

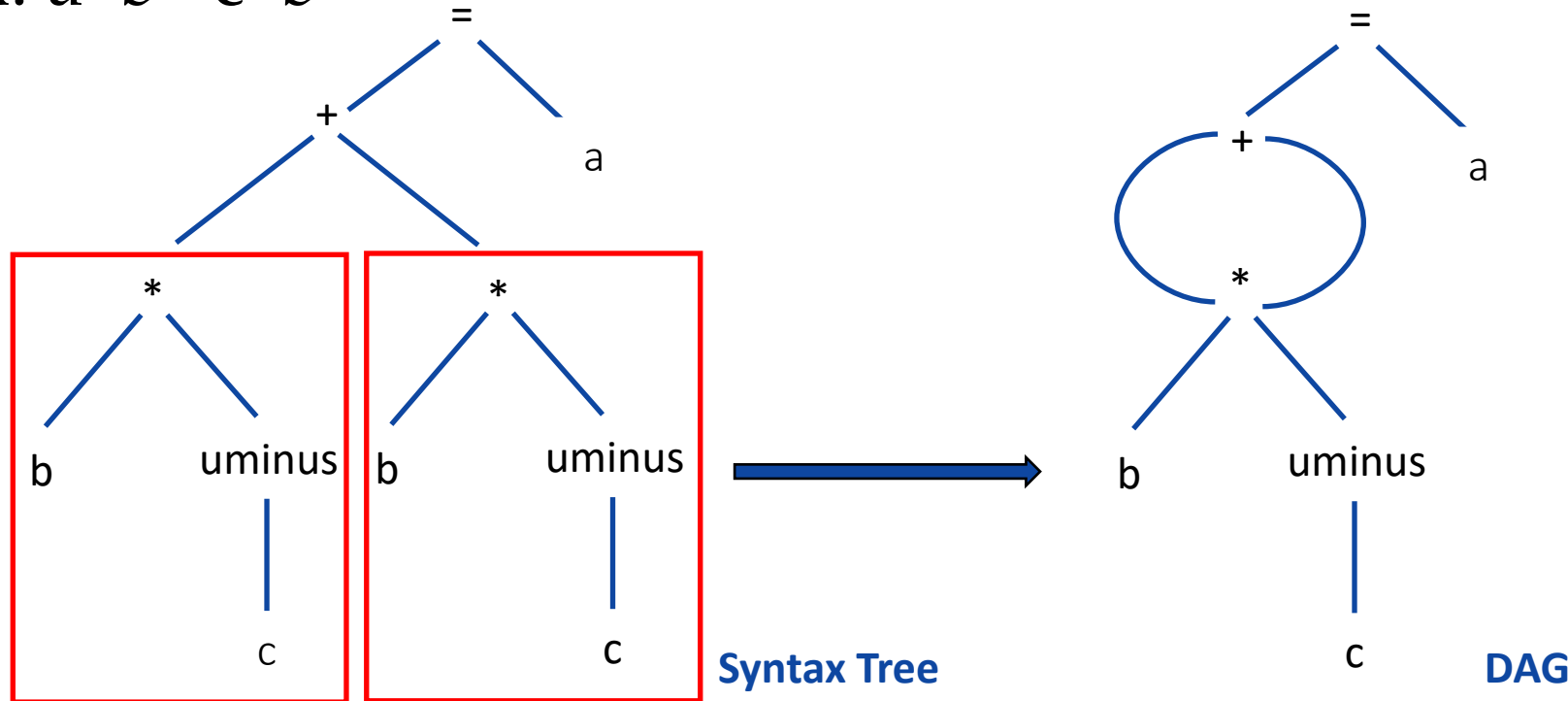
# Module 4 – Intermediate Code Generation

# Different Intermediate Forms

- Abstract syntax tree
- Postfix notation
- Three address code

# Abstract syntax tree & DAG

- A syntax tree depicts the natural hierarchical structure of a source program.
- A DAG (Directed Acyclic Graph) gives the same information but in a more **compact** way because **common sub-expressions** are identified.
- Ex:  $a = b^* - c + b^*$ .



