Algorithms in Invariant Theory

With 0 Figures

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Contents

1 Introduction **1**

1.1 Symmetric Polynomials 1

1.2 Gröbner Bases 3

Preface

A core interest of mine (at the time of writing this) is algorithms in the context of commutative algebra and algebraic geometry. As such, Bernd Sturmfel's text *Algorithms in Invariant Theory* [Str08] is a good resource for first learning this stuff. Perhaps in the future I'll include notes and some source code

Chapter 1

Introduction

1.1 Symmetric Polynomials

Solution: [Str08] 1.1.5: From [Page 23, 2.11'][Mac98], we have the following identities

$$
\begin{aligned}
1 \cdot \sigma_1 &= p_1, \\
2 \cdot \sigma_2 &= p_1 \sigma_1 - p_2, \\
3 \cdot \sigma_3 &= p_1 \sigma_2 - p_2 \sigma_1 + p_3, \\
&\vdots \\
k \cdot \sigma_k &= \sum_{r=1}^{k} (-1)^{r-1} p_r \sigma_{k-r}.
\end{aligned}
$$

Treating the σ_i as indeterminants, we can re-express the above system of equations:

$$
\begin{aligned}
p_1 &= 1 \cdot \sigma_1, \\
p_2 &= p_1 \sigma_1 - 2 \cdot \sigma_2, \\
p_3 &= p_2 \sigma_1 - p_1 \sigma_2 + 3 \cdot \sigma_3, \\
&\vdots \\
p_k &= \left(\sum_{r=1}^{k-1} (-1)^{r-1} p_{k-r} \sigma_r \right) + (-1)^{k-1} \cdot k \sigma_k.
\end{aligned}
$$

Consider $\sigma_1, -\sigma_2, \sigma_3, \ldots, (-1)^n\sigma_{n-1}, \sigma_n$ as indeterminants. Thus, we obtain the following matrix equation:

$$\begin{pmatrix} 1 & 0 & 0 & \cdots & 0 \\ p_1 & 2 & 0 & \cdots & 0 \\ p_2 & p_1 & 3 & \cdots & 0 \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ p_{k-1} & p_{k-2} & \cdots & \cdots & (-1)^k \cdot k \end{pmatrix} \begin{pmatrix} \sigma_1 \\ -\sigma_2 \\ \vdots \\ (-1)^k\sigma_{k-1} \\ \sigma_k \end{pmatrix} = \begin{pmatrix} p_1 \\ p_2 \\ \vdots \\ p_{k-1} \\ p_k \end{pmatrix}.$$

Then, Cramer's rule yields that

$$\begin{aligned}
\sigma_k &= \det\begin{pmatrix} 1 & 0 & 0 & \cdots & p_1 \\ p_1 & 2 & 0 & \cdots & p_2 \\ p_2 & p_1 & 3 & \cdots & p_3 \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ p_{k-1} & p_{k-2} & \cdots & \cdots & p_k \end{pmatrix} \det\begin{pmatrix} 1 & 0 & 0 & \cdots & 0 \\ p_1 & 2 & 0 & \cdots & 0 \\ p_2 & p_1 & 3 & \cdots & 0 \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ p_{k-1} & p_{k-2} & \cdots & \cdots & (-1)^k \cdot k \end{pmatrix}^{-1} \\
&= \frac{(-1)^k}{k!}\det\begin{pmatrix} 1 & 0 & 0 & \cdots & p_1 \\ p_1 & 2 & 0 & \cdots & p_2 \\ p_2 & p_1 & 3 & \cdots & p_3 \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ p_{k-1} & p_{k-2} & \cdots & \cdots & p_k \end{pmatrix} \\
&= \frac{(-1)^k}{k!} \cdot (-1)^k \det\begin{pmatrix} p_1 & 1 & 0 & \cdots & 0 \\ p_2 & p_1 & 2 & \cdots & 0 \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ p_{k-1} & p_{k-2} & \cdots & p_1 & k-1 \\ p_k & p_{k-1} & \cdots & \cdots & p_1 \end{pmatrix} = \frac{1}{k!}\det\begin{pmatrix} p_1 & 1 & 0 & \cdots & 0 \\ p_2 & p_1 & 2 & \cdots & 0 \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ p_{k-1} & p_{k-2} & \cdots & p_1 & k-1 \\ p_k & p_{k-1} & \cdots & \cdots & p_1 \end{pmatrix}.
\end{aligned}$$

