Using Algebraic Geometry

With 0 Figures

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Contents

Preface

At the time of writing this, I am starting my PhD at The Ohio State University. Currently a large part of my interests in algebra are about algorithms as they relate to polynomials and algebraic geometry. I've been doing a bunch of problems from *Ideals, Varieties, and Algorithms* [CLO15]. However, it seems that *Using Algebraic Geometry* [CLO05] Ex. moves through the material faster as it assumes you know more algebra. So I've moved onto working through this book as well as trying to comprehend Sturmfel's *Algorithms in Invariant Theory* [Str08].

Chapter 1

Introduction

1.1 Polynomials and Ideals

Solution: [CLO05] Ex. 1.1.1:

(a) We have that $x(x-y^2)+y(xy)=x^2-xy^2+xy^2=x^2$.

(b) It suffices to check for generators. We have that $x+(-1)(y^2)=x-y^2, y(x)=xy$, and $y^2=y^2$ showing that $\left\langle x-y^2, xy, y^2 \right\rangle \subseteq \left\langle x, y^2 \right\rangle$. Then $x-y^2+y^2=x$ and $y^2=y^2$ shows the reverse containment and overall the ideals are equal.

(c) We already know from 1. that x^2 lives in $\left\langle x-y^2, xy \right\rangle$. Since $xy=xy$, we overall have that $\left\langle x^2, xy \right\rangle \subseteq \left\langle x-y^2, xy \right\rangle$. It remains to check if $x-y^2 \in \left\langle x^2, xy \right\rangle$. However, notice that every element of $\left\langle x^2, xy \right\rangle$ is divisible by x while $x-y^2$ is clearly not divisible by x. Thus $x-y^2 \notin \left\langle x^2, xy \right\rangle$ and the two ideals are not equal.

$\square$

Solution: [CLO05] Ex. 1.1.2:

Let $f, g \in \langle f_1, \ldots, f_s \rangle$. Then $\exists p_1, \ldots, p_s, q_1, \ldots, q_s$ such that $f = \sum_{i=1}^s p_i \cdot f_i$ and $g = \sum_{i=1}^s q_i \cdot f_i$. Thus $f + g = \sum_{i=1}^s (p_i + q_i) \cdot f_i$ which shows that $f + g \in \langle f_1, \ldots, f_s \rangle$. Then let $p \in k[x_1, \ldots, x_n]$. We have that $p \cdot f = p \sum_{i=1}^s p_i f_i = \sum_{i=1}^s (p \cdot p_i) \cdot f_i$ which shows that $\langle f_1, \ldots, f_s \rangle$ is an ideal. $\square$

Solution: [CLO05] Ex. 1.1.3:

We already know that $\langle f_1, \ldots, f_s \rangle$ is an ideal by [CLO05] Ex. 1.1.2. Now suppose that J is an ideal containing $\{ f_1, \ldots, f_s \}$. Then, since ideals are closed under addition and scaling, we have that for all $p_1, \ldots, p_s \in k[x_1, \ldots, x_n]$ that $\sum_{i=1}^s p_i \cdot f_i \in J$. Thus, $\langle f_1, \ldots, f_s \rangle \subseteq J$. $\square$

Solution: [CLO05] Ex. 1.1.4:

We claim that $\langle f_1,\ldots,f_s\rangle = \langle g_1,\ldots,g_t\rangle$ if and only if $\{g_1,\ldots,g_t\} \subseteq I$ and $\{f_1,\ldots,f_s\} \subseteq J$. The forward implication is immediate. Then by [CLO05] Ex. 1.1.3, if $\{g_1,\ldots,g_t\} \subseteq I$ then $J \subseteq I$. Similarly, $\{f_1,\ldots,f_s\} \subseteq J \implies I \subseteq J$ and overall $I = J$. This fact was used in [CLO05] Ex. 1.1.1 (b). □

Solution: [CLO05] Ex. 1.1.5:

It suffices to show that $z - x^3 \in \left\langle y - x^2, z - xy \right\rangle$ and and $z - xy \in \left\langle x - y^2, z - x^3 \right\rangle$. Indeed we have that $(z - xy) + x(y - x^2) = z - x^3$ which also yields that $z - xy = z - x^3 - x(y - x^2)$. □

Solution: [CLO05] Ex. 1.1.6:

If $I = \{0\}$ then $I = \langle 0\rangle$. So suppose $I \neq 0$. Let $d \in I$ be of minimal degree. **⟨⟨ $d = \gcd(I)$ but I need infinite Bezout. ⟩⟩** Then we claim that $\langle d\rangle = I$. Since $d \in I$, we have that $\langle d\rangle \subseteq I$. Now let $f \in I$. By Euclidean division, there exists $q, r \in k[x]$ such that $f = qd + r$ where either $r = 0$ or $0 \leq \deg(r) \leq \deg(d) - 1$. If $r = 0$ then $f \in \langle d\rangle$ and we are done. So suppose $r \neq 0$. Then $f, qd \in I \implies r = f - qd \in I$. Thus, $r \in I$ is of degree strictly less than d, contradicting the minimality of the degree of d. So we must have that $r = 0$ and overall $\langle d\rangle = I$. □

Solution: [CLO05] Ex. 1.1.7:

(a) Suppose $f(x) \in \langle x\rangle$. Then $f(x)^m \in \langle x^n\rangle$ so $f(x) \in \sqrt{\langle x^n\rangle}$ Now suppose that $f(x) \in \sqrt{\langle x^n\rangle}$. Then $\exists k$ such that $f(x)^k \in \langle x^n\rangle$. Thus $f(x)^k$ is a multiple of x^n. This implies that $f(x)^k$ is a multiple of x. Then notice that the unique factorization of $f(x)^k$ into irreducibles is the kth power of the factorization of $f(x)$ into irreducibles. Thus x must be a factor of $f(x)$ and so $f(x) \in \langle x\rangle$. Note, this heavily uses the fact that $k[x]$ is a unique factorization domain for all fields k.

(b) We claim that $\sqrt{\langle p(x)\rangle} = \langle (x - a_1)\cdots(x - a_m)\rangle = I$. Suppose $f(x) \in I$. Let $k = \max e_1,\ldots,e_n$. Then $p(x) \mid f(x)^k$ so $f(x) \in \sqrt{\langle p(x)\rangle}$. Now suppose that $f(x) \in \sqrt{\langle p(x)\rangle}$. Then $\exists k$ such that $f(x)^k \in \langle p(x)\rangle$. Thus $f(x)^k$ is a multiple of each $(x - a_i)$. Then notice that the unique factorization of $f(x)^k$ into irreducibles is the kth power of the factorization of $f(x)$ into irreducibles. Thus $f(x)$ is a multiple of each $(x - a_i)$ and so $f(x) \in I$.

(c) Radical ideals are the ideals I such that $\sqrt{I} = I$. Notice that $\mathbb{C}[x]$ is a principal ideal domain and so any such I must be generated by a single polynomial. Since every polynomial in $\mathbb{C}[x]$ splits into linear factors, (b) immediately implies that the only radical ideals of $\mathbb{C}[x]$ are the ones which are of the form $\langle (x - a_1)\cdots(x - a_m)\rangle$ for $a_1,\ldots,a_m \in \mathbb{C}[x]$.

□

Solution: [CLO05] Ex. 1.1.8:

(a) Let $\mathfrak{p}$ be a prime ideal in $k[\overline{x}]$. Clearly $\mathfrak{p} \subseteq \sqrt{\mathfrak{p}}$ always. Let $f(\overline{x}) \in \sqrt{\mathfrak{p}}$. Then $f(\overline{x})^m \in \mathfrak{p}$ for some $m \in \mathbb{Z}_{\geq 1}$. We prove the reverse inclusion by induction on m. If $m = 1$ then $f(\overline{x}) = f(\overline{x})^1 \in \mathfrak{p}$. Now let $m > 1$ and suppose the claim holds for all $k \leq m$. Then suppose $f(\overline{x})^{m+1} \in \mathfrak{p}$. Then $f(\overline{x}) \cdot f(\overline{x})^m \in \mathfrak{p}$/ Either $f(\overline{x}) \in \mathfrak{p}$ or $f(\overline{x})^m \in \mathfrak{p}$ which by induction implies that $f(\overline{x}) \in \mathfrak{p}$. Thus, $f(\overline{x})^m \in \mathfrak{p} \implies f(\overline{x}) \in \mathfrak{p}$ for all $m \in \mathbb{Z}_{\geq 1}$ and so $\sqrt{\mathfrak{p}} \subseteq \mathfrak{p}$. Thus, all prime ideals are radical.

