$ . So suppose  $I \neq 0$ . Let  $d \in I$  be of minimal degree.  $\langle d = \gcd(I) \text{ but I need} \rangle$  infinite Bezout. Then we claim that  $\langle d \rangle = I$ . Since  $d \in I$ , we have that  $\langle d \rangle \subseteq I$ . Now let  $f \in I$ . By Euclidean division, there exists  $q, r \in k[x]$  such that f = qd + r where either r = 0 or  $0 \leq \deg(r) \leq \deg(d) - 1$ . If r = 0 then  $f \in \langle d \rangle$  and we are done. So suppose  $r \neq 0$ . Then  $f, qd \in I \implies r = f - qd \in I$ . Thus,  $r \in I$  is of degree strictly less than d, contradicting the minimality of the degree of d. So we must have that r = 0 and overall  $\langle d \rangle = I$ .

### Solution: [CLO05] Ex. 1.1.7:

- (a) Suppose  $f(x) \in \langle x \rangle$ . Then  $f(x)^m \in \langle x^n \rangle$  so  $f(x) \in \sqrt{\langle x^n \rangle}$  Now suppose that  $f(x) \in \sqrt{\langle x^n \rangle}$ . Then  $\exists k$  such that  $f(x)^k \in \langle x^n \rangle$ . Thus  $f(x)^k$  is a multiple of  $x^n$ . This implies that  $f(x)^k$  is a multiple of x. Then notice that the unique factorization of  $f(x)^k$  into irreducibles is the kth power of the factorization of f(x) into irreducibles. Thus x must be a factor of f(x) and so  $f(x) \in \langle x \rangle$ . Note, this heavily uses the fact that k[x] is a unique factorization domain for all fields k.
- (b) We claim that  $\sqrt{\langle p(x)\rangle} = \langle (x-a_1)\cdots(x-a_m)\rangle = I$ . Suppose  $f(x) \in I$ . Let  $k = \max e_1, \dots, e_n$ . Then  $p(x) \mid f(x)^k$  so  $f(x) \in \sqrt{\langle p(x)\rangle}$ . Now suppose that  $f(x) \in \sqrt{\langle p(x)\rangle}$ . Then  $\exists k$  such that  $f(x)^k \in \langle p(x)\rangle$ . Thus  $f(x)^k$  is a multiple of each  $(x-a_i)$ . Then notice that the unique factorization of  $f(x)^k$  into irreducibles is the kth power of the factorization of f(x) into irreducibles. Thus f(x) is a multiple of each  $(x-a_i)$  and so  $f(x) \in I$ .
- (c) Radical ideals are the ideals I such that  $\sqrt{I} = I$ . Notice that  $\mathbb{C}[x]$  is a principal ideal domain and so any such I must be generated by a single polynomial. Since every polynomial in  $\mathbb{C}[x]$  splits into linear factors, (b) immediately implies that the only radical ideals of  $\mathbb{C}[x]$  are the ones which are of the form  $\langle (x-a_1)\cdots(x-a_m)\rangle$  for  $a_1,\ldots,a_m\in\mathbb{C}[x]$ .

