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UNIVERSITY**

“JnanaSangama”, Belgaum -590014, Karnataka.



LAB REPORT

on

Artificial Intelligence (23CS5PCAIN)

Submitted by

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in partial fulfillment for the award of the degree of
BACHELOR OF ENGINEERING
in
COMPUTER SCIENCE AND ENGINEERING



B.M.S. COLLEGE OF ENGINEERING
(Autonomous Institution under VTU)

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CERTIFICATE

This is to certify that the Lab work entitled “Artificial Intelligence (23CS5PCAIN)” carried out by **Spandana M R (1BM23CS337)**, who is bona fide student of **B.M.S. College of Engineering**. It is in partial fulfillment for the award of **Bachelor of Engineering in Computer Science and Engineering** of the Visvesvaraya Technological University, Belgaum. The Lab report has been approved as it satisfies the academic requirements in respect of an Artificial Intelligence (23CS5PCAIN) work prescribed for the said degree.

Sonika Sharma D Assistant Professor Department of CSE, BMSCE	Dr. Kavitha Sooda Professor & HOD Department of CSE, BMSCE
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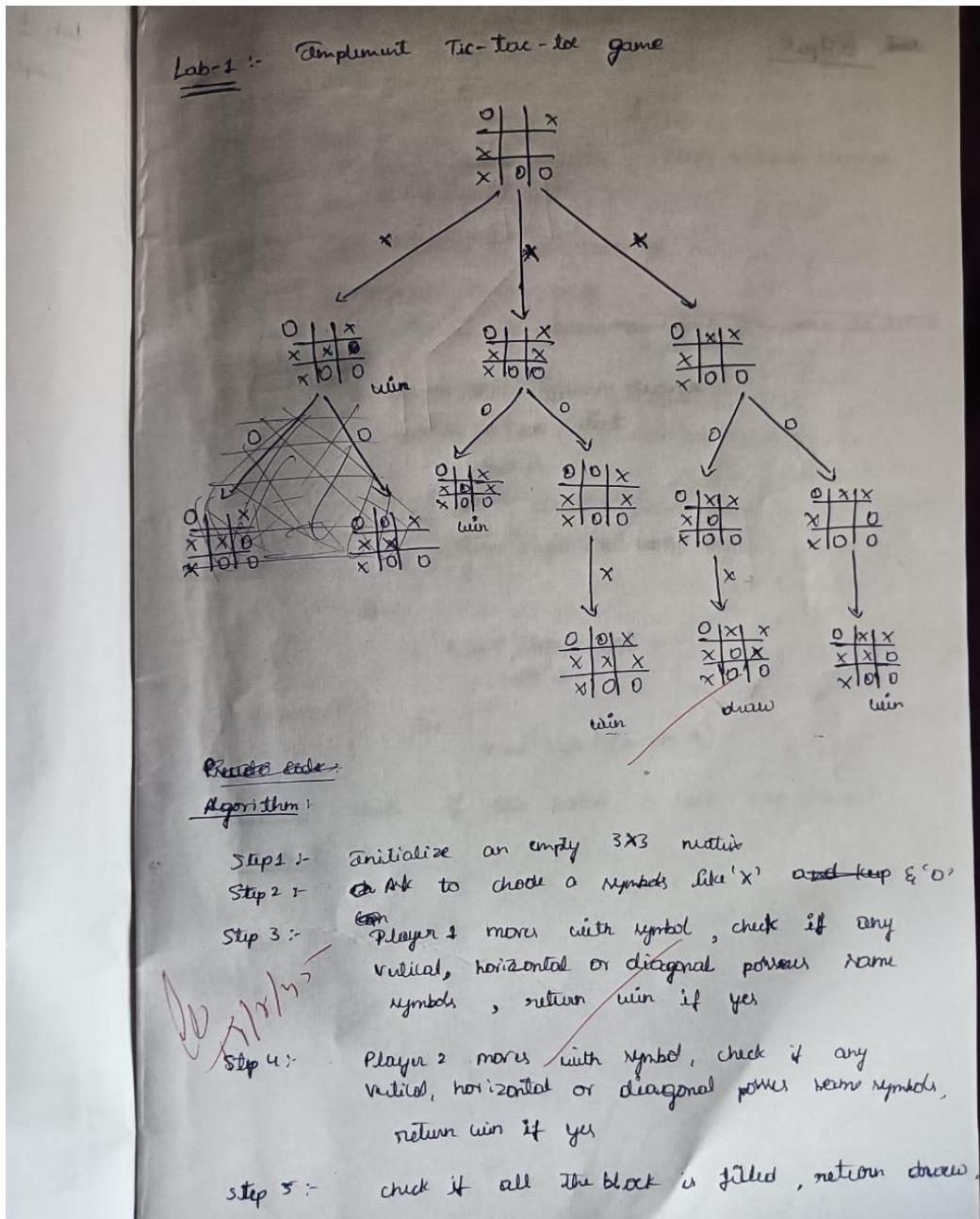
Github Link:

<https://github.com/Spandana-mr/AI>

Program 1

Implement Tic-Tac-Toe Game I
Implement vacuum cleaner agent

Tic-Tac-Toe:



Code:

```
def print_board(board):
    for row in board:
        print(" | ".join(row))
        print("-" * 9)
```

```

def check_winner(board, player):
    # Check rows, columns and diagonals
    for i in range(3):
        if all([cell == player for cell in board[i]]) or \
           all([board[j][i] == player for j in range(3)]):
            return True

    if all([board[i][i] == player for i in range(3)]) or \
       all([board[i][2 - i] == player for i in range(3)]):
        return True

    return False

def is_full(board):
    return all(cell in ['X', 'O'] for row in board for cell in row)

def get_move(player):
    while True:
        try:
            move = input(f"Player {player}, enter your move (row and column: 1 1): ")
            row, col = map(int, move.split())
            if row in [1, 2, 3] and col in [1, 2, 3]:
                return row - 1, col - 1
            else:
                print("Invalid input. Enter numbers between 1 and 3.")
        except ValueError:
            print("Invalid input. Enter two numbers separated by space.")

def play_game():
    board = [[" " for _ in range(3)] for _ in range(3)]
    current_player = "X"

    while True:
        print_board(board)
        row, col = get_move(current_player)

        if board[row][col] != " ":
            print("That spot is taken. Try again.")
            continue

        board[row][col] = current_player

        if check_winner(board, current_player):
            print_board(board)
            print(f"Player {current_player} wins!")

```

```

break

if is_full(board):
    print_board(board)
    print("It's a draw!")
    break

current_player = "O" if current_player == "X" else "X"

if __name__ == "__main__":
    play_game()

```

OUTPUT:

```

| | |
| | |
| | |
-----
Player X, enter your move (row and column: 1 1): 1 1
X | |
| | |
| | |
-----
Player O, enter your move (row and column: 1 1): 1 2
X | O |
| | |
| | |
-----
Player X, enter your move (row and column: 1 1): 1 3
X | O | X
| | |
| | |
-----
Player O, enter your move (row and column: 1 1): 2 2
X | O | X
| O |
| | |
-----
Player X, enter your move (row and column: 1 1): 3 3
X | O | X
| O |
| | X |
-----
Player O, enter your move (row and column: 1 1): 3 1
X | O | X
| O |
O |   | X
-----
Player X, enter your move (row and column: 1 1): 2 3
X | O | X
| O | X
O |   | X
-----
Player X wins!

```

Vacuum Cleaner:

Lab 2- Vacuum cleaner

Algorithm:-

Step 1 :- Start

Step 2 :- construct 1×2 matrix ; keep vacuum cleaner in location A.

Step 3 :- If location A has dirt

if vacuum is in locA
ask if it clean the locA or move to locB

if locA has vacuum cleaner
if locA has dirt
suck it

move right (to locB)

else
if locB has dirt
suck it

move left (to locA)

Step 4 :- check if both locA & locB are clean

Step 5 :- End

~~cost = $2 \times 2^2 = 8$~~

25/02/2022

Output:-

Enter state of A1 (0 for clean, 1 for dirty): 1
 Enter state of A2 (0 for clean, 1 for dirty): 1
 Enter state of B1 (0 for clean, 1 for dirty): 1
 Enter state of B2 (0 for clean, 1 for dirty): 1

Enter starting location (A1, A2, B1, B2): A1

Starting cleaning at A1 and A2 has been cleaned.
 Cleaned A1.

Is A1 clean now? (0 if clean, 1 if dirty): 0

Moving vacuum to A2

Cleaned A2.

Is A2 clean now? (0 if clean, 1 if dirty): 0

Moving vacuum to B2

Cleaned B2.

Is B2 clean now? (0 if clean, 1 if dirty): 0

Moving vacuum to B1

Cleaned B1.

Is B1 clean now? (0 if clean, 1 if dirty): 0

finished cleaning all rooms.

Total cost: 7

{'A1': 0, 'A2': 0, 'B1': 0, 'B2': 0}

Code:

```
def vacuum_simulation():
    cost = 0

    # Get initial states and location
    state_A = int(input("Enter state of A (0 for clean, 1 for dirty): "))
    state_B = int(input("Enter state of B (0 for clean, 1 for dirty): "))
    location = input("Enter location (A or B): ").upper()

    # Vacuum operation loop
    while True:
        if location == 'A':
            if state_A == 1:
                print("Cleaning A.")
                state_A = 0
                cost += 1
            else:
                print("Room A is already clean")
        else:
            if state_B == 1:
                print("Cleaning B.")
                state_B = 0
                cost += 1
            else:
                print("Room B is already clean")
```

```

state_A = 0
cost += 1

elif state_B == 1:
    print("Moving vacuum right")
    location = 'B'
    cost += 1

else:
    print("Turning vacuum off")
    break

elif location == 'B':
    if state_B == 1:
        print("Cleaning B.")
        state_B = 0
        cost += 1

    elif state_A == 1:
        print("Moving vacuum left")
        location = 'A'
        cost += 1

    else:
        print("Turning vacuum off")
        break

print(f'Cost: {cost}')
print(f'{{"A": {state_A}, "B": {state_B}}}')

```

vacuum_simulation()

OUTPUT:

```
Enter state of A (0 for clean, 1 for dirty): 1
Enter state of B (0 for clean, 1 for dirty): 0
Enter location (A or B): A
Cleaning A.
Turning vacuum off
Cost: 1
{'A': 0, 'B': 0}
```

Program 2

Implement 8 puzzle problems using Depth First Search (DFS)
Implement Iterative deepening search algorithm

8 puzzle problems using Depth First Search (DFS)

Lab 3 :- Using BFS solve 8 puzzle without heuristic approach.

2	8	3
1	6	4
7	5	

Initial

1	3	3
6		4
2	6	5

goal

Pseudo code / Algorithm

- S1 :- Given initial & goal state
- S2 :- Construct 3×3 grid with one empty tile
- S3 :- move empty tiles from all possible ways for the given grid.
- S4 :- perform tile movement for each grid from left to right.
- S5 :- Move to next row after completion of all the grid in every row.
- S6 :- write found, if the grid match the goal state.
- S7 :- write no. of states took to reach the goal state.

