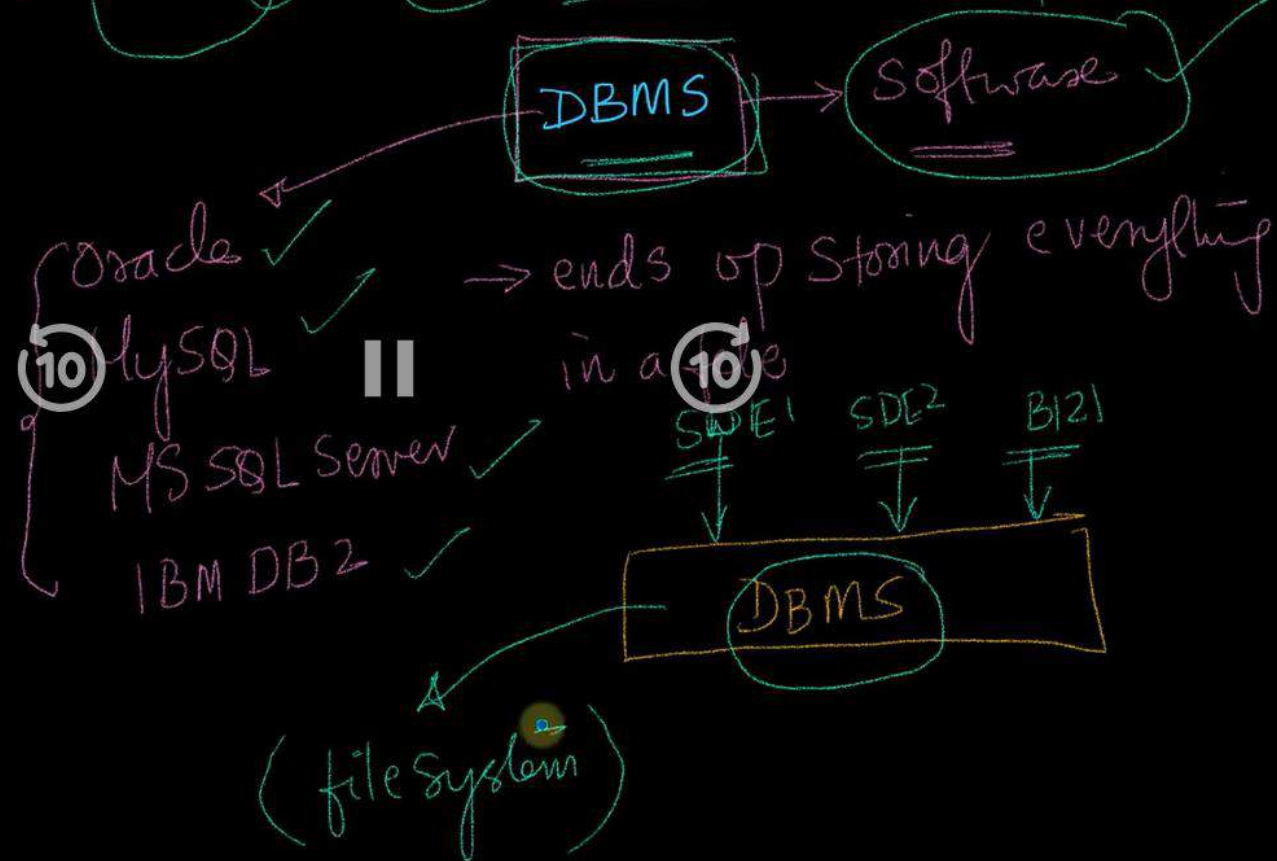


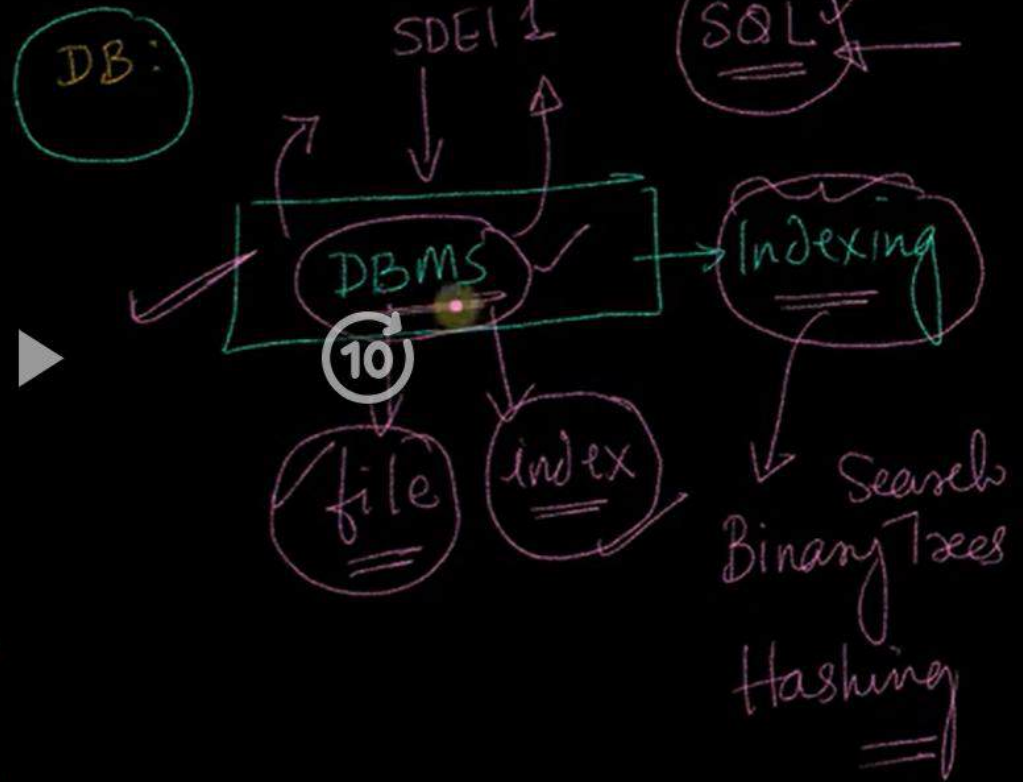
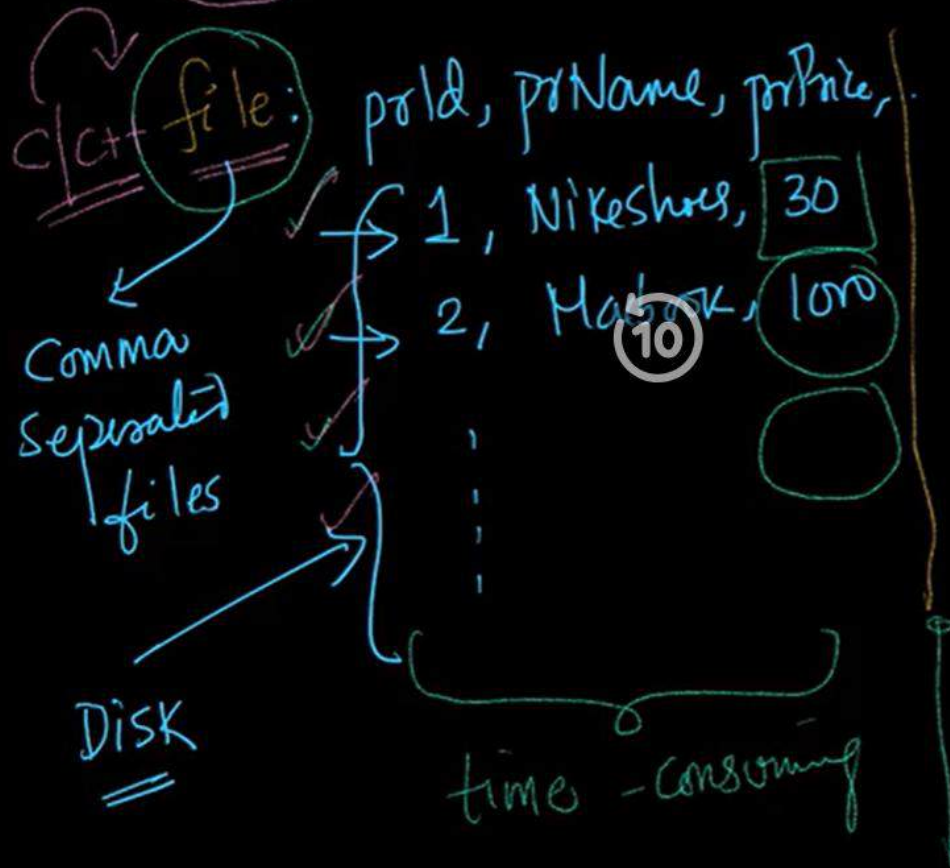
why not simply use files? → C fprintf, fscanf, fopen...