### Solution: [CLO05] Ex. 1.1.8:

- (a) Let  $\mathfrak{p}$  be a prime ideal in  $k[\overline{x}]$ . Clearly  $\mathfrak{p} \subseteq \sqrt{\mathfrak{p}}$  always. Let  $f(\overline{x}) \in \sqrt{\mathfrak{p}}$ . Then  $f(\overline{x})^m \in \mathfrak{p}$  for some  $m \in \mathbb{Z}_{\geq 1}$ . We prove the reverse inclusion by induction on m. If m = 1 then  $f(\overline{x}) = f(\overline{x})^1 \in \mathfrak{p}$ . Now let m > 1 and suppose the claim holds for all  $k \leq m$ . Then suppose  $f(\overline{x})^{m+1} \in \mathfrak{p}$ . Then  $f(\overline{x}) \cdot f(\overline{x})^m \in \mathfrak{p}$  Either  $f(\overline{x}) \in \mathfrak{p}$  or  $f(\overline{x})^m \in \mathfrak{p}$  which by induction implies that  $f(\overline{x}) \in \mathfrak{p}$ . Thus,  $f(\overline{x})^m \in \mathfrak{p} \implies f(\overline{x}) \in \mathfrak{p}$  for all  $m \in \mathbb{Z}_{\geq 1}$  and so  $\sqrt{\mathfrak{p}} \subseteq \mathfrak{p}$ . Thus, all prime ideals are radical.
- (b) Notice that for all fields k that k[x] is a principal ideal domain. Thus, all the prime ideals are the ones generated by a single irreducible polynomial. Also, in k[x] we have that (0) is a prime ideal as well as k[x] is an integral domain. In  $\mathbb{C}[x]$ , these are the ideals generated by x-z for some  $z \in \mathbb{C}$ . In  $\mathbb{R}[x]$ , the primes are the ideals generated by x-r for some  $r \in \mathbb{R}$  or  $x^2+r$  for some positive  $r \in R$ . (( What would be a general condition for  $\mathbb{Q}[x]$ ? ))

Solution: [CLO05] Ex. 1.1.9:

- (a) First, observe that  $\langle x_1, \ldots, x_n \rangle$  is the ideal consisting exactly of polynomials which have no constant term. Let I be an ideal in  $k[x_1, \ldots, x_n]$  such that  $\langle x_1, \ldots, x_n \rangle \subsetneq I$ . Thus there exists  $f(x_1, \ldots, x_n) \in I \setminus \langle x_1, \ldots, x_n \rangle$ . We have by our observation that f has a nonzero constant term z. Then note that the nonconstant terms of f form a polynomial  $g(x_1, \ldots, x_n)$  in  $\langle x_1, \ldots, x_n \rangle$ . Thus, we have that  $z = f(x) g(x) \in I$ . Since I contains a nonzero constant term, we must have that  $I = k[x_1, \ldots, x_n]$ .
- (b) Recall that an ideal I is maximal if and only if R/I is a field. Let  $I = \langle x_1 a_1, \ldots, x_n a_n \rangle$ . Consider the evaluation map  $\operatorname{ev}_{\overline{a}} \colon k[x_1, \ldots, x_n] \to k$  sending  $f(x_1, \ldots, x_n) \mapsto f(a_1, \ldots, a_n)$ . Clearly this map is surjective. Then since for all i we have that  $x_i \equiv a_i \pmod{I}$ , we have that  $f(x_1, \ldots, x_n) \equiv f(a_1, \ldots, a_n) \pmod{I}$  for all  $f(x_1, \ldots, x_n) \in k[x_1, \ldots, x_n]$ . Thus,  $\operatorname{ev}_{\overline{a}}(f) = f(a_1, \ldots, a_n) = 0$  if and only if  $f(x_1, \ldots, x_n) \in I$ . Thus,  $\operatorname{ker}(\operatorname{ev}_{\overline{a}}) = I$  and  $k[x_1, \ldots, x_n]/I$  is a field, meaning  $\langle x_1 a_1, \ldots, x_n a_n \rangle$  is maximal.
- (c) Since  $\mathbb{R}[x]$  is a principal ideal domain, any ideal I strictly containing  $\langle x^2 + 1 \rangle$  is of the form  $\langle g(x) \rangle$  for some  $g(x) \mid x^2 + 1$ . However, since  $x^2 + 1$  is irreducible in  $\mathbb{R}[x]$ , we have that g(x) is either  $z(x^2 + 1)$  for some nonzero  $z \in \mathbb{C}$  or g(x) = z for some nonzero  $z \in \mathbb{C}$ , meaning  $\langle g(x) \rangle = \langle x^2 + 1 \rangle$  or or  $\langle g(x) \rangle = \mathbb{R}[x]$ . Thus,  $\langle x^2 + 1 \rangle$  is maximal. However, in  $\mathbb{C}[x]$ , we have that  $x^2 + 1 = (x + i)(x i)$  and so  $\langle x^2 + 1 \rangle \subseteq \langle x i \rangle \subseteq \mathbb{C}[x]$ .