# Postfix Notation

- Postfix notation is a linearization of a syntax tree.
- In postfix notation the operands occurs first and then operators are arranged.

• Ex:  $(A + B) * (C + D)$

Postfix notation:  $A B + C D + *$

• Ex:  $(A + B) * C$

Postfix notation:  $A B + C *$

• Ex:  $(A * B) + (C * D)$

Postfix notation:  $A B * C D * +$

# Three address code

- Three address code is a sequence of statements of the general form,

$$a := b \text{ op } c$$

- Where  $a$ ,  $b$  or  $c$  are the operands that can be names or constants and  $op$  stands for any operator.
- Example:  $a = b + c + d$

$$t_1 = b + c$$

$$t_2 = t_1 + d$$

$$a = t_2$$

- Here  $t_1$  and  $t_2$  are the temporary names generated by the compiler.
- There are **at most three addresses allowed** (two for operands and one for result). Hence, this representation is called three-address code.

# Different Representation of Three Address Code

- There are three types of representation used for three address code:

1. Quadruples
2. Triples
3. Indirect triples

- Ex:  $x = -a * b + -a * b$

$$t_1 = -a$$

$$t_2 = t_1 * b$$

$$t_3 = -a$$

$$t_4 = t_3 * b$$

$$t_5 = t_2 + t_4$$

$$x = t_5$$

# Quadruple

- The quadruple is a structure with at the most four fields such as op, arg1, arg2 and result.
  - The op field is used to represent the internal code for operator.
  - The arg1 and arg2 represent the two operands.
  - And result field is used to store the result of an expression.
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## Quadruple

$x = -a * b + -a * b$

$t_1 = -a$

$t_2 = t_1 * b$

$t_3 = -a$

$t_4 = t_3 * b$

$t_5 = t_2 + t_4$

$x = t_5$

No.	Operator	Arg1	Arg2	Result
(0)				
(1)				
(2)				
(3)				
(4)				
(5)				

# Triple

- To avoid entering temporary names into the symbol table, we might refer a temporary value by the position of the statement that computes it.
  - If we do so, three address statements can be represented by records with only three fields: op, arg1 and arg2.
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**Quadruple**

No.	Operator	Arg1	Arg2	Result
(0)	uminus	a		t <sub>1</sub>
(1)	*	t <sub>1</sub>	b	t <sub>2</sub>
(2)	uminus	a		t <sub>3</sub>
(3)	*	t <sub>3</sub>	b	t <sub>4</sub>
(4)	+	t <sub>2</sub>	t <sub>4</sub>	t <sub>5</sub>
(5)	=	t <sub>5</sub>		x

**Triple**

No.	Operator	Arg1	Arg2
(0)			
(1)			
(2)			
(3)			
(4)			
(5)			



# Indirect Triple

- In the indirect triple representation the listing of triples has been done. And listing pointers are used instead of using statement.
  - This implementation is called indirect triples.
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**Triple**

No.	Operator	Arg1	Arg2
(0)	uminus	a	
(1)	*	(0)	b
(2)	uminus	a	
(3)	*	(2)	b
(4)	+	(1)	(3)
(5)	=	x	(4)

**Indirect Triple**

	Statement
(0)	(14)
(1)	(15)
(2)	(16)
(3)	(17)
(4)	(18)
(5)	(19)

No.	Operator	Arg1	Arg2
(0)	uminus	a	
(1)	*		b
(2)	uminus	a	
(3)	*		b
(4)	+		
(5)	=	x	

# Exercise

Write quadruple, triple and indirect triple for following:

1.  $-(a*b)+(c+d)$

2.  $a*-(b+c)$

3.  $x=(a+b*c)^(d*e)+f*g^h$

4.  $z=g+a*(b-c)+(x-y)*d$