file system

{ create files/folders
edit files/folders
= ✓



① Querying: find all products priced > \$100.

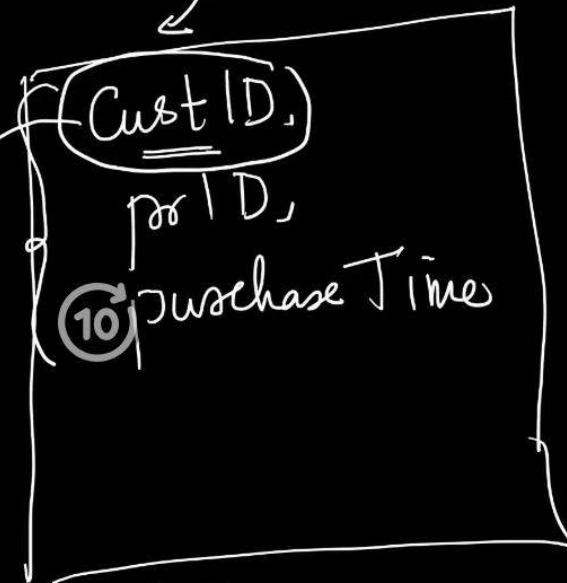
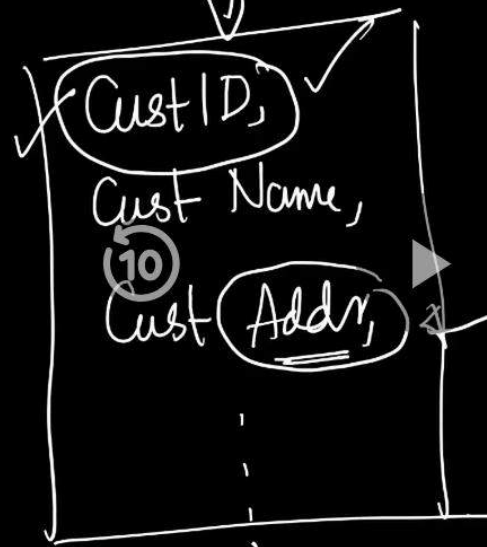


② Redundancy:

repeatedly
storing the
same data

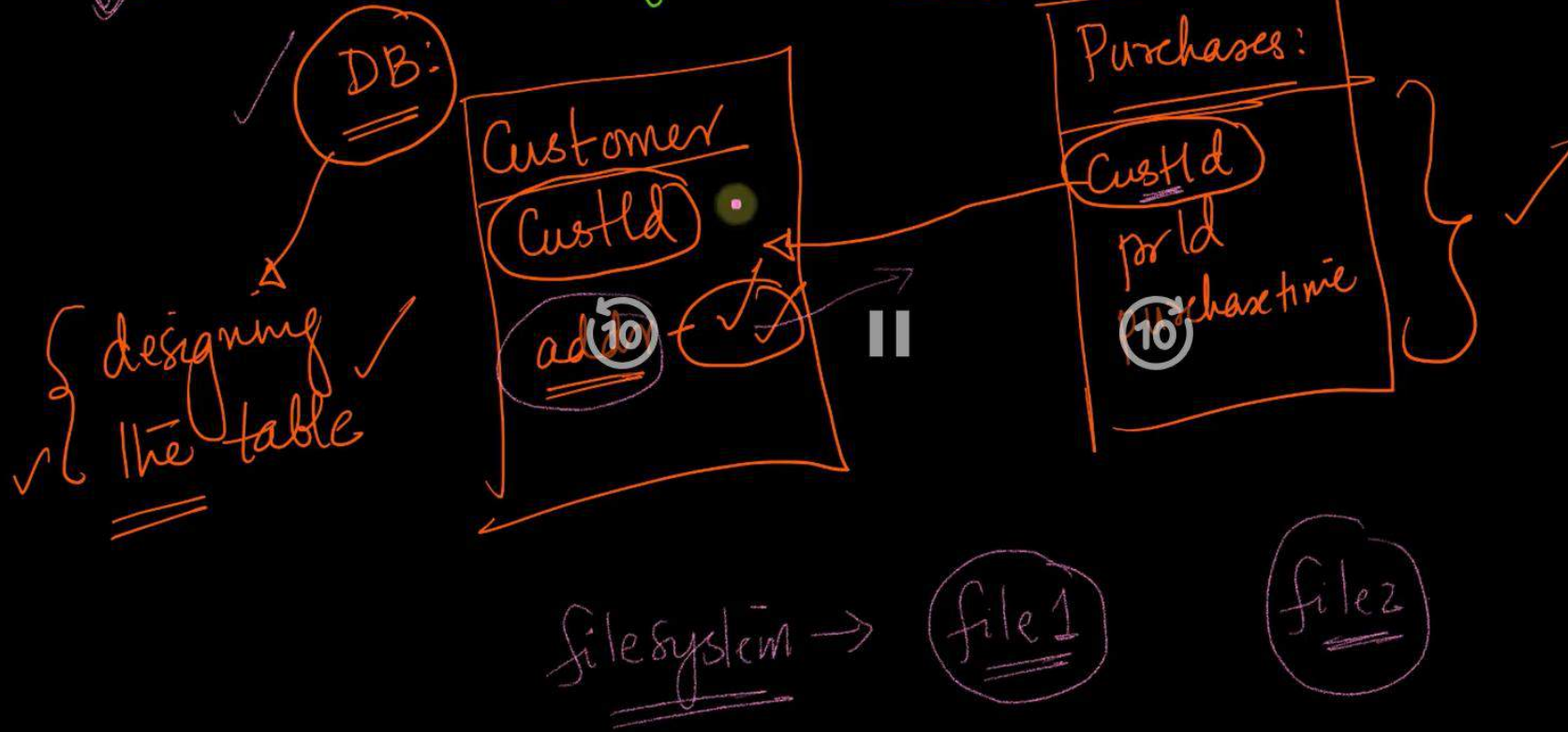
↓
more disk
space

customer table, purchases table

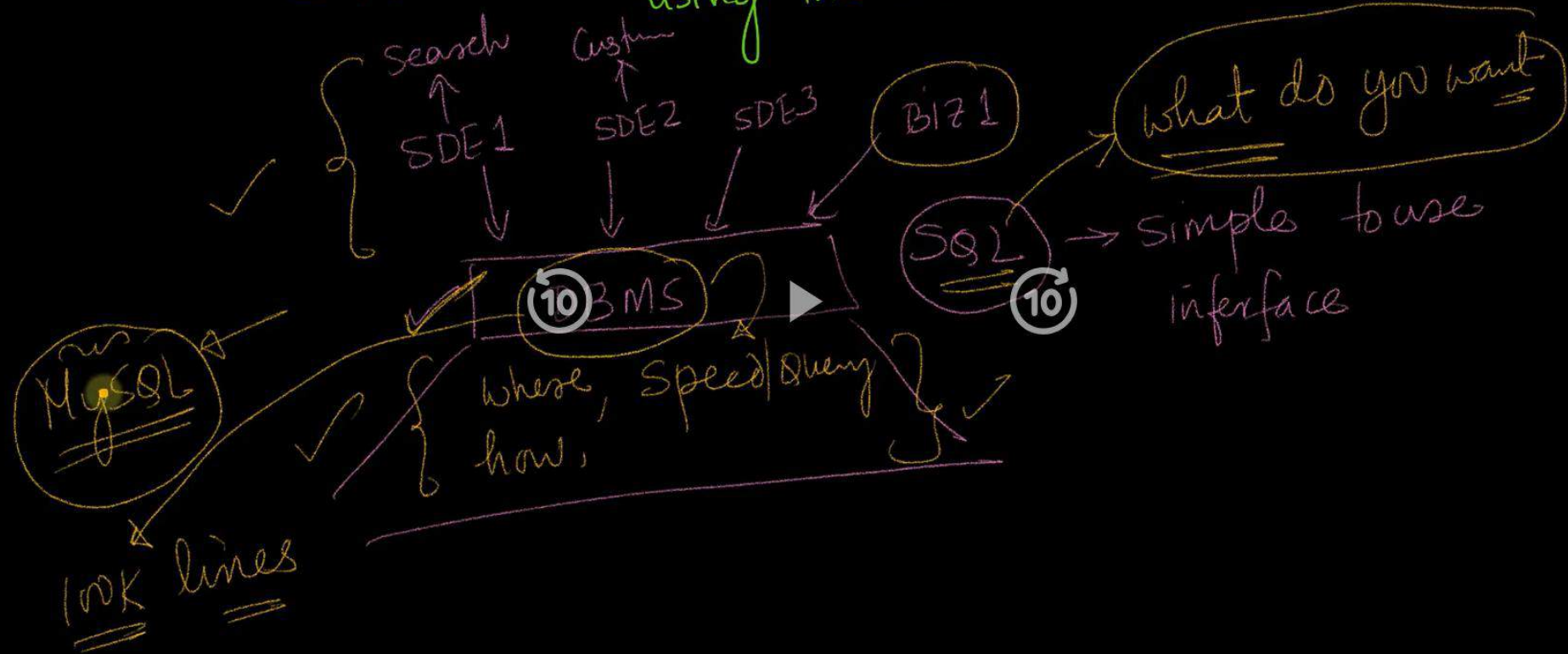


Database ✓
⇒

③ Consistency: change the address



④ Data Independence : Hide low levels details from the engineers
using the DBMS



⑤ Security & Access control: credit card details, personal info.