### Solution: [CLO05] Ex. 1.1.10:

- (a) Since  $x^2 + y^2 (x^2 z^3) = y^2 + z^3$  is an element of I which does not depend on x,  $y^2 + z^3$  is in  $I_1$ .
- (b) For all  $\ell \geq 1$ , we have that  $0 \in I_{\ell}$ . Then, if  $f(x_{\ell+1}, \ldots, x_n)$ ,  $g(x_{\ell+1}, \ldots, x_n)$  are two polynomials in I who do not depend on the first  $\ell$  variables, then so is f+g. Finally, let  $r(x_{\ell}+1, \ldots, x_n) \in k[x_{\ell+1}, \ldots, x_n]$ . Then  $r \cdot f \in I_{\ell}$  since  $r \cdot f \in I$  and still does not depend on any of the first  $\ell$  variables.

Solution: [CLO05] Ex. 1.1.11:

- (a) (( meh ))
- (b) (( meh ))
- (c) We claim that  $I + J = \langle f_1, \dots, f_s, g_1, \dots, g_t \rangle$ . Clearly  $I, J \subseteq \langle f_1, \dots, f_s, g_1, \dots, g_t \rangle$  and thus so is  $I \cup J$ . By (b), this shows that  $I + J \subseteq \langle f_1, \dots, f_s, g_1, \dots, g_t \rangle$ . Then, since  $f_i = f_i + 0$  and  $g_j = 0 + g_j$  for all i, j, we have the reverse inclusion and thus the two ideals are equal.

Solution: [CLO05] Ex. 1.1.12:

- (a) (( meh ))
- (b) Suppose that  $h(\overline{x}) \in IJ$ . Note that IJ is generated by the products  $f(\overline{x}) \cdot g(\overline{x})$  for  $f(\overline{x}) \in I$ , and  $g(\overline{x}) \in J$ . Then  $h(\overline{x})$  consists of sums of terms of the form  $r(\overline{x}) \cdot f(\overline{x}) \cdot g(\overline{x})$  for  $r(\overline{x}) \in k[\overline{x}]$ ,  $f(\overline{x}) \in I$ , and  $g(\overline{x}) \in J$ . Thus, each term is in both I and J and overall so is  $h(\overline{x})$ .

To see an example where  $IJ \subsetneq I \cap J$ , consider  $I = \langle x^2y \rangle$  and  $J = \langle xy^2 \rangle$  in k[x,y]. Then  $I \cap J = \langle x^2y^2 \rangle$  and  $IJ = \langle x^3y^3 \rangle$ . Thus  $IJ \subsetneq I \cap J$  as  $I \cap J$  contains  $x^2y^2$  and IJ does not contain  $x^2y^2$ .

### 1.2 Gröbner Bases

Solution:  $\langle\langle$  [CLO05] Ex. 1.3.11  $\rangle\rangle$  :

### 1.3 Affine Varieties

**Solution:** ⟨⟨ [CLO05] Ex. 1.4.9 ⟩⟩ :

### Chapter 2

## **Solving Polynomial Equations**

### 2.1 Solving Polynomial Systems by Elimination

Solution:  $\langle\langle \text{ [CLO05] Ex. 2.1.1 }\rangle\rangle$ :

Solution:  $\langle\langle \text{ [CLO05] Ex. 2.1.2 }\rangle\rangle$ :

Solution: [CLO05] Ex. 2.1.3:

We may freely rewrite the polynomial as  $p(z)=z^n-a_{n-1}z^{n-1}-\cdots-a_0$  We have that  $0=\overline{z}^n-a_{n-1}\overline{z}^{n-1}-\cdots-a_0$  and so  $\overline{z}^n=a_{n-1}\overline{z}^{n-1}+\cdots+a_0$ . Suppose now that  $|\overline{z}|\geq 1$ . Then

$$|\overline{z}|^n = |a_{n-1}\overline{z}^{n-1} + \dots + a_0| \le |a_{n-1}||z|^{n-1} + \dots + a_0 \le |a_{n-1}|\overline{z}^{n-1} + \dots + a_0|\overline{z}^{n-1}.$$

Thus,  $|\overline{z}| \le |a_{n-1}| + \dots + |a_0|$ . However, we assumed that  $|\overline{z}| \ge 1$ . This may not be the case. Thus,  $|\overline{z}| \le B := \max\{1, |a_{n-1}| + \dots + |a_0|\}$ .