(b) Notice that for all fields k that $k[x]$ is a principal ideal domain. Thus, all the prime ideals are the ones generated by a single irreducible polynomial. Also, in $k[x]$ we have that (0) is a prime ideal as well as $k[x]$ is an integral domain. In $\mathbb{C}[x]$, these are the ideals generated by $x - z$ for some $z \in \mathbb{C}$. In $\mathbb{R}[x]$, the primes are the ideals generated by $x - r$ for some $r \in \mathbb{R}$ or $x^2 + r$ for some positive $r \in R$. **⟨⟨ What would be a general condition for $\mathbb{Q}[x]$? ⟩⟩**

$\square$

Solution: [CLO05] Ex. 1.1.9:

(a) First, observe that $\langle x_1, \ldots, x_n \rangle$ is the ideal consisting exactly of polynomials which have no constant term. Let I be an ideal in $k[x_1, \ldots, x_n]$ such that $\langle x_1, \ldots, x_n \rangle \subsetneq I$. Thus there exists $f(x_1, \ldots, x_n) \in I \setminus \langle x_1, \ldots, x_n \rangle$. We have by our observation that f has a nonzero constant term z. Then note that the non-constant terms of f form a polynomial $g(x_1, \ldots, x_n)$ in $\langle x_1, \ldots, x_n \rangle$. Thus, we have that $z = f(x) - g(x) \in I$. Since I contains a nonzero constant term, we must have that $I = k[x_1, \ldots, x_n]$.

(b) Recall that an ideal I is maximal if and only if R/I is a field. Let $I = \langle x_1 - a_1, \ldots, x_n - a_n \rangle$. Consider the evaluation map $\mathrm{ev}_{\overline{a}} \colon k[x_1, \ldots, x_n] \to k$ sending $f(x_1, \ldots, x_n) \mapsto f(a_1, \ldots, a_n)$. Clearly this map is surjective. Then since for all i we have that $x_i \equiv a_i \pmod{I}$, we have that $f(x_1, \ldots, x_n) \equiv f(a_1, \ldots, a_n) \pmod{I}$ for all $f(x_1, \ldots, x_n) \in k[x_1, \ldots, x_n]$. Thus, $\mathrm{ev}_{\overline{a}}(f) = f(a_1, \ldots, a_n) = 0$ if and only if $f(x_1, \ldots, x_n) \in I$. Thus, $\ker(\mathrm{ev}_{\overline{a}}) = I$ and $k[x_1, \ldots, x_n]/I$ is a field, meaning $\langle x_1 - a_1, \ldots, x_n - a_n \rangle$ is maximal.

(c) Since $\mathbb{R}[x]$ is a principal ideal domain, any ideal I strictly containing $\langle x^2 + 1 \rangle$ is of the form $\langle g(x) \rangle$ for some $g(x) \mid x^2 + 1$. However, since $x^2 + 1$ is irreducible in $\mathbb{R}[x]$, we have that $g(x)$ is either $z(x^2 + 1)$ for some nonzero $z \in \mathbb{C}$ or $g(x) = z$ for some nonzero $z \in \mathbb{C}$, meaning $\langle g(x) \rangle = \langle x^2 + 1 \rangle$ or or $\langle g(x) \rangle = \mathbb{R}[x]$. Thus, $\langle x^2 + 1 \rangle$ is maximal. However, in $\mathbb{C}[x]$, we have that $x^2 + 1 = (x + i)(x - i)$ and so $\langle x^2 + 1 \rangle \subsetneq \langle x - i \rangle \subsetneq \mathbb{C}[x]$.

$\square$

Solution: [CLO05] Ex. 1.1.10:

(a) Since $x^2+y^2-(x^2-z^3)=y^2+z^3$ is an element of I which does not depend on x, y^2+z^3 is in I_1.

(b) For all $\ell \geq 1$, we have that $0 \in I_\ell$. Then, if $f(x_{\ell+1},\ldots,x_n), g(x_{\ell+1},\ldots,x_n)$ are two polynomials in I who do not depend on the first ℓ variables, then so is $f+g$. Finally, let $r(x_\ell+1,\ldots,x_n) \in k[x_{\ell+1},\ldots,x_n]$. Then $r \cdot f \in I_\ell$ since $r \cdot f \in I$ and still does not depend on any of the first ℓ variables.