Amazon

Stocking

SDE 1

SDE 2
⑩

SDE 3

Biz 1



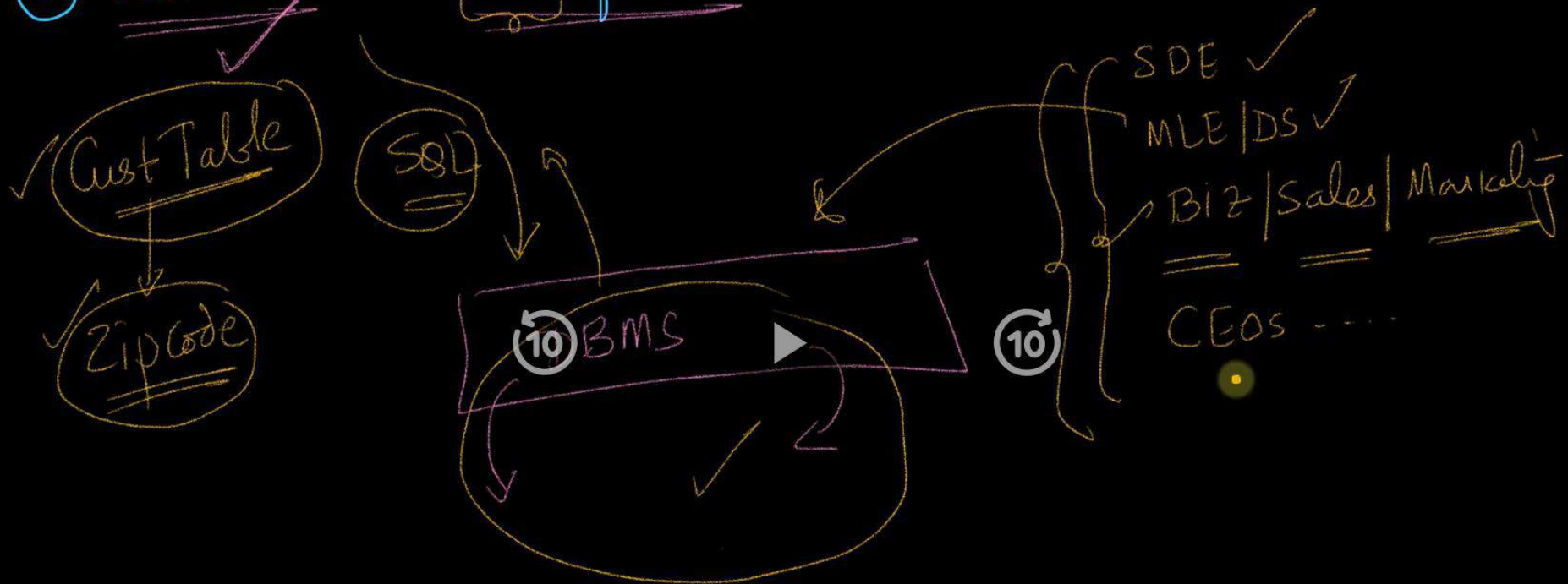
⑩

Sales,

Files → DOS

addr,
ph. no
}

⑥ Abstraction & case of data access: ✓



✓ Tables, Keys & Schema

✓ Relational Databases:

Relation = Table (10)

Table = set of rows & columns

✓ Terminology

✓ (Graph) Databases

Relational Algebra
(10) Set theory

Table:

Cust Table
Relation

Columns / field / attributes

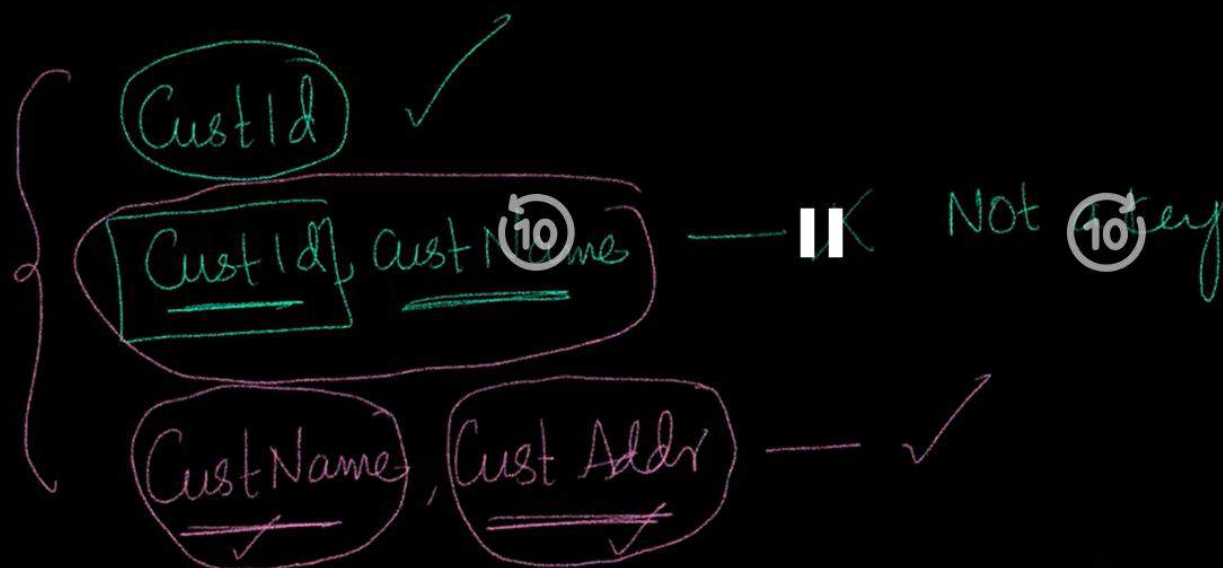
CustId	CustName	CustAddr	CustZipcode
1	abc	202,	94086
2	x4z	Flat,	94021
.	.	.	.
.	.	.	.
.	.	.	.

Records / Tuples

row 1
row 2

instance : set of tuple/rows + Table structure

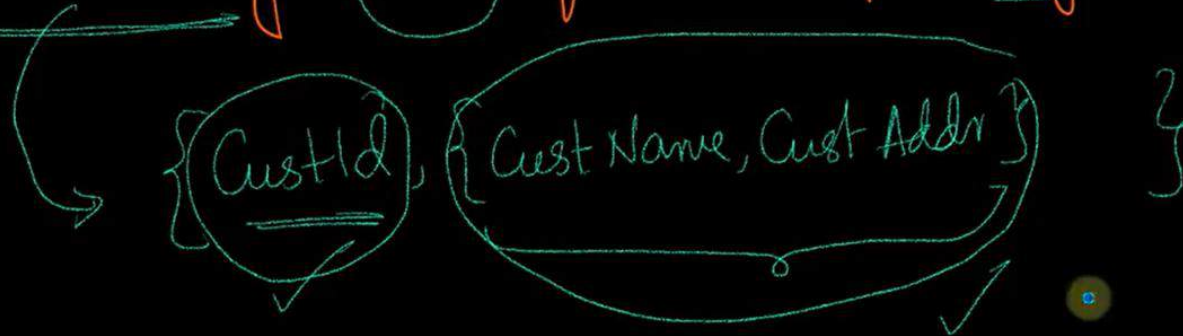
Key: minimum set of attributes / columns to uniquely identify a row / tuple.