Solution:  $\langle\langle \text{ [CLO05] Ex. 2.1.4} \rangle\rangle$ :

Numerically find all roots of  $2z^6 + 2z^5 - z^4 - z^3 - 2z^2 - 2z - 2$ .

### Solution: [CLO05] Ex. 2.1.5:

We apply Buchberger's Criterion. Let  $f(x, y) = x^2 + 2x + 3 + y^5 - y$  and  $g(x, y) = y^6 - y^2 + 2y$ . Then we have that

$$S(f,g) = \frac{x^2y^6}{x^2} \cdot (x^2 + 2x + 3 + y^5 - y) - \frac{x^2y^6}{y^6} \cdot (y^6 - y^2 + 2y) = y^6 \cdot (x^2 + 2x + 3 + y^5 - y) - x^2 \cdot (y^6 - y^2 + 2y).$$

This shows that  $\overline{S(f,g)}^G = 0$  which yields that *G* is a Gröbner basis.

Solution:  $\langle\langle \text{ [CLO05] Ex. 2.1.6} \rangle\rangle$ :

**Solution:** ⟨⟨ [CLO05] Ex. 2.1.7 ⟩⟩ :

### Solution: [CLO05] Ex. 2.1.8:

- (a) Let  $\overline{z}$  be a simple root of p(z), so p(z) = 0 but  $p'(z) \neq 0$ . Then  $N_p(\overline{z}) = \overline{z} \frac{p(\overline{z})}{p'(\overline{z})} = \overline{z}$  meaning  $\overline{z}$  is a fixed point of  $N_p(z)$ .
- (b) Suppose that  $\overline{z}$  is a multiple root of p(z) with multiplicity  $m \ge 2$ . Then we may express  $p(z) = \tilde{p}(z)(z-\overline{z})^m$  such that  $\tilde{p}(\overline{z}) \ne 0$ . Thus, we have that

$$\begin{split} N_p(z) &\coloneqq z - \frac{p(z)}{p'(z)} \\ &= z - \frac{\tilde{p}(z)(z - \overline{z})^m}{\tilde{p}'(z)(z - \overline{z})^m + m\tilde{p}(z)(z - \overline{z})^{m-1}} = z - \frac{\tilde{p}(z)(z - \overline{z})}{\tilde{p}'(z)(z - \overline{z}) + m\tilde{p}(z)} \end{split}$$

Note that  $m\tilde{p}(\bar{z}) \neq 0$ . Thus, we have that

$$\left|N_p(\overline{z})\right| = \left|\overline{z} - \frac{\tilde{p}(\overline{z})(\overline{z} - \overline{z})}{\tilde{p}'(\overline{z})(\overline{z} - \overline{z}) + m\tilde{p}(\overline{z})}\right| = |\overline{z}| \le \mathrm{LC}(p) \cdot B$$

where *B* is the value from [CLO05] Ex. 2.1.3 and LC(p) is the leading coefficient of p(z).

(c) Suppose now that  $\overline{z}$  is a simple root of  $p(\overline{z})$ . Then we may express  $p(z) = \tilde{p}(z)(z - \overline{z})$  such that  $\tilde{p}(\overline{z}) \neq 0$ . We have that

$$p'(z) = \tilde{p}'(z)(z - \overline{z}) + \tilde{p}(z)$$

and evaluation of p'(z) at  $\overline{z}$  is nonzero.