$\square$

1.2 Gröbner Bases

Lemma 1.2.1:

Let $R = \mathbb{C}[x_1, \ldots, x_n]$. Then with the usual grading, let $H(R, z) := \sum_{d=0}^{\infty} \dim_{\mathbb{C}}(R_d) z^d$. We have that

$$H(R, z) := \sum_{d=0}^{\infty} \dim_{\mathbb{C}}(R_d) z^d = \sum_{d=0}^{\infty} \binom{d+n-1}{n-1} = \frac{1}{(1-z)^n}.$$

Solution: To see that $H(R, z) = \sum_{d=0}^{\infty} \binom{d+n-1}{n-1}$, just count the number of monomials of degree d in n variables. The value $\binom{d+n-1}{n-1}$ is the number of non-negative integer solutions to $a_1 + \cdots + a_n = d$. Each solution corresponds to a monomial $x_1^{a_1} \cdots x_n^{a_n}$. Then to see that $H(R, z) = \frac{1}{(1-z)^n}$, consider the product of infinite sums $(1 + z + z^2 + \cdots) \cdots (1 + z + z^2 + \cdots)$ a total of n-times. Then the coefficient of z^d again corresponds to the number of such non-negative integer solutions. Since $\frac{1}{1-z} = 1 + z + z^2 + \cdots$, we obtain the desired equality. □

Solution: [Str08] 1.2.1: Let M be a set of monomial generators for $\operatorname{init}(I)$ and let m be minimally nonstandard. Since m is a monomial and in $\operatorname{init}(I)$, we have that $m' \mid m$ for some monomial $m' \in M$. However, note that $m' \in \operatorname{init}(I)$ and furthermore by the fact that m is minimally nonstandard, we cannot have that m' strictly divides m. Thus, $m' = m$ and $m \in M$. Then by Dickson's Lemma, we have that M is a finite set and thus there are finitely many minimally nonstandard monomials. □

Solution: [Str08] 1.2.2: This is [CLO15, Chapter 2, §7, Theorem 5]. □

Solution: [Str08] 1.2.3: This is [CLO15, Chapter 3, §1, Theorem 2]. □

Solution: [Str08] 1.2.6: First, we recall some definitions. The S-polynomial of f and g is

$$S(f, g) := \frac{\operatorname{lcm}(\textsc{lm}(f), \textsc{lm}(g))}{\textsc{lt}(f)} f - \frac{\operatorname{lcm}(\textsc{lm}(f), \textsc{lm}(g))}{\textsc{lt}(g)} g.$$

For a set of polynomials $\mathcal{F} = \{f_1, \ldots, f_t\} \subseteq k[x_1, \ldots, x_n]$, we write $f \to_{\mathcal{F}} 0$ if there exists $a_1, \ldots, a_t \in k[x_1, \ldots, x_n]$ such that $a_1 f_1 + \cdots + a_t f_t = 0$. Then [CLO15, Chapter 2, §9, Theorem 3] says that a basis $\mathcal{F} = \{f_1, \ldots f_t\}$ is a Gröbner basis for G if and only if $S(f_i, f_j) \to_{\mathcal{F}} 0$ for all $i \neq j$. But [CLO15, Chapter 2, §9, Proposition 4] says that for $f, g \in \mathcal{F}$ with relatively prime initial monomials, we have that $S(f, g) \to_{\mathcal{F}} 0$. This proves the claim. □

Bibliography

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