Simple Key: Key with only one attrib/column
CustId

Compound Key: Key with multiple attrib/columns
Cust Name, Cust Addr
└────────────────┘

Candidate key: Set of all unique keys.



Primary key: one of the candidate keys that the DB-designer chooses to maintain uniqueness.

admin

CustId

NOT NULL

for any row/tuple

atmost one primary key

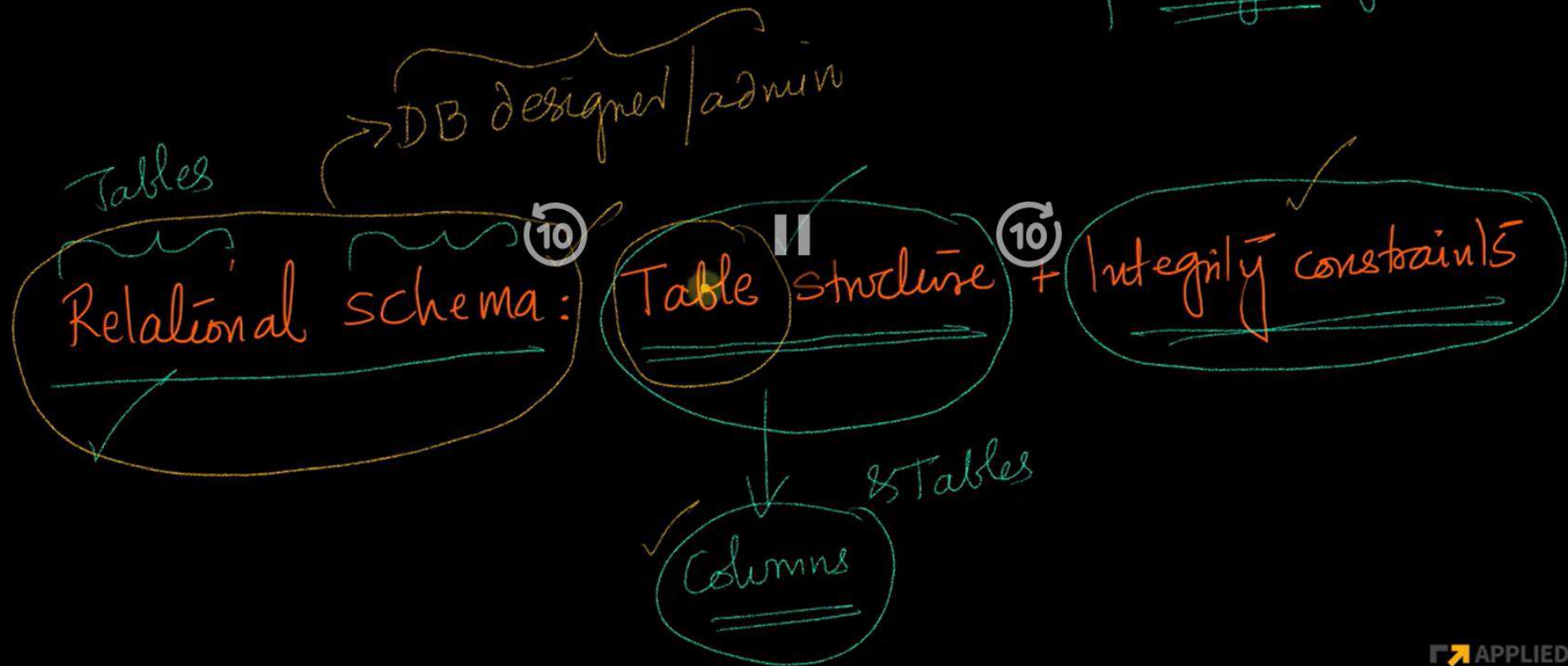
Unique for each row/tuples

Entity Integrity:

Integrity Constraints

(Candidate keys can be NULL)

Alternative / Secondary keys: Candidate keys that are not primary keys



Super Key:

Candidate Key \cup Attributes

$\{ \text{CustId}, \text{CustName} \} = \{ \text{CustId} \} \cup \{ \text{CustName} \}$

Superset



Foreign Key:

✓ avoid inconsistencies

Cust Table

<u>CustId</u>	CustName	CustAddr	CustZip
1			
2			
3			

Purchases Table

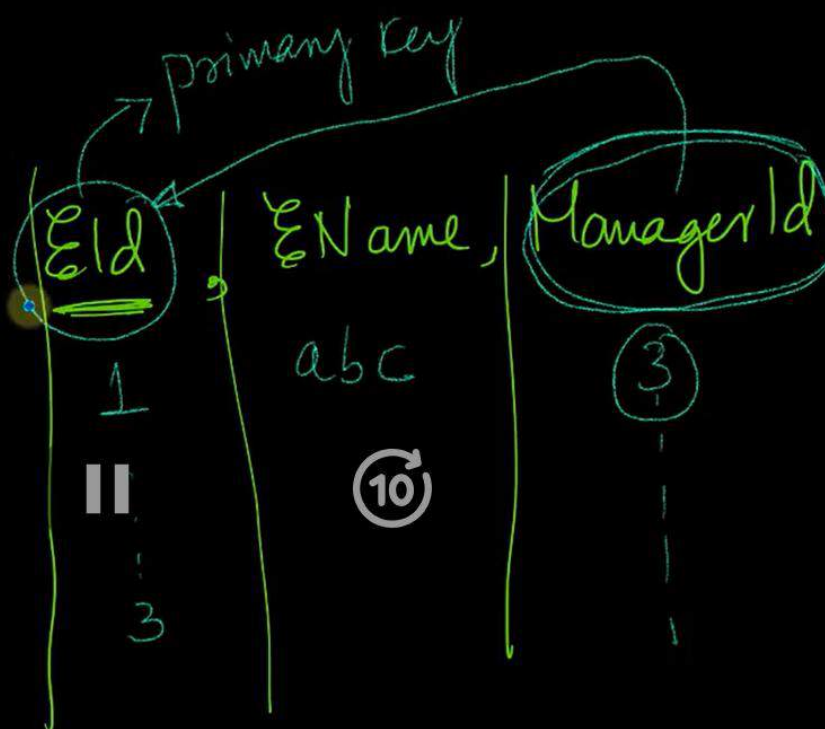
<u>TransId</u>	<u>CustId</u>	PrId	t
12345	1	12	
	2		
	3		
	abc		