(d) Let  $\overline{z}$  be a root of multiplicity m. Following (b), we write  $p(z) = \tilde{p}(z)(z-\overline{z})^m$  such that  $\tilde{p}(\overline{z}) \neq 0$ . Then we have, by differentiating the expression for  $N_p(z)$  from (b), that

$$N_p'(z) = 1 - \frac{(\tilde{p}'(z)(z-\overline{z}) + \tilde{p}(z))(\tilde{p}'(z)(z-\overline{z}) + m\tilde{p}(z)) - (\tilde{p}(z)(z-\tilde{z}))(\tilde{p}''(z)(z-\overline{z}) + \tilde{p}'(z) + m\tilde{p}'(z))}{(\tilde{p}'(z)(z-\overline{z}) + m\tilde{p}(z))^2}.$$

Evaluation at  $z=\overline{z}$  yields that  $\lim_{z\to\overline{z}}N_p'(z)=1-\frac{1}{m}$ .

(e) Let  $\overline{z}$  be a root of multiplicity m. Following (b), we write  $p(z) = \tilde{p}(z)(z-\overline{z})^m$  such that  $\tilde{p}(\overline{z}) \neq 0$ . Then

$$p'(z) = \tilde{p}'(z - \overline{z})^m + m\tilde{p}(z)(z - \overline{z})^{m-1} = (z - \overline{z})^{m-1}(\tilde{p}'(z)(z - \overline{z}) + m\tilde{p}(z)).$$

Notice that  $\tilde{p}'(\overline{z})(\overline{z}-\overline{z})+m\tilde{p}(\overline{z})=m\tilde{p}(\overline{z})\neq 0$ . Thus, a root of multiplicity  $m\geq 1$  of p(z) is a root of multiplicity m-1 of p'(z). This implies that if we have roots  $\overline{z}_1,\ldots,\overline{z}_k$  with multiplicities  $m_1,\ldots,m_k\geq 1$ , then  $\gcd(p(z),p'(z))=(z-\overline{z}_1)^{m_1}\cdots(z-\overline{z}_k)^{m_k}$ . Thus, the polynomial  $p_{\mathrm{red}}(z)=\frac{p(z)}{\gcd(p(z),p'(z))}$  has the same roots of p(z) but all with multiplicity 1 which is the best case for Newton's method.

Solution: [CLO05] Ex. 2.1.9:

(a) Let  $p(z) = z^2 + 1$ . We have that

$$N_p(z) = z - \frac{z^2 + 1}{2z} = \frac{2z^2 - z^2 + 1}{2z} = \frac{z^2 + 1}{2z} = \frac{z^2 + 1}{2z} = \frac{z^2 + 2ixy - y^2 + 1}{2x + 2iy}.$$

If z is real then y = 0 and so  $N_p(x) = \frac{x^2+1}{2x}$  which is always real. Thus, Newton's method will never reach the imaginary roots of  $z^2 + 1$ . However, if we begin with a guess with nonzero imaginary part, then the guess does converge as expected.

(b) ⟨⟨ Just basic arithmetic not worth doing. ⟩⟩

Solution: [CLO05] Ex. 2.1.10:

Let  $\overline{z}$  be a root of p(z). Then  $-\overline{z}^n = a_{n-1}\overline{z}^{n-1} + \cdots + a_0$  and so

$$\begin{split} |\overline{z}|^n &= \left|a_{n-1}\overline{z}^{n-1} + \dots + a_0\right| \\ &\leq \max_i \left\{ \left|a_i\right| \right\} \cdot \left|\overline{z}^{n-1} + \dots + 1\right| \\ &\leq \max_i \left\{ \left|a_i\right| \right\} \cdot \left( |\overline{z}|^{n-1} + \dots + 1\right) \\ &= \max_i \left\{ \left|a_i\right| \right\} \cdot \frac{|\overline{z}|^n + 1}{|\overline{z}| - 1} \leq \max_i \left\{ \left|a_i\right| \right\} \cdot \frac{|\overline{z}|^n}{|\overline{z}| - 1}. \end{split}$$

 $\text{Thus, } |\overline{z}|^n \leq \max_i \left\{ \left. |a_i| \right. \right\} \cdot \frac{|z|^n}{|z|-1} \text{ which implies that } |z|-1 \leq \max_i \left\{ \left. |a_i| \right. \right\}. \text{ Thus, } |z| \leq 1 + \max_i \left\{ \left. |a_i| \right. \right\}. \qquad \qquad \square$ 