$\square$

Solution: [CLO05] Ex. 1.1.11:

(a) **⟨⟨ meh ⟩⟩**

(b) **⟨⟨ meh ⟩⟩**

(c) We claim that $I+J=\langle f_1,\ldots,f_s,g_1,\ldots,g_t\rangle$. Clearly $I,J \subseteq \langle f_1,\ldots,f_s,g_1,\ldots,g_t\rangle$ and thus so is $I \cup J$. By (b), this shows that $I+J \subseteq \langle f_1,\ldots,f_s,g_1,\ldots,g_t\rangle$. Then, since $f_i=f_i+0$ and $g_j=0+g_j$ for all i,j, we have the reverse inclusion and thus the two ideals are equal.

$\square$

Solution: [CLO05] Ex. 1.1.12:

(a) **⟨⟨ meh ⟩⟩**

(b) Suppose that $h(\overline{x}) \in IJ$. Note that IJ is generated by the products $f(\overline{x})\cdot g(\overline{x})$ for $f(\overline{x}) \in I$, and $g(\overline{x}) \in J$. Then $h(\overline{x})$ consists of sums of terms of the form $r(\overline{x})\cdot f(\overline{x})\cdot g(\overline{x})$ for $r(\overline{x}) \in k[\overline{x}]$, $f(\overline{x}) \in I$, and $g(\overline{x}) \in J$. Thus, each term is in both I and J and overall so is $h(\overline{x})$.

To see an example where $IJ \subsetneq I \cap J$, consider $I=\left\langle x^2y \right\rangle$ and $J=\left\langle xy^2 \right\rangle$ in $k[x,y]$. Then $I \cap J=\left\langle x^2y^2 \right\rangle$ and $IJ=\left\langle x^3y^3 \right\rangle$. Thus $IJ \subsetneq I \cap J$ as $I \cap J$ contains x^2y^2 and IJ does not contain x^2y^2.

$\square$

1.2 Gröbner Bases

Solution: ⟨⟨ [CLO05] Ex. 1.3.11 ⟩⟩ :

□

1.3 Affine Varieties

Solution: ⟨⟨ [CLO05] Ex. 1.4.9 ⟩⟩ :

□

Chapter 2

Solving Polynomial Equations

2.1 Solving Polynomial Systems by Elimination

Solution: ⟨⟨ **[CLO05] Ex. 2.1.1** ⟩⟩ :

□

Solution: ⟨⟨ **[CLO05] Ex. 2.1.2** ⟩⟩ :

□

Solution: [CLO05] Ex. 2.1.3:

We may freely rewrite the polynomial as $p(z) = z^n - a_{n-1}z^{n-1} - \cdots - a_0$ We have that $0 = \overline{z}^n - a_{n-1}\overline{z}^{n-1} - \cdots - a_0$ and so $\overline{z}^n = a_{n-1}\overline{z}^{n-1} + \cdots + a_0$. Suppose now that $|\overline{z}| \geq 1$. Then

$$|\overline{z}|^n = \left|a_{n-1}\overline{z}^{n-1} + \cdots + a_0\right| \leq |a_{n-1}||z|^{n-1} + \cdots + a_0 \leq |a_{n-1}|\overline{z}^{n-1} + \cdots |a_0|\overline{z}^{n-1}.$$

Thus, $|\overline{z}| \leq |a_{n-1}| + \cdots + |a_0|$. However, we assumed that $|\overline{z}| \geq 1$. This may not be the case. Thus, $|\overline{z}| \leq B := \max\{1, |a_{n-1}| + \cdots + |a_0|\}$. □

Solution: ⟨⟨ **[CLO05] Ex. 2.1.4** ⟩⟩ :

Numerically find all roots of $2z^6 + 2z^5 - z^4 - z^3 - 2z^2 - 2z - 2$. □

Solution: [CLO05] Ex. 2.1.5:

We apply Buchberger's Criterion. Let $f(x,y)=x^2+2x+3+y^5-y$ and $g(x,y)=y^6-y^2+2y$. Then we have that

$$S(f,g)=\frac{x^2y^6}{x^2}\cdot(x^2+2x+3+y^5-y)-\frac{x^2y^6}{y^6}\cdot(y^6-y^2+2y)=y^6\cdot(x^2+2x+3+y^5-y)-x^2\cdot(y^6-y^2+2y).$$

This shows that $\overline{S(f,g)}^G=0$ which yields that G is a Gröbner basis. $\square$

Solution: **⟨⟨ [CLO05] Ex. 2.1.6 ⟩⟩** :

$\square$

Solution: **⟨⟨ [CLO05] Ex. 2.1.7 ⟩⟩** :

$\square$

Solution: [CLO05] Ex. 2.1.8:

(a) Let $\overline{z}$ be a simple root of $p(z)$, so $p(z)=0$ but $p'(z)\neq 0$. Then $N_p(\overline{z})=\overline{z}-\frac{p(\overline{z})}{p'(\overline{z})}=\overline{z}$ meaning $\overline{z}$ is a fixed point of $N_p(z)$.