Output :-

Step 0 :-

2	8	3
1	6	4
7	5	

Step 1 :-

2	8	3
1	0	4
7	6	5

Step 2 :-

2	0	3
1	8	4
7	6	5

Step 3 :-

0	2	3
1	3	4
7	6	5

Step 4 :-

1	2	3
0	8	4
7	6	5

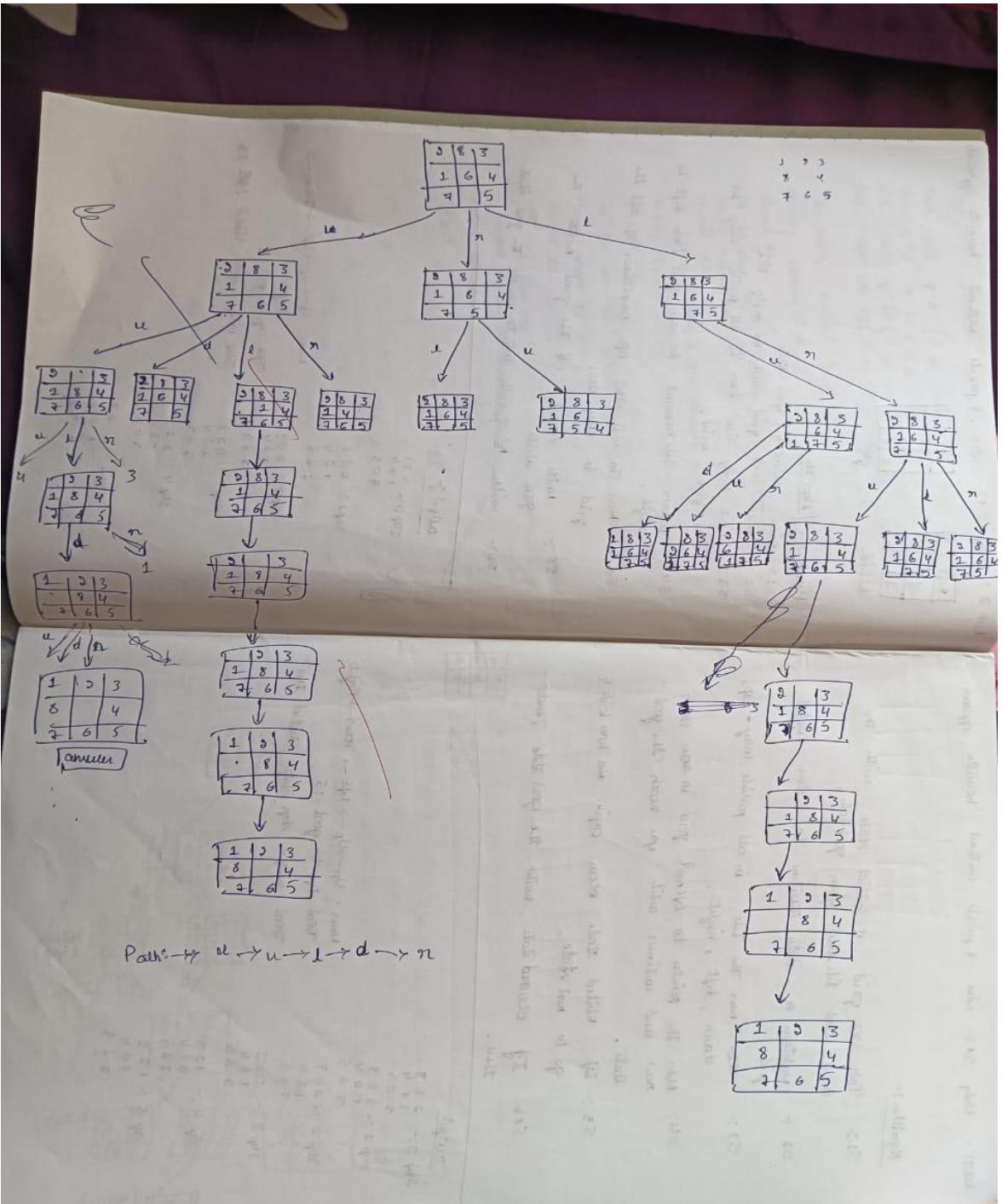
Step 5 :-

1	2	3
8	0	4
7	6	5

Movs: Up \rightarrow Up \rightarrow Left \rightarrow Down \rightarrow Right

Total steps to goal : 5

Total unique states visited : 35



Lab3:- Using DFS solve 8 puzzle without heuristic approach.

Algorithm :-

S1:- Take 3x3 grid of initial state with one empty tile and given goal state.

S2:- Maintain a data structure to store visited grid.

S3:- Move the tile in all possible way - up, down, left, right.

S4:- Move the pointer to leftmost grid in each row and continue until you reach the goal state.

S5:- If visited state occur again no back & go to next state.

S6:- If occurred state, match the goal state, end there.

Output

Step 0:-
2 8 3
1 6 4
7 0 5

Step 1:-
9 8 3
1 0 4
7 6 5

Step 2:-
0 0 3
1 8 4
7 6 5

Step 3:-
0 2 3
1 8 4
7 6 5

Step 4:-
1 2 3
0 8 4
7 6 5

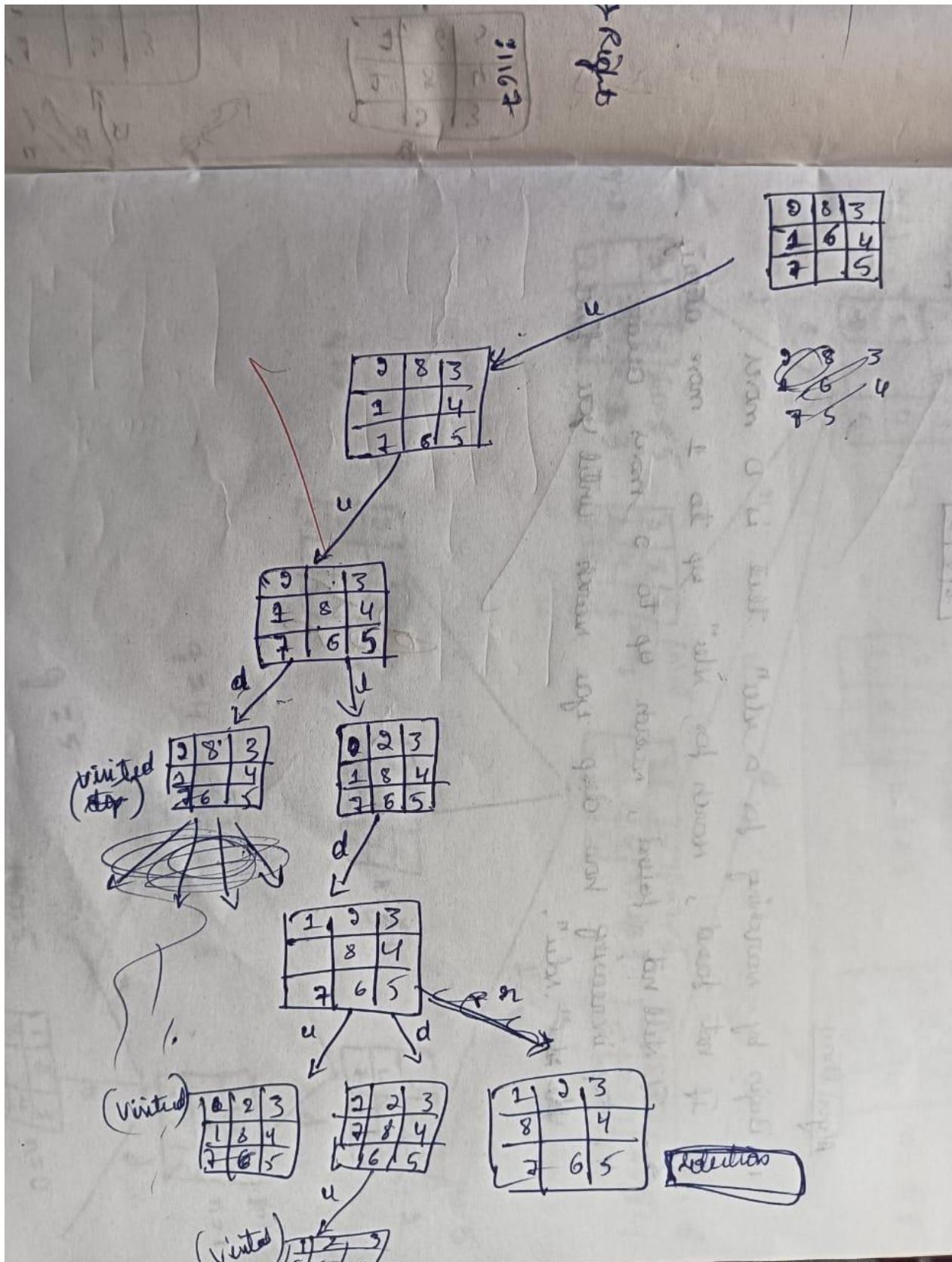
Step 5:-
1 0 3
8 0 4
7 6 5

Note: Up → Up → Left → Down → Right

Total steps to goal : 5

Total unique steps states visited : 16

3	2	1
0	6	4
7	5	8



Code:

```
Usig DFS 8 puzzel without heuristic
# Goal state
goal = ((1, 2, 3),
         (8, 0, 4),
         (7, 6, 5))

# Moves: Up, Down, Left, Right
moves = [(-1, 0), (1, 0), (0, -1), (0, 1)]

def get_neighbors(state):
    # Find the empty tile (0)
    for i in range(3):
        for j in range(3):
            if state[i][j] == 0:
                x, y = i, j
                break

    neighbors = []
    for dx, dy in moves:
        nx, ny = x + dx, y + dy
        if 0 <= nx < 3 and 0 <= ny < 3:
            # Swap empty tile with adjacent tile
            new_state = [list(row) for row in state]
            new_state[x][y], new_state[nx][ny] = new_state[nx][ny], new_state[x][y]
            neighbors.append(tuple(tuple(row) for row in new_state))
    return neighbors

def dfs_limited(start, depth_limit):
    stack = [(start, [start])]
    visited = set([start])

    while stack:
        current, path = stack.pop()

        if current == goal:
            return path

        if len(path) - 1 >= depth_limit: # already reached depth limit
            continue

        for neighbor in get_neighbors(current):
            if neighbor not in visited:
                visited.add(neighbor)
                stack.append((neighbor, path + [neighbor]))

    return None
```

```

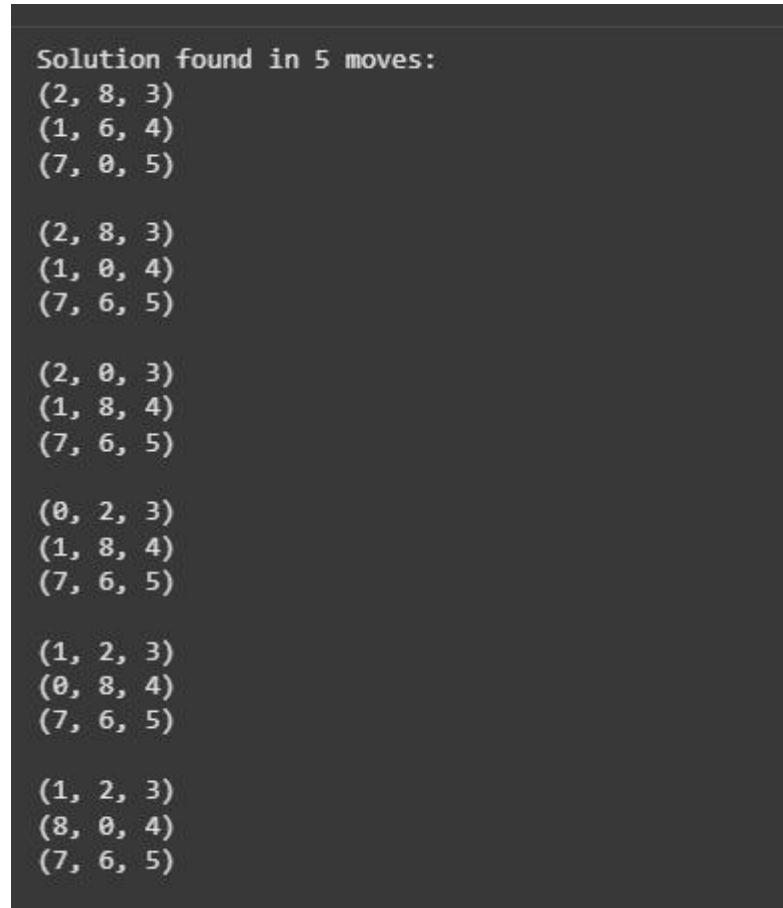
# Example start state
start = ((2, 8, 3),
          (1, 6, 4),
          (7, 0, 5))

solution_path = dfs_limited(start, depth_limit=5)

if solution_path:
    print(f"Solution found in {len(solution_path) - 1} moves:")
    for state in solution_path:
        for row in state:
            print(row)
        print()
else:
    print("No solution found within 5 moves.")