### 2.2 Finite Dimensional Algebras

Solution:  $\langle\langle \text{ [CLO05] Ex. 2.2.1 }\rangle\rangle$ :

### Solution: [CLO05] Ex. 2.2.2:

It is clear that  $\langle p_i(x_i) \rangle \subseteq I \cap k[x_i]$ . Now suppose that  $f(x_i) \in I \cap k[x_i]$ . Then  $\deg(f(x_i))$  must be  $\geq m_i$ . If not, then by the minimality of  $m_i$  we would arrive at a contradiction. Now by the division algorithm, write  $f(x_i) = q(x_i)p_i(x_i) + r(x_i)$  where  $\deg(r_{x_i}) < m_i$ . Then  $r(x_i) = f(x_i) - q(x_i)p(x_i) \in I$  and so  $r(x_i)$  must be 0 since if not, we would arrive at a contradiction of the minimality of  $m_i$ .

This gives us an algorithm to compute  $p_i(x_i)$ . Let I be a zero dimensional ideal and G a Gröbner basis for I. Then we know there exists  $m_i$  such that  $\{1, [x_i], \ldots, [x^{m_i}]\}$  is linearly dependent in  $k[\overline{x}]/I$ . In fact, we may use the Finiteness Theorem to set  $m_i$  to the smallest integer such that  $x_i^{m_i} = \operatorname{LT}(g)$  for some  $g \in G$ . Since  $k[x_1, \ldots, x_n]/I$  is a vector space, we can check linear independence in the usual way. See code/ch2/2\_2\_2.sage for a SageMath implementation of this.

### Solution: [CLO05] Ex. 2.2.3:

Let  $0 \neq f(x) \in \sqrt{\langle p(x) \rangle}$ . Then there exists  $m \geq 1$  such that  $f^m \in \langle p(x) \rangle$  and so  $p(x) \mid f(x)^m$ . In particular, each linear factor  $(x - \overline{z})$  of p(x) divides  $f(x)^n$  and so divides f(x) as  $(x - \overline{z})$  is irreducible. Thus,  $p_{\text{red}}(x) \mid f(x)$  and so  $f(x) \in \langle p_{\text{red}}(x) \rangle$ . Conversely, suppose  $f(x) \in \langle p_{\text{red}}(x) \rangle$  so that  $\langle p_{\text{red}} \rangle \mid f(x)$ . Label the roots of p(x) as  $\overline{z}_1, \dots, \overline{z}_r$ , each  $\overline{z}_i \in \overline{k}$ . Then for each  $i, (x - \overline{z}_i) \mid f(x)$  Let  $m_i$  be the multiplicity of  $z_i$  in p(x) and  $m = \max\{m_1, \dots, m_r\}$ . Then  $p(x) \mid f(x)^m$  and so  $f(x) \in \sqrt{\langle p(x) \rangle}$ 

#### Solution: [CLO05] Ex. 2.2.4:

We use the algorithm from [CLO05] Ex. 2.2.2 implemented in code/ch2/2\_2\_2sage. See code/ch2/2\_2\_2sage for the code in action.

### Solution: $\langle\langle \text{ [CLO05] Ex. 2.2.5} \rangle\rangle$ :

Then  $\sqrt{I} = I + \langle x(x-1), y(y-2) \rangle$ . Since  $I \subseteq \sqrt{I}$ , we see that  $\dim \mathbb{C}[x,y]/I \ge \dim \mathbb{C}[x,y]/\sqrt{I}$ . A quick SageMath computation confirms this:  $\dim \mathbb{C}[x,y]/I = 9$  and  $\dim \mathbb{C}[x,y]/\sqrt{I} = 2$ . See code/ch2/2\_2\_5. sage for the code in action. Then, since  $I \subseteq \sqrt{I}$  we have that  $V(\sqrt{I}) \subseteq V(I)$ . Notice that

$$y^{4}x + 3x^{3} - y^{4} - 3x^{2} = y^{4}(x - 1) + 3x^{2}(x - 1) = (y^{4} + 3x^{2})(x - 1)$$
$$x^{2}y - 2x^{2} = x^{2}(y - 2)$$
$$2y^{4}x - x^{3} - 2y^{4} + x^{2} = 2y^{4}(x - 1) - x^{2}(x - 1) = (2y^{4} - x^{2})(x - 1).$$