(b) Suppose that $\overline{z}$ is a multiple root of $p(z)$ with multiplicity $m\geq 2$. Then we may express $p(z)=\tilde{(p)}(z)(z-\overline{z})^m$ such that $\tilde{p}(\overline{z})\neq 0$. Thus, we have that

$$\begin{aligned}N_p(z) &:= z-\frac{p(z)}{p'(z)}\\ &= z-\frac{\tilde{p}(z)(z-\overline{z})^m}{\tilde{p}'(z)(z-\overline{z})^m+m\tilde{p}(z)(z-\overline{z})^{m-1}} = z-\frac{\tilde{p}(z)(z-\overline{z})}{\tilde{p}'(z)(z-\overline{z})+m\tilde{p}(z)}\end{aligned}$$

Note that $m\tilde{p}(\overline{z})\neq 0$. Thus, we have that

$$\left|N_p(\overline{z})\right| = \left|\overline{z}-\frac{\tilde{p}(\overline{z})(\overline{z}-\overline{z})}{\tilde{p}'(\overline{z})(\overline{z}-\overline{z})+m\tilde{p}(\overline{z})}\right| = |\overline{z}| \leq \textsc{lc}(p)\cdot B$$

where B is the value from [CLO05] Ex. 2.1.3 and $\textsc{lc}(p)$ is the leading coefficient of $p(z)$.

(c) Suppose now that $\overline{z}$ is a simple root of $p(\overline{z})$. Then we may express $p(z)=\tilde{p}(z)(z-\overline{z})$ such that $\tilde{p}(\overline{z})\neq 0$. We have that

$$p'(z)=\tilde{p}'(z)(z-\overline{z})+\tilde{p}(z)$$

and evaluation of $p'(z)$ at $\overline{z}$ is nonzero.

(d) Let $\overline{z}$ be a root of multiplicity m. Following (b), we write $p(z)=\tilde{p}(z)(z-\overline{z})^m$ such that $\tilde{p}(\overline{z})\neq 0$. Then we have, by differentiating the expression for $N_p(z)$ from (b), that

$$N_p'(z)=1-\frac{(\tilde{p}'(z)(z-\overline{z})+\tilde{p}(z))(\tilde{p}'(z)(z-\overline{z})+m\tilde{p}(z))-(\tilde{p}(z)(z-\tilde{z}))(\tilde{p}''(z)(z-\overline{z})+\tilde{p}'(z)+m\tilde{p}'(z))}{(\tilde{p}'(z)(z-\overline{z})+m\tilde{p}(z))^2}.$$

Evaluation at $z=\overline{z}$ yields that $\lim_{z\to\overline{z}}N_p'(z)=1-\frac{1}{m}$.

(e) Let $\overline{z}$ be a root of multiplicity m. Following (b), we write $p(z)=\tilde{p}(z)(z-\overline{z})^m$ such that $\tilde{p}(\overline{z})\neq 0$. Then

$$p'(z)=\tilde{p}'(z-\overline{z})^m+m\tilde{p}(z)(z-\overline{z})^{m-1}=(z-\overline{z})^{m-1}(\tilde{p}'(z)(z-\overline{z})+m\tilde{p}(z)).$$

Notice that $\tilde{p}'(\overline{z})(\overline{z}-\overline{z})+m\tilde{p}(\overline{z})=m\tilde{p}(\overline{z})\neq 0$. Thus, a root of multiplicity $m\geq 1$ of $p(z)$ is a root of multiplicity $m-1$ of $p'(z)$. This implies that if we have roots $\overline{z}_1,\ldots,\overline{z}_k$ with multiplicities $m_1,\ldots,m_k\geq 1$, then $\gcd(p(z),p'(z))=(z-\overline{z}_1)^{m_1}\cdots(z-\overline{z}_k)^{m_k}$. Thus, the polynomial $p_{\text{red}}(z)=\frac{p(z)}{\gcd(p(z),p'(z))}$ has the same roots of $p(z)$ but all with multiplicity 1 which is the best case for Newton's method.

