Faculty of Science

Cryptographic Library for FPGAs using SME

A Bachelor Project Defense

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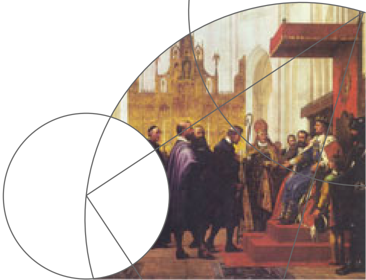
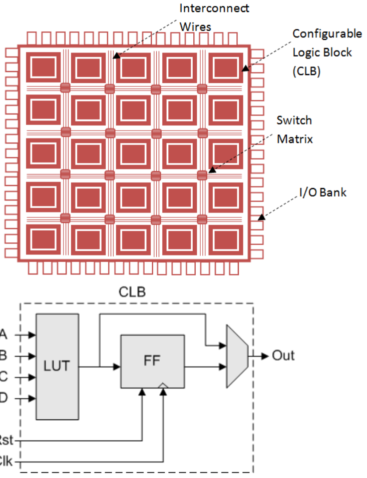


Agenda

- Why is FPGA's interesting?
- Presentation of MD5 and AES
- Results
- suggestions for future work

Motivation - Why use FPGAs?

- Architecture
- Configurable not only in computation but also in interface
- Often fast as the overhead from generality is ommitted
- Often lower power consumption than CPU's
- Difficult to program (Xilinx Vivado 😩)



MD5

- Cryptograpic hash function
- Merkle-Damgaard construction
- Four different compression functions, 64 rounds

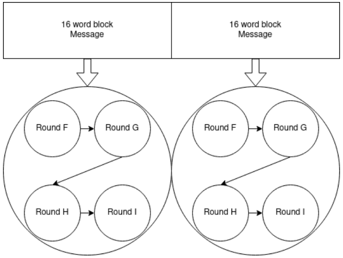


Figure: MD5 round

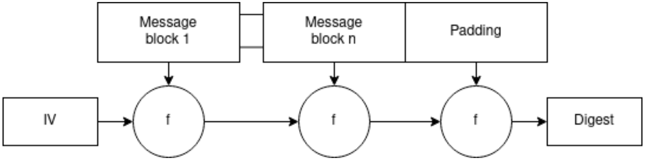


Figure: Merkle-Damgaard construction

MD5 - Optimizations

- Naive; 1 simple process, 4 busses
- Pipelined version; clocked process for preprocessing and each compression stage
- Same idea for SHA-2

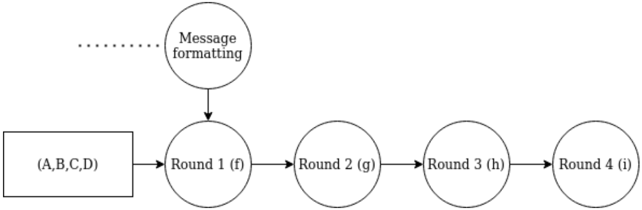


Figure: MD5 pipeline

MD5 - Pipeline

- Stalls on large messages due to data dependency
- Solution is enabling multiple inputs

	Independent message blocks								
clock	0	1	2	3	4	5	6	7	8
	P_1	M_1	F_1	G_1	H_1	I_1	C_1		
		P_2	M_2	F_2	G_2	H_2	I_2	C_2	

	Dependent message blocks											
clock	0	1	2	3	4	5	6	7	8	9	10	11
	P_1	M_1	F_1	G_1	H_1	I_1	C_1					
		P_2	M_2	-	-	-	-	F_2	G_2	H_2	I_2	C_2

AES

- The algorithm is Rijndael. It is a Block Cipher and a Substitution-permutation (SP) network.
- Four steps, or lookups:

$$W_i \begin{cases} K_i & \text{if } i < N \\ W_{i-N} \oplus \text{SubWords}(W_{i-1} \lll 8) \oplus \text{rcon}_{i/N} & \text{if } i \geq N \text{ and } i \equiv 0 (\text{mod } N) \\ W_{i-N} \oplus \text{SubWords}(W_{i-1}) & \text{if } i \geq N,\ N > 6, \text{ and } i \equiv 4 \ (\text{mod } N) \\ W_{i-N} \oplus W_{i-1} & \text{otherwise} \end{cases}$$

$$T_0[a] = \begin{bmatrix} S[a] \cdot 02_{16} \\ S[a] \\ S[a] \\ S[a] \cdot 03_{16} \end{bmatrix} T_1[a] = \begin{bmatrix} S[a] \cdot 03_{16} \\ S[a] \cdot 02_{16} \\ S[a] \\ S[a] \end{bmatrix} T_2[a] = \begin{bmatrix} S[a] \\ S[a] \cdot 03_{16} \\ S[a] \cdot 02_{16} \\ S[a] \end{bmatrix} T_3[a] = \begin{bmatrix} S[a] \\ S[a] \\ S[a] \cdot 03_{16} \\ S[a] \cdot 02_{16} \end{bmatrix}$$

$$e_j = T_0[a_{0,3}] \oplus T_1[a_{1,2}] \oplus T_2[a_{2,1}] \oplus T_3[a_{3,0}] \oplus k_j$$

AES - Optimization

- Fast naive version
- Pipelined by splitting up rounds
- no data dependecy

Results from pipelining

- $hi(x) = x + 2 \cdot blocks$
- $lo(x) = 2 + x \cdot blocks$
- $C(x) = x + 2 \cdot blocks$
- Streaming AXI
- Zedboard bus: 2132 MBps

More processes is better throughput (to a certain extend)

MD5: Is easily optimised beyond the memory limit. 20x faster than Naive version.

