

STUDENT: **The complexity of problems often depends on what?**  
 TEACHER:  $\hookrightarrow$  "" (the type of reduction being used )

STUDENT: **What would create a conflict between a problem X and problem C within the context of reduction?**  
 TEACHER:  $\hookrightarrow$  "" (if every problem in C can be reduced to X )

STUDENT: **An algorithm for X which reduces to C would us to do what?**  
 TEACHER:  $\hookrightarrow$  "" (solve any problem in C )

STUDENT: **A problem set that that is hard for the expression NP can also be stated how?**  
 TEACHER:  $\hookrightarrow$  "" (NP-hard )

STUDENT: **What does the complexity of problems not often depend on?**  
 TEACHER:  $\hookrightarrow$  "" (CANNOTANSWER )

STUDENT: **What would not create a conflict between a problem X and problem C within the context of reduction?**  
 TEACHER:  $\hookrightarrow$  "" (CANNOTANSWER )

STUDENT: **What problem in C is harder than X?**  
 TEACHER:  $\hookrightarrow$  "" (CANNOTANSWER )

STUDENT: **How is a problem set that is hard for expression QP be stated?**  
 TEACHER:  $\hookrightarrow$  "" (CANNOTANSWER )

### Section:Computational complexity theory31

Context: This motivates the concept of a problem being hard for a complexity class. A problem X is hard for a class of problems C if every problem in C can be reduced to X. Thus no problem in C is harder than X, since an algorithm for X allows us to solve any problem in C. Of course, the notion of hard problems depends on the type of reduction being used. For complexity classes larger than P, polynomial-time reductions are commonly used. In particular, the set of problems that are hard for NP is the set of NP-hard problems. CANNOTANSWER