$\square$

Solution: [CLO05] Ex. 2.1.9:

(a) Let $p(z) = z^2 + 1$. We have that

$$N_p(z) = z - \frac{z^2+1}{2z} = \frac{2z^2 - z^2 + 1}{2z} = \frac{z^2+1}{2z} = \frac{x^2 + 2ixy - y^2 + 1}{2x + 2iy}.$$

If z is real then $y = 0$ and so $N_p(x) = \frac{x^2+1}{2x}$ which is always real. Thus, Newton's method will never reach the imaginary roots of $z^2 + 1$. However, if we begin with a guess with nonzero imaginary part, then the guess does converge as expected.

(b) **⟨⟨ Just basic arithmetic not worth doing. ⟩⟩**

□

Solution: [CLO05] Ex. 2.1.10:

Let $\overline{z}$ be a root of $p(z)$. Then $-\overline{z}^n = a_{n-1}\overline{z}^{n-1} + \cdots + a_0$ and so

$$\begin{aligned}
|\overline{z}|^n &= \left|a_{n-1}\overline{z}^{n-1} + \cdots + a_0\right| \\
&\leq \max_i \{\,|a_i|\,\} \cdot \left|\overline{z}^{n-1} + \cdots + 1\right| \\
&\leq \max_i \{\,|a_i|\,\} \cdot \left(|\overline{z}|^{n-1} + \cdots + 1\right) \\
&= \max_i \{\,|a_i|\,\} \cdot \frac{|\overline{z}|^n + 1}{|\overline{z}| - 1} \leq \max_i \{\,|a_i|\,\} \cdot \frac{|\overline{z}|^n}{|\overline{z}| - 1}.
\end{aligned}$$

Thus, $|\overline{z}|^n \leq \max_i \{\,|a_i|\,\} \cdot \frac{|z|^n}{|z|-1}$ which implies that $|z| - 1 \leq \max_i \{\,|a_i|\,\}$. Thus, $|z| \leq 1 + \max_i \{\,|a_i|\,\}$. □

2.2 Finite Dimensional Algebras

Solution: **⟨⟨ [CLO05] Ex. 2.2.1 ⟩⟩** :

□

Solution: [CLO05] Ex. 2.2.2:

It is clear that $\langle p_i(x_i)\rangle \subseteq I \cap k[x_i]$. Now suppose that $f(x_i) \in I \cap k[x_i]$. Then $\deg(f(x_i))$ must be $\geq m_i$. If not, then by the minimality of m_i we would arrive at a contradiction. Now by the division algorithm, write $f(x_i) = q(x_i)p_i(x_i) + r(x_i)$ where $\deg(r_{x_i}) < m_i$. Then $r(x_i) = f(x_i) - q(x_i)p(x_i) \in I$ and so $r(x_i)$ must be 0 since if not, we would arrive at a contradiction of the minimality of m_i.

This gives us an algorithm to compute $p_i(x_i)$. Let I be a zero dimensional ideal and G a Gröbner basis for I. Then we know there exists m_i such that $\{1, [x_i], \ldots, [x^{m_i}]\}$ is linearly dependent in $k[\overline{x}]/I$. In fact, we may use the Finiteness Theorem to set m_i to the smallest integer such that $x_i^{m_i} = \textsc{lt}(g)$ for some $g \in G$. Since $k[x_1, \ldots, x_n]/I$ is a vector space, we can check linear independence in the usual way. See `code/ch2/2_2_2.sage` for a SageMath implementation of this. □

Solution: [CLO05] Ex. 2.2.3:

Let $0 \neq f(x) \in \sqrt{\langle p(x)\rangle}$. Then there exists $m \geq 1$ such that $f^m \in \langle p(x)\rangle$ and so $p(x) \mid f(x)^m$. In particular, each linear factor $(x - \overline{z})$ of $p(x)$ divides $f(x)^n$ and so divides $f(x)$ as $(x - \overline{z})$ is irreducible. Thus, $p_{\text{red}}(x) \mid f(x)$ and so $f(x) \in \langle p_{\text{red}}(x)\rangle$. Conversely, suppose $f(x) \in \langle p_{\text{red}}(x)\rangle$ so that $\langle p_{\text{red}}\rangle \mid f(x)$. Label the roots of $p(x)$ as $\overline{z}_1, \ldots, \overline{z}_r$, each $\overline{z}_i \in \overline{k}$. Then for each i, $(x - \overline{z}_i) \mid f(x)$ Let m_i be the multiplicity of z_i in $p(x)$ and $m = \max\{m_1, \ldots, m_r\}$. Then $p(x) \mid f(x)^m$ and so $f(x) \in \sqrt{\langle p(x)\rangle}$ □