```

OUTPUT:



The image shows a terminal window with a dark background and light-colored text. It displays the output of a Python script that performs a depth-limited search for a solution to a puzzle. The output starts with a message indicating a solution was found in 5 moves. It then lists the states of the puzzle at each step of the search, starting from the initial state and moving through several intermediate states before reaching a goal state.

```

Solution found in 5 moves:
(2, 8, 3)
(1, 6, 4)
(7, 0, 5)

(2, 8, 3)
(1, 0, 4)
(7, 6, 5)

(2, 0, 3)
(1, 8, 4)
(7, 6, 5)

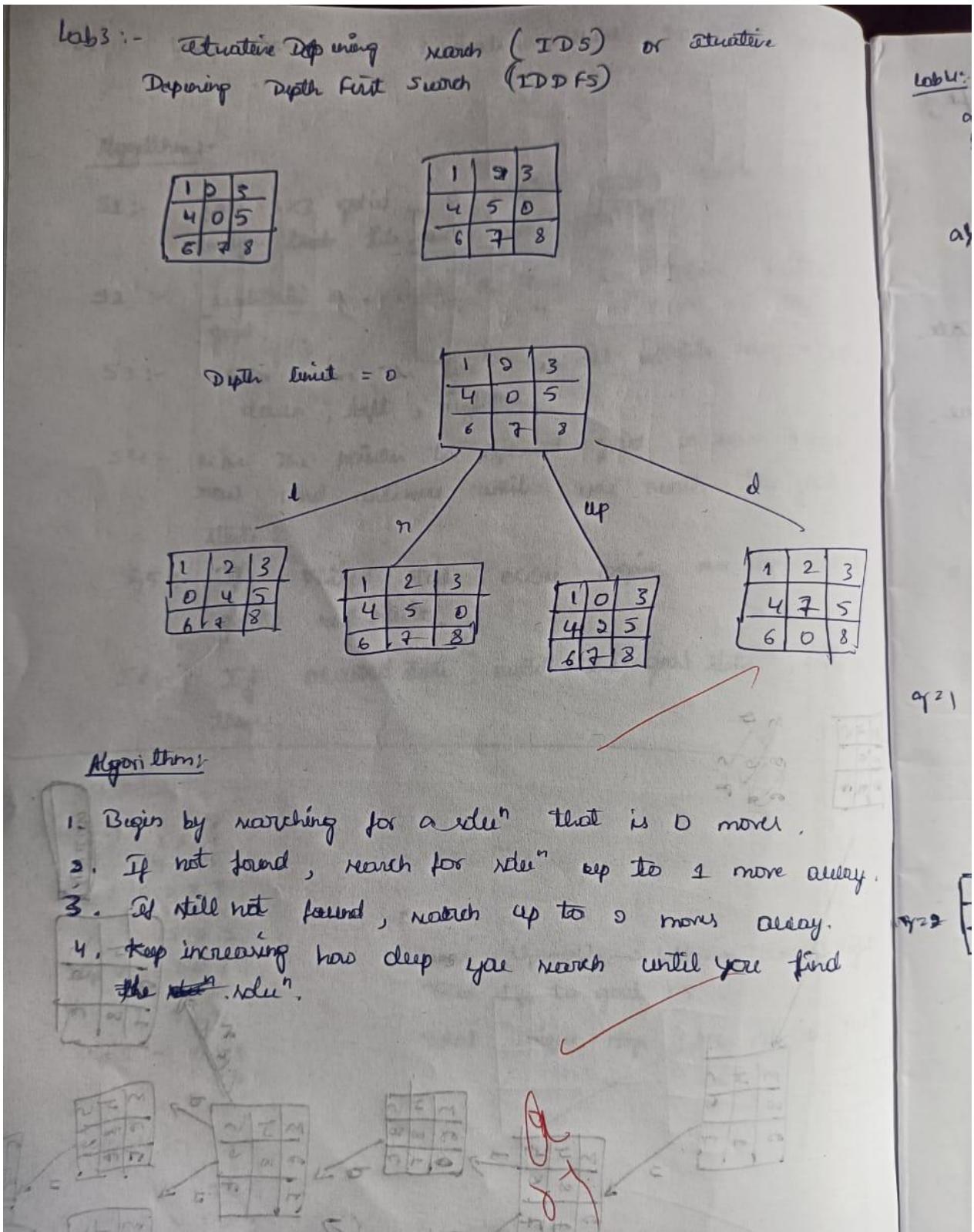
(0, 2, 3)
(1, 8, 4)
(7, 6, 5)

(1, 2, 3)
(0, 8, 4)
(7, 6, 5)

(1, 2, 3)
(8, 0, 4)
(7, 6, 5)

```

Iterative Deepening Search(IDS) or Iterative Deepening Depth First Search(IDDFS)



Code:

```
From copy import deepcopy
GOAL_STATE = [
    [1, 2, 3],
    [4, 5, 0],
    [6, 7, 8]
]

# Possible moves of the blank (0) tile: up, down, left, right
MOVES = [(-1, 0), (1, 0), (0, -1), (0, 1)]

def find_blank(state):
    for i in range(3):
        for j in range(3):
            if state[i][j] == 0:
                return i, j

def is_goal(state):
    return state == GOAL_STATE

def get_neighbors(state):
    neighbors = []
    x, y = find_blank(state)

    for dx, dy in MOVES:
        nx, ny = x + dx, y + dy
        if 0 <= nx < 3 and 0 <= ny < 3:
            new_state = deepcopy(state)
            # Swap blank with neighbor
            new_state[x][y], new_state[nx][ny] = new_state[nx][ny], new_state[x][y]
            neighbors.append(new_state)

    return neighbors
```

```

neighbors.append(new_state)

return neighbors

def dfs(state, depth, limit, path, visited):
    if is_goal(state):
        return path + [state]
    if depth == limit:
        return None

    for neighbor in get_neighbors(state):
        # To avoid cycles, do not revisit states in current path
        if neighbor not in visited:
            result = dfs(neighbor, depth + 1, limit, path + [state], visited + [neighbor])
            if result is not None:
                return result
    return None

def iterative_deepening_search(initial_state, max_depth=50):
    for depth_limit in range(max_depth):
        print(f"Searching with depth limit = {depth_limit}")
        result = dfs(initial_state, 0, depth_limit, [], [initial_state])
        if result is not None:
            return result
    return None

def print_state(state):
    for row in state:
        print(' '.join(str(x) for x in row))
    print()

```

```

if __name__ == "__main__":
    initial_state = [
        [1, 2, 3],
        [4, 0, 5],
        [6, 7, 8]
    ]

    solution = iterative_deepening_search(initial_state)

    if solution:
        print(f"Solution found in {len(solution)-1} moves!")
        for step, state in enumerate(solution):
            print(f"Step {step}:")
            print_state(state)
    else:
        print("No solution found.")

```

OUTPUT

```

Searching with depth limit = 0
Searching with depth limit = 1
Solution found in 1 moves!
Step 0:
1 2 3
4 0 5
6 7 8

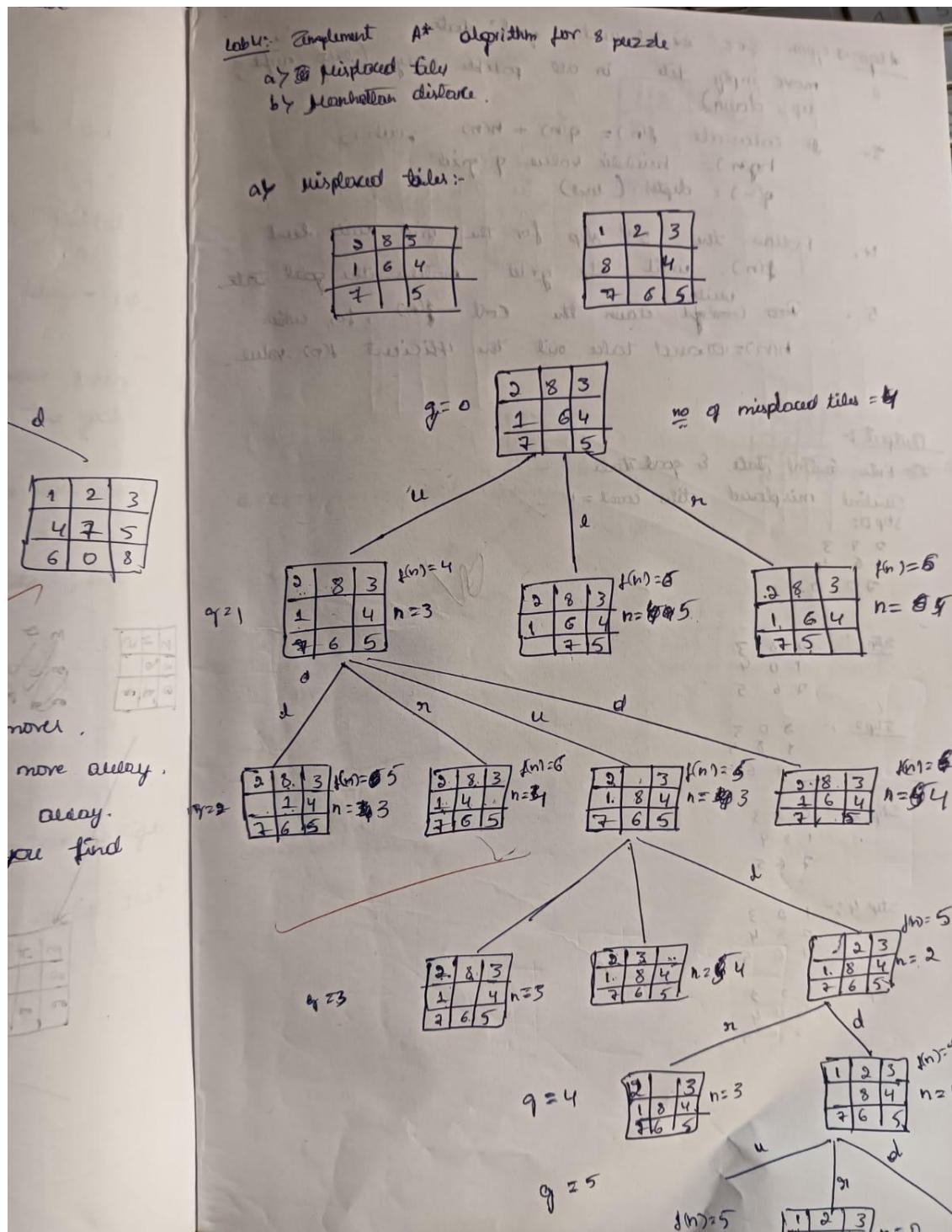
Step 1:
1 2 3
4 5 0
6 7 8