✓ { Integrity Constraints

Self referential

Employees Table:

same table



✓ Integrity Constraints ✓ → various types

① Entity Integrity : Primary Keys

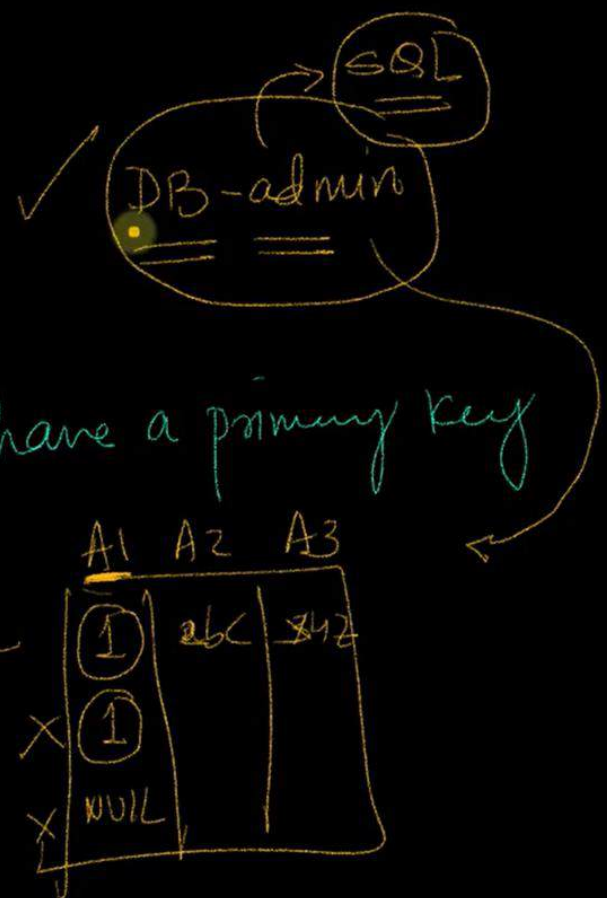
ER-diagrams

10

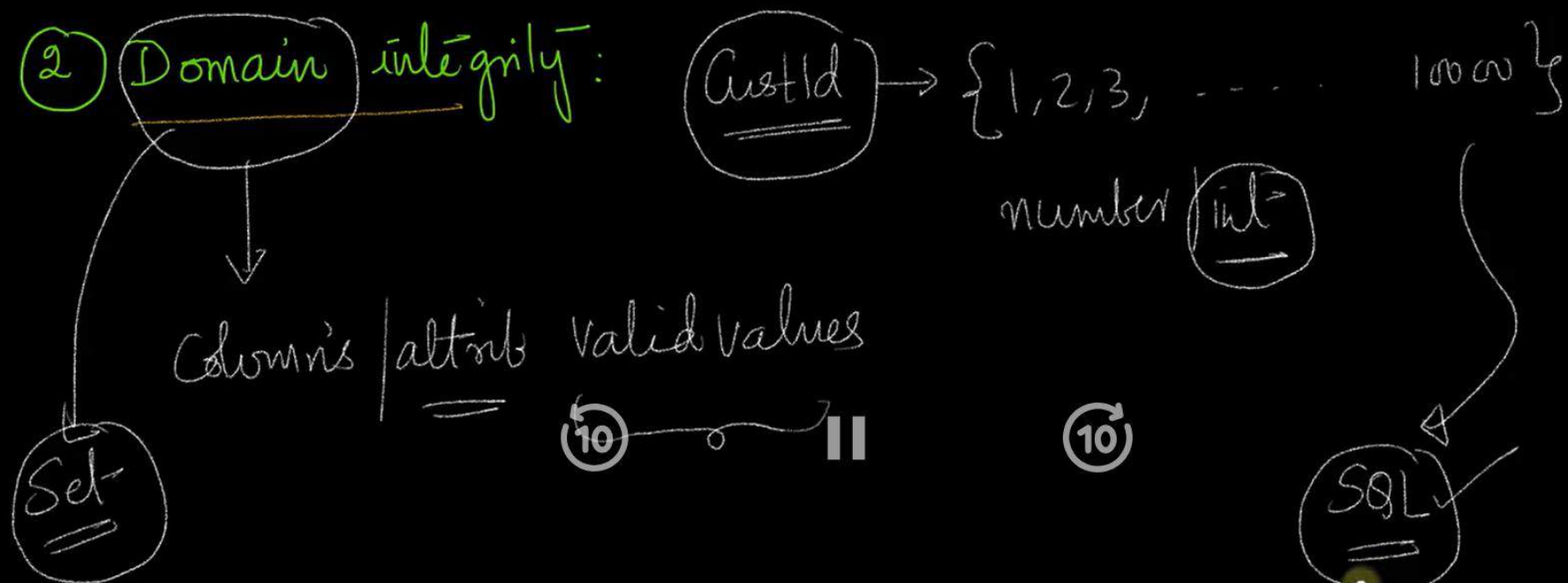
every table must have a primary key

10

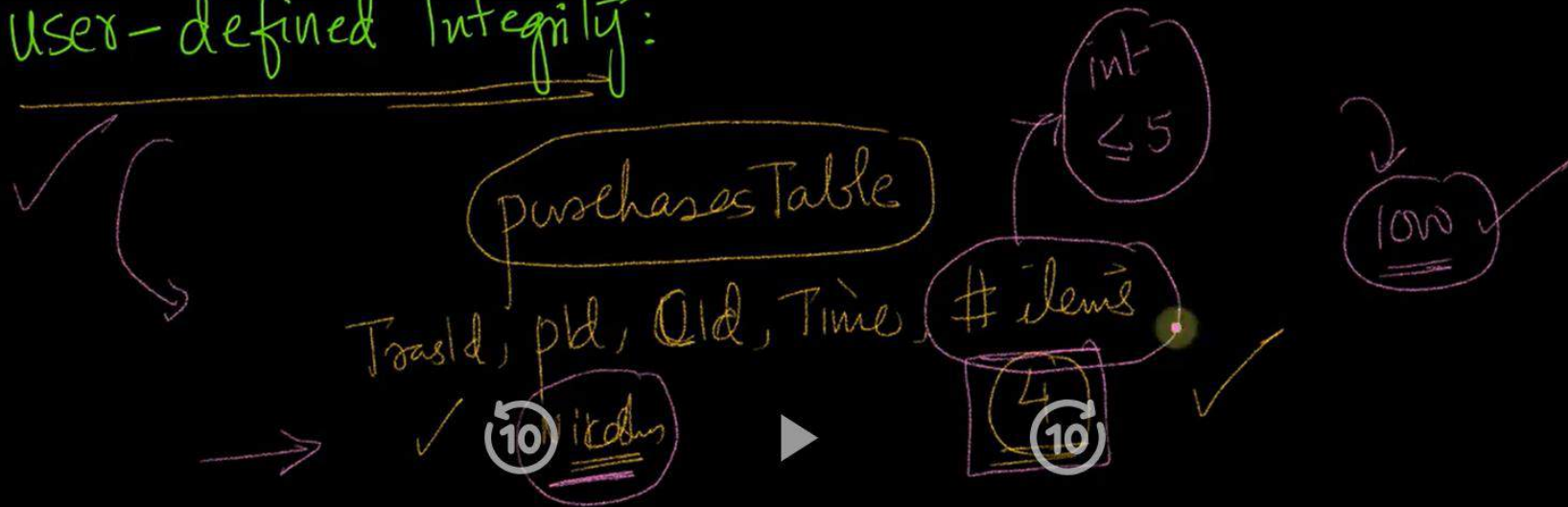
- ✓ Unique
- ✓ NOT NULL



② Domain integrity:



③ user-defined Integrity:



④ Referential Integrity: foreign keys

✓ Cust Table

<u>Cid</u>	CName	CAddr
1	abc	123456...
2	x4z	7891011...

✓ Referenced Table

foreign key

✓ Purchases Table

<u>TransId</u>	<u>Cid</u>	PId	time
123	1	12	12:30
456	2	11	11:30

✓ Referencing Table

Changes to Referenced Table:

Cust Table

<u>Cid</u>	CName	CAddr
1	abc	123456...
2	xyz	7891011...
3	mno	---

Purchases Table

<u>TransId</u>	<u>Cid</u>	PId	time
123	1	12	12:30
456	2	11	11:30
	NULL		

part of the primary key
referencing
Table

✓ Insert → No changes

Delete

on delete no action (1) → violates

on delete Cascade (2)
(dangerous/avoid)

on delete set null (3)

Changes to Referenced Table:

Cust Table

<u>Cid</u>	CName	CAddr
1a	abc	123456...
2	x4z	7891011...

==