Solution: [CLO05] Ex. 2.2.4:

We use the algorithm from [CLO05] Ex. 2.2.2 implemented in `code/ch2/2_2_2sage`. See `code/ch2/2_2_2sage` for the code in action. □

Solution: ⟨⟨ **[CLO05] Ex. 2.2.5** ⟩⟩ :

Then $\sqrt{I} = I + \langle x(x-1), y(y-2)\rangle$. Since $I \subseteq \sqrt{I}$, we see that $\dim \mathbb{C}[x,y]/I \geq \dim \mathbb{C}[x,y]/\sqrt{I}$. A quick SageMath computation confirms this: $\dim \mathbb{C}[x,y]/I = 9$ and $\dim \mathbb{C}[x,y]/\sqrt{I} = 2$. See `code/ch2/2_2_5.sage` for the code in action. Then, since $I \subseteq \sqrt{I}$ we have that $V(\sqrt{I}) \subseteq V(I)$. Notice that

$$\begin{aligned} y^4x + 3x^3 - y^4 - 3x^2 &= y^4(x-1) + 3x^2(x-1) = (y^4 + 3x^2)(x-1) \\ x^2y - 2x^2 &= x^2(y-2) \\ 2y^4x - x^3 - 2y^4 + x^2 &= 2y^4(x-1) - x^2(x-1) = (2y^4 - x^2)(x-1). \end{aligned}$$

Thus, $(1,2)$ and $(0,0)$ are the only two points in $V(I)$. Since it is evident that $V(\sqrt{I})$ contains these two points, we see in this case that $V(\sqrt{I}) = V(I)$. □

Solution: [CLO05] Ex. 2.2.6:

A grevlex Gröbner basis for I is $\{y^4 - 16y^2, x^3 - x^2, -2x^2\}$. Thus, by the Finiteness Theorem, we know that the for monomials $x^a y^b$ in $\mathbb{C}[x,y]/I$ we must have that $0 \leq a \leq 1$ and $0 \leq b \leq 3$. See `code/ch2/2_2_6.sage` for the code in action to compute the table. □

Solution: [CLO05] Ex. 2.2.7:

We implement the algorithm described in ⟨⟨ **[CLO05] Ex. 1.3.11** ⟩⟩ . See `/code/ch2/2_2_7.sage` for the code in action. □

Solution: [CLO05] Ex. 2.2.8:

(a) See `code/ch2/2_2_8.sage` for the code in action.

(b) Since each of the I_j are maximal ideals and $I_j \subseteq \sqrt{I_j}$, we must have that $I = \sqrt{I_j}$. Thus $I(V(I_j)) = I_j$ and we must have that $I_j = \sqrt{I_j}$. Since each I_j is radical and $I = \bigcap_{j=1}^{5} I_j$, we have by [CLO05] Ex. 2.2.7 that I is radical.

□

Solution: [CLO05] Ex. 2.2.9:

(a) Let $f(\overline{x}) \in I + \langle p \rangle$ and let $1 \le j \le d$. Then $f(\overline{x}) = g(\overline{x}) + h(\overline{x})p(x_1)$ for some $g(\overline{x}) \in I$ and $h(\overline{x}) \in k[\overline{x}]$. We have that $(x_1 - a_j) \mid p(x_1)$ and so $h(\overline{x})p(x_1) \in \left\langle x_1 - a_j \right\rangle$. Thus, $f(\overline{x}) = g(\overline{x}) + h(\overline{x})p(x_1) \in I + \left\langle x_1 - a_j \right\rangle$. As j was arbitrary, we have that $f(\overline{x}) \in \bigcap_j (I + \left\langle x_1 - a_j \right\rangle)$.