```

Program 3

Implement A* search algorithm

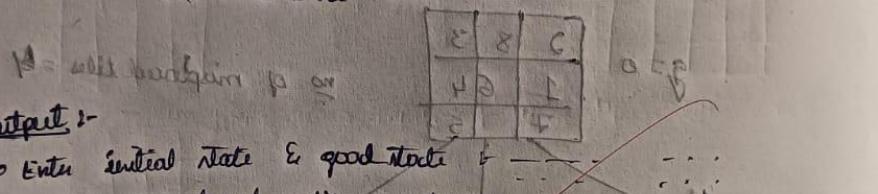
1. Misplace Tiles



Algorithm:- Given 3×3 initial and goal state
 1. move empty tile in all possible way (left, right, up, down)

2. calculate $f(n) = g(n) + h(n)$ where,
 $h(n)$ = heuristic value of grid
 $g(n)$ = depth (level)

3. Follow the 3rd step for the grid with least $f(n)$, until the grid matches the goal state.
4. Repeat with the cost $f(n)$, for which $h(n)=0$ and take out the efficient $f(n)$ value.



Output :-

1. Enter initial state & goal state
 Initial misplaced tiles count = 4

Step 0:

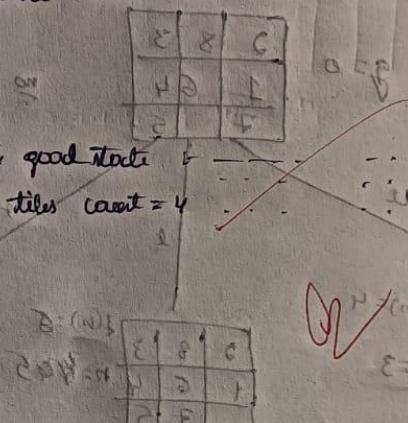
0	8	3
1	6	4
7	0	5

$d = 0$

$\Delta B = N$

Step 1:

0	8	3
1	0	4
7	6	5



Step 2 :-

0	8	3
1	8	4
7	6	5

$\Delta B = N$

$\Delta C = N$

$\Delta E = N$

Step 3 :-

0	8	3
1	8	4
7	6	5

3	8	C
H	0	1
2	F	

Step 4 :-

1	8	3
0	8	4
7	6	5

$\Delta B = N$

$\Delta C = N$

$\Delta E = N$

Step 5 :-

1	8	3
8	0	4
7	6	5

$H = P$

```

Code:
import heapq

# Goal state

goal = ((1, 2, 3),
         (8, 0, 4),
         (7, 6, 5))

# Moves: Up, Down, Left, Right

moves = [(-1, 0), (1, 0), (0, -1), (0, 1)]


# Heuristic: Misplaced tiles

def misplaced_tiles(state):
    count = 0

    for i in range(3):
        for j in range(3):
            if state[i][j] != 0 and state[i][j] != goal[i][j]:
                count += 1

    return count


# Find blank position (0)

def find_blank(state):
    for i in range(3):
        for j in range(3):
            if state[i][j] == 0:
                return i, j


# Generate neighbors

def get_neighbors(state):
    neighbors = []

    x, y = find_blank(state)

    for dx, dy in moves:
        nx, ny = x + dx, y + dy

```

```

if 0 <= nx < 3 and 0 <= ny < 3:
    new_state = [list(row) for row in state]
    new_state[x][y], new_state[nx][ny] = new_state[nx][ny], new_state[x][y]
    neighbors.append(tuple(tuple(row) for row in new_state))

return neighbors

# A* Search

def astar(start):
    pq = []
    heapq.heappush(pq, (misplaced_tiles(start), 0, start, []))
    visited = set()

    while pq:
        f, g, state, path = heapq.heappop(pq)
        if state == goal:
            return path + [state]
        if state in visited:
            continue
        visited.add(state)
        for neighbor in get_neighbors(state):
            if neighbor not in visited:
                new_g = g + 1
                new_f = new_g + misplaced_tiles(neighbor)
                heapq.heappush(pq, (new_f, new_g, neighbor, path + [state]))


    return None

```

```

# Example usage

start_state = ((2, 8, 3),
               (1, 6, 4),
               (7, 0, 5))

```

```
solution = astar(start_state)
```

```
# Print solution path
```

```
for step in solution:
```

```
    for row in step:
```

```
        print(row)
```

```
        print("-----")
```

OUTPUT

```
(2, 8, 3)  
(1, 6, 4)  
(7, 0, 5)
```

```
-----  
(2, 8, 3)  
(1, 0, 4)  
(7, 6, 5)
```

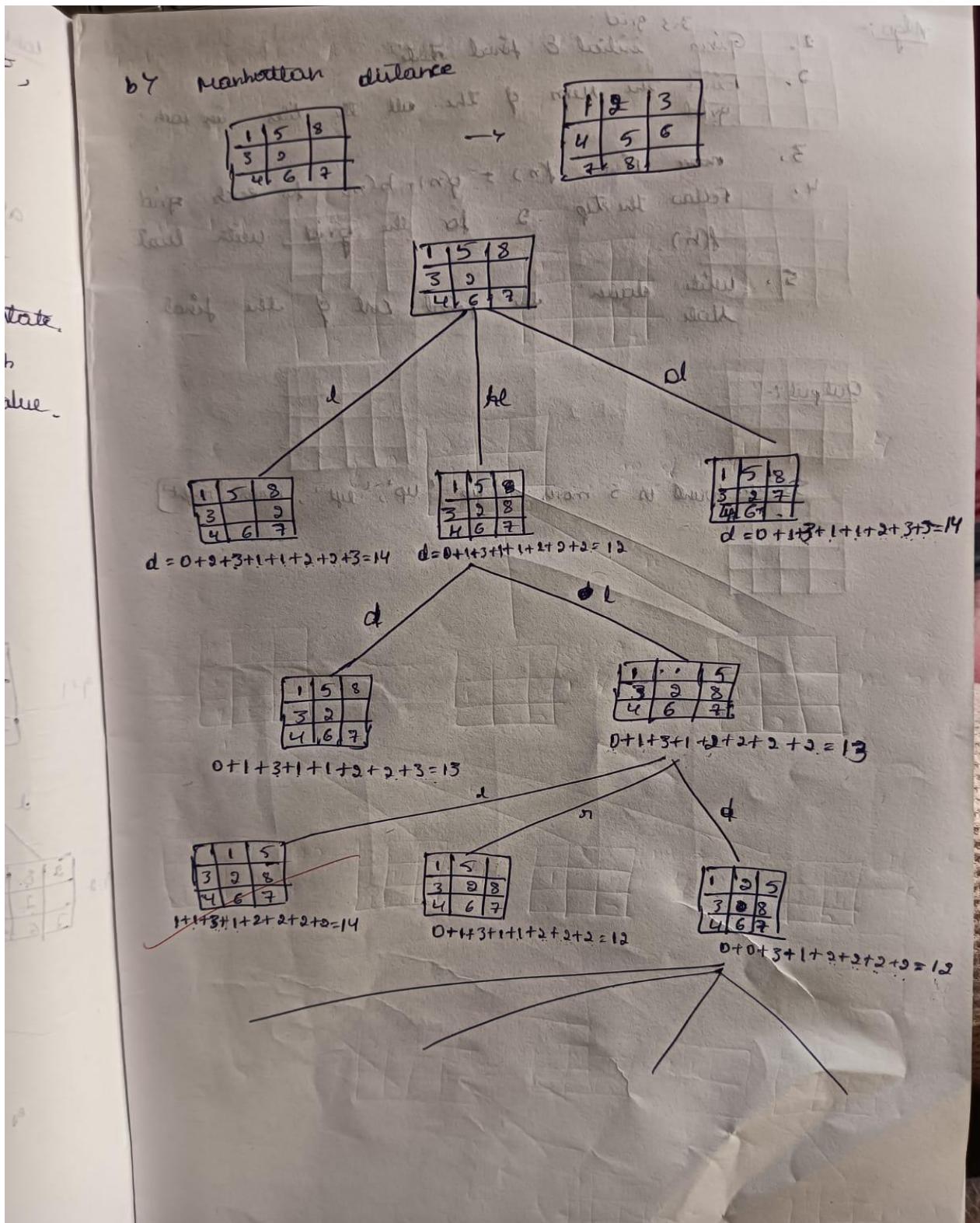
```
-----  
(2, 0, 3)  
(1, 8, 4)  
(7, 6, 5)
```

```
-----  
(0, 2, 3)  
(1, 8, 4)  
(7, 6, 5)
```

```
-----  
(1, 2, 3)  
(0, 8, 4)  
(7, 6, 5)
```

```
-----  
(1, 2, 3)  
(8, 0, 4)  
(7, 6, 5)
```

Manhattan:



- Algo:-
1. Given initial & final state
 2. Find the sum of the all the tiles in each grid
 3. Then find $f(n) = g(n) + h(n)$ for each grid
 4. Follow the step 2 for the grid with least $f(n)$.
 5. Write down the total cost of the final state.