Purchases Table

<u>TransId</u>	Cid	PId	time
123	1a	12	12:30
456	2	11	11:30
-	1a	-	-

Modify/edit

No action

Cascade ✓

Set NULL

most widely ✓

time Consuming

Changes to Referencing Table:

referenced
Cust Table

<u>Cid</u>	<u>CName</u>	<u>CAddr</u>
1	abc	123456...
2	xyz	7891011...
3	MND	...

Purchases Table

<u>TransId</u>	<u>Cid</u>	<u>PId</u>	<u>time</u>
123	1a	12	12:30
456	2	11	11:30
789	3	10	10:30

Insert

Delete → do nothing

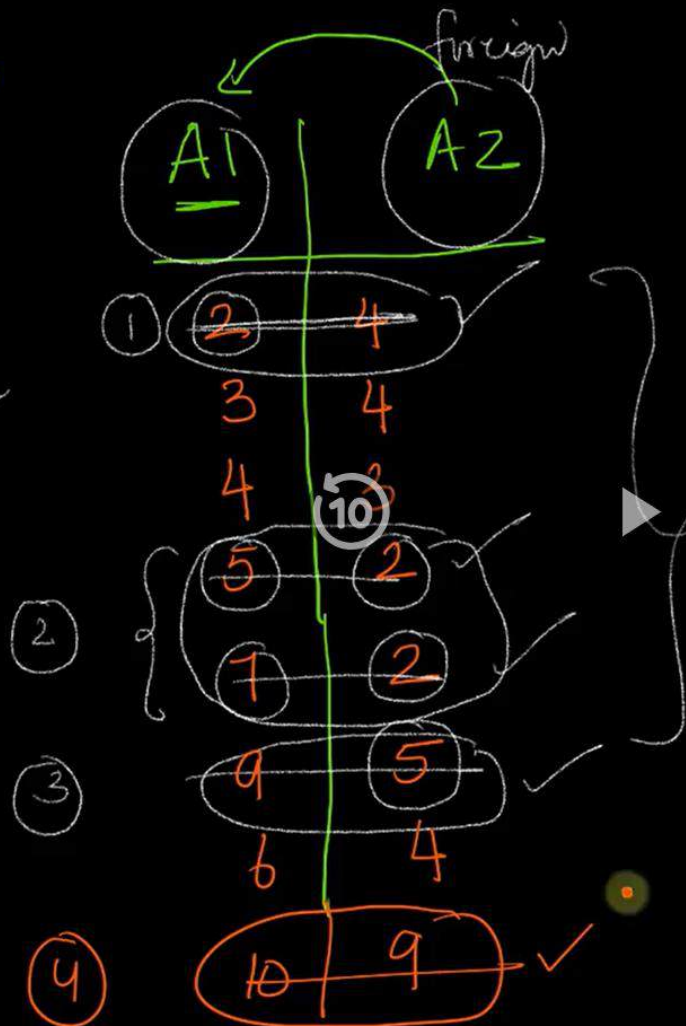
Check for any violation

Edit → Check for any violations

error ✓

Example

Relation/
Table



ON DELETE CASCADE

2,4 ← row/tuple

(10) multi-level cascade

(Q1) $R(A_1, A_2, A_3, \dots, A_n)$

(b) Number of super keys if candidate keys are $\{A_1, A_2\}$



(Q1) Table $R(A_1, A_2, A_3, \dots, A_n)$

(a) Number of Super keys if candidate keys are $\{A_1\}$

n
↓
 2^n

Candidate \cup other attrib

Counting / Sets

$\{A_1 \cup \{\phi, A_2, A_3, \dots, A_2, A_3\}\}$

$(n-1)$ elements

powerset

2^{n-1}

(Q1) $R(A_1, A_2, A_3, \dots, A_n)$

Inclusion-Exclusion

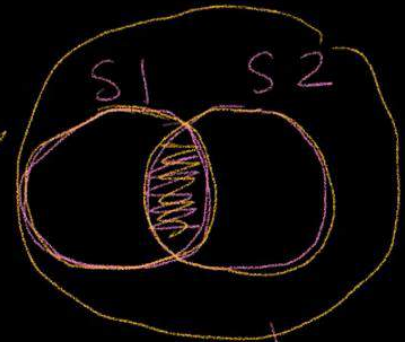
(b) Number of super keys if candidate keys are $\{A_1, A_2\}$