Thus, (1,2) and (0,0) are the only two points in V(I). Since it is evident that  $V(\sqrt{I})$  contains these two points, we see in this case that  $V(\sqrt{I}) = V(I)$ .

### Solution: [CLO05] Ex. 2.2.6:

A grevlex Gröbner basis for I is  $\{y^4 - 16y^2, x^3 - x^2, -2x^2\}$ . Thus, by the Finiteness Theorem, we know that the for monomials  $x^ay^b$  in  $\mathbb{C}[x,y]/I$  we must have that  $0 \le a \le 1$  and  $0 \le b \le 3$ . See code/ch2/2\_2\_6. sage for the code in action to compute the table.

### Solution: [CLO05] Ex. 2.2.7:

We implement the algorithm described in (( [CLO05] Ex. 1.3.11 )). See /code/ch2/2\_2\_7. sage for the code in action.

### Solution: [CLO05] Ex. 2.2.8:

- (a) See code/ch2/2\_2\_8. sage for the code in action.
- (b) Since each of the  $I_j$  are maximal ideals and  $I_j \subseteq \sqrt{I_j}$ , we must have that  $I = \sqrt{I_j}$ . Thus  $I(V(I_j)) = I_j$  and we must have that  $I_j = \sqrt{I_j}$ . Since each  $I_j$  is radical and  $I = \bigcap_{j=1}^5 I_j$ , we have by [CLO05] Ex. 2.2.7 that I is radical.

### Solution: [CLO05] Ex. 2.2.9:

- (a) Let  $f(\overline{x}) \in I + \langle p \rangle$  and let  $1 \le j \le d$ . Then  $f(\overline{x}) = g(\overline{x}) + h(\overline{x})p(x_1)$  for some  $g(\overline{x}) \in I$  and  $h(\overline{x}) \in k[\overline{x}]$ . We have that  $(x_1 a_j) \mid p(x_1)$  and so  $h(\overline{x})p(x_1) \in \langle x_1 a_j \rangle$ . Thus,  $f(\overline{x}) = g(\overline{x}) + h(\overline{x})p(x_1) \in I + \langle x_1 a_j \rangle$ . As j was arbitrary, we have that  $f(\overline{x}) \in \bigcap_j (I + \langle x_1 a_j \rangle)$ .
- (b) Let  $f(\overline{x}) \in p_j \cdot (I + \langle x_1 a_j \rangle)$ . Then  $f(\overline{x}) = p_j(x_1) \cdot (g(\overline{x}) + h(\overline{x})(x_1 a_j))$  for some  $g(\overline{x}) \in I$  and  $h(\overline{x}) \in k[\overline{x}]$ . We have that  $p_j(x_1)g(\overline{x}) \in I$  and  $p_j(x_1)h(\overline{x})(x_1 a_j) = h(\overline{x})p(x_1) \in \langle p \rangle$ . Thus,  $f(\overline{(x)}) = p_j(x_1)g(\overline{x}) + h(\overline{x})p(x_1) \in I + \langle p \rangle$ .
- (c) Let  $i \neq j$  and let  $d = \gcd(p_i, p_j)$ . We must have that  $d \mid p_i p_j = a_i a_j$  which is nonzero as  $a_i \neq a_j$ . Thus, d itself is a unit in  $k[\overline{x}]$  and  $p_i$  and  $p_j$  are coprime. By Bezout's Lemma, there exists polynomials  $h_1, \ldots, h_d \in k[\overline{x}]$  such that  $1 = \sum_{i=1}^d h_i(\overline{x}) p_i(x_1)$ .

(d)

## **Bibliography**

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