Output:-

Solution found in 5 moves: ['up', 'up', 'left', 'down', 'right']

$$f_1 = C_1 + G_1 + H_1 + F_1 + D_1 = 0 \quad f_2 = C_2 + G_2 + H_2 + F_2 + D_2 = 0 \quad f_3 = C_3 + G_3 + H_3 + F_3 + D_3 = 0$$

Code:

```
import heapq
goal = ((1, 2, 3),
        (8, 0, 4),
        (7, 6, 5))
```

```
moves = [(-1, 0), (1, 0), (0, -1), (0, 1)] # Up, Down, Left, Right
```

```
def manhattan_distance(state):
```

```
    distance = 0
```

```
    for i in range(3):
```

```
        for j in range(3):
```

```
            value = state[i][j]
```

```

if value != 0:
    # goal position of this tile
    goal_x = (value - 1) // 3
    goal_y = (value - 1) % 3
    distance += abs(i - goal_x) + abs(j - goal_y)

return distance

def find_blank(state):
    for i in range(3):
        for j in range(3):
            if state[i][j] == 0:
                return i, j

def get_neighbors(state):
    neighbors = []
    x, y = find_blank(state)
    for dx, dy in moves:
        nx, ny = x + dx, y + dy
        if 0 <= nx < 3 and 0 <= ny < 3:
            new_state = [list(row) for row in state]
            new_state[x][y], new_state[nx][ny] = new_state[nx][ny], new_state[x][y]
            neighbors.append(tuple(tuple(row) for row in new_state))

    return neighbors

def astar(start):
    pq = []
    heapq.heappush(pq, (manhattan_distance(start), 0, start, [])) # (f, g, state, path)
    visited = set()
    while pq:
        f, g, state, path = heapq.heappop(pq)
        if state == goal:
            return path + [state]
        if state in visited:

```

```

        continue

visited.add(state)

for neighbor in get_neighbors(state):

    if neighbor not in visited:

        new_g = g + 1

        new_f = new_g + manhattan_distance(neighbor)

        heapq.heappush(pq, (new_f, new_g, neighbor, path + [state]))


return None

start_state = ((2, 8, 3),
               (1, 6, 4),
               (7, 0, 5))

solution = astar(start_state)

if solution is None:

    print("No solution found.")

else:

    print("Solution path:")

    for step in solution:

        for row in step:

            print(row)

            print("-----")

```

OUTPUT:

Solution path:

(2, 8, 3)

(1, 6, 4)

(7, 0, 5)

(2, 8, 3)

(1, 0, 4)

(7, 6, 5)

(2, 0, 3)

(1, 8, 4)

(7, 6, 5)

(0, 2, 3)

(1, 8, 4)

(7, 6, 5)

(1, 2, 3)

(0, 8, 4)

(7, 6, 5)

(1, 2, 3)

(8, 0, 4)

(7, 6, 5)

Program 4

Implement Hill Climbing search algorithm to solve N-Queens problem

Lab Study Hill Climbing To solve 4 Queens :-

on each

b grid
b least
final
right
ret = 0

Step:

- 1 : Start
- 2 : An initial state of the 4 Queens arranged in 4*4 grid is arranged.
- 3 : The next states change the position of the Queens such that it doesn't get attacked. Fill all the positions till so the cost is minimum.
- 4 : Calculate the cost for each step / ~~not written~~ and consider the least cost.
- 5 : continue step 4 till the queens are arranged such that there is no attack to anyone from any side.
- 6 : Stop.

Output:-

~~order initial state : 3, 1, 2, 0
(All 16 counts with their costs).~~

initial state $x_1 = 3, x_2 = 1, x_3 = 2, x_4 = 0$

$\text{cost} = 3$

$\text{cost} = 1$

$\text{cost} = 2$

$\text{cost} = 0$

$\text{cost} = 6$

Code:

```
import random
```

```
def compute_cost(state):
    n = len(state)
    cost = 0
    for i in range(n):
        for j in range(i + 1, n):
            if state[i] == state[j]:          # same row
                cost += 1
            elif abs(state[i] - state[j]) == abs(i - j): # same diagonal
                cost += 1
    return cost
```

```
def get_neighbors(state):
    neighbors = []
    n = len(state)
    for col in range(n):      # pick a column
        for row in range(n):  # try moving queen in this column to another row
            if row != state[col]:
                new_state = state.copy()
                new_state[col] = row
                neighbors.append(new_state)
    return neighbors
```

```
def print_board(state):
```

```
    n = len(state)
```

```

for r in range(n):
    line = ""
    for c in range(n):
        line += "Q " if state[c] == r else ". "
    print(line)
print("")

def hill_climb(initial_state, max_sideways=50):
    current = initial_state
    current_cost = compute_cost(current)
    steps = 0
    sideways_moves = 0

    print("Initial State (cost={}):".format(current_cost))
    print_board(current)

    while True:
        neighbors = get_neighbors(current)
        costs = [compute_cost(n) for n in neighbors]
        min_cost = min(costs)

        if min_cost > current_cost:
            # no better neighbor -> stop
            break

        # pick one of the best neighbors randomly
        best_neighbors = [n for n, c in zip(neighbors, costs) if c == min_cost]
        next_state = random.choice(best_neighbors)
        next_cost = compute_cost(next_state)

```

```

# handle sideways moves

if next_cost == current_cost:

    if sideways_moves >= max_sideways:

        break

    else:

        sideways_moves += 1

else:

    sideways_moves = 0


current = next_state

current_cost = next_cost

steps += 1


print("Step {} (cost={}):".format(steps, current_cost))
print_board(current)

if current_cost == 0:

    print("Solution found in {} steps ✅".format(steps))

    return current


print("Local minimum reached (cost={}) ❌".format(current_cost))

return current


initial_state = [3, 1, 2, 0]

final = hill_climb(initial_state, max_sideways=10)

```

OUTPUT:

Initial State (cost=2):

. . . Q
. Q . .
. . Q .
Q . . .

Step 1 (cost=2):

. . . Q
Q Q . .
. . Q .
. . . .

Step 2 (cost=1):

. . . Q
Q . . .
. . Q .
. Q . .

Step 3 (cost=1):

. . Q Q
Q . . .
. . . .
. Q . .

Step 4 (cost=0):

. . Q .
Q . . .
. . . Q
. Q . .

Solution found in 4 steps ✓

Program 5

Simulated Annealing to Solve 8-Queens problem

By Simulated Annealing

algo. description

Algorithm :-

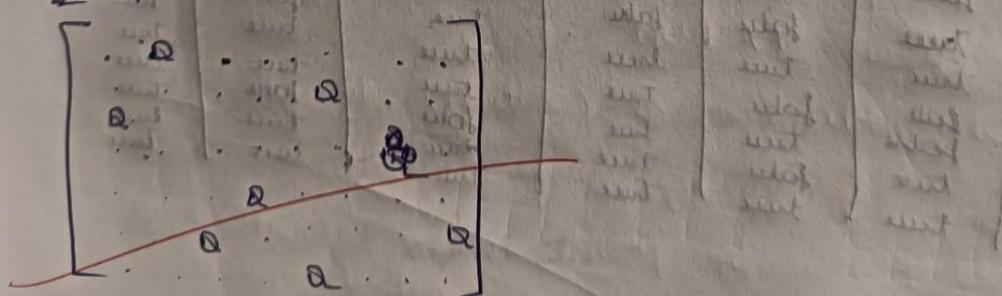
1. current \leftarrow initial state
2. $T \leftarrow$ a large positive value
3. while $T > 0$ do
4. next \leftarrow a random neighbour of current
5. $\Delta E \leftarrow$ current cost $-$ next cost
6. if $\Delta E > 0$ then
current \leftarrow next
7. else
current \leftarrow next with probability $e^{\frac{-\Delta E}{T}}$
8. end if
9. decrease T
10. end while
11. return current.

Output

But position found : [2, 0, 6, 4, 7, 1, 3, 5]

No. of non attacking pairs = 28

board:



Code:

```
from scipy.optimize import dual_annealing
import numpy as np

def queens_max(x):
    cols = np.round(x).astype(int)
    n = len(cols)

    if len(set(cols)) < n:
        return 1e6

    attacks = 0

    for i in range(n):
        for j in range(i + 1, n):
            if abs(i - j) == abs(cols[i] - cols[j]):
                attacks += 1

    return attacks

n = 8
bounds = [(0, n - 1)] * n
result = dual_annealing(queens_max, bounds)

best_cols = np.round(result.x).astype(int).tolist()
not_attacking = n

print(f"The best position found is: {best_cols}")
print(f"The number of queens that are not attacking each other is: {not_attacking}")
```

OUTPUT:

```
The best position found is: [7, 4, 6, 1, 3, 5, 0, 2]
The number of queens that are not attacking each other is: 8
```

Program 6

Create a knowledge base using propositional logic and show that the given query entails the knowledge base or not.

Lab 6

22/9/25

Propositional logic

→ Truth table for connectives.

P	Q	$\neg P$	and $P \wedge Q$	or $P \vee Q$	xor $P \Leftrightarrow Q$
false	false	true	false	false	true
false	true	true	false	true	false
true	false	false	false	true	false
true	true	false	true	true	false

→ Propositional inference : Enumeration method.

$$\alpha = A \vee B$$

$$KB = (A \vee C) \wedge (B \vee \neg C)$$

Checking that $KB \models \alpha$

A	B	C	$A \vee C$	$B \vee \neg C$	KB	α
false	false	false	false	true	false	false
false	false	true	true	true	false	false
false	true	false	true	false	false	true
false	true	true	true	true	true	true
true	false	false	true	true	true	true
true	false	true	true	true	true	true
true	true	false	true	true	true	true
true	true	true	true	true	true	true

Algorithms

- Start
- The KB (Knowledge base), α (propositional logic), a query and a query is given.
- Calculate the values for both KB & α based on the given condition.
- Check in truth table if both the KB & α values are true.
- If both values $\neq B$ and α are false or either of them will false they are ignored.

For

false
false
true
true

⇒

⇒

7. Assume a P value where is true. Once it's true
 once then recur with the rest. Once E false
 8. Combine the results with and operation.
 9. If both branches are false then the whole check
 succeed.
 10. Stop iteration as soon as it is true.

true

false

false

false

true

Output

A	B	C	KB	alpha	KB \ alpha
true	true	true	true	true	true
true	true	false	false	true	false
true	false	true	true	true	true
true	false	false	true	true	true
false	true	true	true	true	true
false	true	false	false	true	true
false	false	true	false	false	false
false	false	false	false	true	false
false	false	false	false	false	false

Models where KB & alpha are true:

{'A': true, 'B': true, 'C': true}

{'A': true, 'B': false, 'C': true}

{'A': true, 'B': false, 'C': false}

{'A': false, 'B': true, 'C': true}

Does KB entail alpha? true if false has no

error $a: s \wedge (s \vee t), b: \neg(s \wedge t) \Rightarrow c = (t \vee \neg t)$

S	T	A	B	C
false	false	true	false	true
false	true	false	false	true
true	false	false	false	true
true	true	false	true	true

$\Rightarrow a \text{ entails } b: \text{false}$

$\Rightarrow a \text{ entails } c: \text{true}$ (but $s: \text{false}$ & $t: \text{false}$)

22/3/22

Code:

```
import itertools
from sympy import symbols, sympify

A, B, C = symbols('A B C')

alpha_input = input("Enter alpha (example: A | B): ")
kb_input = input("Enter KB (example: (A | C) & (B | ~C)): ")

alpha = sympify(alpha_input, evaluate=False)
kb = sympify(kb_input, evaluate=False)

GREEN = "\033[92m"
RESET = "\033[0m"

print(f"\nTruth Table for \alpha = {alpha_input}, KB = {kb_input}\n")
print(f"{'A':<6} {'B':<6} {'C':<6} {'\alpha':<10} {'KB':<10}")

entailed = True

for values in itertools.product([False, True], repeat=3):
    subs = {A: values[0], B: values[1], C: values[2]}
    alpha_val = alpha.subs(subs)
    kb_val = kb.subs(subs)

    alpha_str = f"\033[92m{alpha_val}\033[0m" if kb_val else str(alpha_val)
    kb_str = f"\033[92m{kb_val}\033[0m" if kb_val else str(kb_val)

    print(f"\n{str(values[0]):<6} {str(values[1]):<6} {str(values[2]):<6}"
          f"\n{alpha_str:<10} {kb_str:<10}")
```

```

if kb_val and not alpha_val:
    entailed = False

if entailed:
    print(f"\n KB |= α holds (KB entails α)\n")
else:
    print(f"\n KB does NOT entail α\n")

```

OUTPUT:

```

Enter alpha (example: A | B): A|B
Enter KB (example: (A | C) & (B | ~C)): (A | C) & (B | ~C)

Truth Table for α = A|B, KB = (A | C) & (B | ~C)

A   B   C   α      KB
False False False False  False
False False True  False  False
False True  False True   False
False True  True   True   TrueTrue
True  False False True   TrueTrue
True  False True  True   False
True  True  False True   TrueTrue
True  True  True   True   TrueTrue

KB |= α holds (KB entails α)

```

Program 7

Implement unification in first order logic

Lab 7:

Algorithm:

Step 1: If ψ_1 or ψ_2 is a variable or constant then:

- a) If ψ_1 or ψ_2 are identical, then return NIL.
- b) Else if ψ_1 is a variable,
 - a. Then if ψ_1 occurs in ψ_2 , then return FAILURE.
 - b. Else return $\{(\psi_2 / \psi_1)\}$
- c) Else if ψ_2 is a variable,
 - a. If ψ_2 occurs in ψ_1 , then return FAILURE.
 - b. Else return $\{(\psi_1 / \psi_2)\}$
- d) Else return FAILURE.