$\checkmark A_1 \rightarrow 2^{n-1}$ super keys
 $\checkmark A_2 \rightarrow 2^{n-1}$ super keys

||

$\{A_1, A_2\}$ ✓
 A_1, A_2, A_3 ✓

A_2, A_1 ✓
 A_2, A_1, A_3 ✓



A_1, A_2
 \downarrow
 2^{n-2}

$$2^{n-1} + 2^{n-1} - 2^{n-2}$$

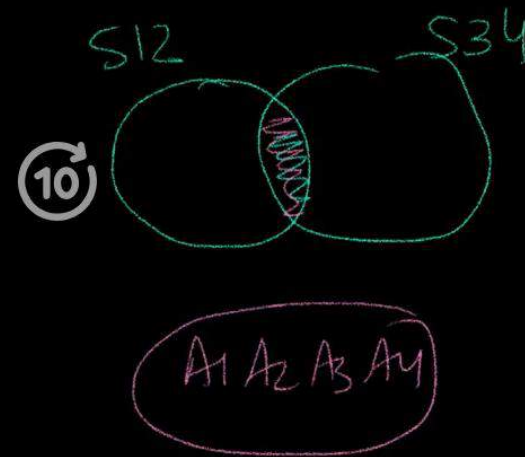
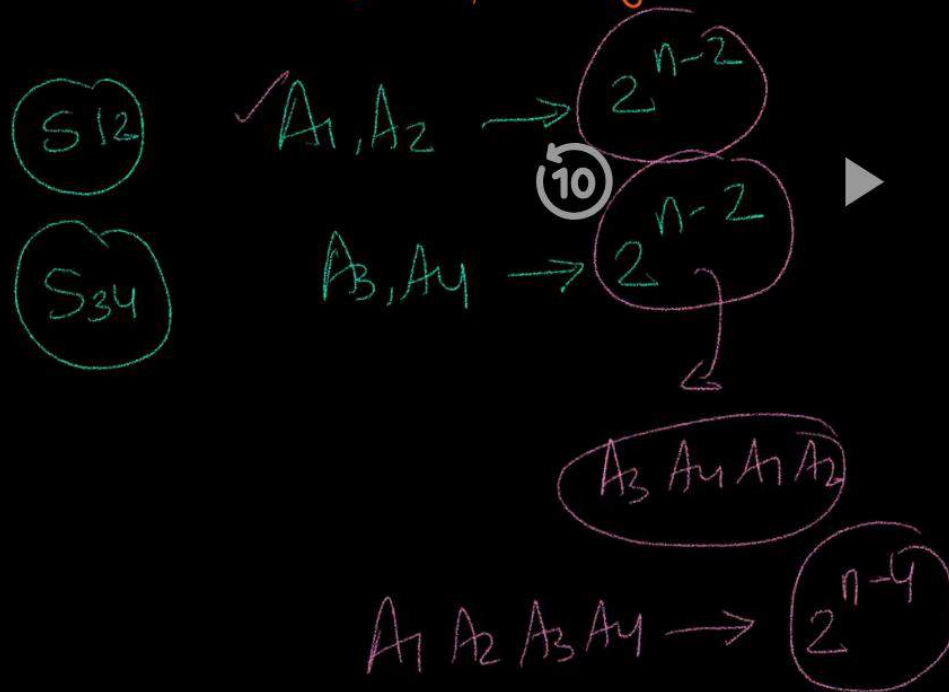
$$= 2 \cdot 2^{n-1} - 2^{n-2} = 2^n - 2^{n-2}$$

$$|S_1 \cup S_2| = |S_1| + |S_2| - |S_1 \cap S_2|$$

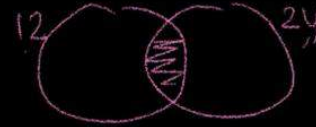
(Q1) $R(\underline{A_1}, A_2, A_3, \dots, A_n)$

$$2^{n-2} + 2^{n-2} - 2^{n-4}$$

(c) Number of Super keys if candidate keys are $\{\underline{A_1, A_2}, A_3, A_4\}$



(Q1) $R(A_1, A_2, A_3, \dots, A_n)$



(d) Number of super keys if candidate keys are $\{\{A_1, A_2\}, \{A_2, A_4\}\}$

$A_1, A_2, A_4 \rightarrow 2^{n-3}$ (10)



(10) 2^{n-2}
 2^{n-2}

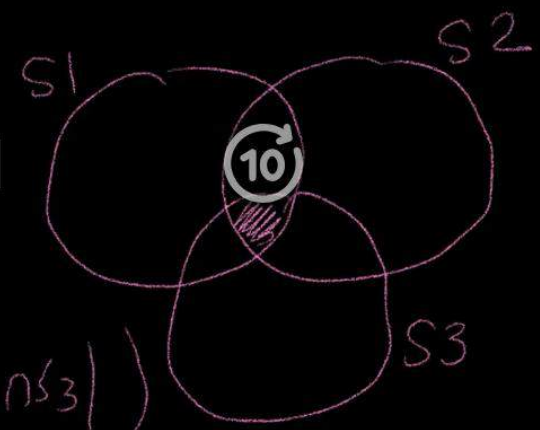
$2^{n-2} + 2^{n-2} - 2^{n-3}$

$$(Q1) R(A_1, A_2, A_3, \dots, A_n) \left(\underline{3 \cdot 2^{n-1}} - \{ \underline{3 \cdot 2^{n-2}} \} + 2^{n-3} \right)$$

(c) Number of Super keys if candidate keys are $\{ \underline{A_1}, \underline{A_2}, \underline{A_3} \}$

$$\left\{ \begin{aligned} & |S_1 \cup S_2 \cup S_3| \\ &= |S_1| + |S_2| + |S_3| \\ &\quad - (|S_1 \cap S_2| + |S_2 \cap S_3| + |S_1 \cap S_3|) \\ &\quad + (|S_1 \cap S_2 \cap S_3|) \end{aligned} \right\} \quad \begin{array}{c} \text{S1} \\ \text{S2} \\ \text{S3} \end{array}$$

||



② Given a relation with n attributes, how many super keys can be present?