(b) Let $f(\overline{x}) \in p_j \cdot (I + \left\langle x_1 - a_j \right\rangle)$. Then $f(\overline{x}) = p_j(x_1) \cdot (g(\overline{x}) + h(\overline{x})(x_1 - a_j))$ for some $g(\overline{x}) \in I$ and $h(\overline{x}) \in k[\overline{x}]$. We have that $p_j(x_1)g(\overline{x}) \in I$ and $p_j(x_1)h(\overline{x})(x_1 - a_j) = h(\overline{x})p(x_1) \in \langle p \rangle$. Thus, $f(\bar{(}x)) = p_j(x_1)g(\overline{x}) + h(\overline{x})p(x_1) \in I + \langle p \rangle$.

(c) Let $d = \gcd(p_1, \ldots, p_d)$. Then as $d \mid p_1$ and $d \mid p_2$, we have that $d \mid \prod_{j \neq 1,2}(x_1 - a_j)$. Continuing on inductively, we have that for all $c \le d$ that $d \mid \prod_{j \notin [c]}(x_1 - a_j)$. In particular, this means that $d \mid \prod_{j \notin [d]}(x_1 - a_j) = 1$. Thus, d itself is a unit in $k[\overline{x}]$ and p_i and p_j are coprime. By Bezout's Lemma, there exists polynomials $h_1, \ldots, h_d \in k[\overline{x}]$ such that $1 = \sum_{j=1}^d h_j(\overline{x})p_j(x_1)$.

(d) Now let $h(\overline{x}) \in \bigcap_{j=1}^d (I + \left\langle x_1 - a_j \right\rangle)$. As all the p_j are coprime, we have that there exist polynomials $h_1, \ldots, h_d \in k[\overline{x}]$ such that $1 = \sum_{j=1}^d h_j(\overline{x})p_j(x_1)$. Thus, $h = \sum_{j=1}^d h_j(\overline{x})p_j(x_1)h(\overline{x})$. Then for all $1 \le j \le d$, we have that as $p_j(x_1)h(\overline{x}) \in p_j \cdot (I + \left\langle x_1 - a_j \right\rangle) \subseteq I + \langle p \rangle$. Thus, each summand of $\sum_{j=1}^d h_j(\overline{x})p_j(x_1)h(\overline{x})$ is in $I + \langle p \rangle$ and so overall $h \in I + \langle p \rangle$.

$\square$

Solution: [CLO05] Ex. 2.2.10:

(a) Let $\overline{f}^G = \sum_{j=1}^d c_j(f)x^{\alpha(j)}$ and $\overline{g}^G = \sum_{j=1}^d c_j(g)x^{\alpha(j)}$. Then by combining like terms, we have that $\overline{f}^G + \overline{g}^G = \sum_{j=1}^d (c_j(f)+c_j(g))x^{\alpha(j)}$. On the other hand, we have that $\overline{f+g}^G = \sum_{j=1}^d c_j(f+g)x^{\alpha(j)}$. Since $\overline{f}^G + \overline{g}^G = \overline{f+g}^G$ and each of the $x^{\alpha(j)}$ are linearly independent, we may equate coefficients and conclude that $c_j(f)+c_j(g) = c_j(f+g)$. For $\lambda \in k$, $\overline{\lambda f}^G = \sum_{j=1}^d c_j(\lambda f)x^{\alpha(j)}$. Now notice that $\overline{\lambda f}^G = \lambda \overline{f}^G$ as we are working over a field. Thus, we have by equating coefficients that $c_j(\lambda f) = \lambda c_j(f)$. Thus, c_j is a linear function $A \to k$.

(b) Let $\alpha_j \in A^*$ be the linear map $\alpha_j(f) = c_j(f)$. Notice that for all $1 \le i, j \le d$ we have that $\alpha_j(x^{\alpha(i)}) = c_j(x^{\alpha(i)}) = \delta_{i,j}$. Suppose there exists $\lambda_1, \dots, \lambda_d \in k$ such that $\lambda_1\alpha_1 + \cdots + \lambda_d\alpha_d = 0$. Then for all $1 \le i \le d$ we have that
$$0 = \left(\sum_{j=1}^d \lambda_j \alpha_j\right)(x^{\alpha(i)}) = \sum_{j=1}^d \lambda_j \alpha_j(x^{\alpha(i)}) = \lambda_i$$
and so for all $1 \le i \le d$, $\lambda_i = 0$ meaning that $\{\alpha_1, \dots, \alpha_d\}$ is linearly independent. Since we know that $d = \dim A = \dim A^*$, we have that $\{\alpha_1, \dots, \alpha_d\}$ is a basis for A^*.

(c) This was proven in **(b)**.

$\square$

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