Step 2: If the initial predicate symbol in ψ_1 and ψ_2 are not same, then return FAILURE.

Step 3: If ψ_1 and ψ_2 have a different no. of arguments, then return FAILURE.

Step 4: Set substitution set (SUBST) to NIL.

Step 5: For $i = 1$ to the no. of elements

- a. Call unify ψ_1^i with the i^{th} element of ψ_1 and i^{th} element of ψ_2 , and put the result into s .
- b. If $s = \text{failure}$ then return FAILURE.
- c. If $s \neq \text{NIL}$ then do,
 - a. Apply s to the remainder of both ψ_1 and ψ_2 .
 - b. SUBST = APPEND (s , SUBST).

Step 6: Return SUBST

Output:

Unification succeeded; { 'x': 'B', 'y': 'A' }

solve the following:

1) Find Most General Unifier (M.G.U) of $\{P(b, x, f(g(z)))$
and $P(z, f(y), f(y))\}$

$$\Rightarrow \{P(b, x, f(g(z))) \quad ? \\ P(z, f(y), f(y)) \quad ?\}$$

both have arguments &
predicates but same

then:

NIL.

FAILURE.

FAILURE.

d_1, d_2 are

arguments,

$$so, b = z$$

$$x = f(y)$$

$$f(g(z)) = f(y)$$

$$\therefore g(z) = y$$

$$\Rightarrow z = b$$

$$x = f(y)$$

$$y = g(z)$$

$$x = f(y) \text{ but } y = g(z)$$

$$\{x \neq f(g(z))\}$$

$$P(z, f(y), f(g(z))) \Rightarrow b = z, x = f(y)$$

$$P(z, f(y), f(g(z))) \Rightarrow f(y) = f(g(y))$$

$$\text{Unifier} = \{b/z, x/f(y), y/g(z)\}$$

2) Find M.G.U of $\{Q(a, g(x, a), f(y))$ and

$$\Rightarrow Q(a, g(f(b), a), x)\}$$

a

$$\Rightarrow Q(a, g(x, a), f(y))$$

$$a (a, g(f(b), a), b)$$

$$Q(a, g(f(b), a), f(y))$$

$$Q(a, g(f(b), a), b)$$

$$x = f(b)$$

$$f(y) = x$$

$$f(y) = f(b)$$

$$\therefore y = b$$

$$Q(a, g(f(b), a), f(b))$$

~~$$Q(a, g(f(b), a), f(b))$$~~

$$\text{Unifier} : \{x/f(b)\}$$

3) Find MGU of $\{ p(f(a), g(y)), p(x, x) \}$

$$p(f(a), g(y))$$

$$p(x, x)$$

If $x = f(a)$ and $x = g(y)$ then $x = f(a) = g(y) \Rightarrow f(a) \neq g(y)$
 it is not possible
 \therefore NO Unifiers

4) Unify $\{ \text{prime}(11) \text{ and } \text{prime}(y) \}$

$\text{prime}(11)$ both have same predicate & same no. of arguments.
 $\text{prime}(y)$

$\therefore y = 11$
 $\text{prime}(11)$ sub: $\{ y/u \}$
 $\text{prime}(u)$ $(s)u = v$ and $(v)t = x$

5) Unify $\{ \text{knows}(\text{John}, x), \text{knows}(y, \text{mother}(y)) \}$

$\text{knows}(\text{John}, x)$ \leftarrow
 $\text{knows}(y, \text{mother}(y))$

Let $y = \text{John}$ and $x = \text{mother}(y)$
 $x = \text{mother}(\text{John})$
 $\therefore \text{knows}(\text{John}, \text{mother}(\text{John})) = \text{knows}(\text{John}, \text{mother}(\text{John}))$

6) Unify $\{ \text{knows}(\text{John}, x), \text{knows}(y, \text{Bill}) \}$

$\text{knows}(\text{John}, x)$ Let $y = \text{John}$
 $\text{knows}(y, \text{Bill})$ $x = \text{Bill}$

$\text{knows}(\text{John}, \text{Bill}) = \text{knows}(\text{John}, \text{Bill})$

Code:

```
def occurs_check(var, expr):
    if var == expr:
        return True
    if isinstance(expr, tuple):
        return any(occurs_check(var, sub) for sub in expr[1:]) # Skip function symbol
    return False

def substitute(expr, subst):
    if isinstance(expr, str):
        # Follow substitution chain until fully resolved
        while expr in subst:
            expr = subst[expr]
        return expr
    # If it's a function term: (f, arg1, arg2, ...)
    return (expr[0],) + tuple(substitute(sub, subst) for sub in expr[1:])

def unify(Y1, Y2, subst=None):
    if subst is None:
        subst = {}
    Y1 = substitute(Y1, subst)
    Y2 = substitute(Y2, subst)

    # Case 1: identical
    if Y1 == Y2:
        return subst

    # Case 2: Y1 is variable
    if isinstance(Y1, str):
        if occurs_check(Y1, Y2):
            return "FAILURE"
        subst[Y1] = Y2
```

```

return subst

# Case 3: Y2 is variable
if isinstance(Y2, str):
    if occurs_check(Y2, Y1):
        return "FAILURE"
    subst[Y2] = Y1
return subst

# Case 4: function mismatch
if Y1[0] != Y2[0] or len(Y1) != len(Y2):
    return "FAILURE"

# Case 5: unify arguments
for a, b in zip(Y1[1:], Y2[1:]):
    subst = unify(a, b, subst)
    if subst == "FAILURE":
        return "FAILURE"

return subst

expr1 = ("p", "X", ("f", "Y"))
expr2 = ("p", "a", ("f", "b"))

output = unify(expr1, expr2)
print(output)

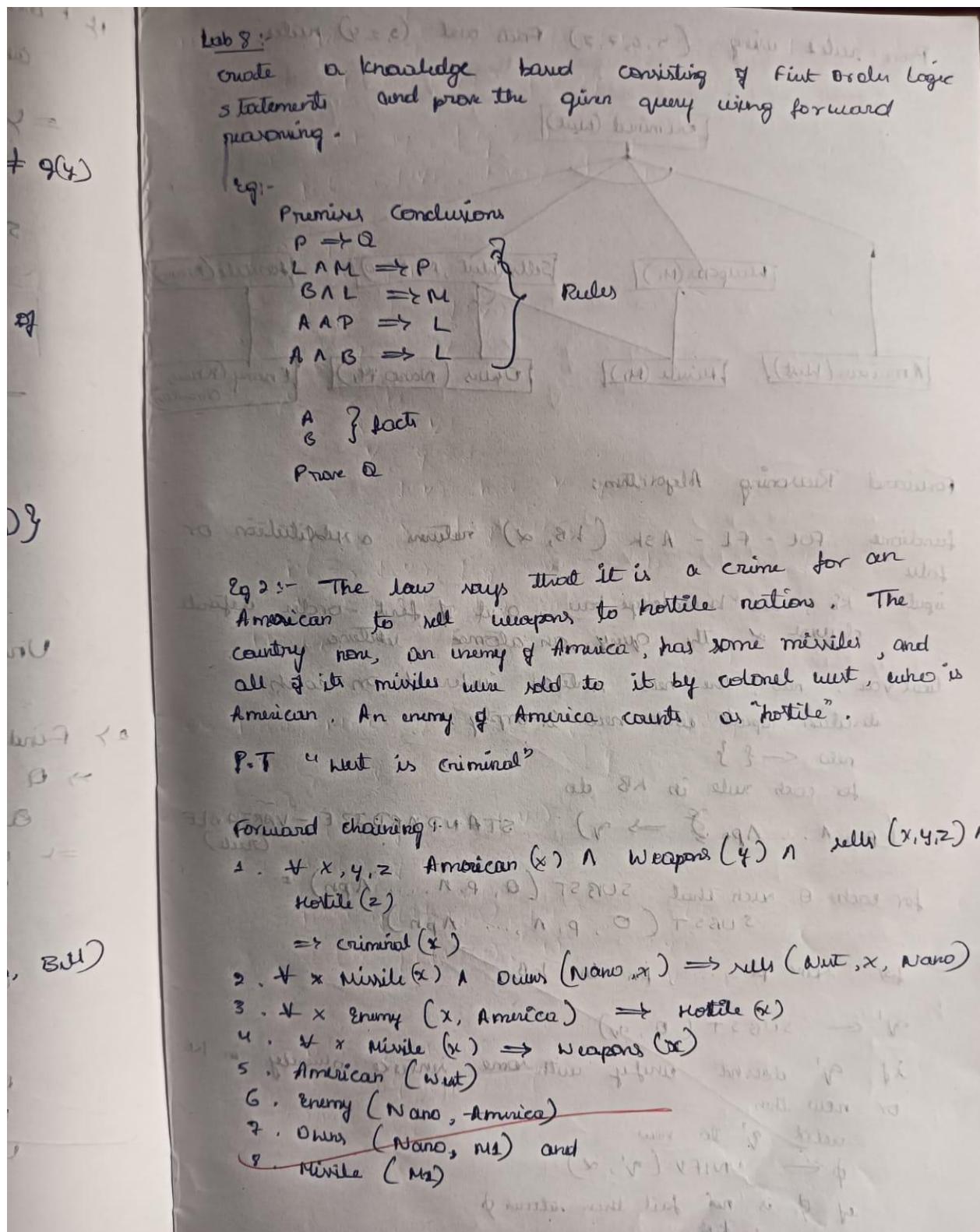
```