SHA: Worst. Hard to optimize because of message expansion. 2x faster.

AES: Reached limit with current approach? Potentially other approaches can reach higher. 4x faster.

ChaCha: Starting to reach limit of board (10000 FF on board). 34x faster.

MD5

Version	f_{max} (Mhz)	$clocks_{hi}$	$TP(MBps)_{hi}$	$clocks_{lo}$	$TP(MBps)_{lo}$	LUT	FF
Naive	2.38	b	152	b	152	11607	2304
$Proc_4$	9.50	hi(6)	266	lo(6)	101	10247	5226
$Proc_8$	19.00	hi(10)	532	lo(10)	122	10087	7538
$Proc_{16}$	33.50	hi(18)	937	lo(18)	119	10206	12162
$Proc_{32}$	65.00	hi(34)	1817	lo(34)	123	10149	21347
$Proc_{64}$	**115.00**	**hi(66)**	**3209**	**lo(66)**	**112**	**10350**	**39718**

SHA

Version	f_{max} (Mhz)	$clocks_{hi}$	$TP(MBps)_{hi}$	$clocks_{lo}$	$TP(MBps)_{lo}$	LUT	FF
Naive	2.1	b	134.4	b	134.4	24330	2560
$Proc_4$	**8.0**	**hi(6)**	**223.9**	**lo(6)**	**85.3**	**24466**	**8938**
$Proc_8$	8.0	hi(10)	223.8	lo(10)	51.2	24756	14066

AES

Version	f_{max} (Mhz)	clocks	TP(MBps)	LUT	FF	BRAM
Naive	22	b	352	10612	3195	0
TBox	25	b	400	16458	3195	0
$Proc_4$	68	C(3)	544	16474	2817	0
$Proc_{11}$	**208**	**C(12)**	**1663**	**15659**	**4383**	**0**
$Proc_{22}$	217	C(24)	1662	15454	7401	0
$BRAM_{11}$	195	C(31)	1556	10012	10398	72

ChaCha

Version	f_{max} (Mhz)	clocks	TP(MBps)	LUT	FF
Naive	1.25	b	80	14670	3457
$Proc_{11}$	40.00	C(9)	1279	14736	16898
$Proc_{22}$	82.00	C(20)	2557	17565	32420
$Proc_{44}$	**85.00**	**C(40)**	**2715**	**17612**	**62436**

Results compared to CPU

Comparing to CPU over GPU

- GPU is the standard approach for hardware acceleration.
- CPU is more approachable and making a GPU version would require higher development time.
- already reached some board limitations.

MD5: ~4.5 times faster than any comparable CPU version.

AES: Proximity of the C# version, but cannot compete AES-NI.

SHA: Only half the speed of i5, but faster than ARM processor. Potential for improvement.

ChaCha: Percentagewise best but not quite speed of i5. Reaches the bandwidth limit of the Zynq board.

MD5

	Naive	Proc_{64}	C#	C	OpenSLL_{low}	OpenSLL_{high}
Pi	152	3210	287	256	42	293
i5	152	3210	604	622	81	691

AES

	Naive	Proc_{11}	C#	C	OpenSLL_{low}	OpenSLL_{high}
Pi	400	1963	70	198	72	89
i5	400	1699	1963	340	847	5722

SHA

	Naive	Proc_{4}	C#	OpenSLL_{low}	OpenSLL_{high}
Pi	134	224	163	42	165
i5	134	224	438	61	461

ChaCha

	Naive	Proc	OpenSLL_{low}	OpenSLL_{high}
Pi	80	2715	84	307
i5	80	2715	388	3092

Power usage

Why use TDP?

- Selling point for FPGAs
- Only possibility in the current stage.
- Pretty diffuse concept as no standard way to measure it (Intel boost)

All of our versions seems significantly more power efficient than the CPUs:

FPGA		1.765W	
i5	65.0W	~6.500W	(i5 is 39 times bigger).
Pi	7.5W	~0.750W	

Future work

Critical work

- The Dependency routing needs to be fixed in the hashing functions.
- Make hashes able to switch between messages to circumvent the stalling.
- Make a useful interface to expose our implementations.

Optimizing work

- Investigate SHA hopefully getting the performance to reasonable levels
- Test if different approaches of AES approaches could yield better performance. Sugestions:
  1. Naive but pipelined
  2. Stateful BRAM.

Comparing work

- Compare to other research papers results which often are written in HDL to see if SME can actually provide comparable results. This would require better a better FPGA.
- Test against a GPU.

Questions?