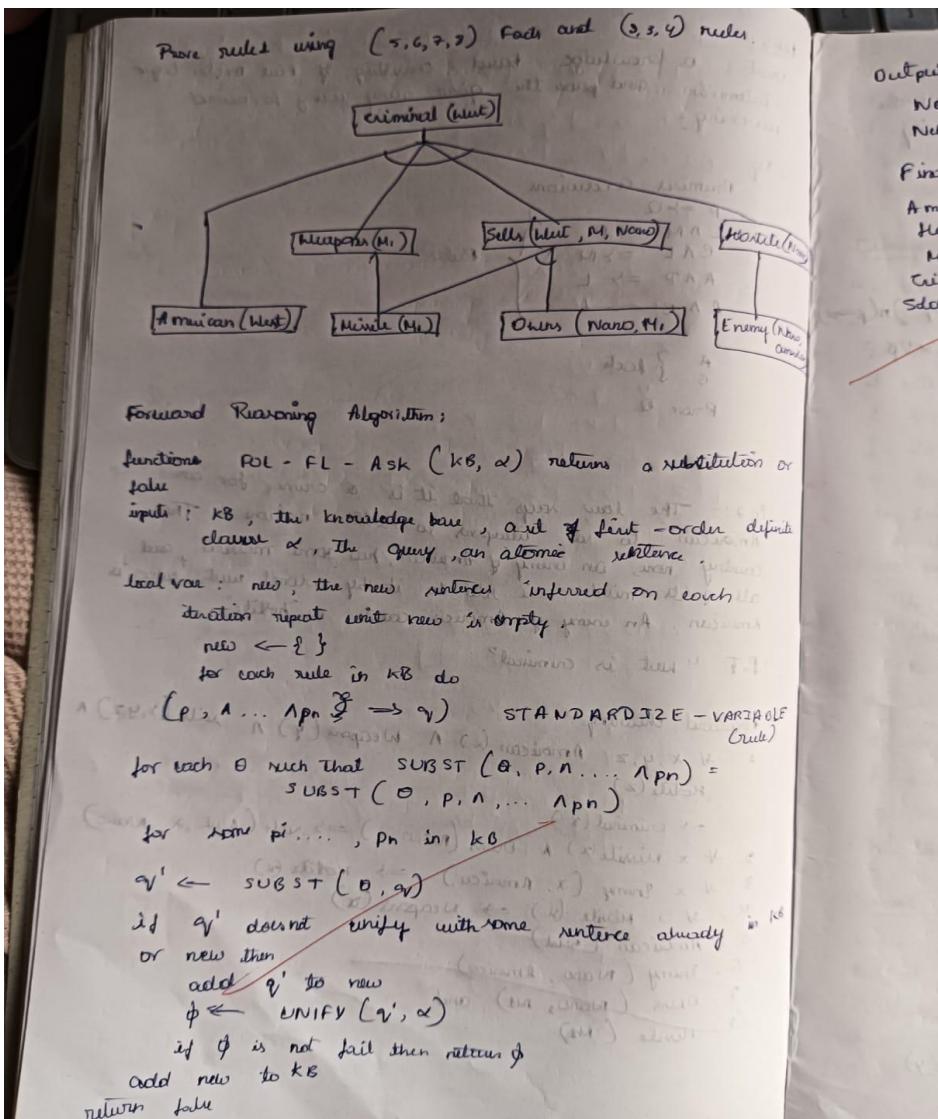
OUTPUT:

```
{'X': 'a', 'Y': 'b'}
```

Program 8

Create a knowledge base consisting of first order logic statements and prove the given query using forward reasoning.





Output:

New fact inferred : criminal (west)

New fact inferred : SdolWeapons (West, None)

Final facts :

- American (West)
- Hostile (None)
- Missile (None)
- Criminal (West)
- Sdol Weapons (West, None)

Sdol Weapons (West, None)

27/10/22

Code:

```
from collections import deque

class KnowledgeBase:

    def __init__(self):
        self.facts = set()
        self.rules = []
        self.inferred = set()

    def add_fact(self, fact):
        if fact not in self.facts:
            print(f"Adding fact: {fact}")
            self.facts.add(fact)
            return True
        return False

    def add_rule(self, premises, conclusion):
        self.rules.append((premises, conclusion))

    def forward_chain(self):
        agenda = deque(self.facts)

        while agenda:
            fact = agenda.popleft()
            if fact in self.inferred:
                continue
            self.inferred.add(fact)

            for (premises, conclusion) in self.rules:
                if all(p in self.inferred for p in premises):
                    if conclusion not in self.facts:
```

```

print(f"Inferred new fact: {conclusion} from {premises} => {conclusion}")

self.facts.add(conclusion)

agenda.append(conclusion)

```

```

if conclusion == 'Criminal(West)':
    print("\n Goal Reached: West is Criminal")
    return True

```

```
return False
```

```
kb = KnowledgeBase()
```

```

kb.add_fact('American(West)')
kb.add_fact('Enemy(Nono, America)')
kb.add_fact('Missile(M1)')
kb.add_fact('Owns(Nono, M1)')

```

```
kb.add_rule(premises=['Missile(M1)'], conclusion='Weapon(M1)')
```

```
kb.add_rule(premises=['Missile(M1)', 'Owns(Nono, M1)'], conclusion='Sells(West, M1, Nono)')
```

```

kb.add_rule(premises=['Enemy(Nono, America)'], conclusion='Hostile(Nono)')
kb.add_rule(premises=['American(West)', 'Weapon(M1)', 'Sells(West, M1, Nono)', 'Hostile(Nono)'],
            conclusion='Criminal(West)')
kb.forward_chain()

```

OUTPUT:

```

Adding fact: American(West)
Adding fact: Enemy(Nono, America)
Adding fact: Missile(M1)
Adding fact: Owns(Nono, M1)
Inferred new fact: Weapon(M1) from ['Missile(M1)'] => Weapon(M1)
Inferred new fact: Hostile(Nono) from ['Enemy(Nono, America)'] => Hostile(Nono)
Inferred new fact: Sells(West, M1, Nono) from ['Missile(M1)', 'Owns(Nono, M1)'] => Sells(West, M1, Nono)
Inferred new fact: Criminal(West) from ['American(West)', 'Weapon(M1)', 'Sells(West, M1, Nono)', 'Hostile(Nono)'] => Criminal(West)

```

```

 Goal Reached: West is Criminal
True

```

Program 9

Create a knowledge base consisting of first order logic statements and prove the given query using Resolution

Lab-9 :-

create a knowledge base consisting of first order logic statements and prove the given query using Resolution

Algo:-

1. Eliminate biconditionals & implications:
 - a. Eliminates \leftrightarrow replacing $\alpha \leftrightarrow \beta$ with $(\alpha \rightarrow \beta) \wedge (\beta \rightarrow \alpha)$.
 - b. Eliminates \rightarrow replacing $\alpha \rightarrow \beta$ with $\neg \alpha \vee \beta$.
2. Now \neg increases.
 - a. $\neg(\forall x p) = \exists x \neg p,$
 - b. $\neg(\exists x p) = \forall x \neg p.$
 - c. $\neg(\alpha \vee \beta) = \neg \alpha \wedge \neg \beta$

Proof by Resolution :-

Create KB or Premises

- $\neg \rightarrow \text{John likes all kind of food}$
- $\neg \rightarrow \text{Apple and vegetables are food.}$
- $\neg \rightarrow \text{Anything anyone eats and not killed is food}$
- $\neg \rightarrow \text{Anil eats peanut \& still alive.}$
- $\neg \rightarrow \text{Harvy eats everything that Anil eats.}$
- $\neg \rightarrow \text{Anyone who is alive implies not killed.}$
- $\neg \rightarrow \text{Anyone who is not killed implies alive.}$

Prove by Resolution that :-

John likes peanut.

Representation in FOL

- a. $\forall x ; \text{food}(x) \rightarrow \text{likes}(\text{John}, x)$
- b. $\text{food}(\text{Apple}) \wedge \text{food}(\text{vegetables})$
- c. $\forall x \forall y : \text{eats}(x, y) \wedge \neg \text{killed}(x) \rightarrow \text{food}(y)$

- d. $\neg \neg$
e. -
f. +
g. +
h. -
- Eliminate
- a. \forall
 - b. fo
 - c. \forall
 - d. \neg
 - e. \neg
 - f. \neg
 - g. +
 - h. -

- More
- a. $\forall x$
 - b. fo
 - c. $\forall x$
 - d. \neg
 - e. $\forall x$
 - f. \neg
 - g. +
 - h. \neg

- Residue
- a. $\forall x$
 - b. food
 - c. $\forall y \neg$
 - d. \neg
 - e. $\forall N$

int. order
using

- d. eats (Anil, Peanuts) \wedge alive (Anil)
- e. $\forall x : \text{eats}(\text{Anil}, x) \rightarrow \text{eats}(\text{Harry}, x)$
- f. $\forall x : \neg \text{killed}(x) \rightarrow \text{alive}(x)$
- g. $\forall x : \text{alive}(x) \rightarrow \neg \text{killed}(x)$
- h. likes (John, Peanuts)

Elimination Implication

- a. $\forall x \rightarrow \text{food}(x) \vee \text{likes}(\text{John}, x)$
- b. food (Apple) \wedge food (Vegetable)
- c. $\forall x \forall y \rightarrow [\text{eats}(x, y) \wedge \neg \text{killed}(x)] \vee \text{food}(y)$
- d. eats (Anil, Peanuts) \wedge alive (Anil)
- e. $\forall x \rightarrow \text{eats}(\text{Anil}, x) \vee \text{eats}(\text{Harry}, x)$
- f. $\forall x \rightarrow [\neg \text{killed}(x)] \vee \text{alive}(x)$
- g. $\forall x \rightarrow \text{alive}(x) \vee \neg \text{killed}(x)$
- h. likes (John, Peanuts)

More negation (\neg) inwards and rewrite

- a. $\forall x \rightarrow \text{food}(x) \vee \text{likes}(\text{John}, x)$
- b. food (Apple) \wedge food (Vegetable)
- c. $\forall x \forall y \rightarrow \neg \text{eats}(x, y) \vee \text{killed}(x) \vee \text{food}(y)$
- d. eats (Anil, Peanuts) \wedge alive (Anil)
- e. $\forall x \rightarrow \text{eats}(\text{Anil}, x) \vee \text{eats}(\text{Harry}, x)$
- f. $\forall x [\neg \text{killed}(x)] \vee \text{alive}(x)$
- g. $\forall x \rightarrow \text{alive}(x) \vee \neg \text{killed}(x)$
- h. likes (John, Peanuts)

Rename Variable or standardize variables

- a. $\forall x \rightarrow \text{food}(x) \vee \text{likes}(\text{John}, x)$
- b. food (Apple) \wedge food (Vegetable)
- c. $\forall y \forall z \rightarrow \text{eats}(y, z) \vee \text{killed}(y) \vee \text{food}(z)$
- d. eats (Anil, Peanuts) \wedge alive (Anil)
- e. $\forall N \rightarrow \neg \text{eats}(\text{Anil}, N) \vee \text{eats}(\text{Harry}, N)$

(y)

- f. $\neg q \text{ killed } (q) \rightarrow v \text{ alive } (q)$
 g. $\neg k \rightarrow \text{alive } (k) \rightarrow \text{killed } (k)$
 h. $\text{Likes } (\text{John}, \text{Peanuts}) \rightarrow \text{killed } (x) \rightarrow \neg v \text{ alive } (x)$

Lemma:
NEE

Drop variable:

- a. $\rightarrow \text{food}(x) \vee \text{Likes } (\text{John}, x)$
 b. $\text{Food } (\text{Apple})$
 c. $\text{Food } (\text{Vegetables})$
 d. $\rightarrow \text{eats } (y, z) \vee \text{killed } (y) \vee \text{Food } (z)$
 e. $\text{eats } (\text{Anil}, \text{Peanuts}) \rightarrow \text{killed } (x) \rightarrow \neg v \text{ alive } (x)$
 f. $\text{alive } (\text{Anil})$
 g. $\rightarrow \text{eats } (\text{Anil}, w) \vee \text{eats } (\text{Harry}, w)$
 h. $\text{killed } (g) \vee \text{alive } (g)$
 i. $\rightarrow \text{alive } (k) \vee \rightarrow \text{killed } (k)$
 j. $\text{Likes } (\text{John}, \text{Peanuts}) \rightarrow \text{killed } (x) \rightarrow \neg v \text{ alive } (x)$

$\rightarrow \text{Likes } (\text{John}, \text{Peanuts}) \quad \rightarrow \text{Food } (x) \vee \text{Likes } (\text{John}, x)$

Drop variable: $\{ \text{Peanuts} / x \}$

$\rightarrow \text{Food } (\text{Peanuts}) \quad \rightarrow \text{eats } (y, z) \vee \text{killed } (y) \vee \text{Food } (z)$

Drop variable: $\{ \text{Peanuts} / z \}$

$\rightarrow \text{Food } (v) \vee \text{Food } (w) \quad \{ \text{Peanuts} / z \}$

$\rightarrow \text{eats } (y, \text{peanuts}) \vee \text{killed } (y) \quad \rightarrow \text{eats } (\text{Anil}, \text{Peanuts})$

Drop variable: $\{ \text{Anil} / y \}$

$\rightarrow \text{killed } (\text{Anil}) \quad \rightarrow \text{alive } (k) \vee \rightarrow \text{killed } (k)$

Drop variable: $\{ \text{Anil} / k \}$

$\rightarrow \text{alive } (\text{Anil}) \quad \rightarrow \text{alive } (\text{Anil})$

Drop variable: $\{ \text{Anil} / \text{Anil} \}$

$\rightarrow \{ \} \text{ Hence Proved}$

Drop variable: $\{ \} \text{ Hence Proved}$

Code:

```
from itertools import combinations
```

```
def get_clauses():

    n = int(input("Enter number of clauses in Knowledge Base: "))

    clauses = []

    for i in range(n):

        clause = input(f'Enter clause {i+1}: ')

        clause_set = set(clause.replace(" ", "").split("v"))

        clauses.append(clause_set)

    return clauses
```

```
def resolve(ci, cj):

    resolvents = []

    for di in ci:

        for dj in cj:

            if di == ('~' + dj) or dj == ('~' + di):

                new_clause = (ci - {di}) | (cj - {dj})

                resolvents.append(new_clause)

    return resolvents
```

```
def resolution_algorithm(kb, query):

    kb.append(set(['~' + query]))

    derived = []

    clause_id = {frozenset(c): f'C{i+1}' for i, c in enumerate(kb)}
```

```
    step = 1

    while True:

        new = []

        for (ci, cj) in combinations(kb, 2):
```

```

resolvents = resolve(ci, cj)

for res in resolvents:
    if res not in kb and res not in new:
        cid_i, cid_j = clause_id[frozenset(ci)], clause_id[frozenset(cj)]
        clause_name = f'R{step}'
        derived.append((clause_name, res, cid_i, cid_j))
        clause_id[frozenset(res)] = clause_name
        new.append(res)
        print(f'[Step {step}] {clause_name} = Resolve({cid_i}, {cid_j}) → {res or "{}"}')
        step += 1

# If empty clause found → proof complete
if res == set():
    print("\n✓ Query is proved by resolution (empty clause found).")
    print("\n--- Proof Tree ---")
    print_tree(derived, clause_name)
    return True

if not new:
    print("\n✗ Query cannot be proved by resolution.")
    return False

kb.extend(new)

def print_tree(derived, goal):
    tree = {name: (parents, clause) for name, clause, *parents in [(r[0], r[1], r[2:][0], r[2:][1]) for r in derived]}

def show(node, indent=0):
    if node not in tree:
        print(" " * indent + node)
    return

```

```

parents, clause = tree[node]
print(" " * indent + f"{node}: {set(clause) or '{}'}")
for p in parents:
    show(p, indent + 4)

show(goal)

```

OUPUT:

```

==== FOL Resolution Demo with Proof Tree ====
Enter number of clauses in Knowledge Base: 3
Enter clause 1: P
Enter clause 2: ~P v Q
Enter clause 3: ~Q
Enter query to prove: Q
[Step 1] R1 = Resolve(C1, C2) → {'Q'}
[Step 2] R2 = Resolve(C2, C4) → {'~P'}
[Step 3] R3 = Resolve(C1, R2) → {}

✓ Query is proved by resolution (empty clause found).

```

```

--- Proof Tree ---
R3: {}
  C1
    R2: {'~P'}
      C2
        C4
True

```

Program 10

Implement Alpha-Beta Pruning.

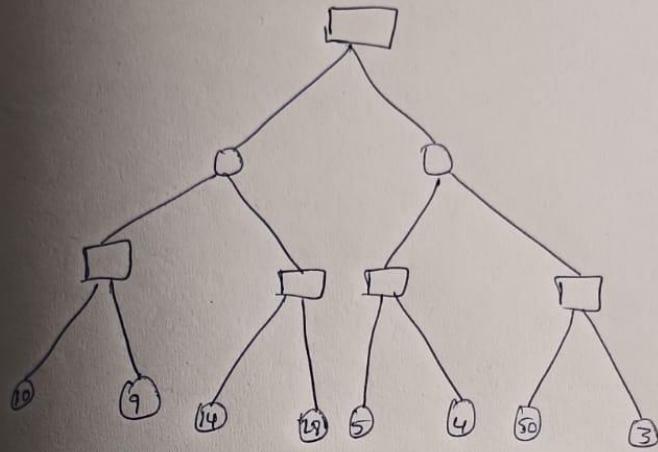
Lab 10

Adversarial Search: Implement Alpha-Beta Pruning

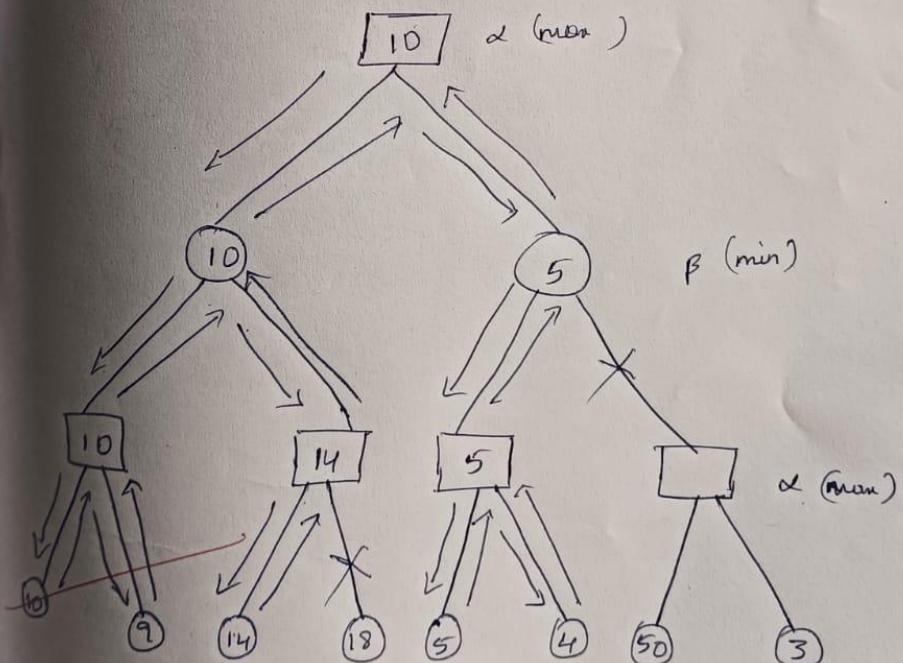
Algo:

1. Start at the root node (current game starts) at current player is either max or min.
2. Initialize
 - $\alpha = -\infty$
 - $\beta = +\infty$
3. If terminal node (end of game):
→ Return the utility (score) of that node.
4. If it's a max player:
 - Set value = $-\infty$
 - For each child of this node:
 - Compute each child ~~of this node~~ value =
 - AlphaBeta (child, depth - 1, α , β , False)
 - Update value = max (value, child_value)
 - Update α = max (value, child_value)
 - If $\alpha \geq \beta$, then break → (prune remaining branches)
 - Return value
5. If it's a min player:
 - Set value = $+\infty$
 - For each child_value =
 - Alpha Beta (child, depth - 1, α , β , True)
 - Update value = min (value, child_value)
 - Update β = min (β , value)

"If $\alpha \geq \beta$, then \rightarrow (prune remaining branches)
Return Value"



Solu:-



Was
50 10 13

Code:

class Node:

```
def __init__(self, name):
    self.name = name
    self.children = []
    self.value = None
    self.pruned = False
```

def alpha_beta(node, depth, maximizing, values, alpha, beta, index):

Terminal node

if depth == 3:

```
    node.value = values[index[0]]
    index[0] += 1
    return node.value
```

if maximizing:

best = float('-inf')

for i in range(2): # 2 children

```
        child = Node(f'{node.name} {i}')
        node.children.append(child)
```

```
        val = alpha_beta(child, depth + 1, False, values, alpha, beta, index)
```

```
        best = max(best, val)
```

```
        alpha = max(alpha, best)
```

if beta <= alpha:

```
            node.pruned = True
```

```
            break
```

```
    node.value = best
```

```
    return best
```

else:

```
    best = float('inf')
```

```

for i in range(2):
    child = Node(f"{{node.name}} {i}")
    node.children.append(child)
    val = alpha_beta(child, depth + 1, True, values, alpha, beta, index)
    best = min(best, val)
    beta = min(beta, best)
    if beta <= alpha:
        node.pruned = True
        break
    node.value = best
return best

def print_tree(node, indent=0):
    prune_mark = "[PRUNED]" if node.pruned else ""
    val = f"= {node.value}" if node.value is not None else ""
    print(" " * indent + f"{{node.name}} {val} {prune_mark}")
    for child in node.children:
        print_tree(child, indent + 4)

# --- main ---
print("== Alpha-Beta Pruning with Tree ==")
values = list(map(int, input("Enter 8 leaf node values separated by spaces: ").split()))

root = Node("R")
alpha_beta(root, 0, True, values, float('-inf'), float('inf'), [0])

print("\n--- Game Tree ---")
print_tree(root)

print("\nOptimal Value at Root:", root.value)

```

OUTPUT:

```
--- Alpha-Beta Pruning with Tree ---
Enter 8 leaf node values separated by spaces: 3 5 6 9 1 2 0 7

--- Game Tree ---
R = 5
R0 = 5
R00 = 5
    R000 = 3
    R001 = 5
R01 = 6 [PRUNED]
    R010 = 6
R1 = 2 [PRUNED]
R10 = 9
    R100 = 9
    R101 = 1
R11 = 2
    R110 = 2
    R111 = 0

Optimal Value at Root: 5
```